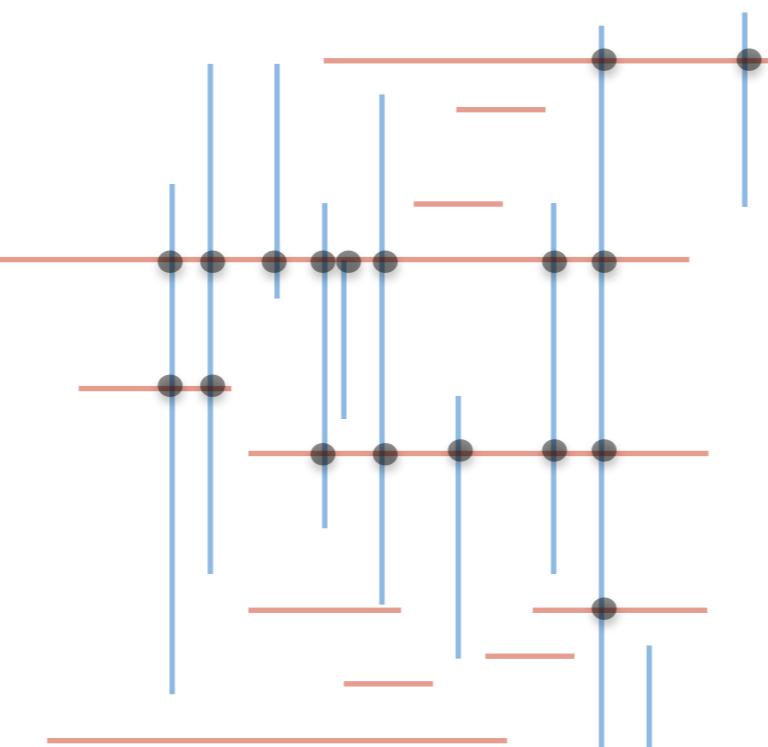


Line segment intersection:

(I) **Orthogonal** line segment intersection

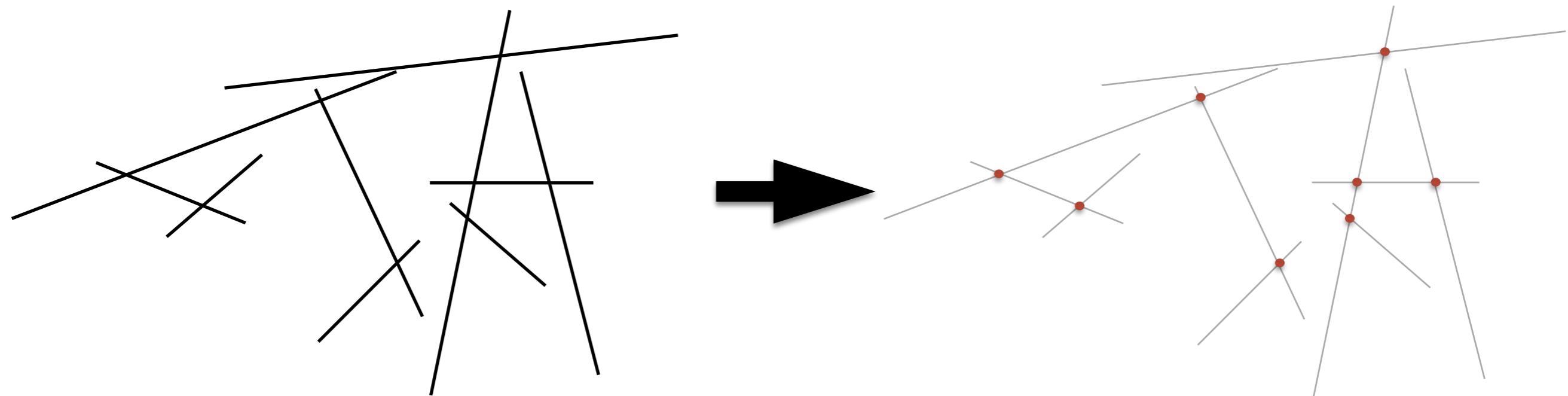


Line segment intersection

- The problem (what)
- Applications (why)
- Algorithms (how)
 - A special case: Orthogonal line segments
 - Next time: General case: Bentley-Otman line sweep algorithm

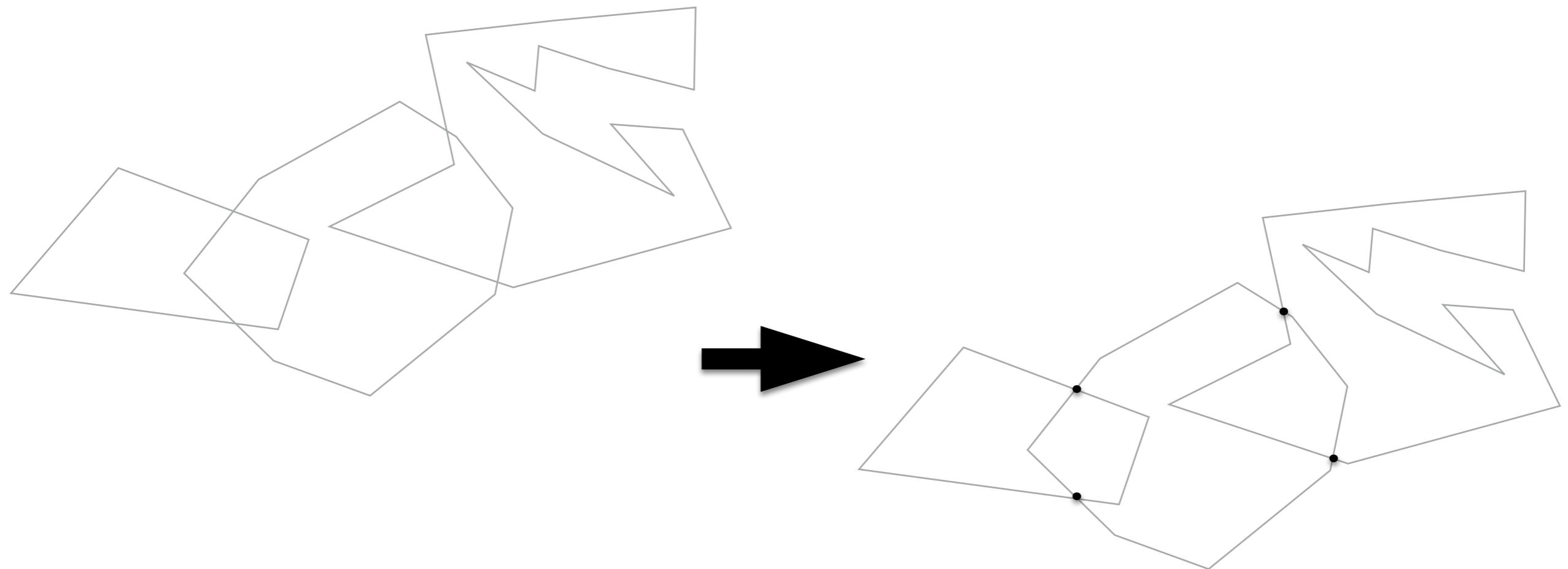
Line segment intersection

Problem: Given a set of line segments in 2D, find all their pairwise intersections.



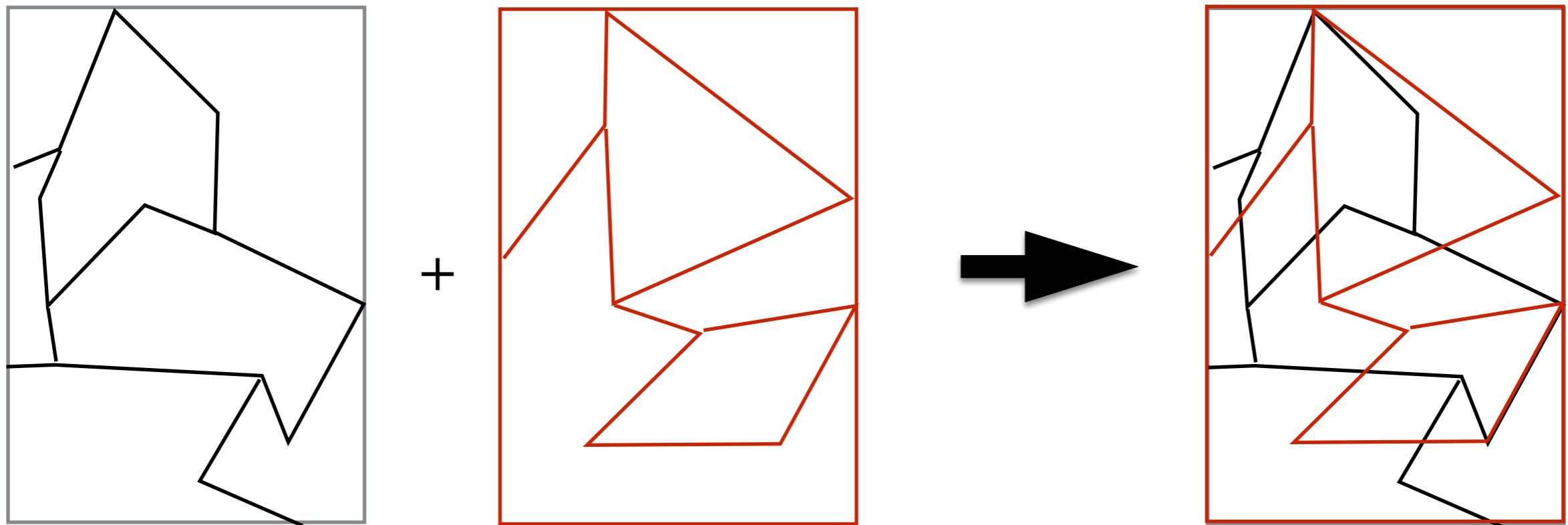
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Line segment intersection

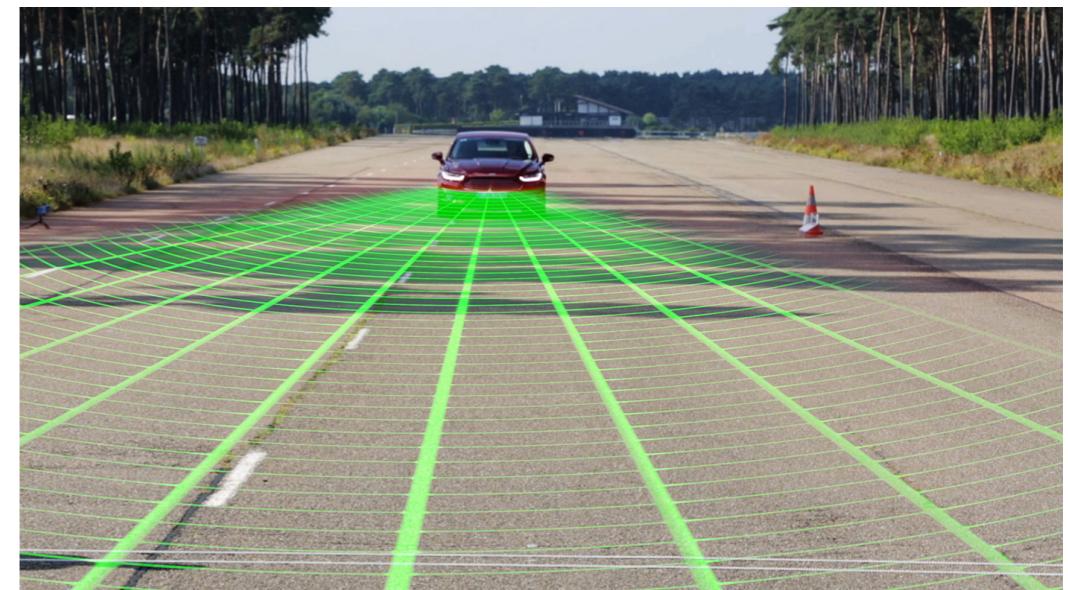
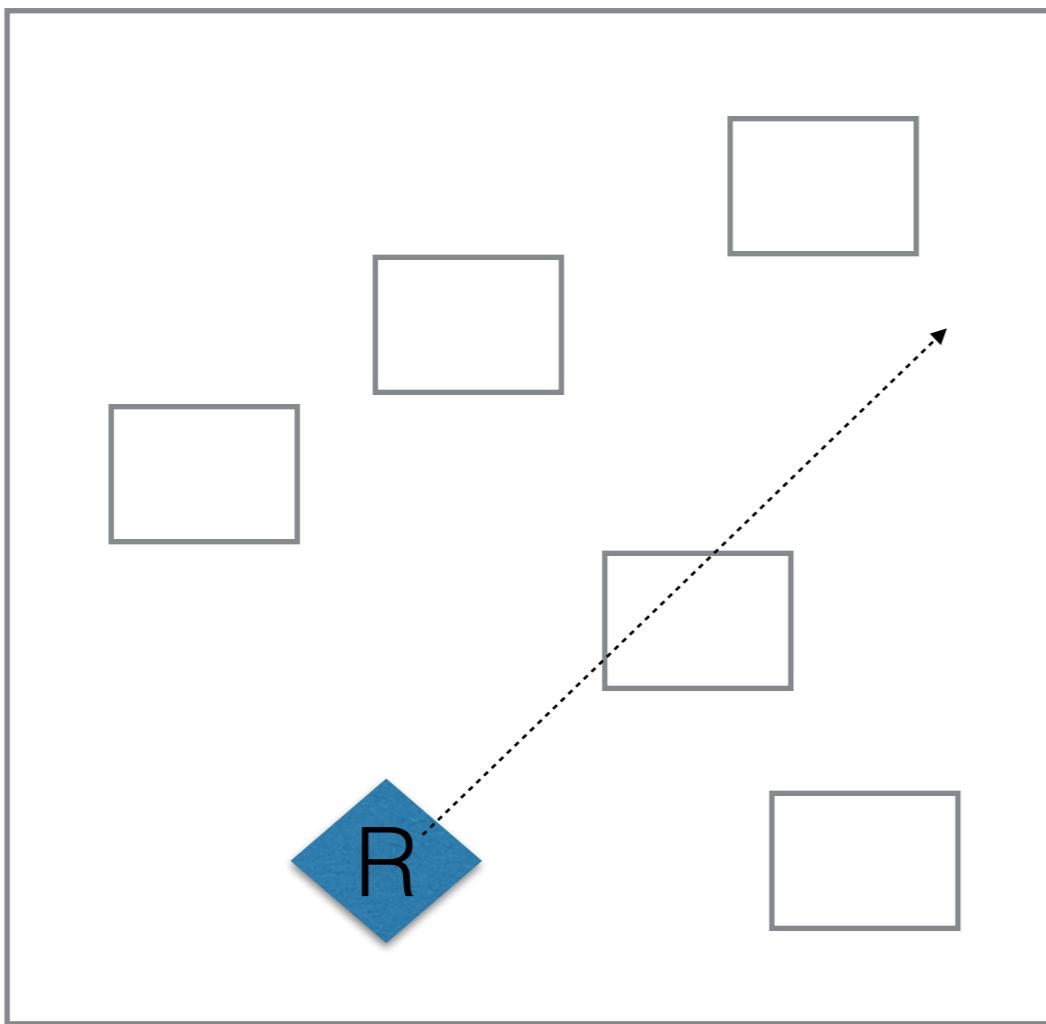
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Line segment intersection: Applications

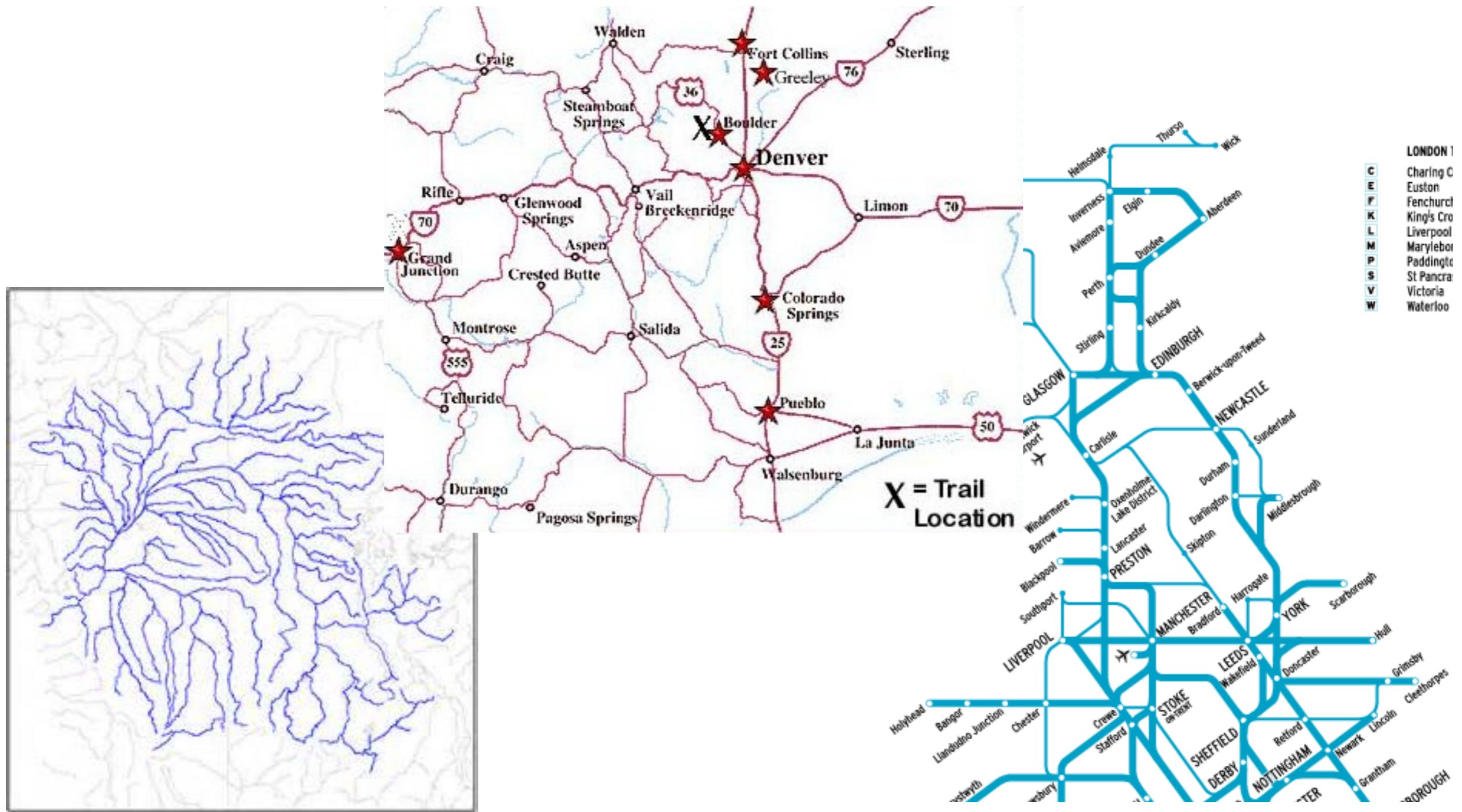
Applications

Motion planning and collision detection in autonomous systems/robotics



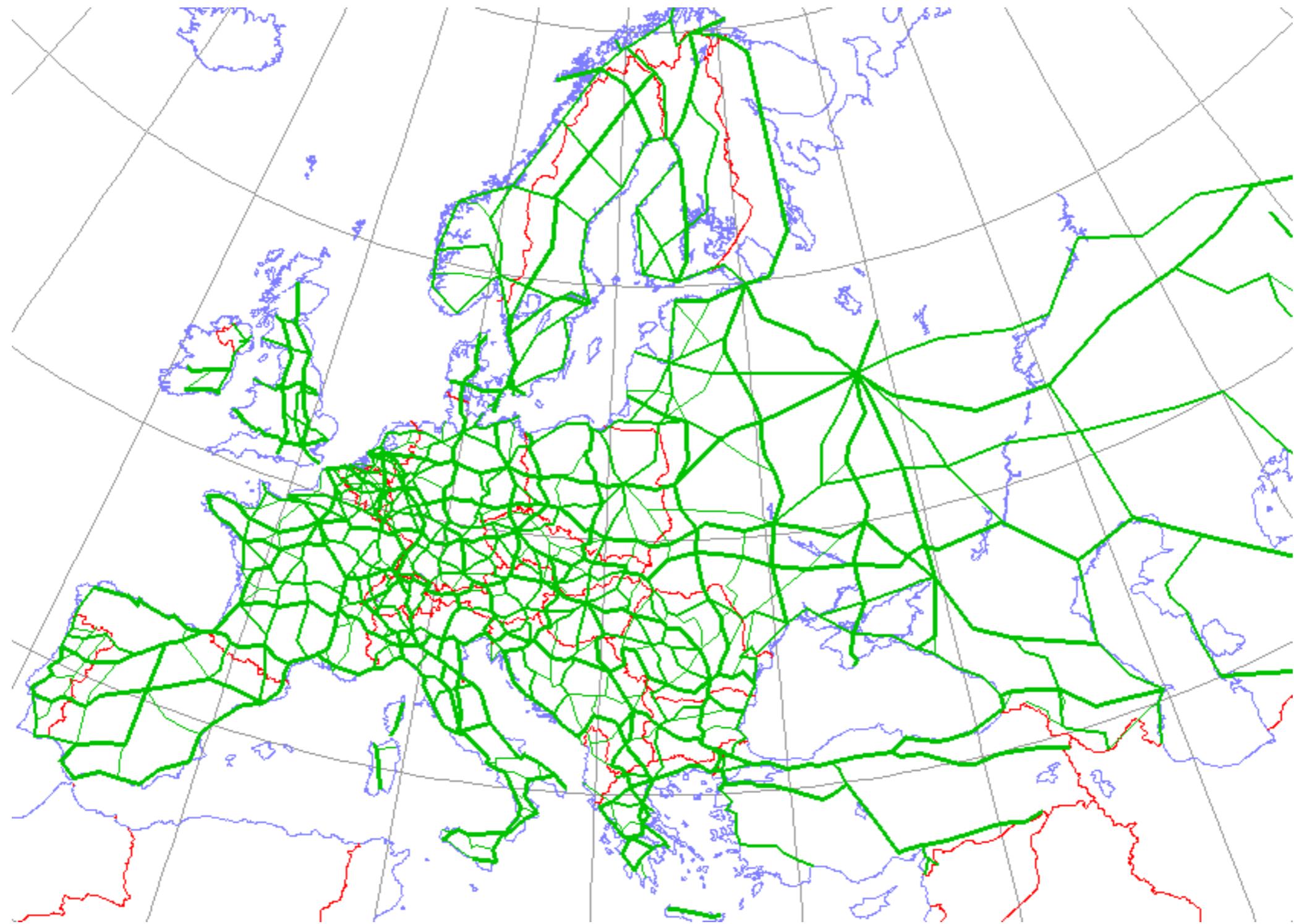
Applications

Geographic data: River networks, road networks, railways, ..



