# Imperative Programming in OCaml CS496

## Imperative Features in OCaml

OCaml (as seen so far) is purely functional

every expression is evaluated solely for its value

This lack of side-effects has an important consequence

- purely functional languages are said to enjoy referential transparency
  - ► This means that the order in which subexpressions are evaluated, in some large expression, is irrelevant

As a result, one

- ▶ can use standard algebraic equations (eg. a + b = b + a) to reason about programs
- can easily parallelize

## Imperative Features in OCaml

However sometimes imperative features are needed

- variable assignment and destructive update of data structures (specially for efficiency reasons)
- ▶ I/O: communication with some external device

Therefore, OCaml is enriched with expressions that are evaluated solely for their effects

- ► Reference types, assignment
- Mutable fields in records
- Arrays
- ▶ I/O: communication with some external device:

scanf, printf, ...

Many more

#### References

Records with Mutable Fields

Arrays

#### **Examples**

```
1  # let x=2
2  val x : int = 2
3
4  # let x=ref 2;;
5  val x : int ref = {contents = 2}
```

- ▶ ref 2 denotes a location in memory (i.e. a memory address)
- ref is similar to malloc in C

#### **Examples**

```
1 # let x = ref 2 in !x;;
_{2} - : int = 2
3 # let x=ref 2 in x:=!x+1; !x;;
4 - : int = 3
5 # let x=ref 2;;
6 val x : int ref = {contents = 2}
7 # x;;
8 - : int ref = {contents = 2}
9 # !x;;
10 - : int = 2
# x:=!x+1;;
12 - : unit = ()
13 # !x;;
-: int = 3
15
  * x := ! x + 1;;
-: unit = ()
17 # !x;;
18 - : int = 4
```

# Modeling a Counter Object

- Hidden state: value of counter
- First we declare the type of such an object, namely a record type

▶ Models the public interface of the object

#### Modeling a Counter Object

Then we define the object c itself:

#### and interact with it

```
1 # c.get ();;
2 - : int = 0
3 # c.inc ();;
4 - : unit = ()
5 # c.get ();;
6 - : int = 1
7 # c.set 4;;
8 - : unit = ()
9 # c.get ();;
10 - : int = 4
```

# Modeling a Function that Creates Counter Objects

```
# let newCounter n =
     let s = ref n
2
   in { get = (fun () -> !s);
3
          set = (fun x \rightarrow s:=x);
          inc = (fun () -> s:=!s+1);;
5
6 val newCounter : int -> counter = <fun>
7 # let c1 = newCounter 1;;
8 val c1 : counter = {get = <fun>; set = <fun>; inc = <</pre>
      \hookrightarrow fun>}
9 # let c2 = newCounter 2;;
10 val c2 : counter = {get = <fun>; set = <fun>; inc = <
     \hookrightarrow fun>}
# c1.get();;
12 - : int = 1
# c2.get();;
-: int = 2
# c1.inc();;
-: unit = ()
17 # c2.get();;
-: int = 2
```

# Modeling a Counter Object with this

```
1 # let newCounter n =
2 let s = ref n
3 in let rec this =
\{ get = (fun () -> !s); \}
set = (fun x -> s:=x);
   inc = (fun () \rightarrow s:=this.get ()+1)
7 in this;;
8 val newCounter : int -> counter = <fun>
9 # let c= newCounter 4;;
val c : counter = {get = <fun>; set = <fun>; inc = <</pre>
    \hookrightarrow fun>}
# c.get ();;
12 - : int = 4
13 # c.inc ();;
-: unit = ()
# c.get();;
16 - : int = 5
```

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# Example – Linked List of Integers

▶ option allows a field to be None, representing a null reference

# Example – Linked List of Integers

```
let create () =
   { head = None;
      size=0}
3
   let add x ll =
     11.head <- Some (ref {data=x; next=ll.head});</pre>
     11.size < -11.size + 1
8
   let string_of_list ll =
     let rec string_of_node = function
10
       | None -> ""
11
       | Some r -> string_of_node (!r.next) ^
12
           ⇔ string_of_int (!r.data)
     in string_of_node (ll.head)
13
```

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## Arrays.

```
1 # let a = [|1;2;3|];;
2 val a : int array = [|1; 2; 3|]
3 # a.(1);;
4 - : int = 2
5 # a.(1) <-4;;
6 - : unit = ()
7 # a;;
8 - : int array = [|1; 4; 3|]</pre>
```