## CSCE 221 Cover Page Homework #1 Due July 12 at midnight to eCampus

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Please list all sources in the table below including web pages which you used to solve or implement the current homework. If you fail to cite sources you can get a lower number of points or even zero, read more: Aggie Honor System Office

Type of sources		
People	I worked with Aaron Ingram	
	I worked with Callen McCauley	
Web pages (provide URL)	https://stackoverflow.com/questions/8480640/how-to-throw-a-c-exception	http://www.cplu
Printed material	"Data Structures and Algorithms" textbook for the class	
Other Sources		

I certify that I have listed all the sources that I used to develop the solutions/codes to the submitted work.

"On my honor as an Aggie, I have neither given nor received any unauthorized help on this academic work."

Your Name Laura Austin Date July 12, 2017

## 1.

part (a) Ensure that if the vector is not sorted, then exception is thrown.

```
[laustin254]@linux2 ~/HWl> (19:42:49 07/12/17)
:: ./a.out BinarySearch
vector tested is: < 2 1 4 8 16 32 64 128 256 512 1024 2048 >
terminate called after throwing an instance of 'std::invalid_argument'
   what(): Input vector not sorted
Aborted
```

part (b) Searching both the ascending and descending vectors.

```
[laustin254]@linux2 ~/HWl> (19:38:55 07/12/17)
:: ./a.out BinarySearch
vector tested is: < 2048 1024 512 256 128 64 32 16 8 4 2 1 >
found element 1 after 11 comparisons.

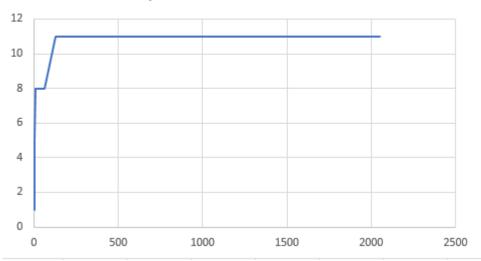
[laustin254]@linux2 ~/HWl> (19:38:55 07/12/17)
:: g++ -std=c++11 BinarySearch.cpp

[laustin254]@linux2 ~/HWl> (19:39:19 07/12/17)
:: ./a.out BinarySearch
vector tested is: < 1 2 4 8 16 32 64 128 256 512 1024 2048 >
found element 2048 after 11 comparisons.
```

## part (c)

Range $[1,n]$	Target for	# comp. for	Target for	# comp. for	Result of the
	incr. values	incr. values	decr. values	decr. values	formula in
					item 5 (e)
[1,1]	1	2	1	2	
[1,2]	2	5	1	5	
[1,4]	4	5	1	5	
[1,8]	8	8	1	8	
[1,16]	16	8	1	8	
[1,32]	32	8	1	8	
[1,64]	64	8	1	8	
[1,128]	128	11	1	11	
[1,256]	256	11	1	11	
[1,512]	512	11	1	11	
[1,1024]	1024	11	1	11	
[1,2048]	2048	11	1	11	

## Comparison #'s vs. Vector size 2^k



part (e)  $f(n) = 2 + \log_2 n$ 

part (f) Changing vector to 2<sup>k</sup> - 1.

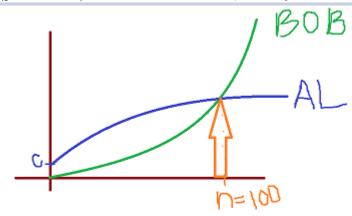
Range $[1,n]$	Target for	# comp. for	Target for	# comp. for	Result of the
	incr. values	incr. values	decr. values	decr. values	formula in
					item 5 (e)
[1,1]	1	2	1	2	
[1,3]	3	5	1	5	
[1,7]	7	5	1	5	
[1,15]	15	8	1	8	
[1,31]	31	8	1	8	
[1,63]	63	8	1	8	
[1,127]	127	8	1	8	
[1,255]	255	11	1	11	
[1,511]	511	11	1	11	
[1,1023]	1023	11	1	11	
[1,2047]	2047	11	1	11	

 $\overline{\text{part (g) } O(logn)}$ 

2. (problem 4.7) 8nlogn will be faster than  $2n^2$ , however up until  $n_0$ ,  $2n^2$  will be faster. To find  $n_0$ , setting both equations equal to each other will determine when they equal, and at what value of n, that this occurs.  $8nlogn = 2n^2 \Longrightarrow 4logn = n$ . Assuming that this is a logarithm base 2, then  $n/log_2n = 4 \Longrightarrow 16/log_216 = 16/4 = 4$ . So n = 16. For every value after n = 16, algorithm A will be better, so for all n > 16.

3. (problem 4.21) The worst case run time in terms of n is:  $O(n^2)$  In terms of in, then it'd be  $O(n^2)$ . This is not a linear algorithm because  $n^2$  is quadratic.

4. (problem 4.39) The cutoff mark n = 100, basically means that the following graph applies:



since nlogn does not  $crossn^2$ , we know that Al's f(n) formula must be ADDED to some constant, in other words, Al's f(n) = nlogn + C, whereas Bob's  $f(n) = n^2$ . Also that is an image I drew in paint. copyright.

$$\begin{split} 5.s &= n. \\ f(n) &= 2(n-1) + 1 \Longrightarrow O(n) \\ f(n) &= 2(n-2) + 1 \Longrightarrow O(n) \\ f(n) &= 2(n-1)(2)(n) + 1 = 4n(n-1) + 1 \Longrightarrow O(n^2) \\ f(n) &= (2+2)(n-1) + 2 = 4n - 4 + 2 = 4n - 2 \Longrightarrow O(n) \end{split}$$