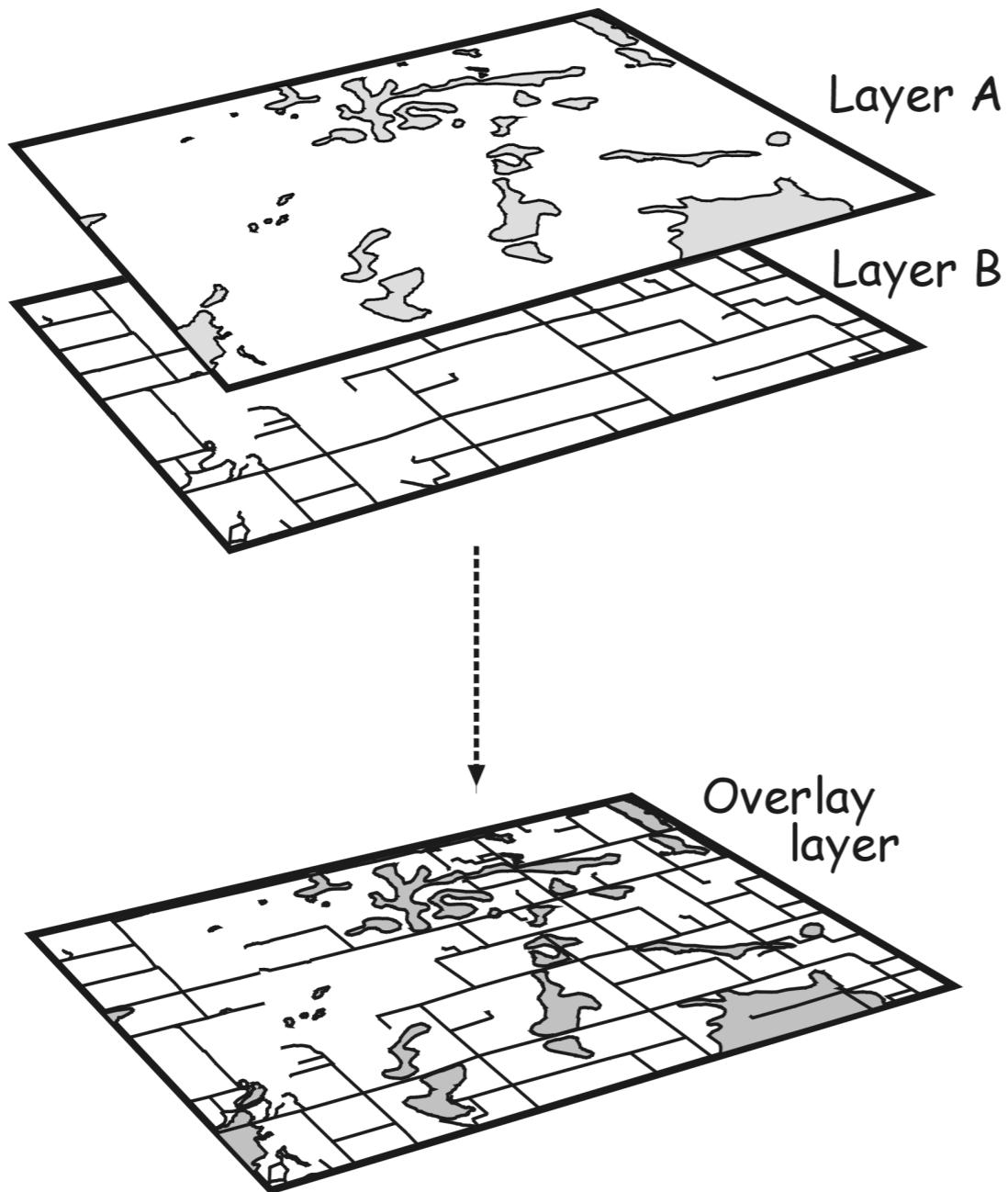
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Applications

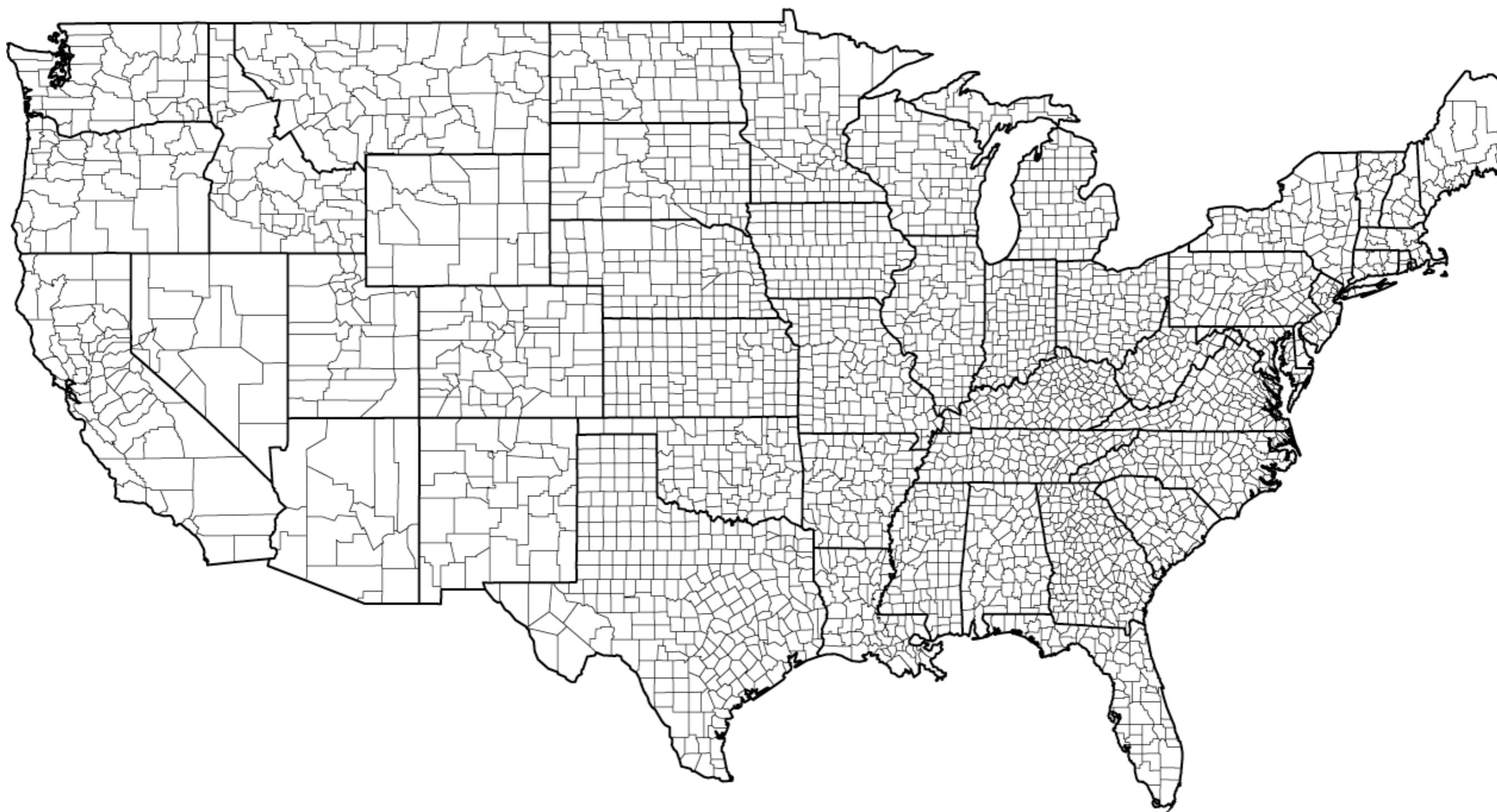
Map overlay in GIS



from: www.geo.hunter.cuny.edu/aierulli/gis2/lectures/Lecture2/fig9-30_raster_overlay.gif

Applications

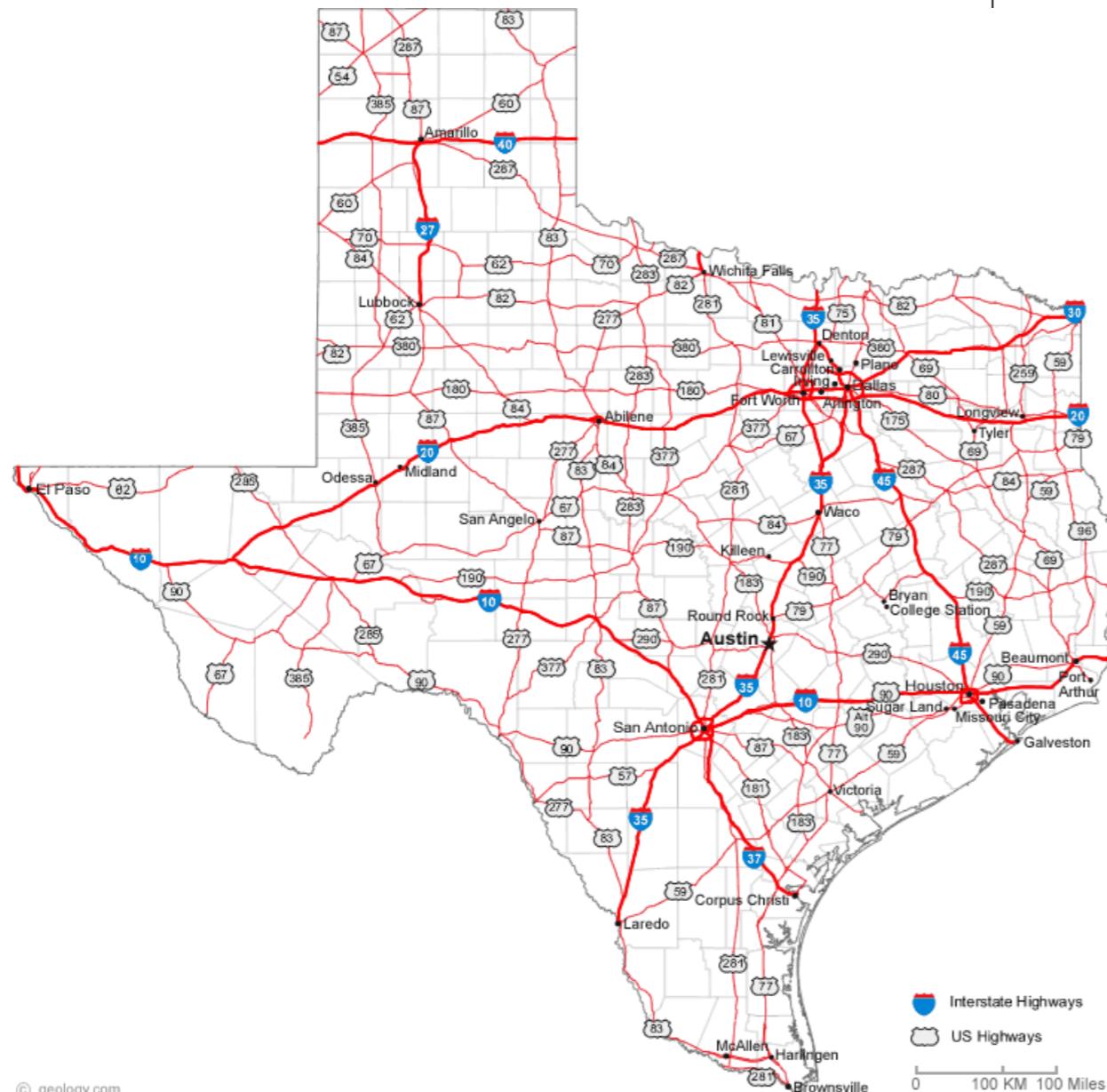
Segment data in GIS: river network, road networks, counties, etc



Counties of Texas

Applications

Map overlay in GIS



Computing line segment intersection: Algorithms

Naive

Notation

- n : size of the input (number of segments)
- k : size of output (number of intersections)

Problem: Given a set of n line segments in 2D, find all their pairwise intersections.

Naive

Notation

- n : size of the input (number of segments)
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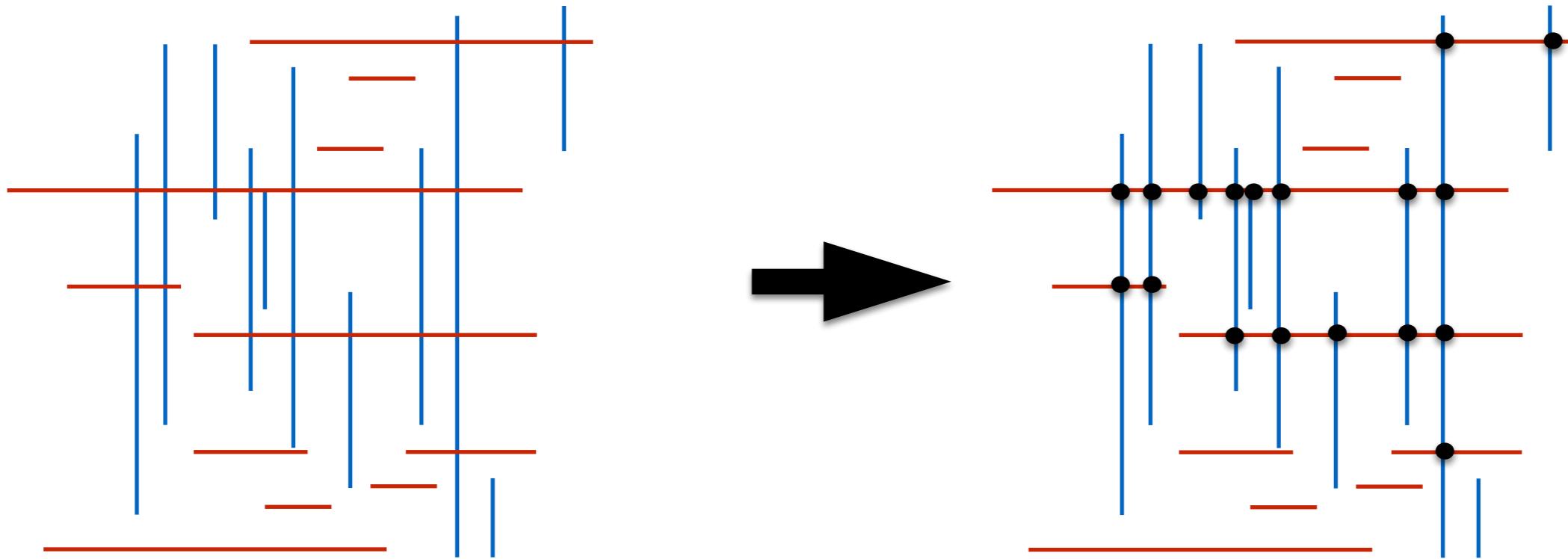
Problem: Given a set of n line segments in 2D, find all their pairwise intersections.

To think:

- Give upper and lower bounds for k , draw examples that achieve these bounds.
- Give a straightforward algorithm that computes all intersections and analyze its running time. Give scenarios when this algorithm is efficient/inefficient.
- What is your intuition of an upper bound for this problem? (how fast would you hope to be able to solve it?)

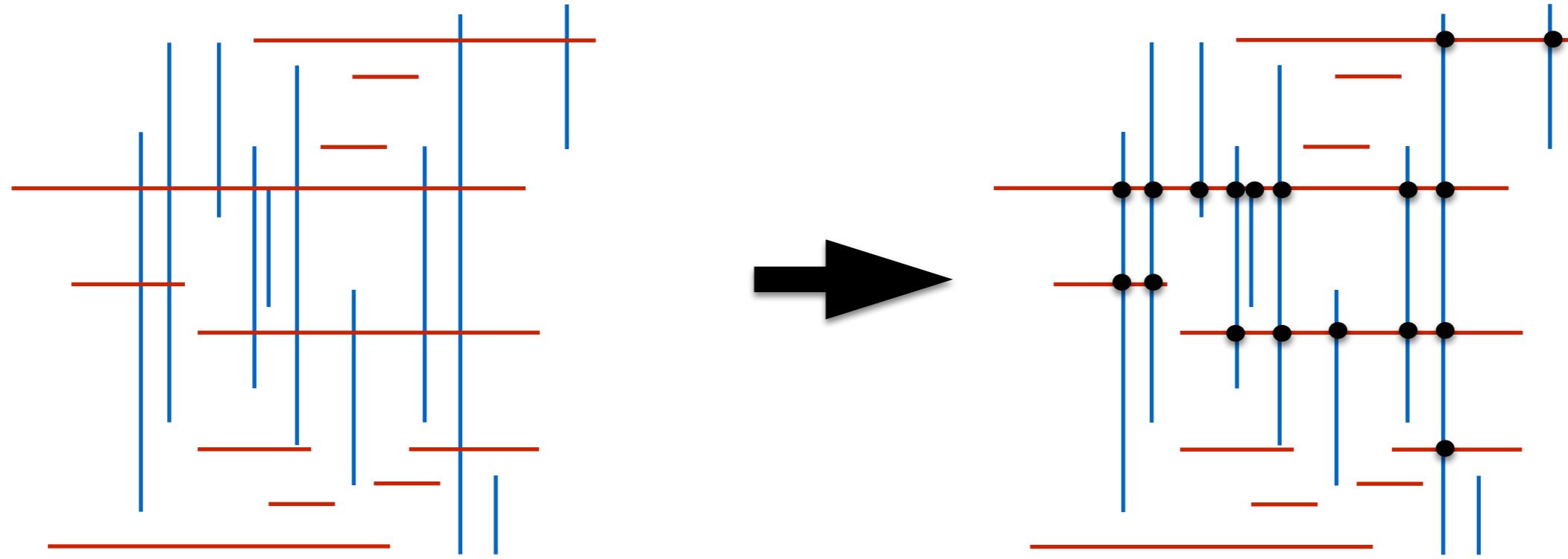
A special case: Orthogonal line segment intersection

Problem: Given a set of **orthogonal** line segments in 2D, find all their pairwise intersections.



A special case: Orthogonal line segment intersection

Problem: Given a set of **orthogonal** line segments in 2D, find all their pairwise intersections.



To think:

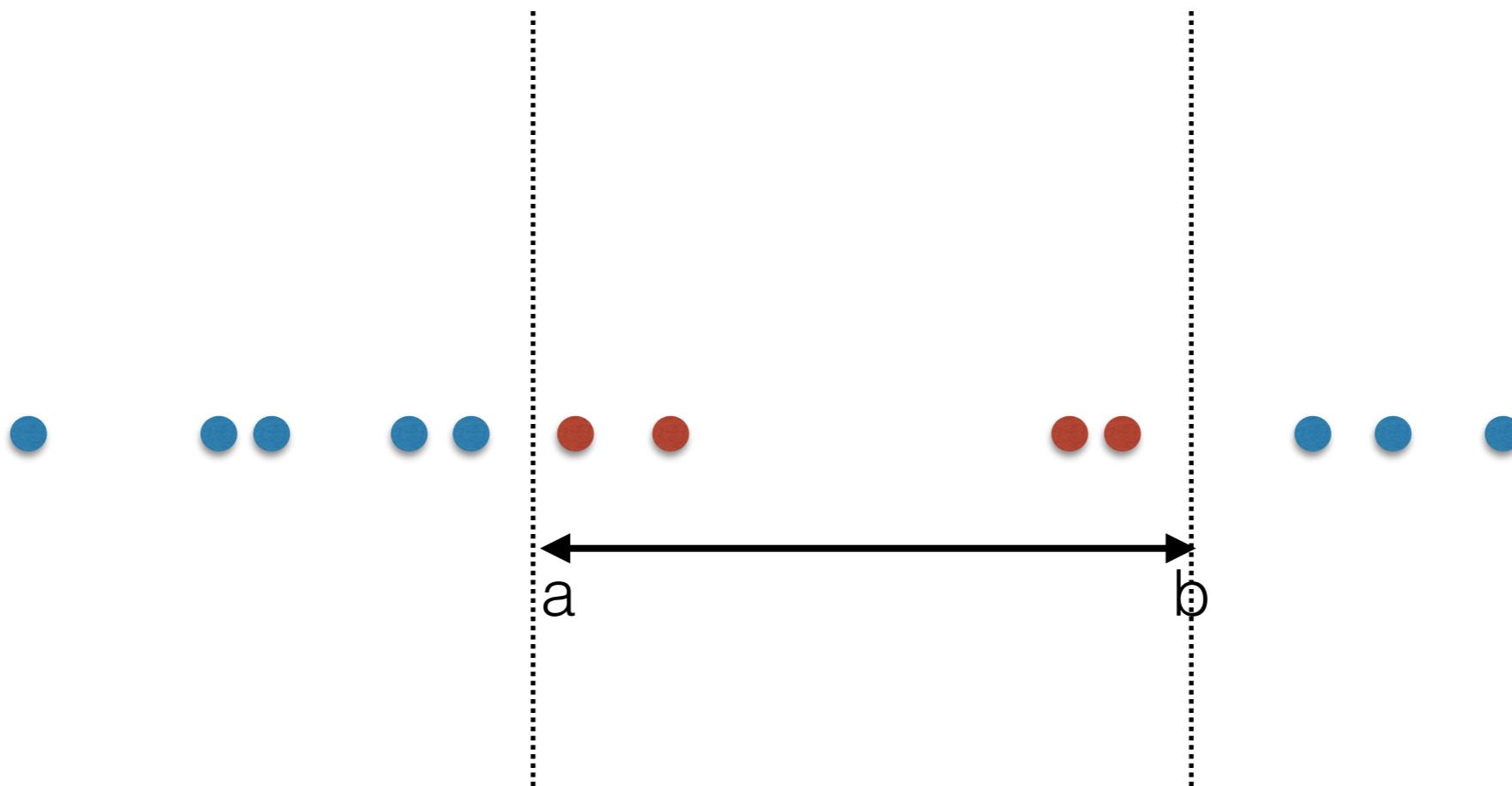
- Come up with a straightforward algorithm and analyze its time
- Can you do better?

detour: range searching

1D Range Searching

- Given a set of values $P = \{x_1, x_2, x_3, \dots x_n\}$
- Pre-process it in order to answer

rangeSearch(a, b): return all elements in P in interval (a, b)

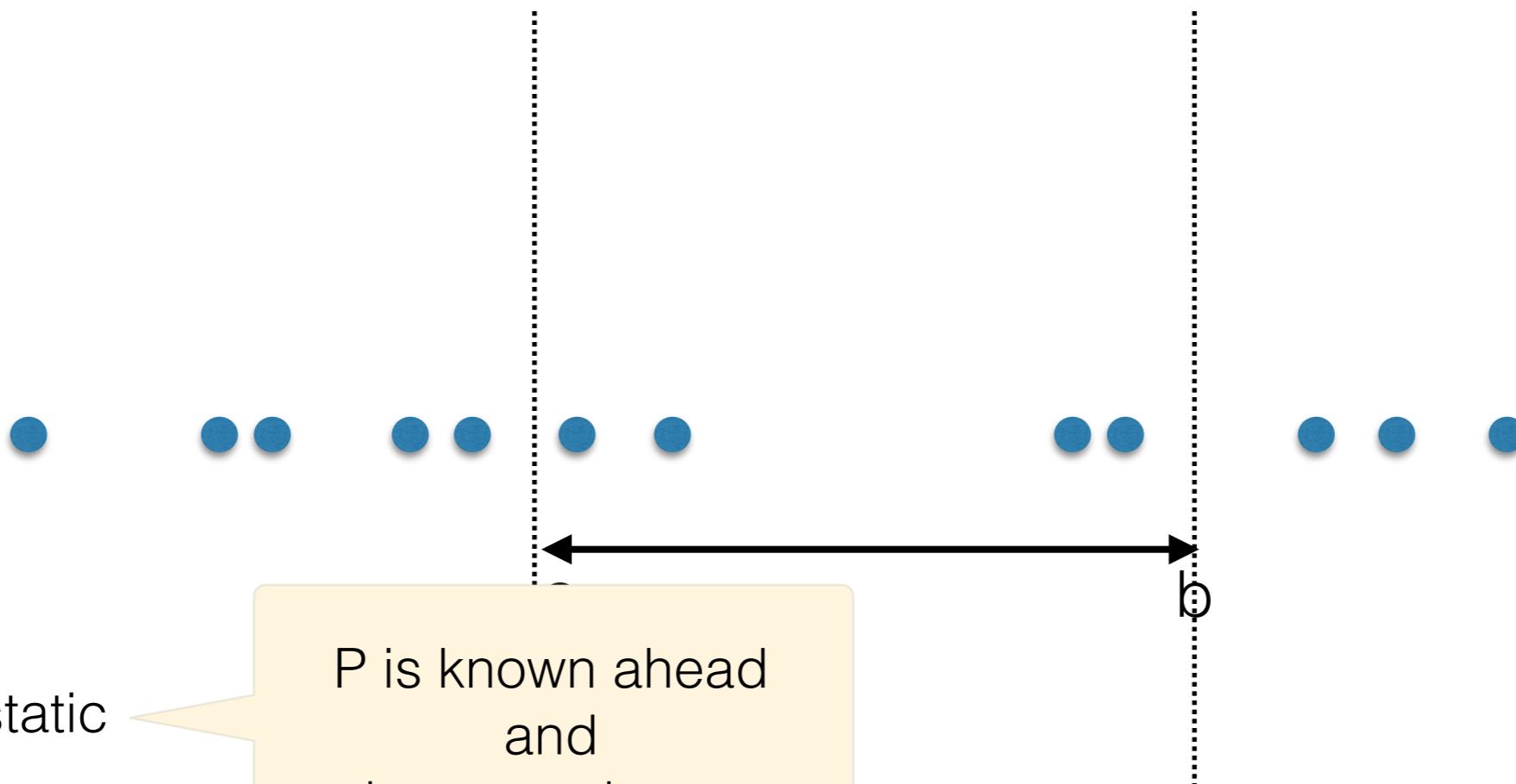


1D Range Searching

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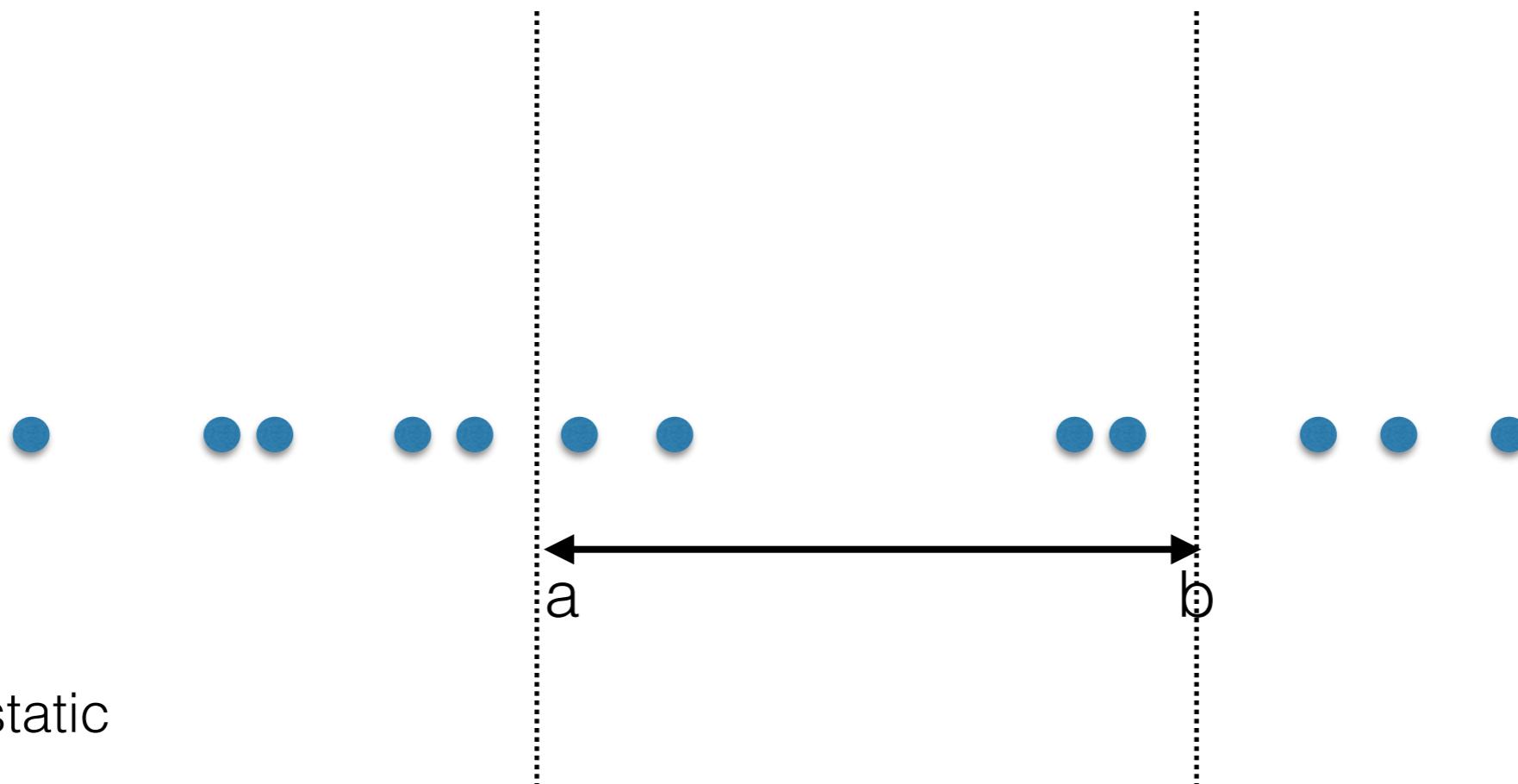


- If P is static

1D Range Searching

- Given a set of values $P = \{x_1, x_2, x_3, \dots, x_n\}$
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rangeSearch(a, b): return all elements in P in interval (a, b)

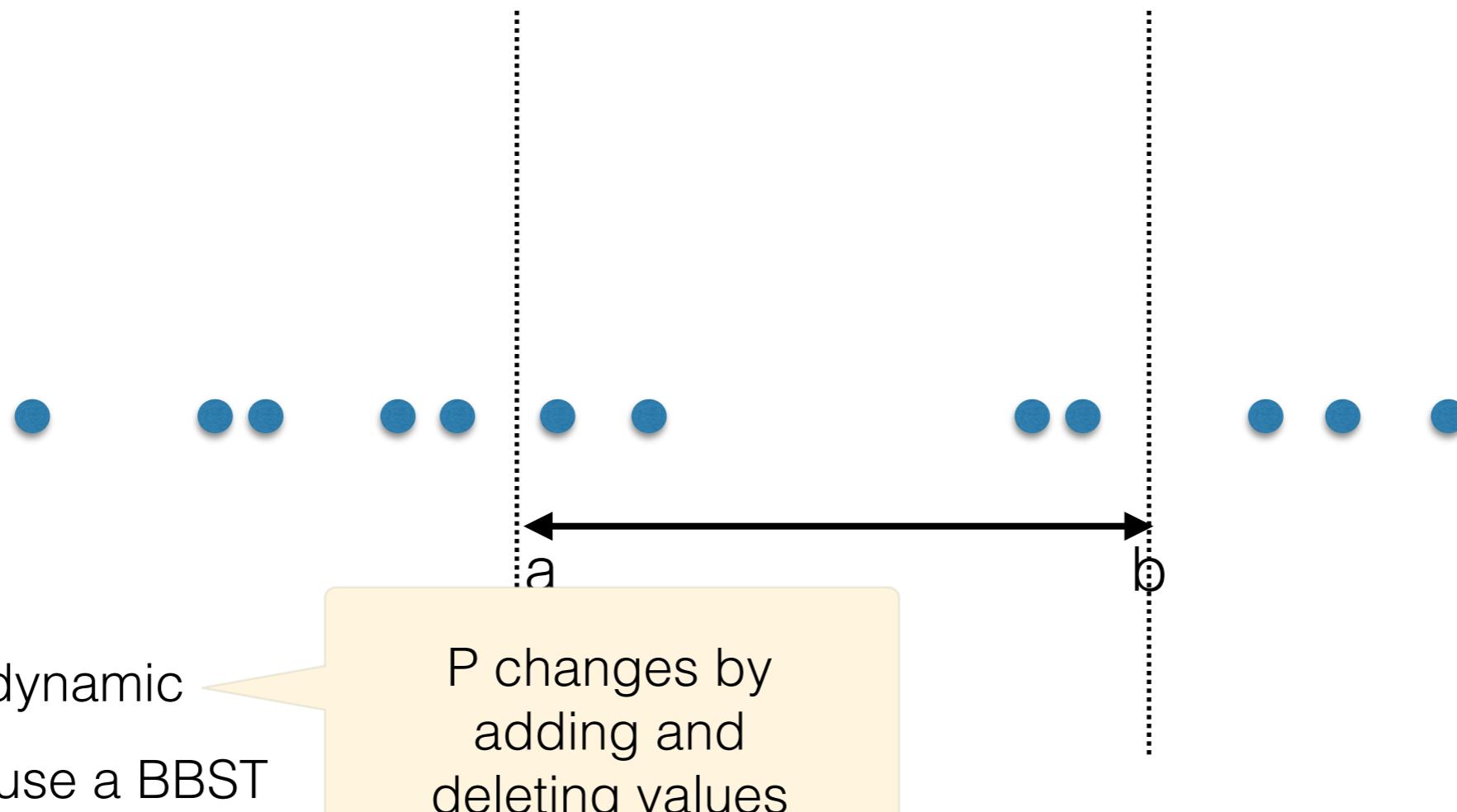


1D Range Searching

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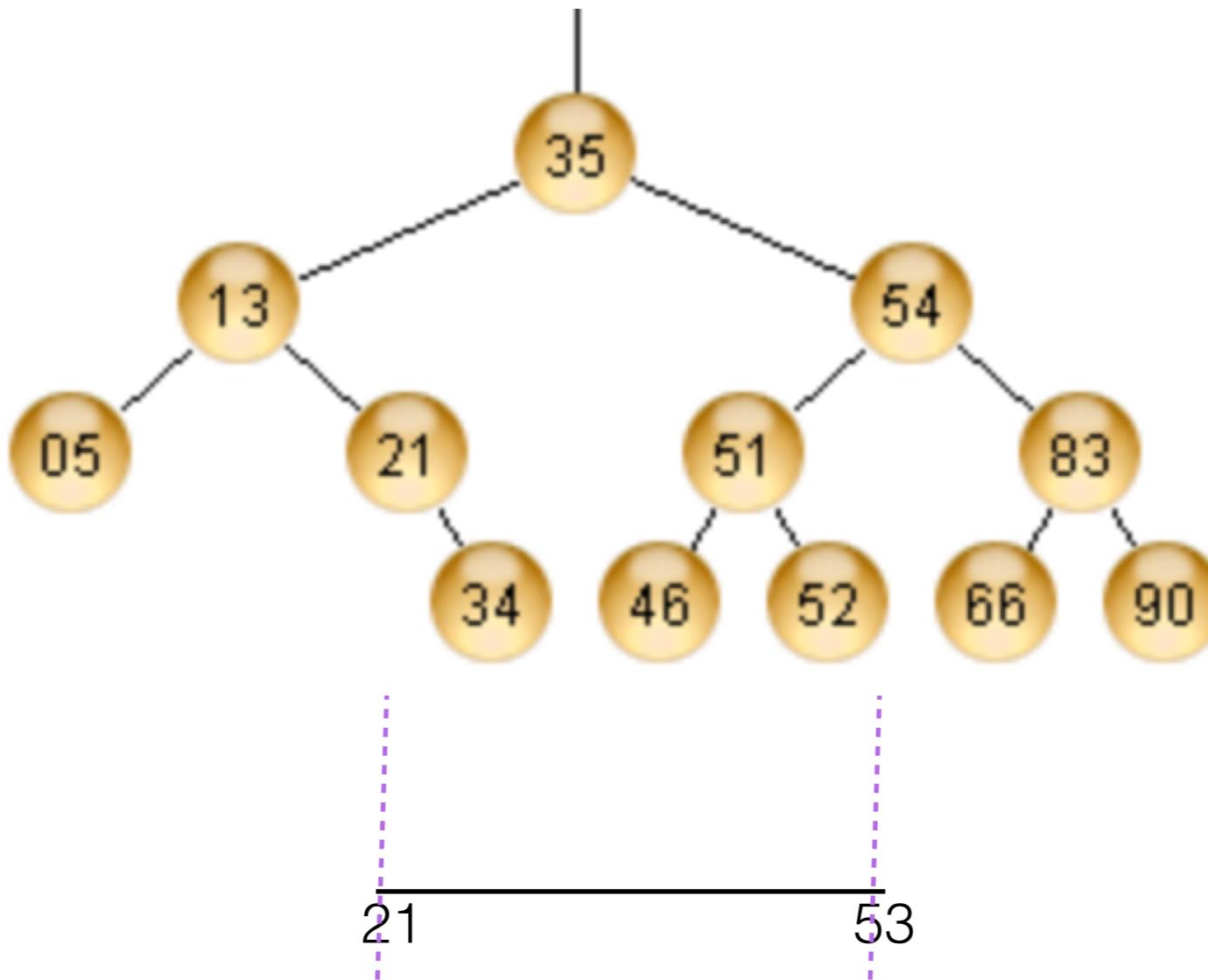
- Pre-process it in order to answer

`rangeSearch(a,b)`: return all elements in P in interval (a,b)



1D range searching with Binary Search Trees

Example: range_search(21, 53): return 21, 34, 35, 46, 51, 52

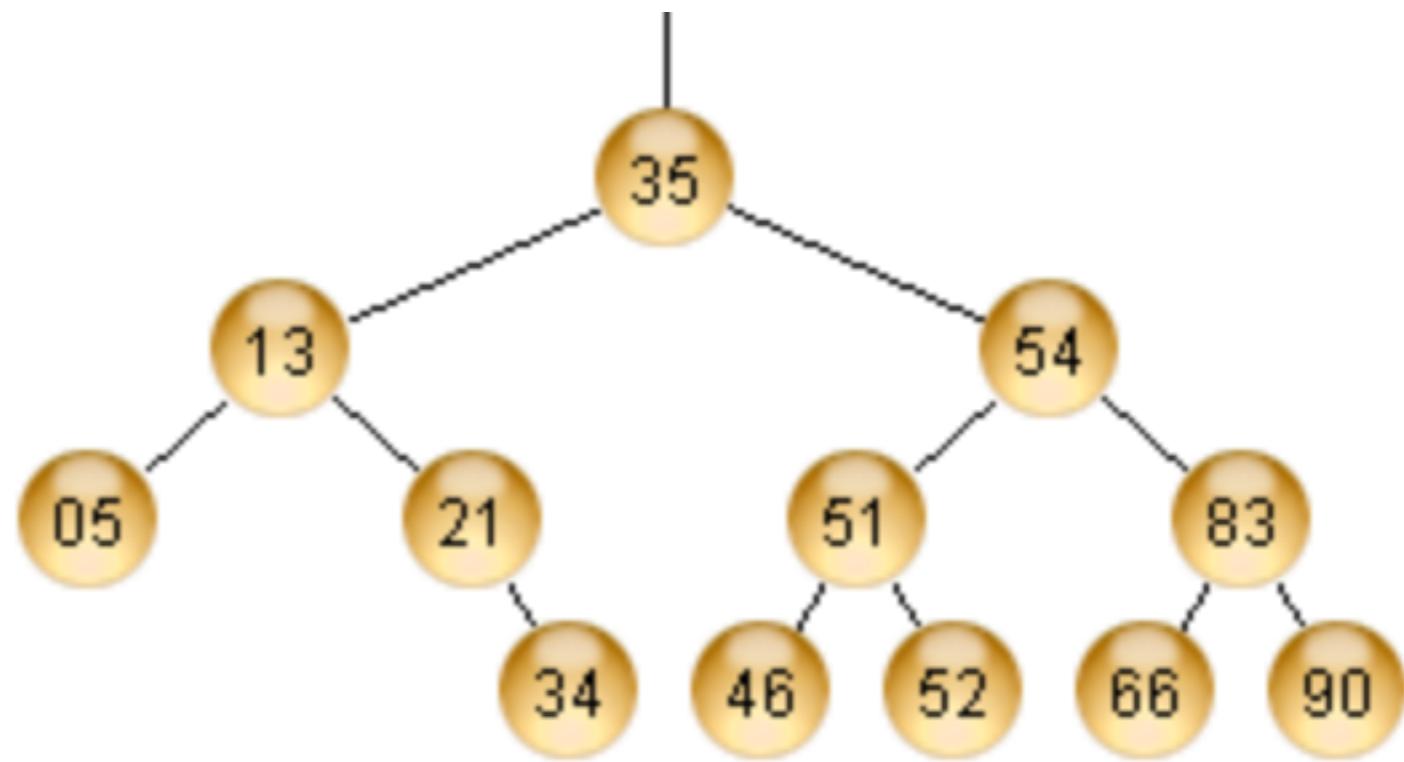


Balanced Binary Search Trees

- crash course -

Binary Search Trees (BST)

- Operations
 - insert
 - delete
 - search
 - successor, predecessor
 - traversals (in order, ..)
 - min, max

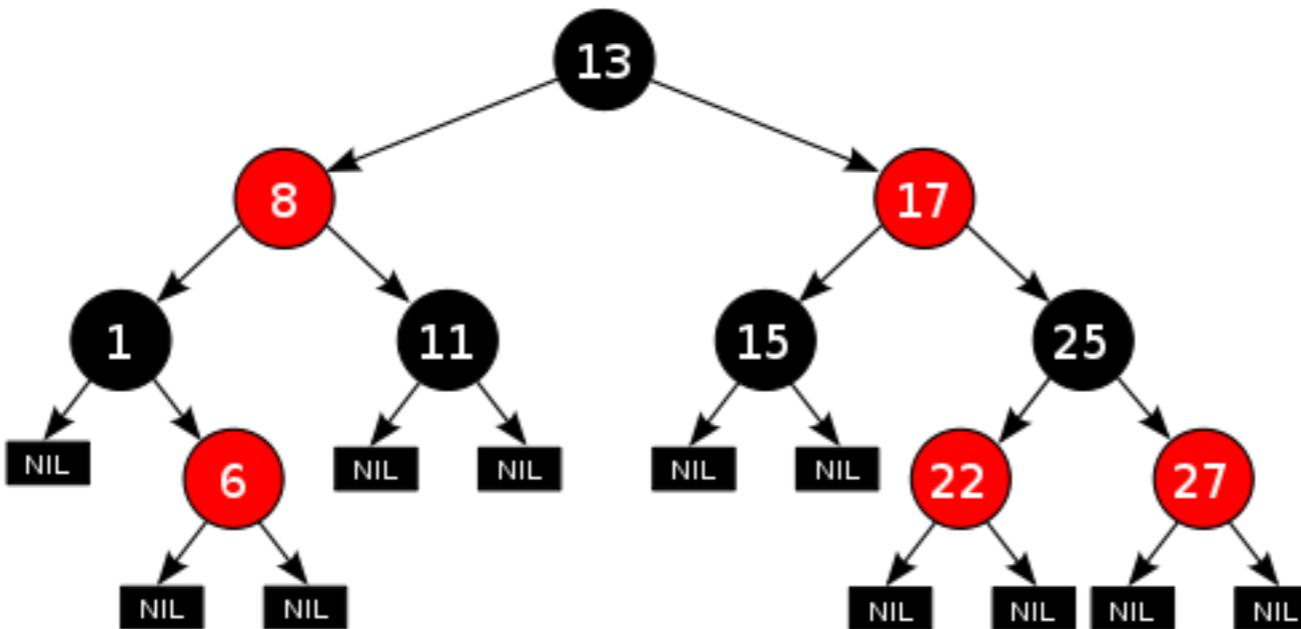


Balanced Binary Search Trees (BBST)

- Binary search trees + invariants that constrain the tree to be balanced
 - $h = O(\lg n)$
- These invariants have to be maintained when inserting and deleting
 - we can think of the tree as self-balancing
- BBST variants
 - red-black trees
 - AVL trees
 - B-trees
 - (a,b) trees
 - ...

Example: Red-Black trees

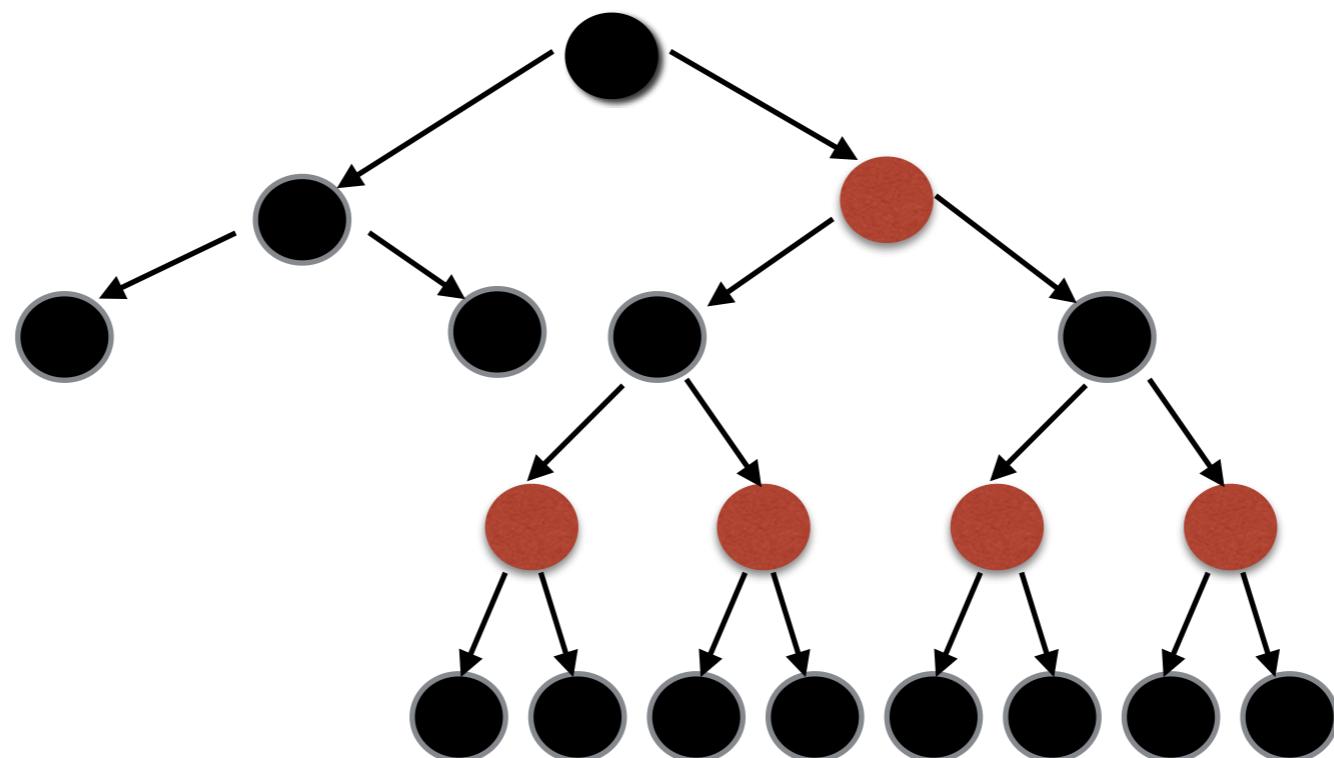
- Binary search tree, and
 - Each node is Red or Black
 - The children of a Red node must be Black
 - The number of Black nodes on any path from the root to any node that does not have two children must be the same



- easier to conceptualize the tree as containing explicit NULL leaves, all Black
- the number of Black nodes on any root-to-leaf path must be the same

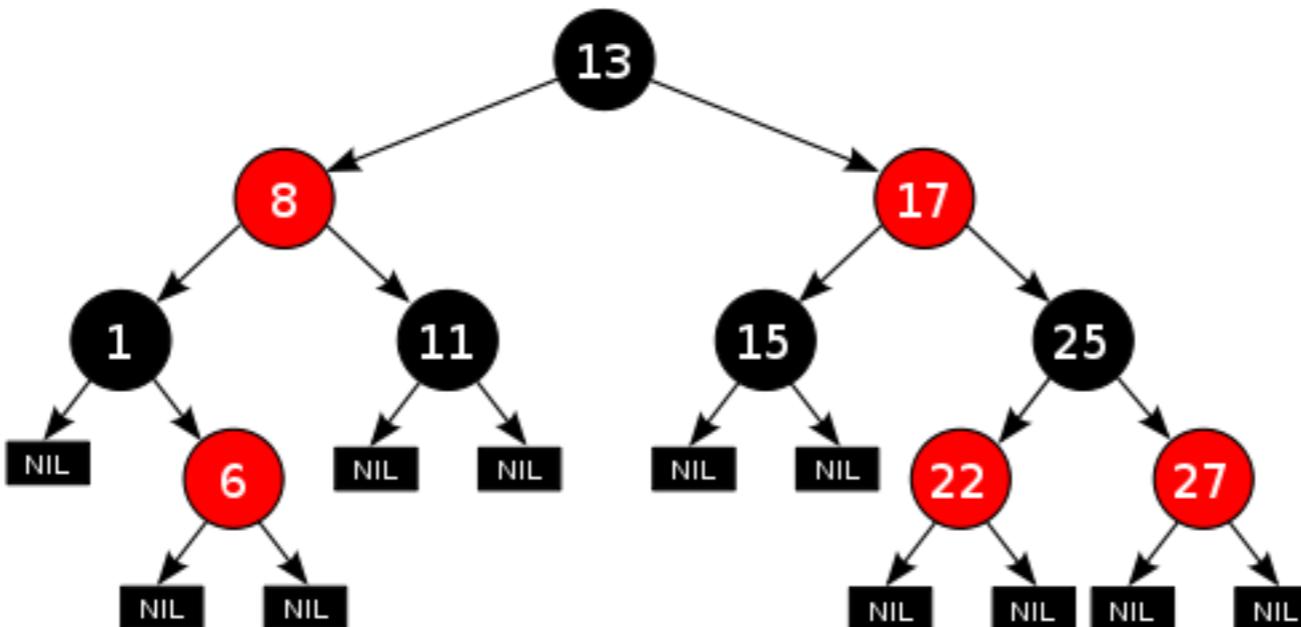
Example: Red-Black trees

- Theorem:
 - A Red-Black tree of n nodes has height $\Theta(\lg n)$.



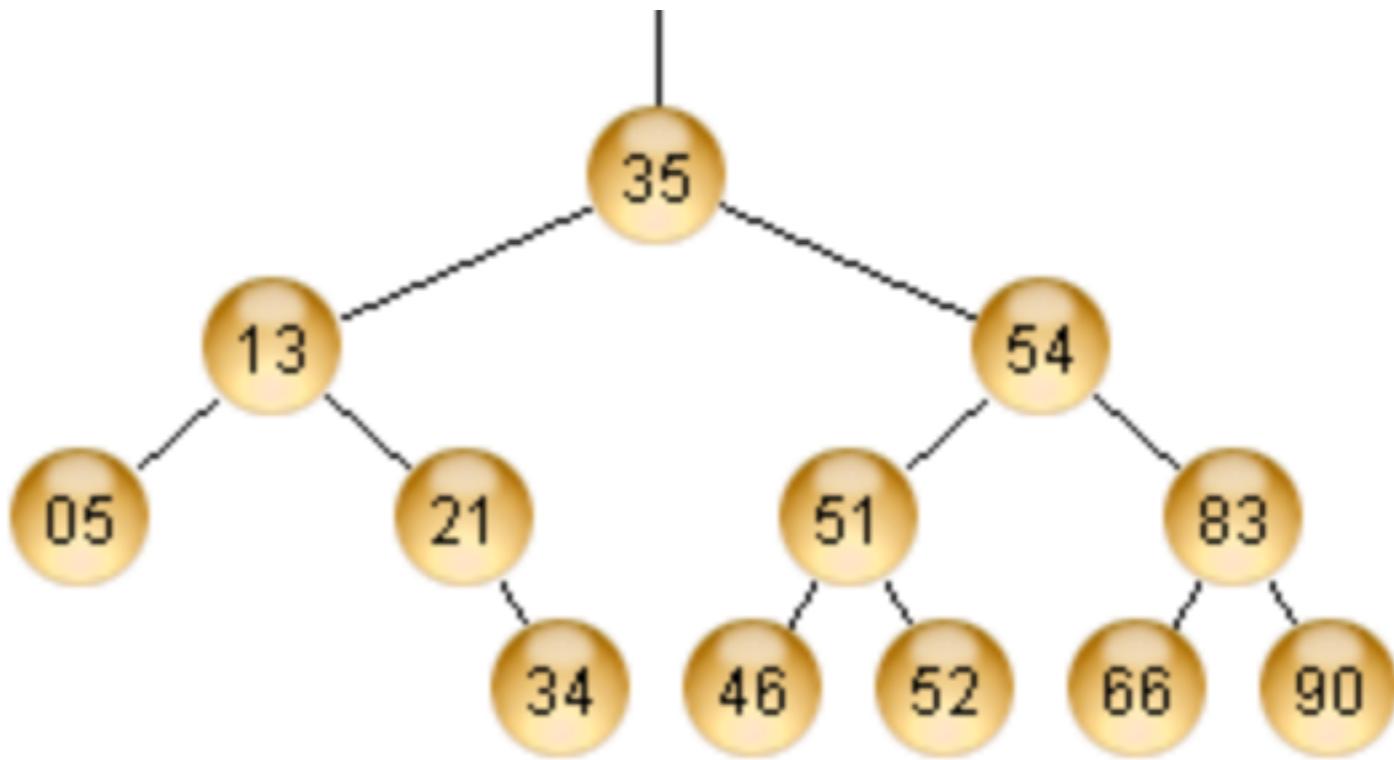
Example: Red-Black trees

- Theorem:
 - After an insertion or a deletion, the RB tree invariants can be maintained in additional $O(\lg n)$ time. This is done by performing rotations and recoloring nodes on the path from the inserted/deleted node to the root.



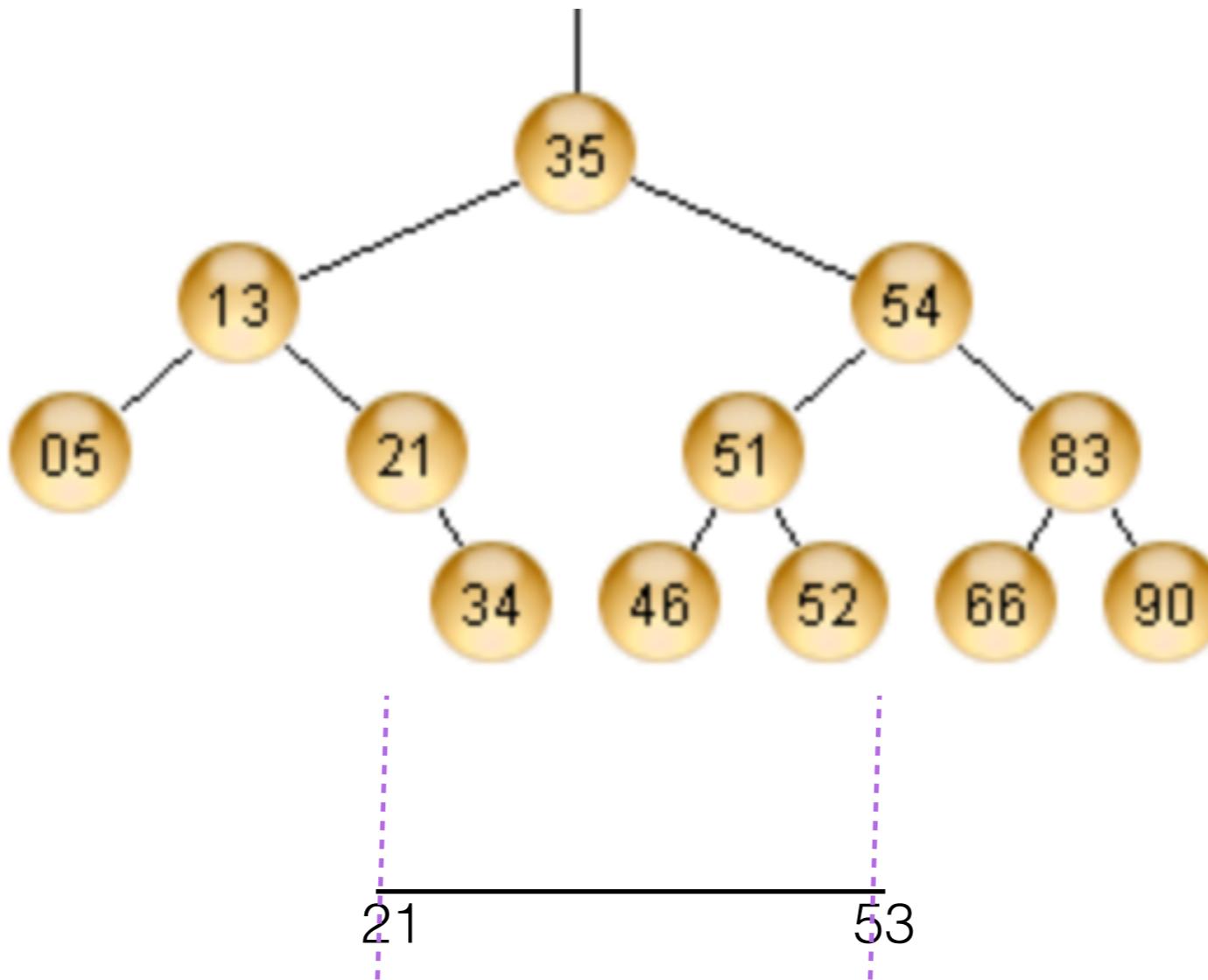
Binary Search Trees

- Operations
 - insert
 - delete
 - search
 - successor, predecessor
 - traversals (in order, ..)
 - min, max
 - range search (1D)



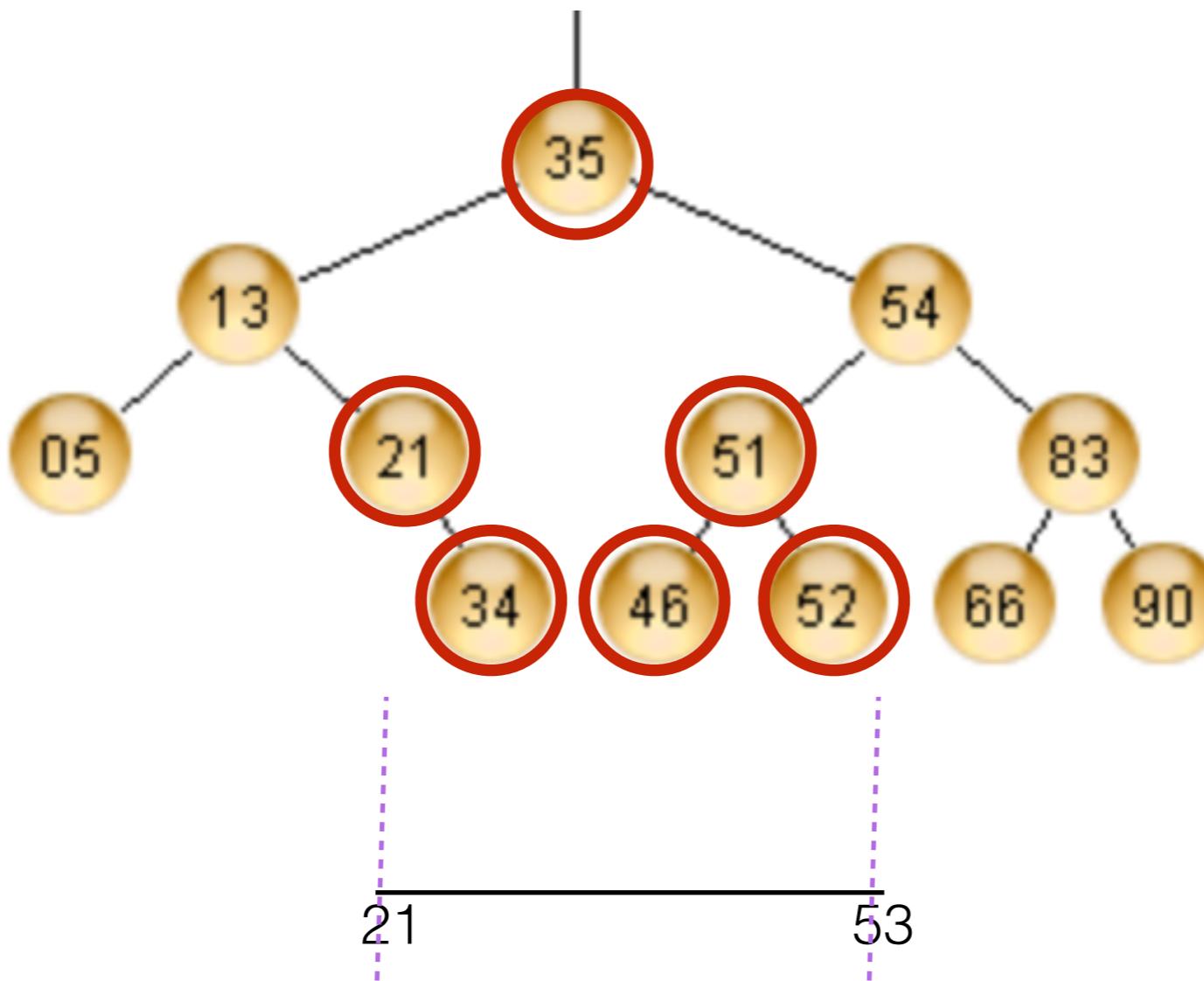
1D range searching with Binary Search Trees

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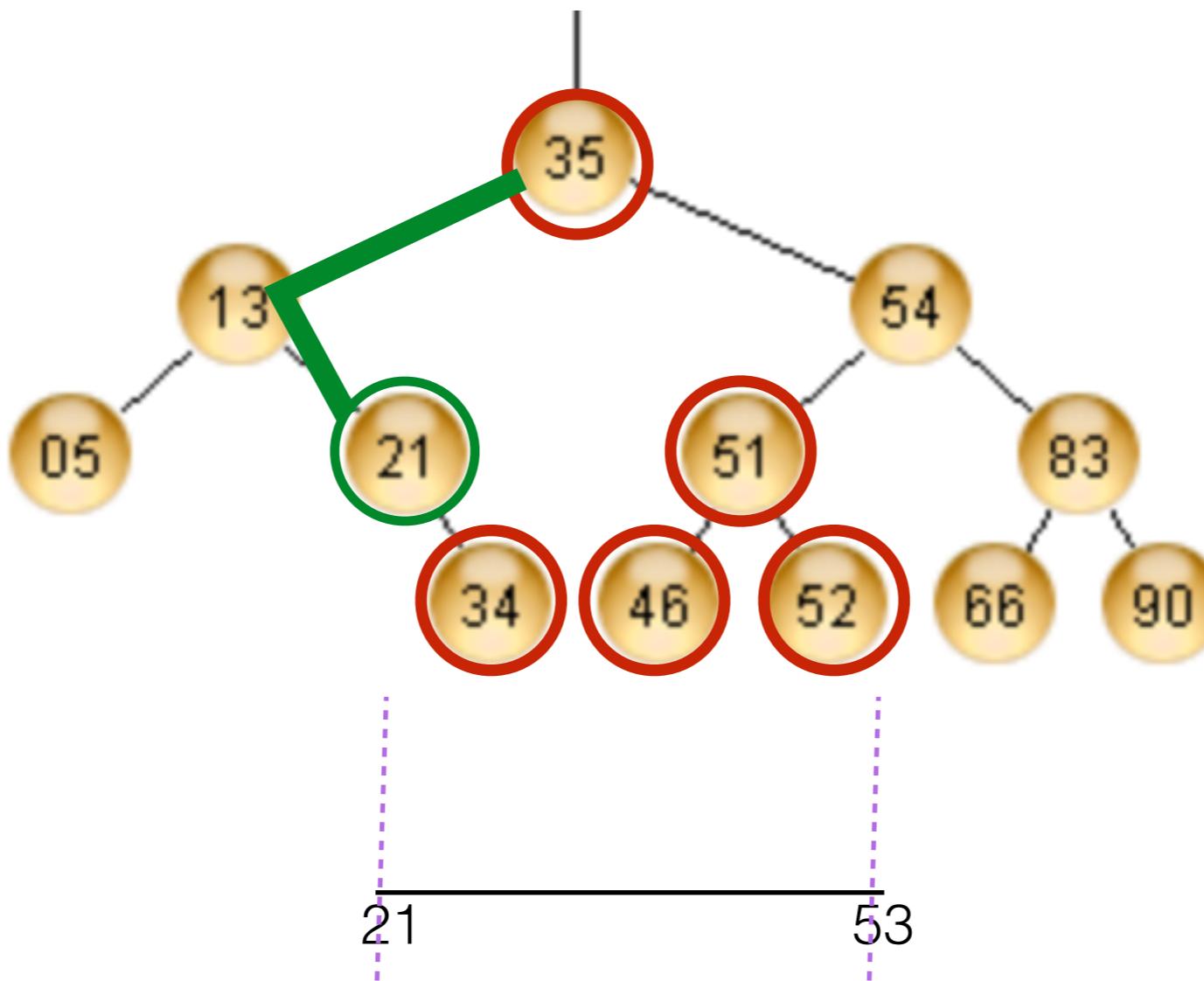
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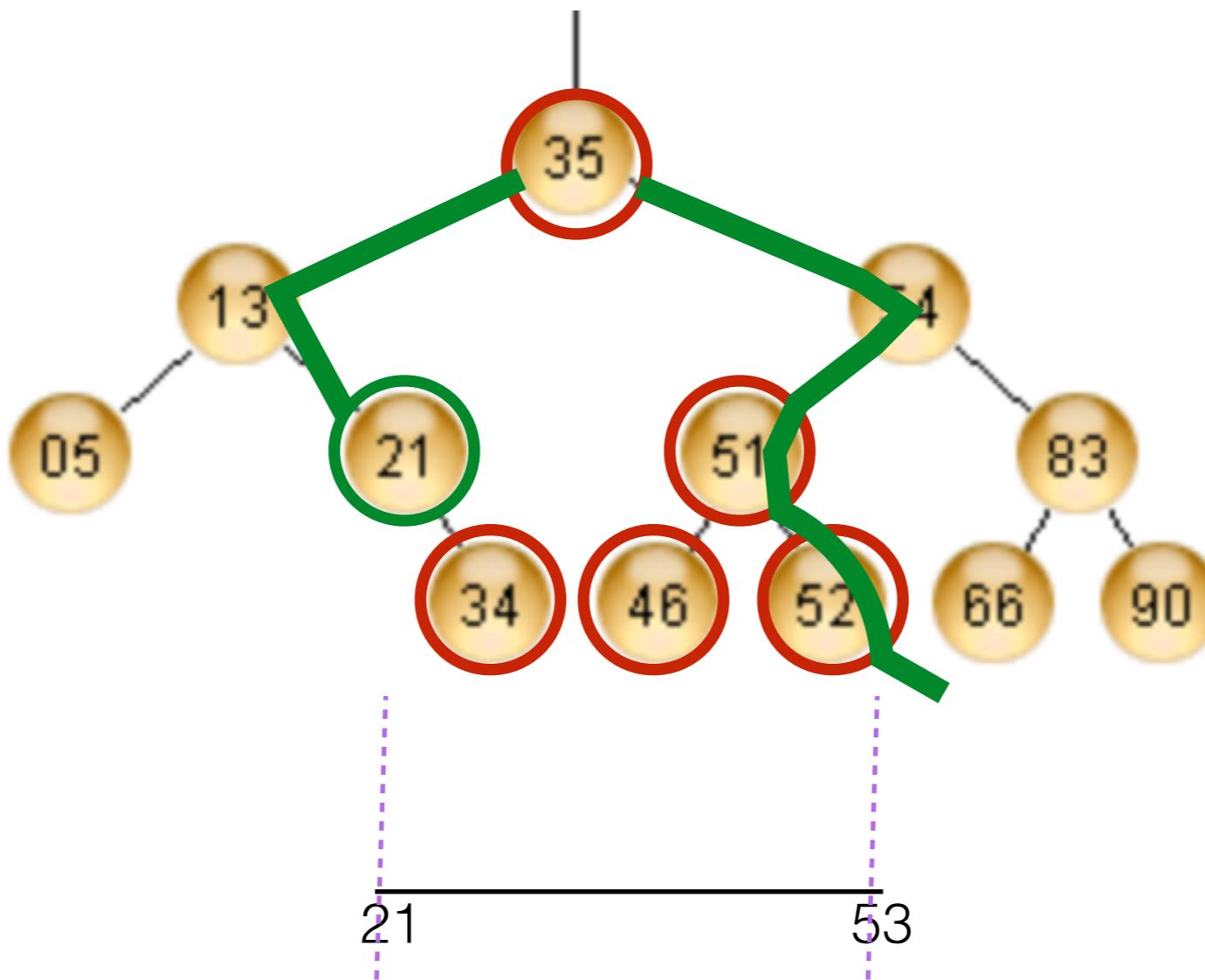
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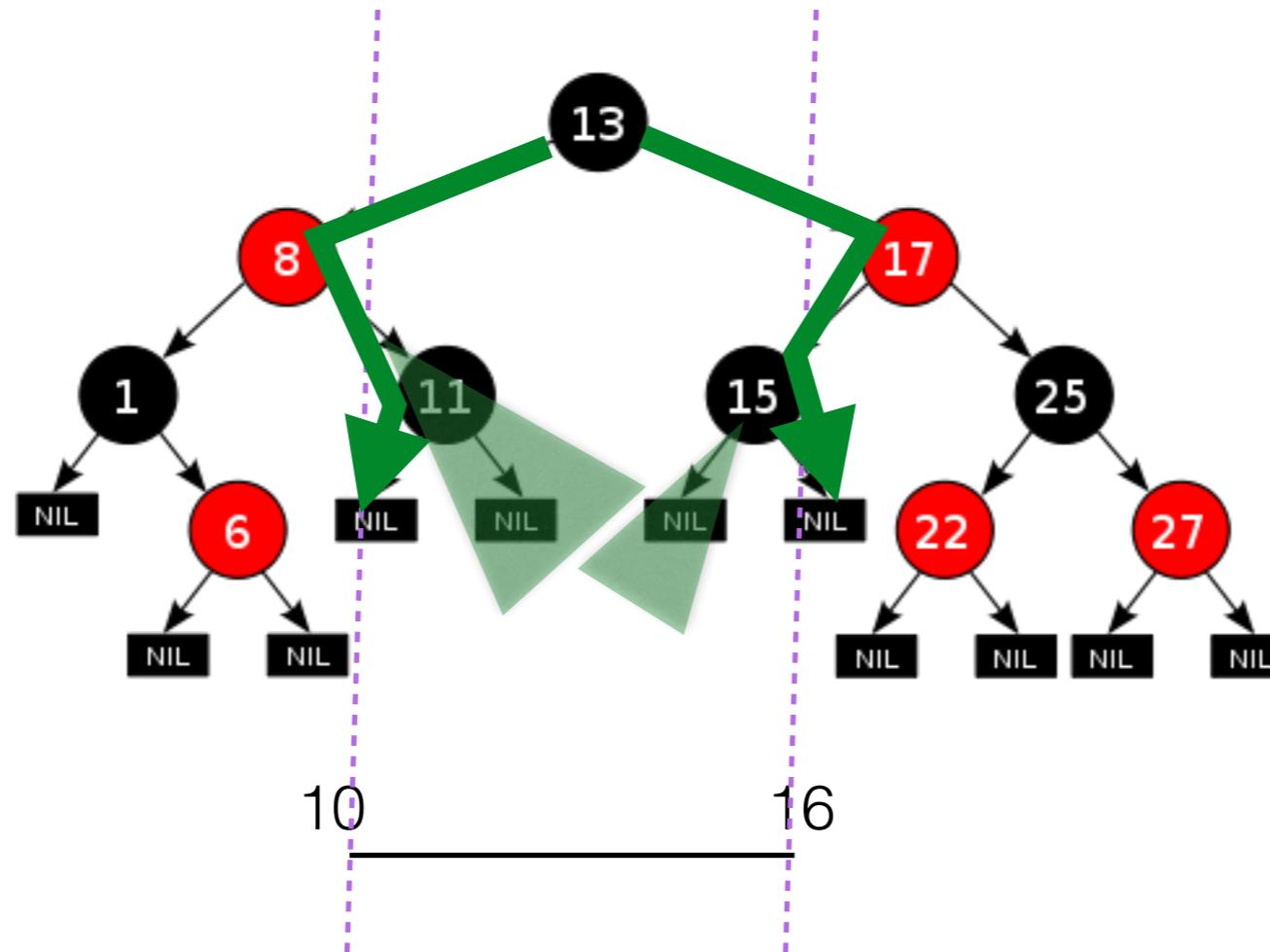
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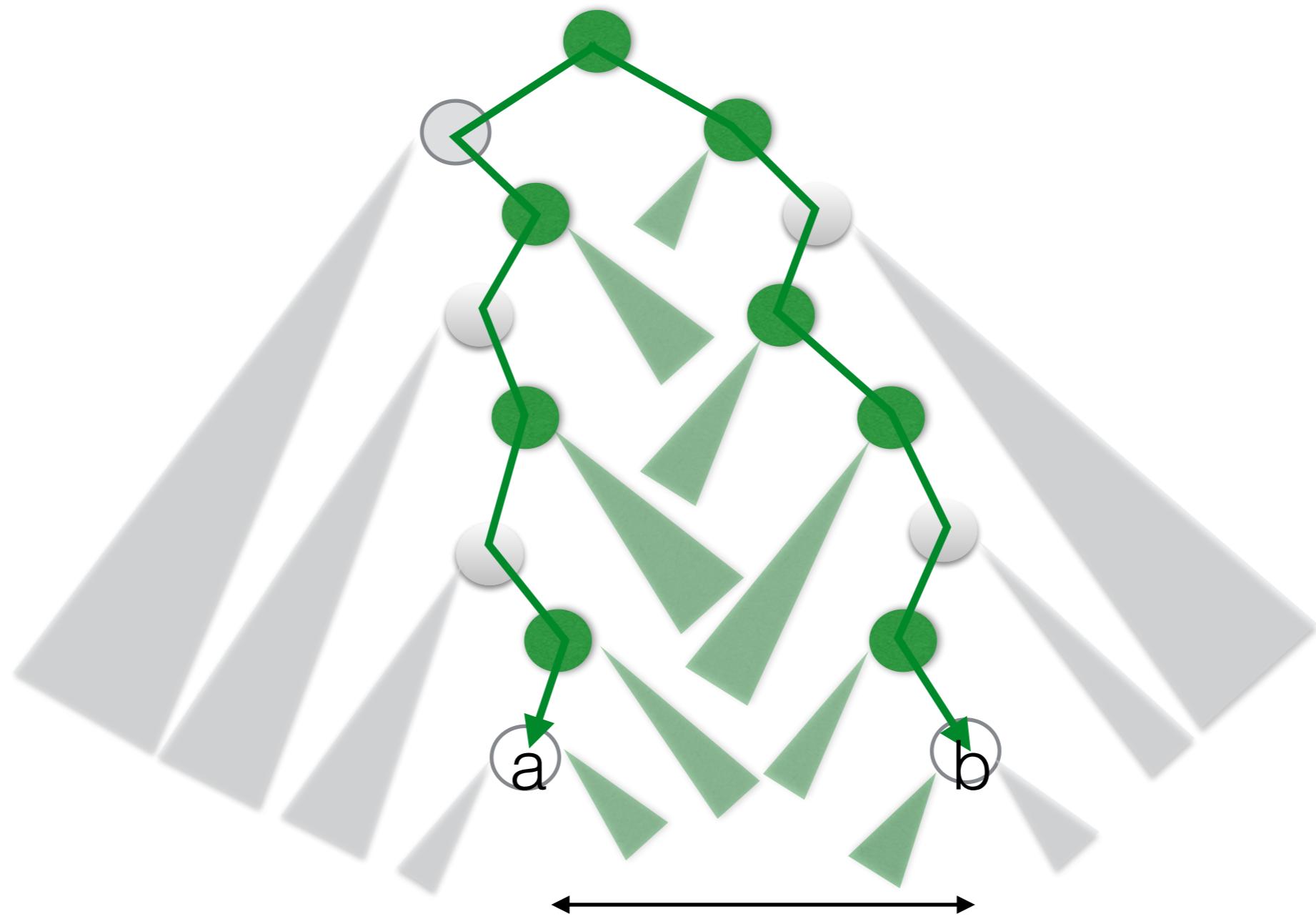
1D Range Searching with Red-Black Trees

Example: range_search(10, 16): return 11, 13, 15



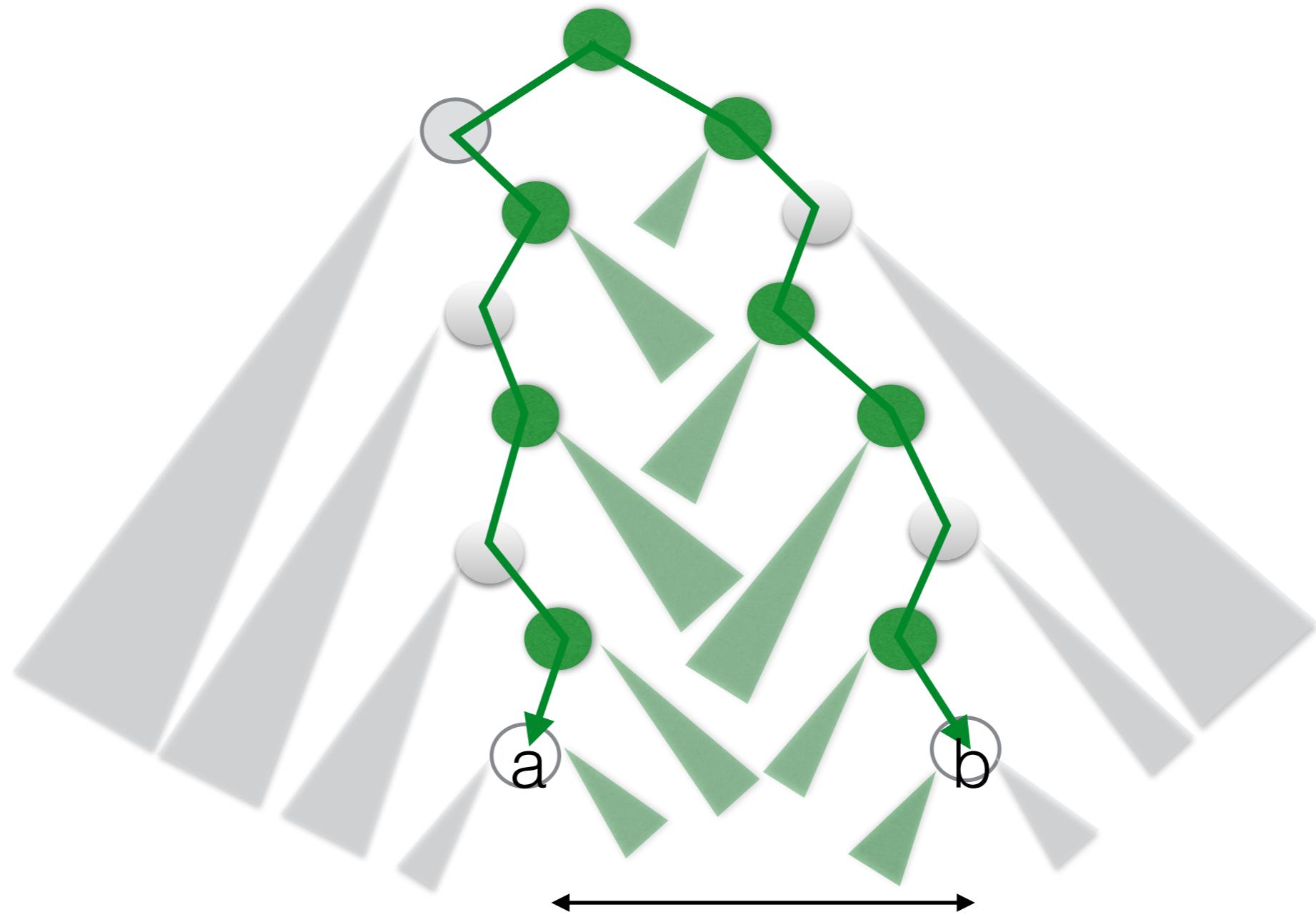
1D range searching with Binary Search Trees

- Range search (a, b):



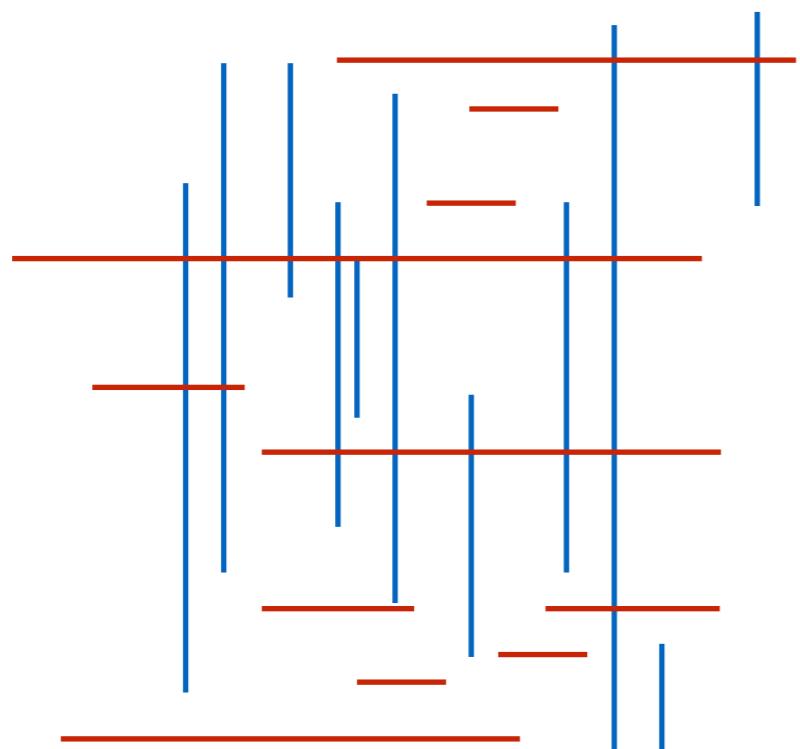
1D range searching with Binary Search Trees

- Range search (a, b):
- Can be answered in $O(\lg n+k)$, where $k = O(n)$ is the size of output

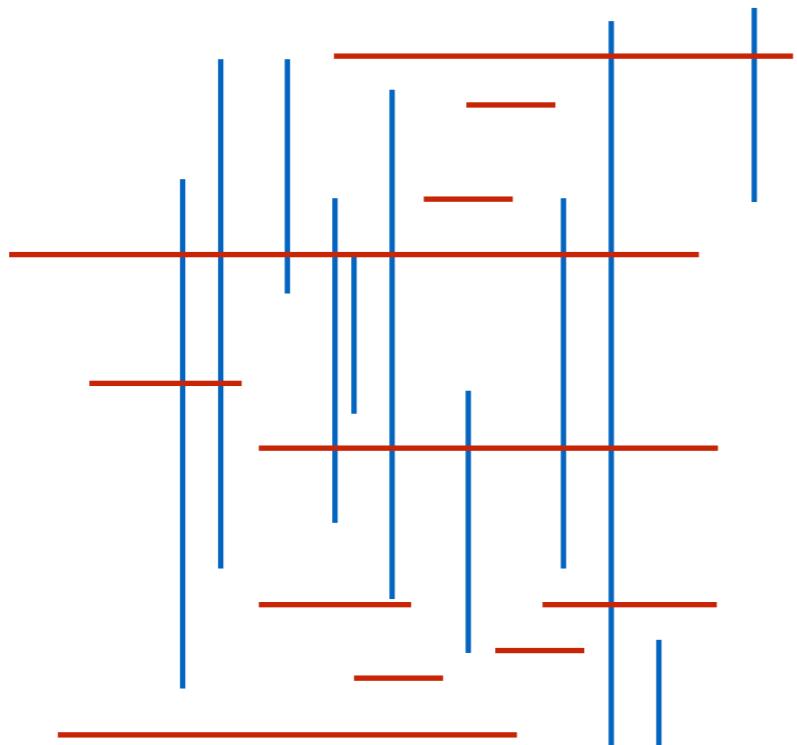


Balanced Binary Search Trees
- end -

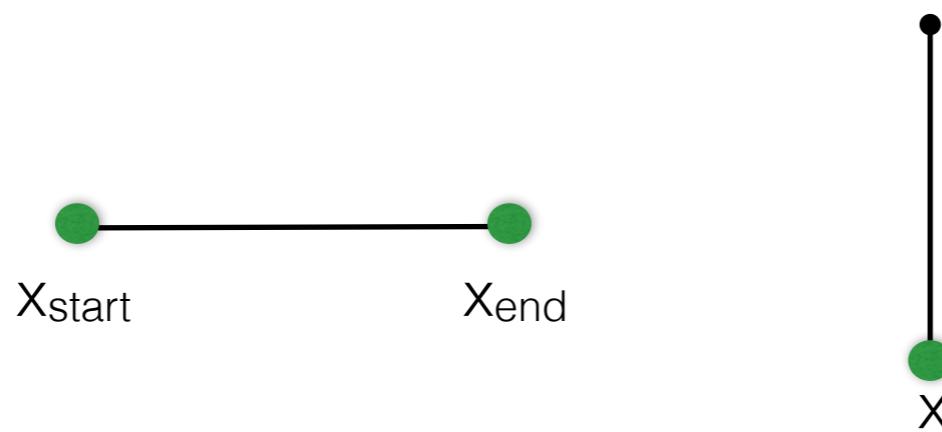
The line sweep technique



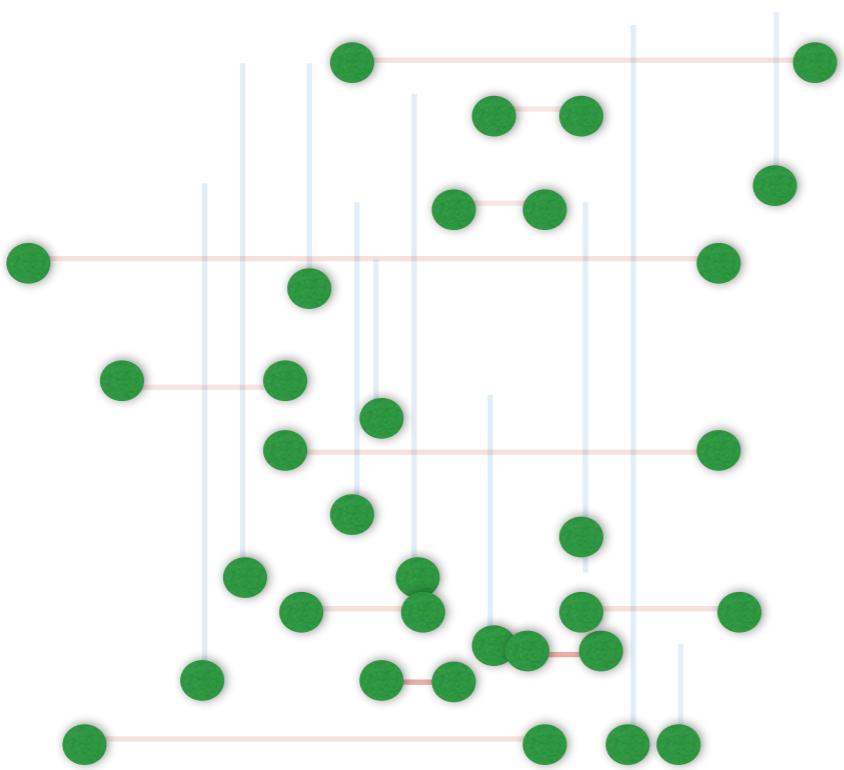
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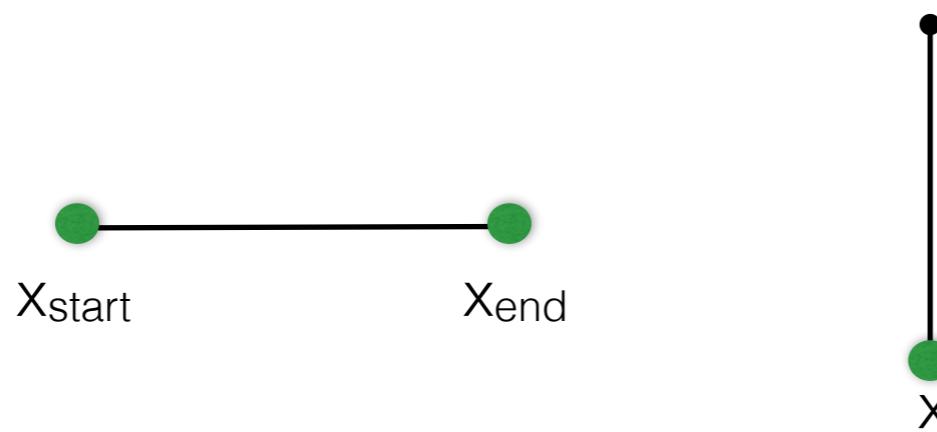
- Let X be the set of x -coordinates of all segments: these are the “events”



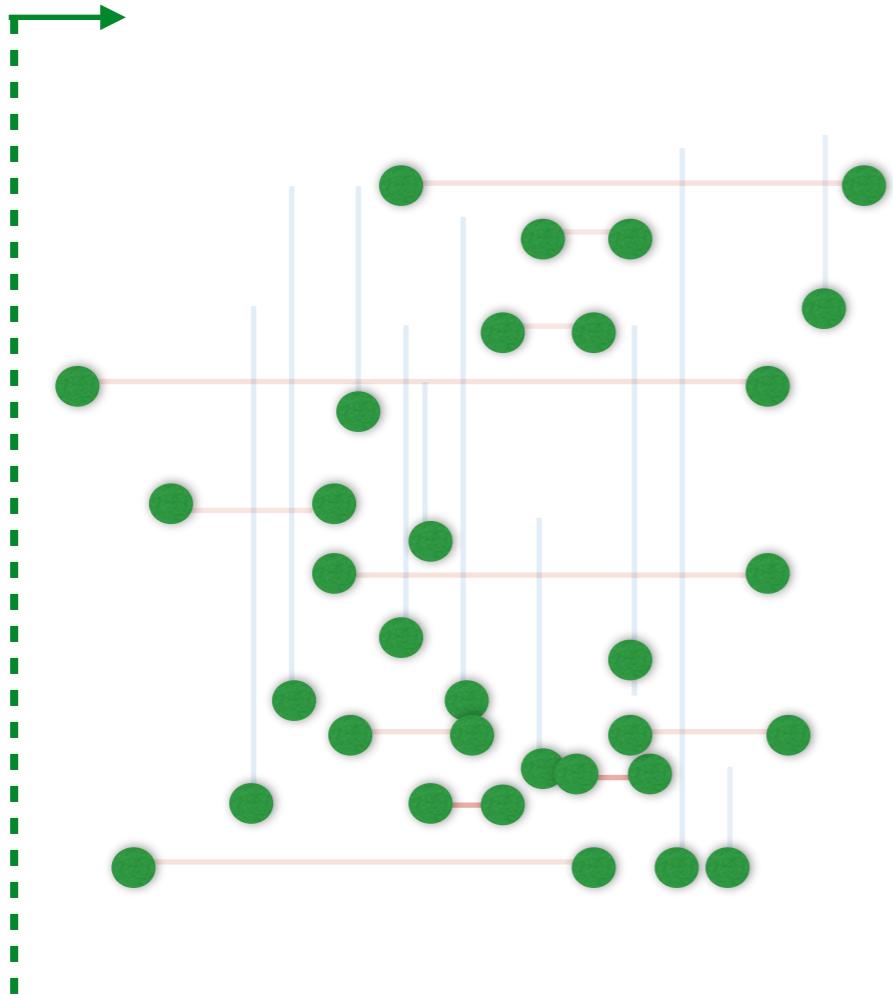
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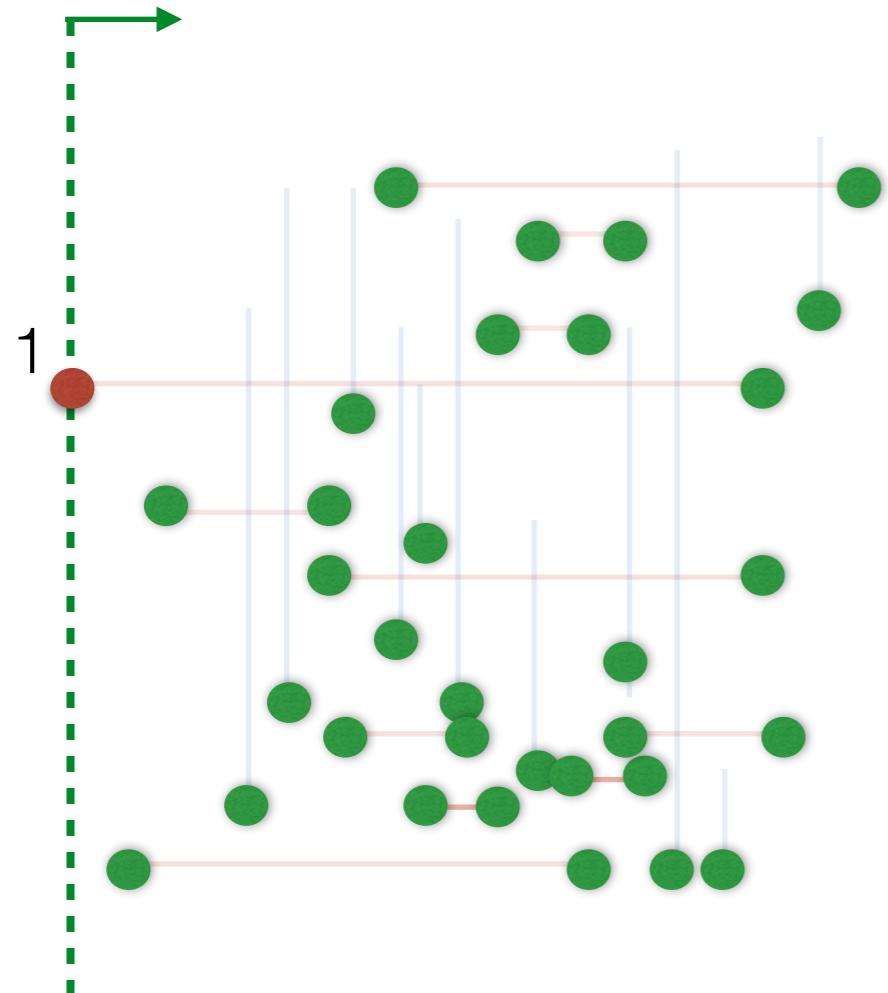


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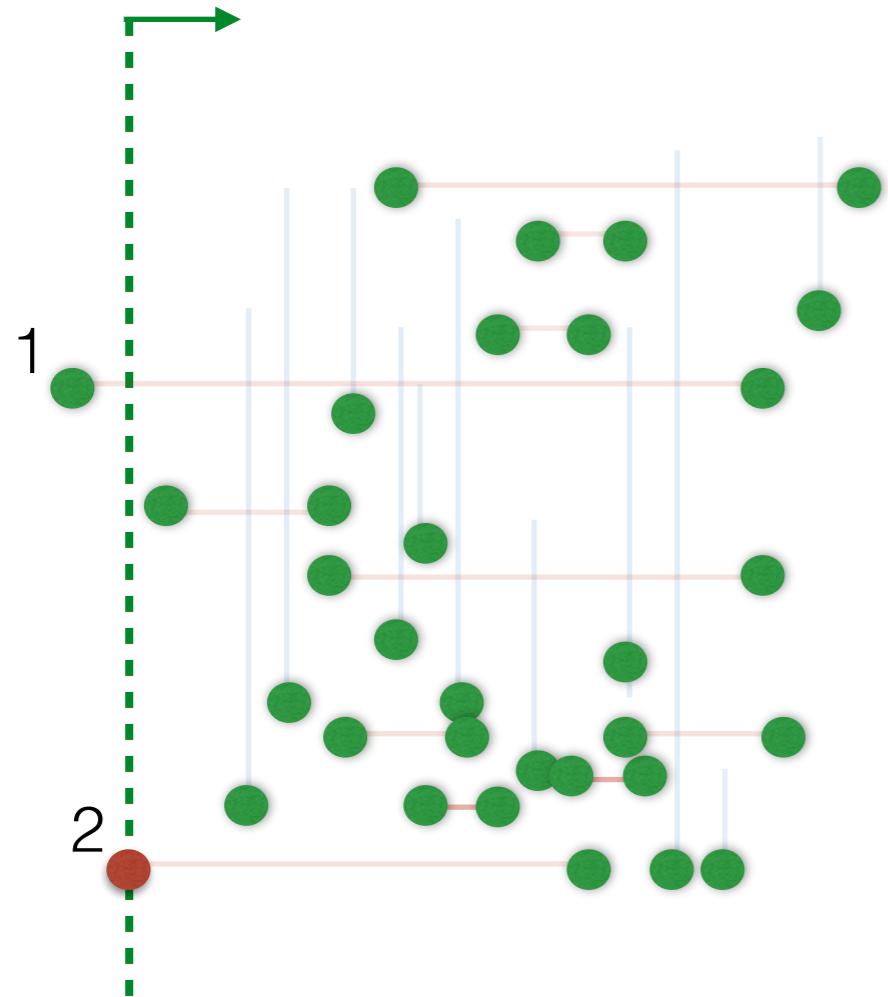
- Traverse the **events** in order

The line sweep technique



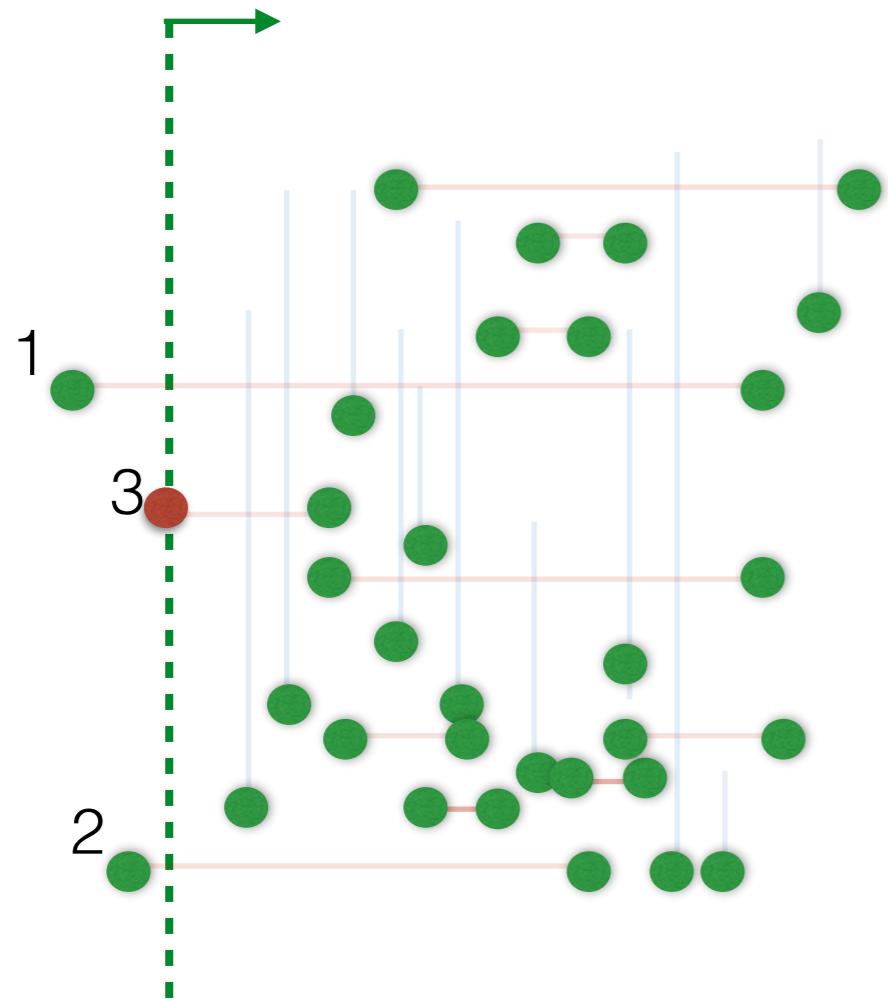
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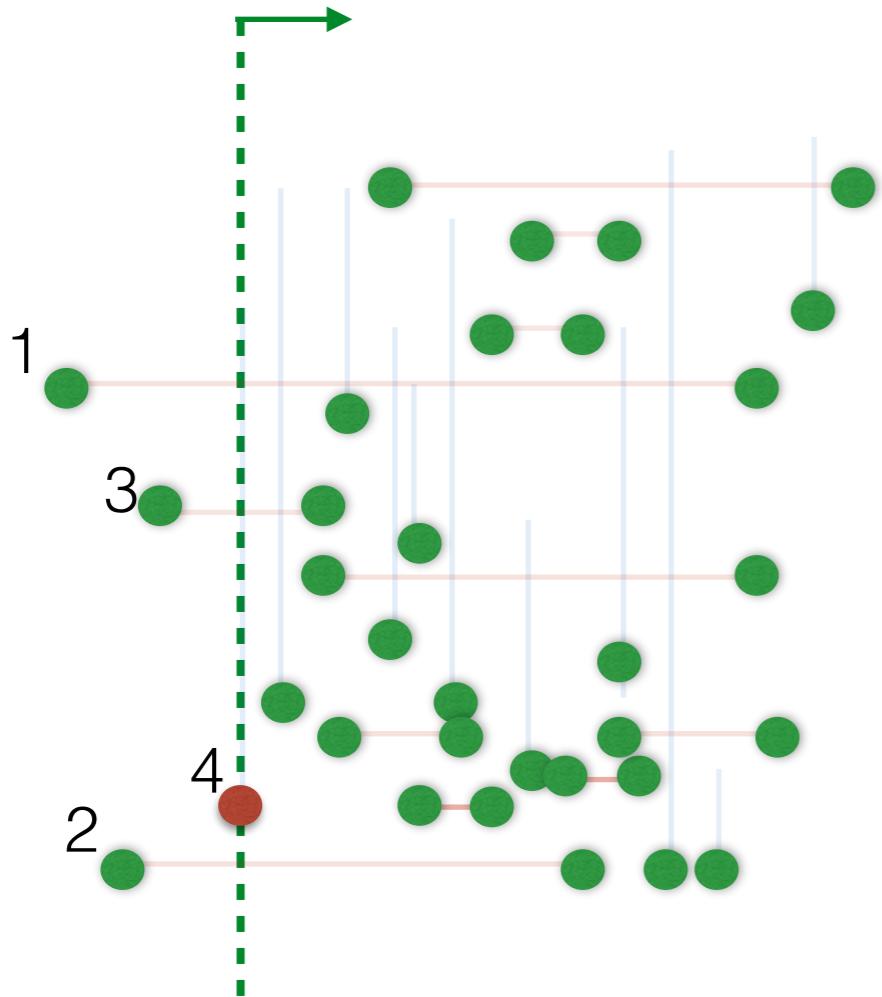
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The line sweep technique

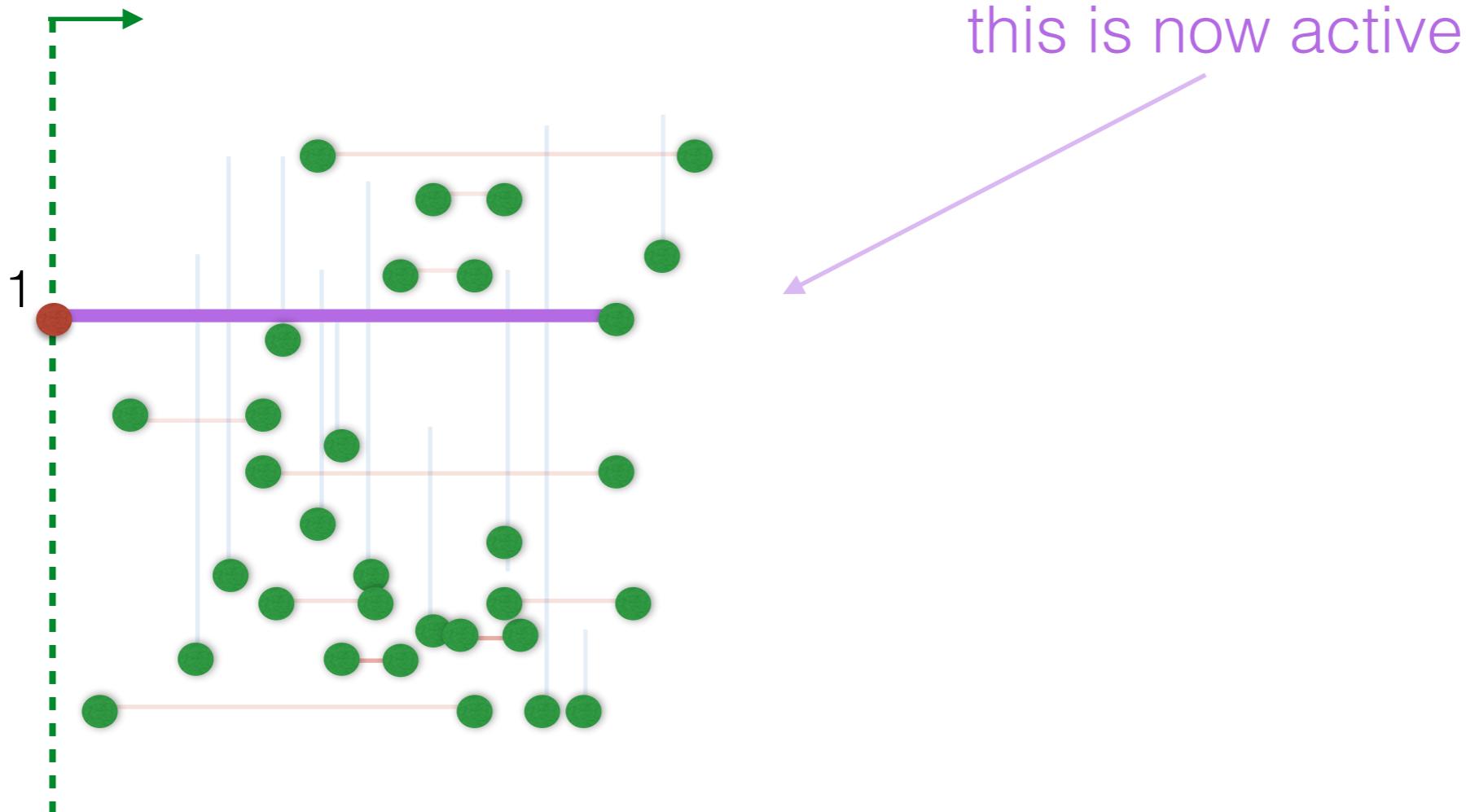


And so on.

How to process events?

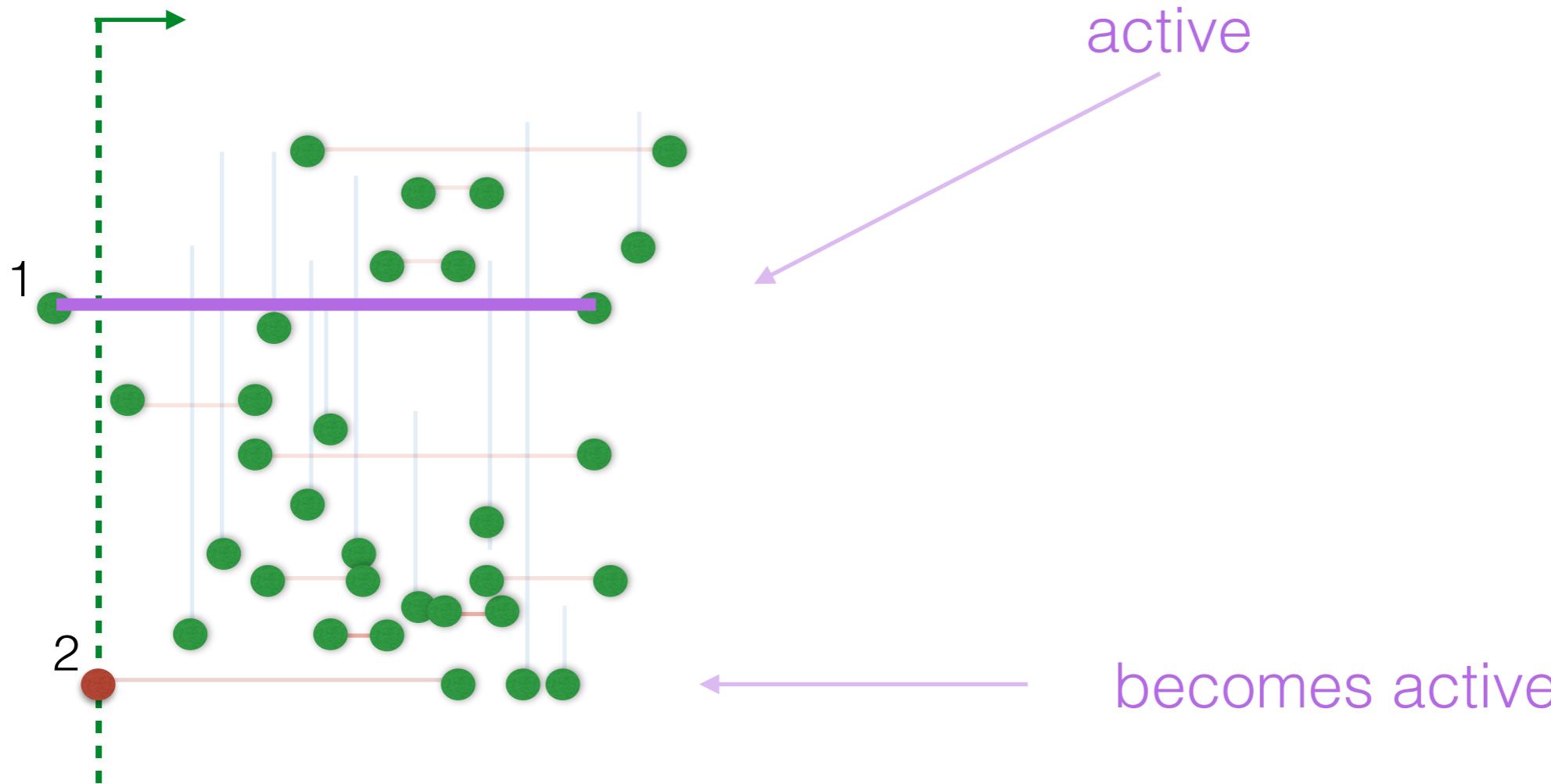
When reach a start event: segment becomes active

The line sweep technique



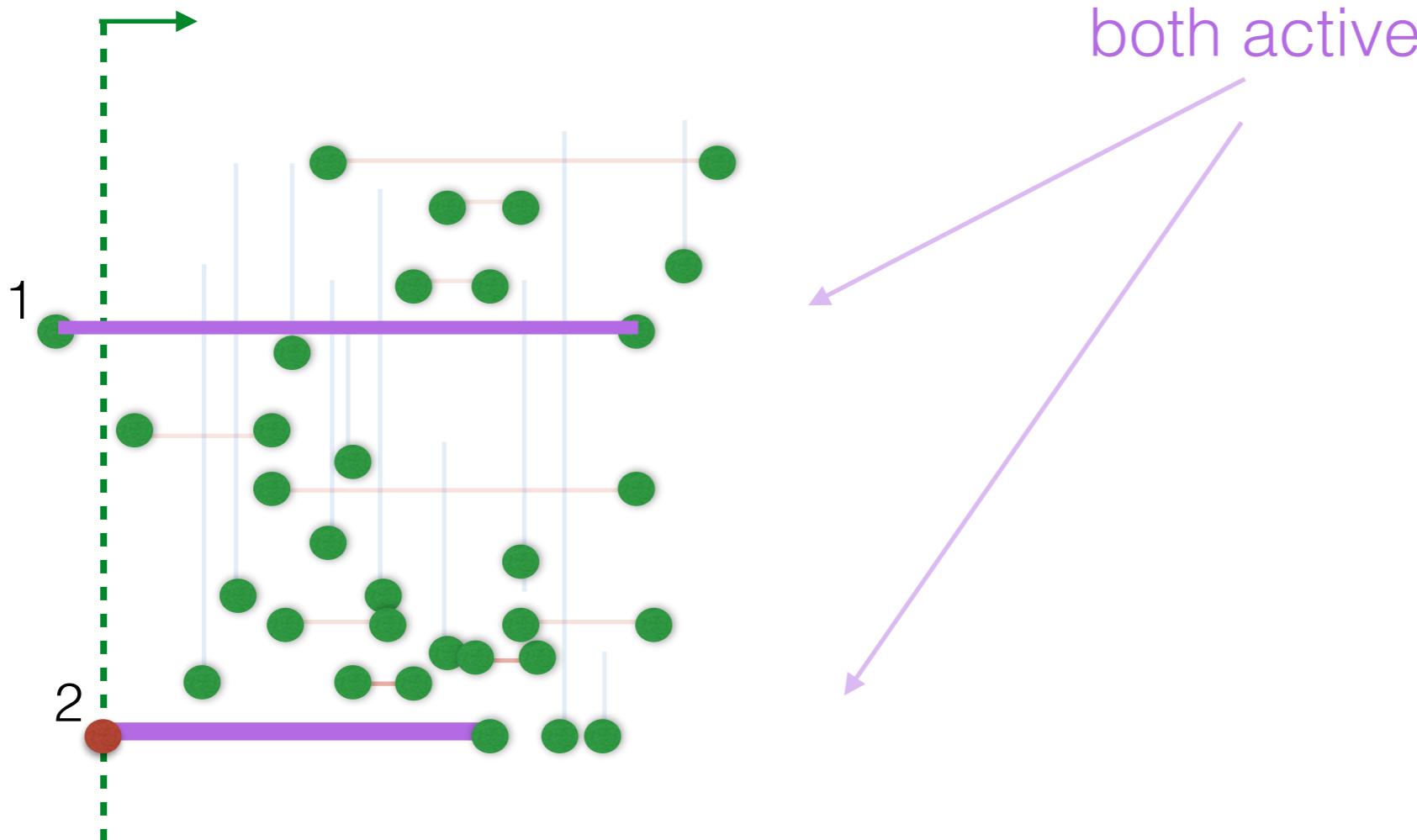
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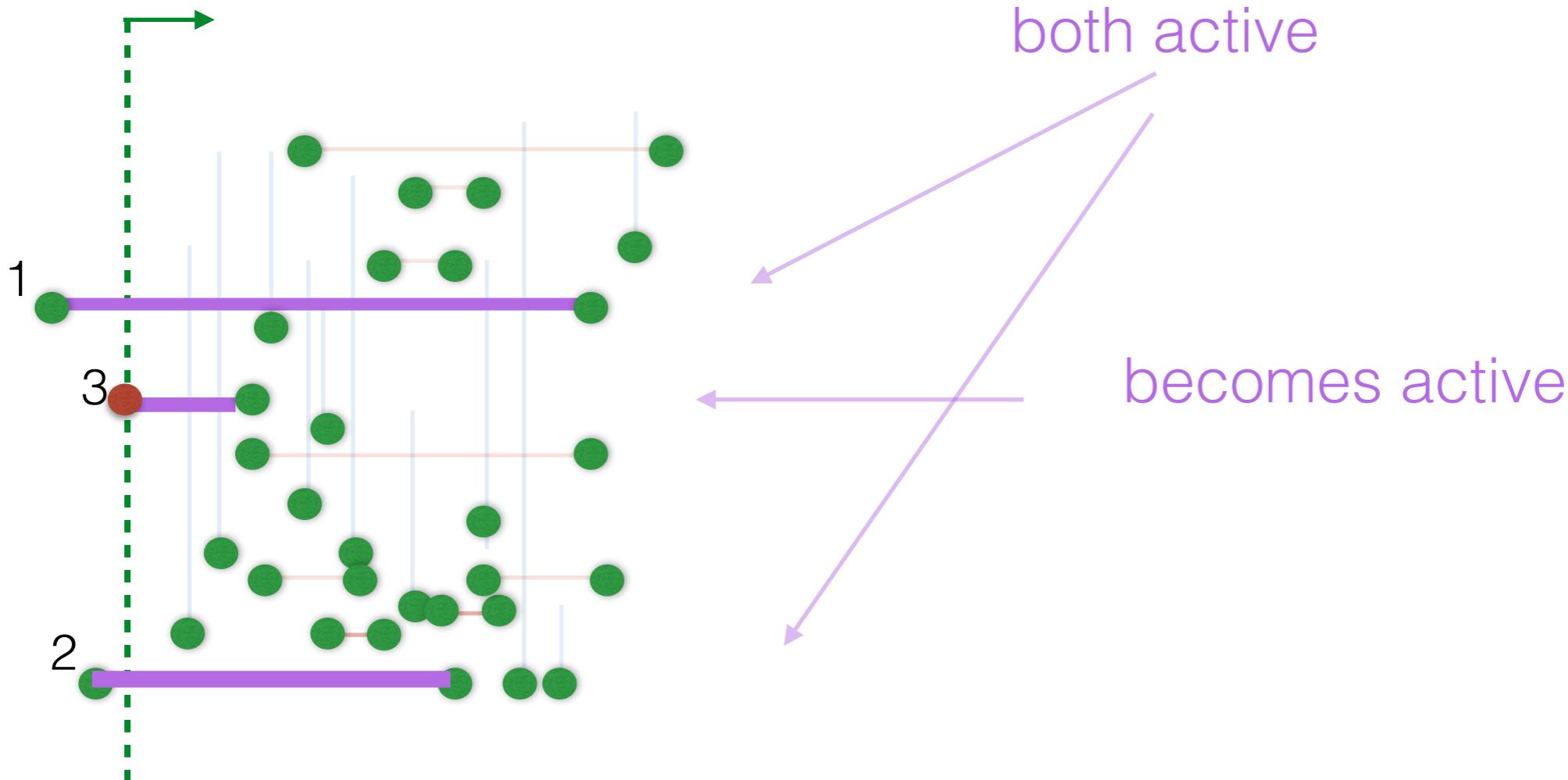
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The line sweep technique



When reach a start event: segment becomes active

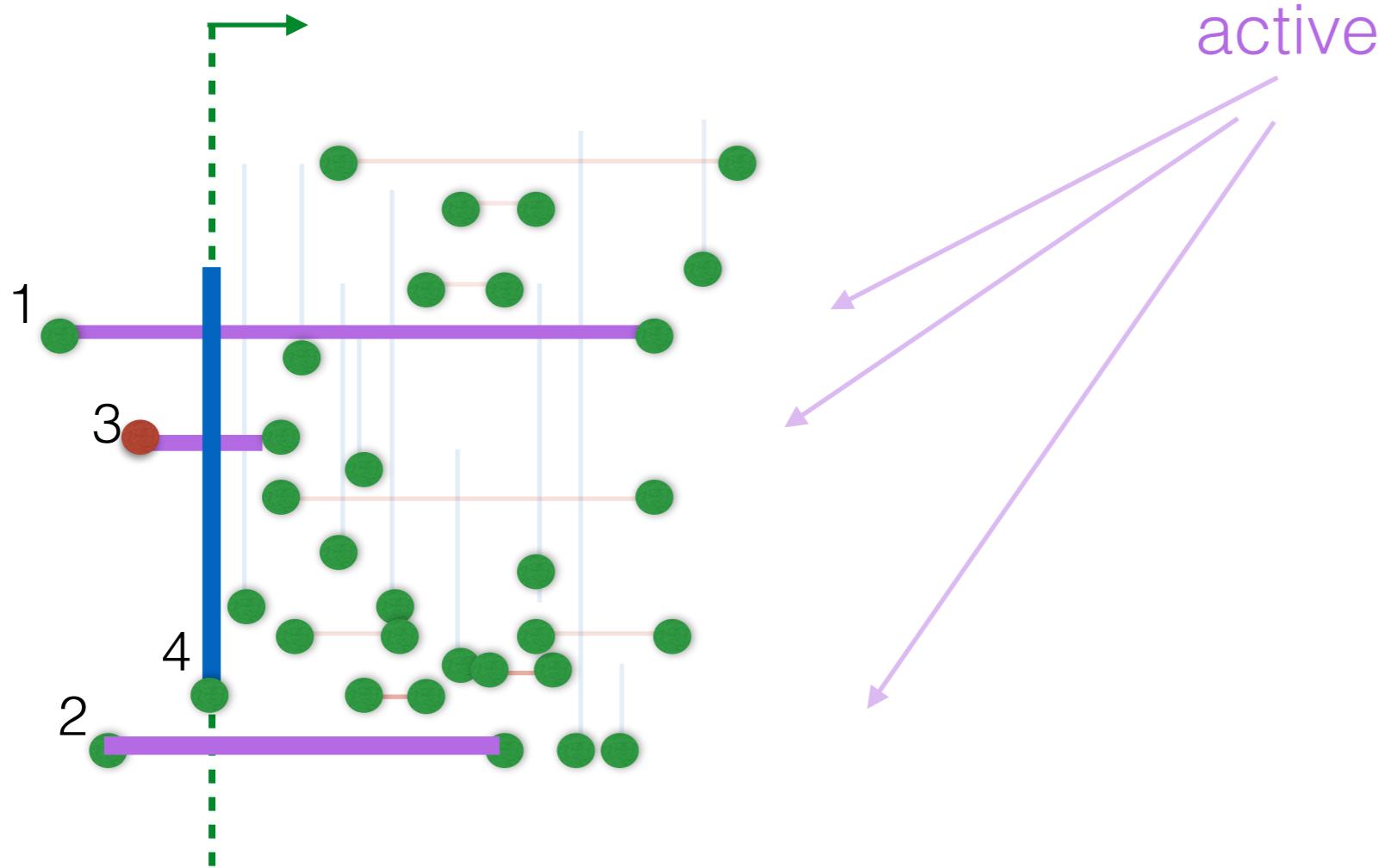
The line sweep technique



When reach a start event: segment becomes active

Let's see what happens when we reach an event
corresponding to a vertical segment

The line sweep technique

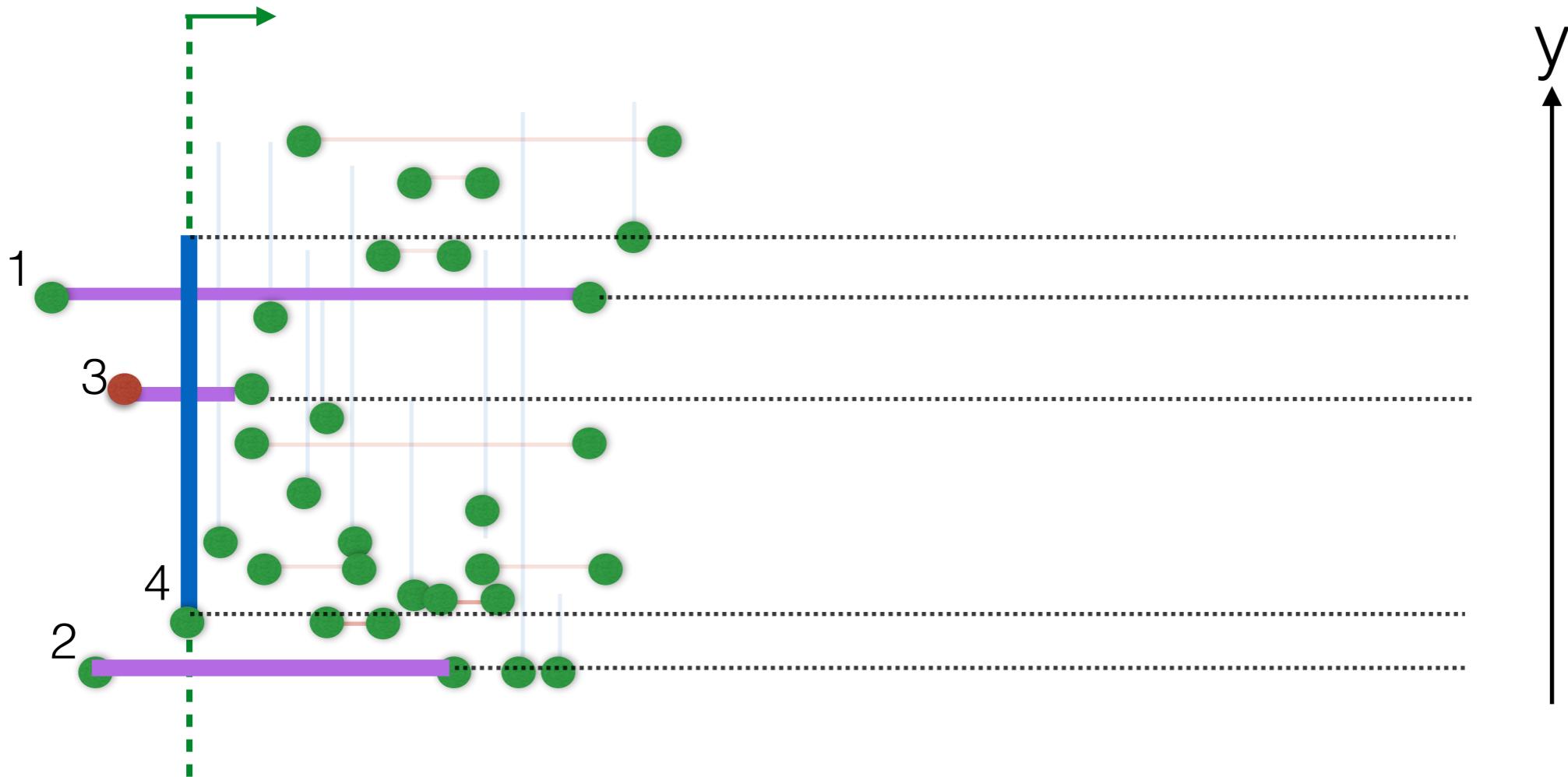


When reach an event corresponding to a vertical segment:

Claim: All horizontal segments that it intersects must be active

But, not all active segments intersect the vertical segment

The line sweep technique

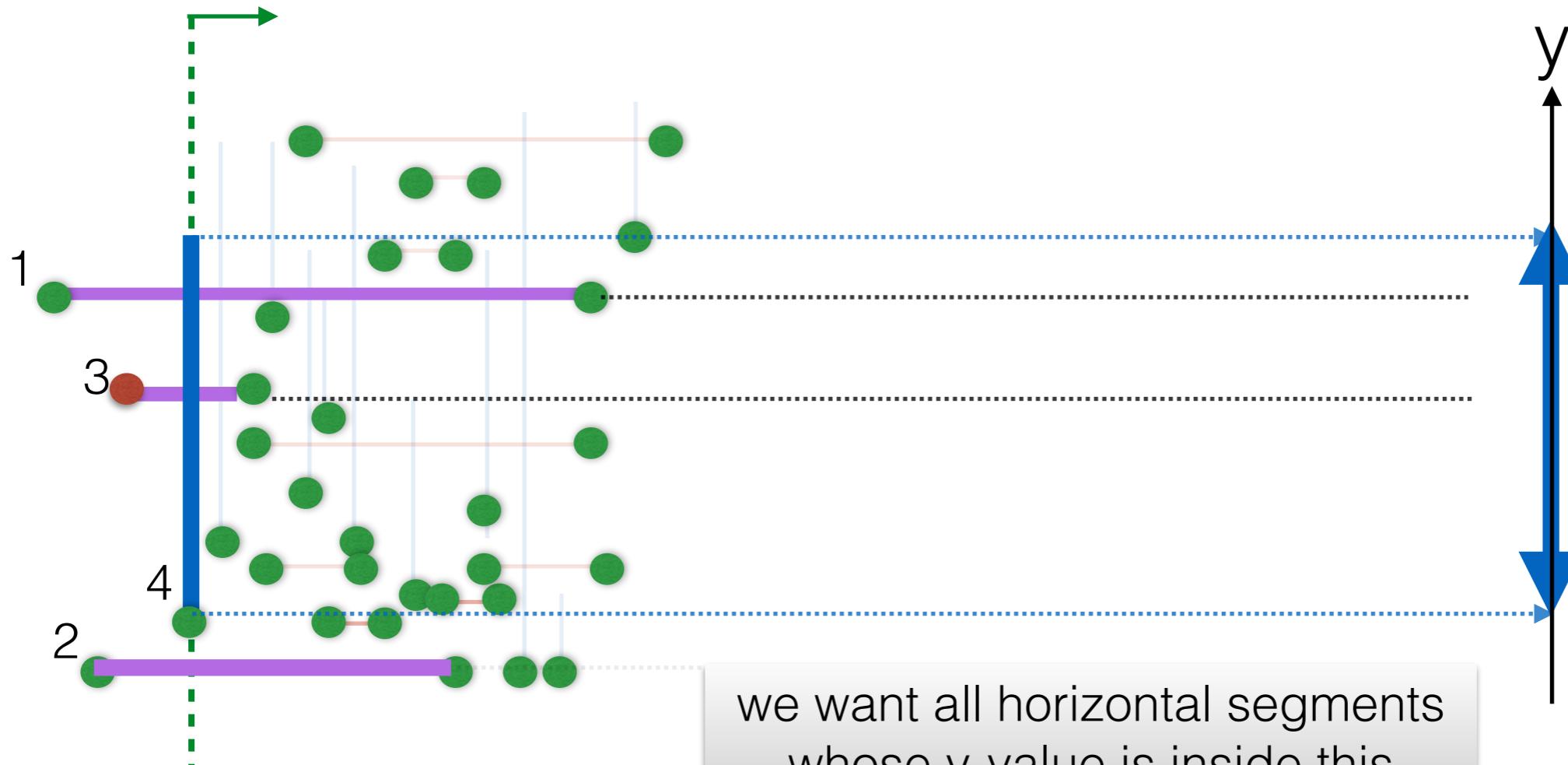


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The line sweep technique

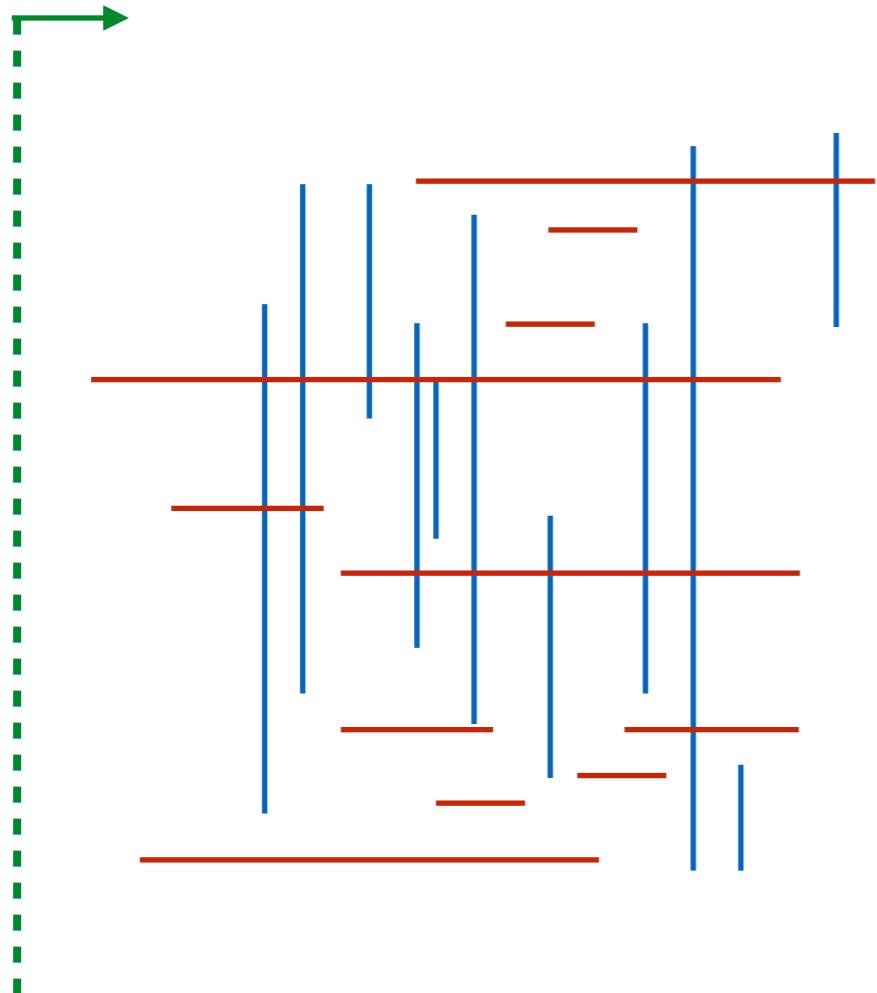


When reach an event corresponding to a vertical segment:

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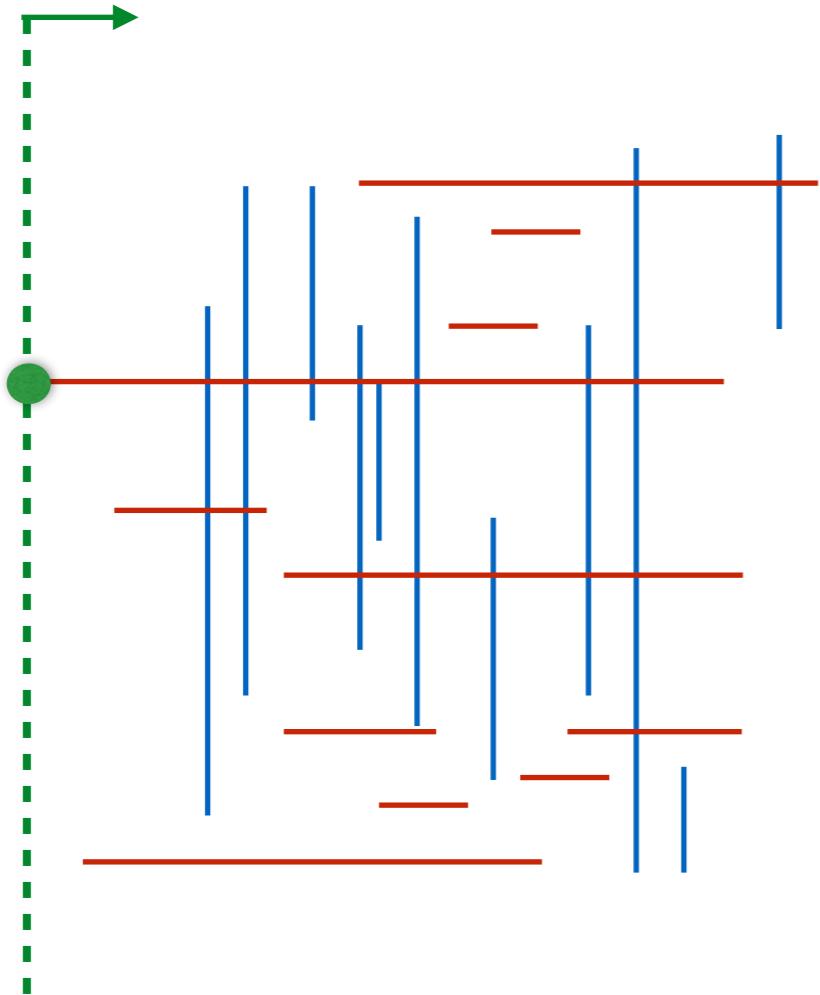
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Orthogonal line segment intersection



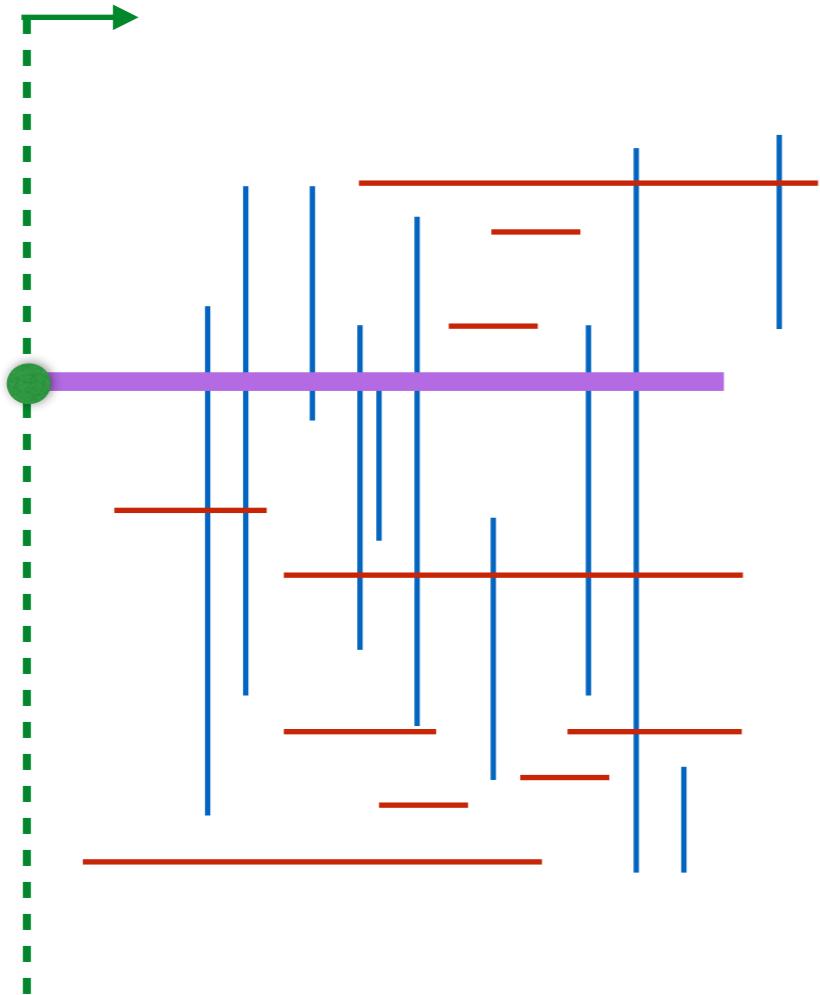
- Let X be the set of x-coordinates of all segments //the events
- Initialize $AS = \{\}$
- Sort X and traverse the events in sorted order; let x be the next event in X
 - if x is start of horizontal segment (x, x', y) :
//segment becomes active
insert segment (x, x', y) in AS
 - if x is end of horizontal segment (x, x', y) :
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delete segment (x, x', y) from AS
 - if x corresponds to a vertical segment (y, y', x) :
//All active segments start before x and end after x . We need those whose y -value is in $[y, y']$
search AS for all segments with y -value in given range $[y, y']$ and report intersections

Orthogonal line segment intersection



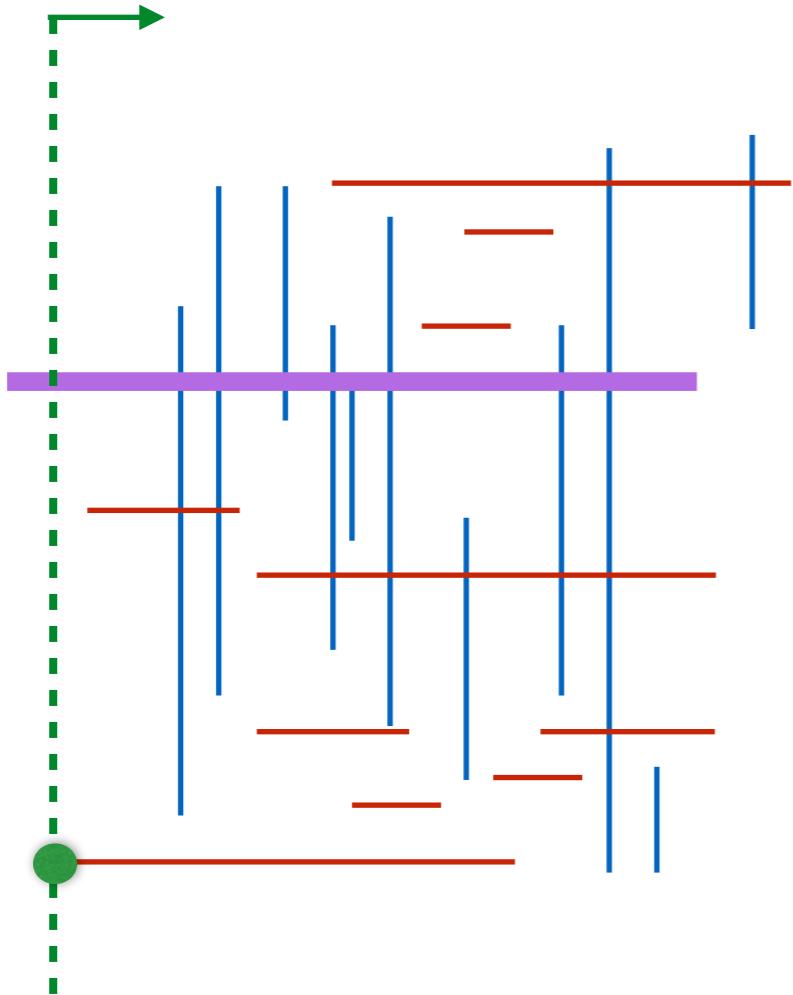
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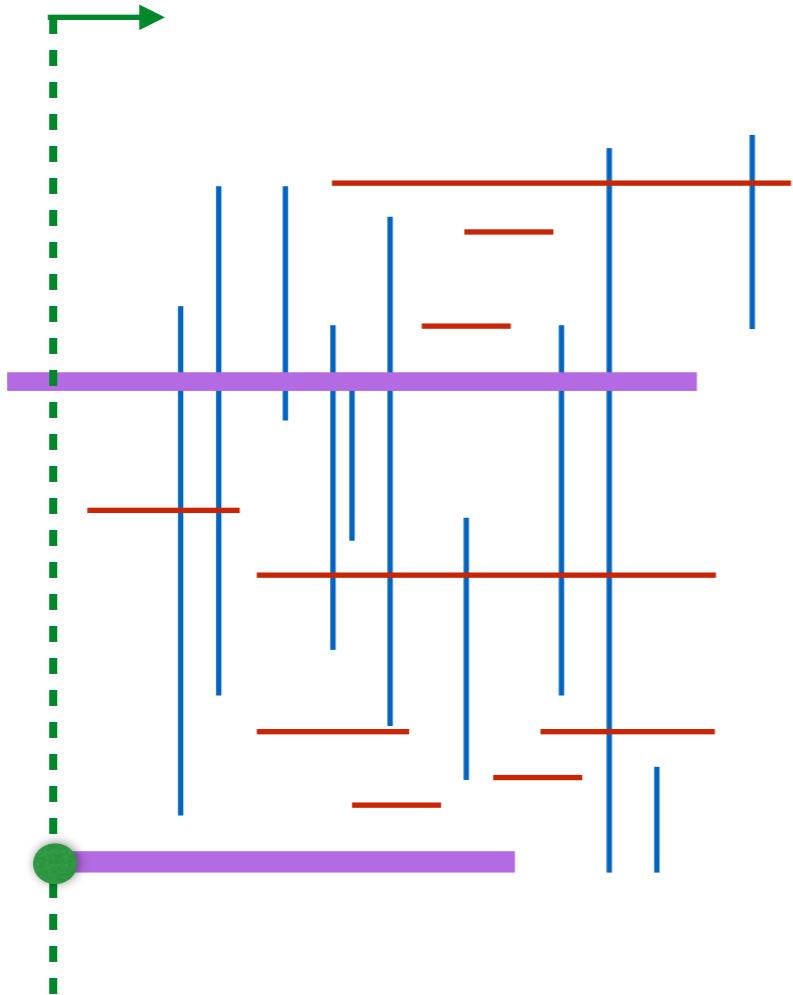
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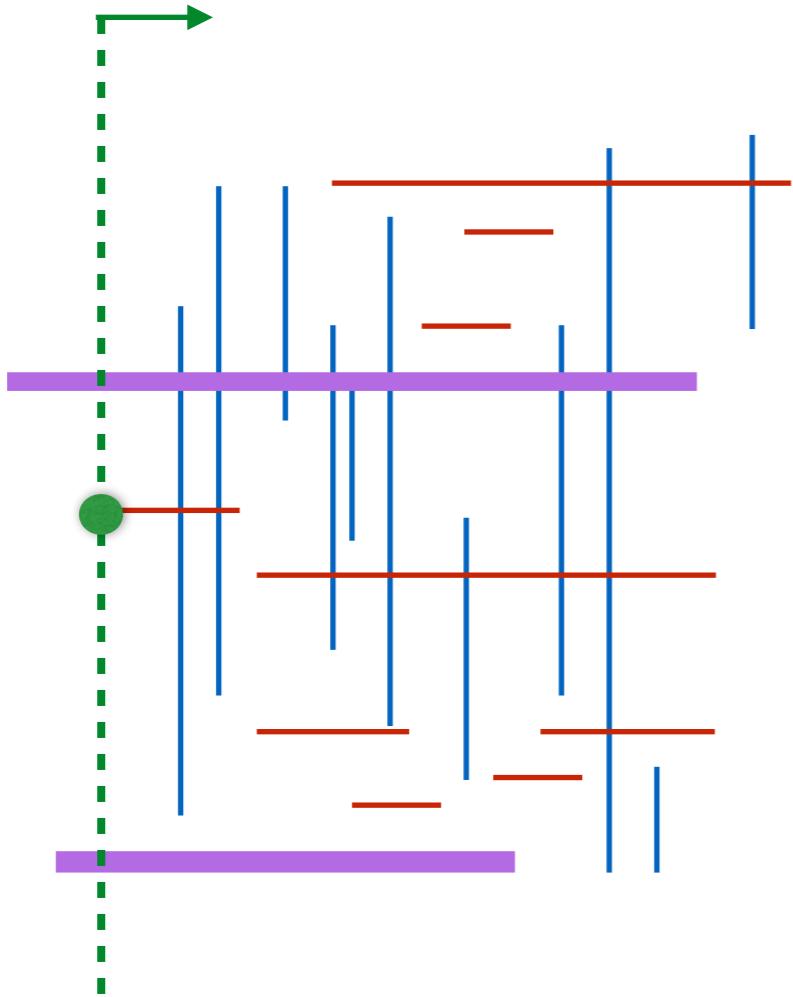
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Orthogonal line segment intersection



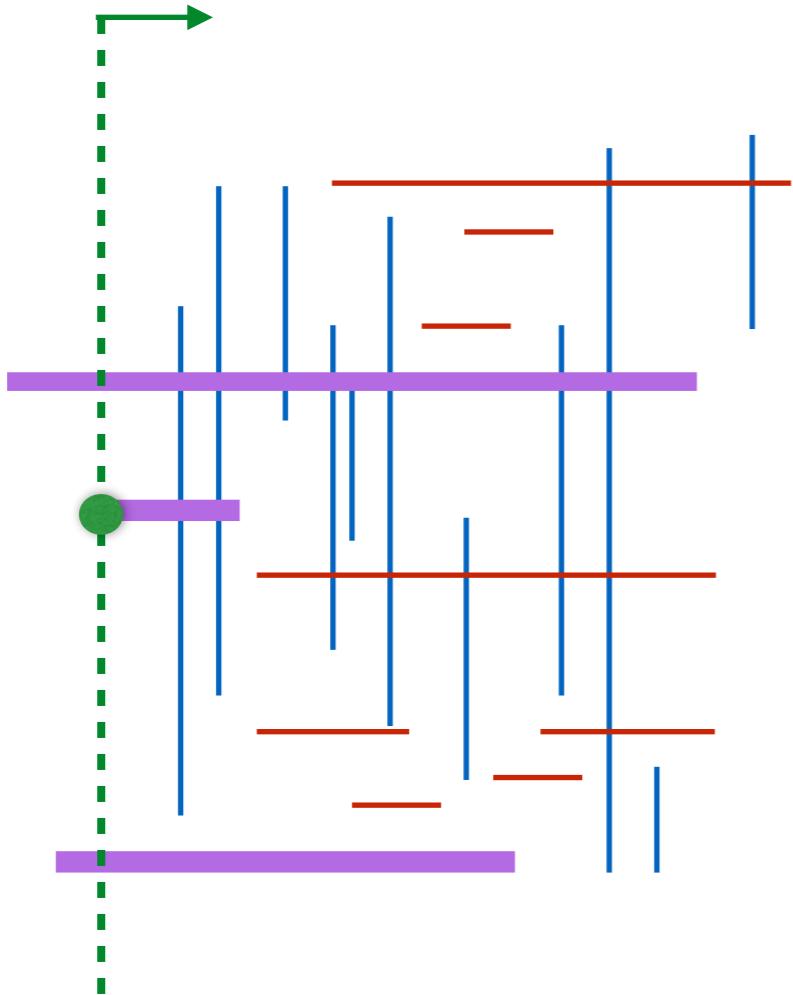
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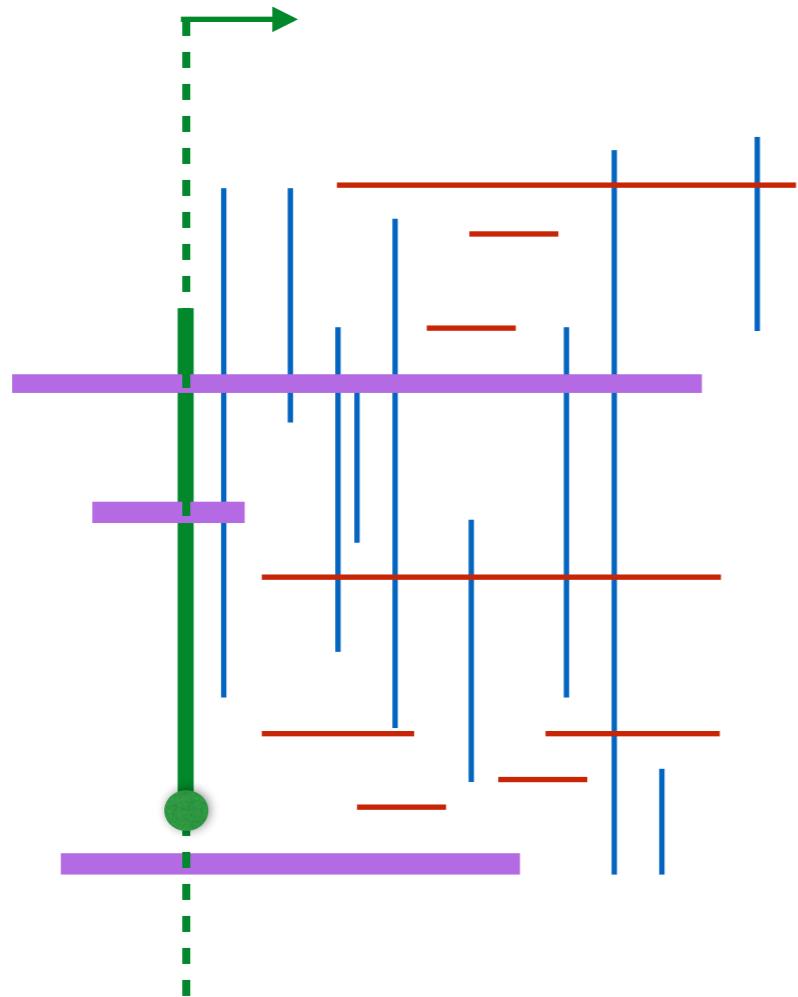
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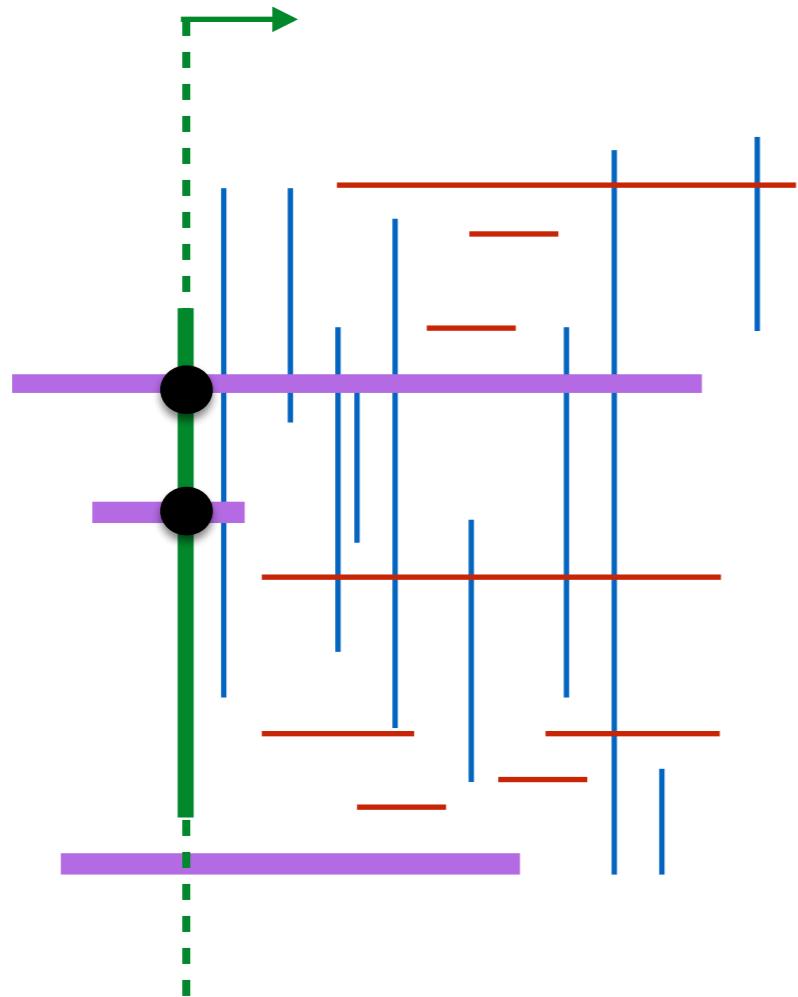
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Orthogonal line segment intersection



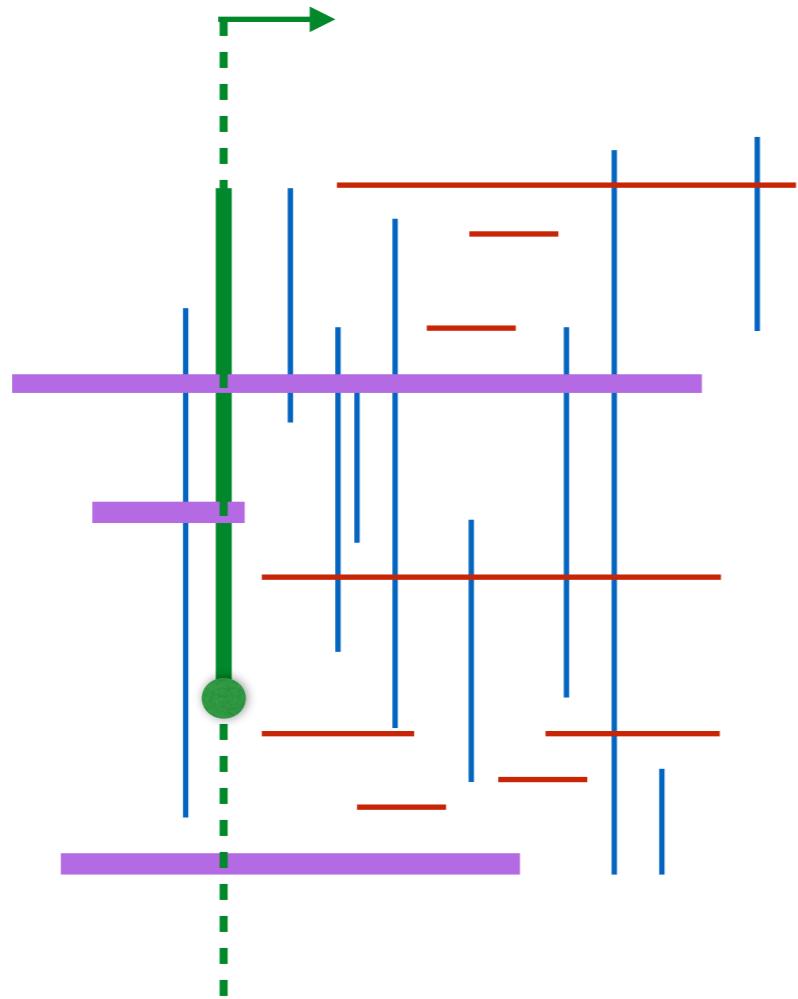
- Let X be the set of x -coordinates of all segments //the events
- Initialize $AS = \{\}$
- Sort X and traverse the events in sorted order; let x be the next event in X
 - if x is start of horizontal segment (x, x', y) :
//segment becomes active
insert segment (x, x', y) in AS
 - if x is end of horizontal segment (x, x', y) :
//segment stops being active
delete segment (x, x', y) from AS
 - if x corresponds to a vertical segment (y, y', x) :
//All active segments start before x and end after x . We need those whose y is in $[y, y']$
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Orthogonal line segment intersection



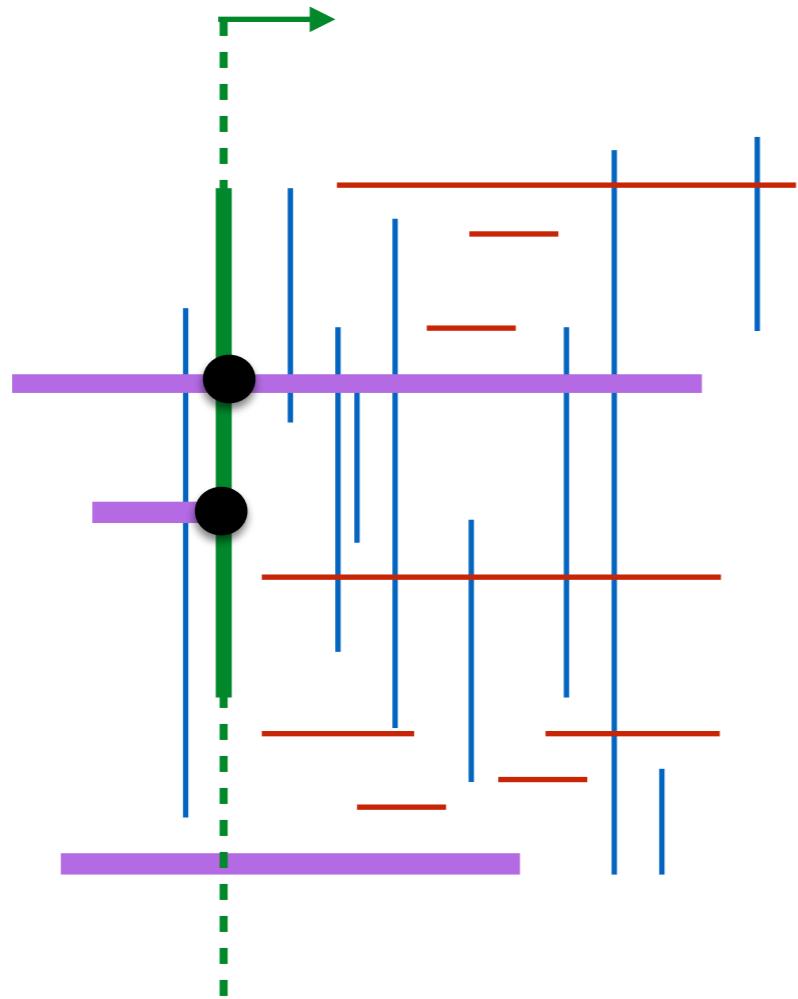
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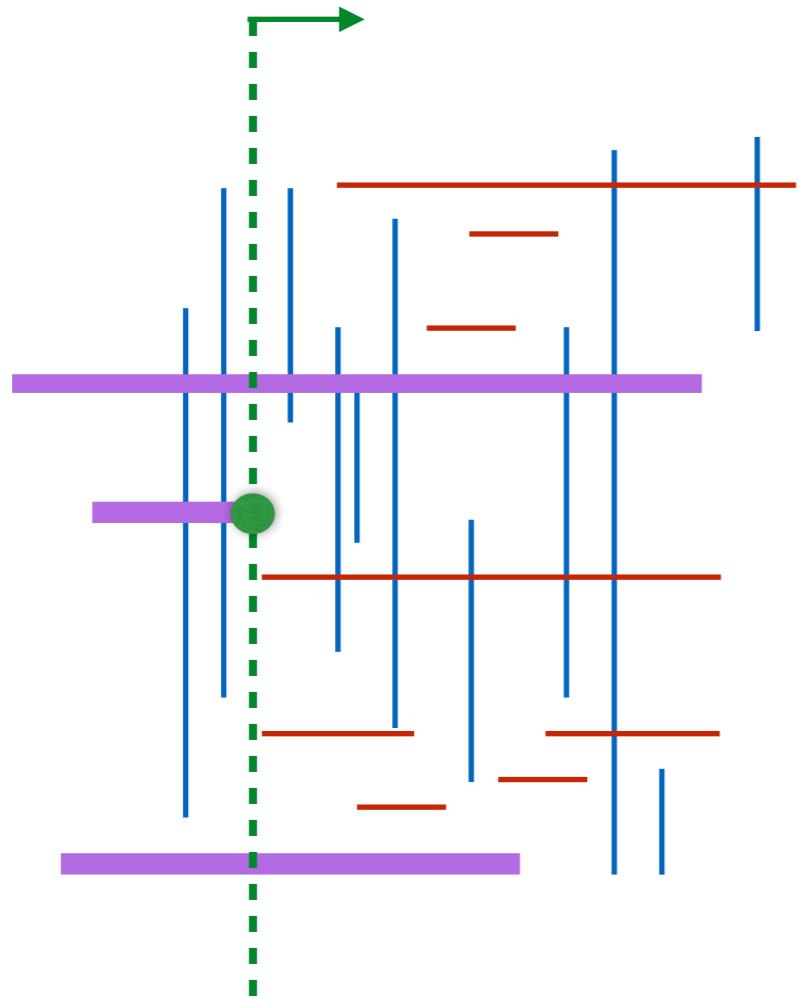
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Orthogonal line segment intersection



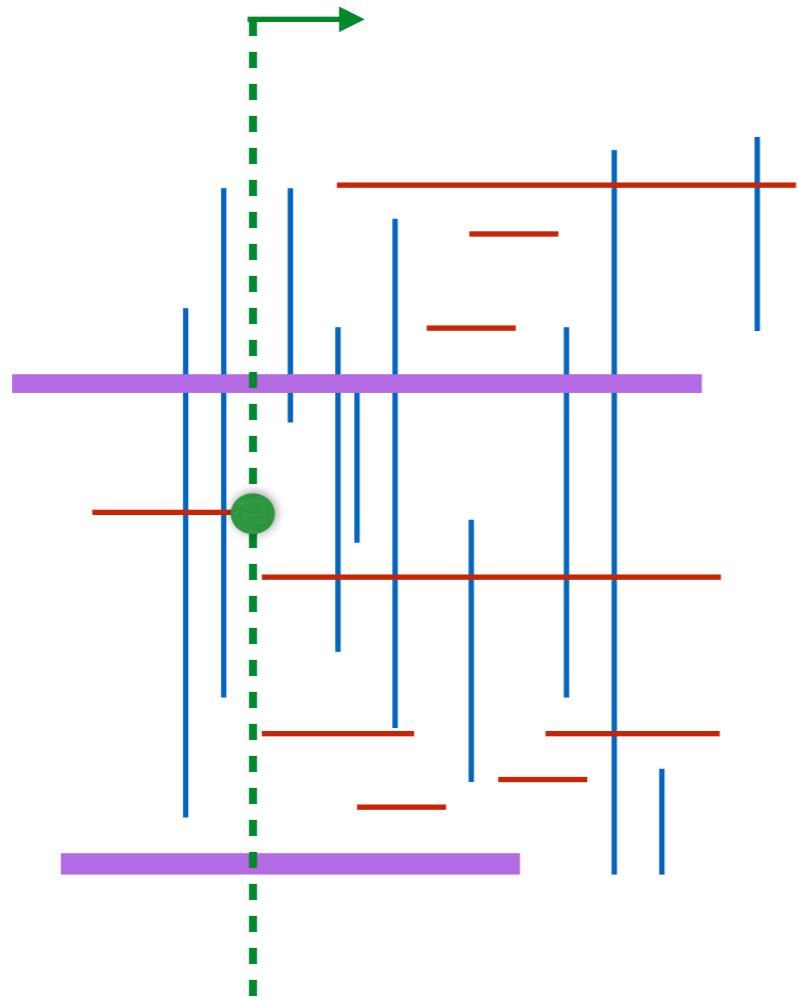
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Orthogonal line segment intersection



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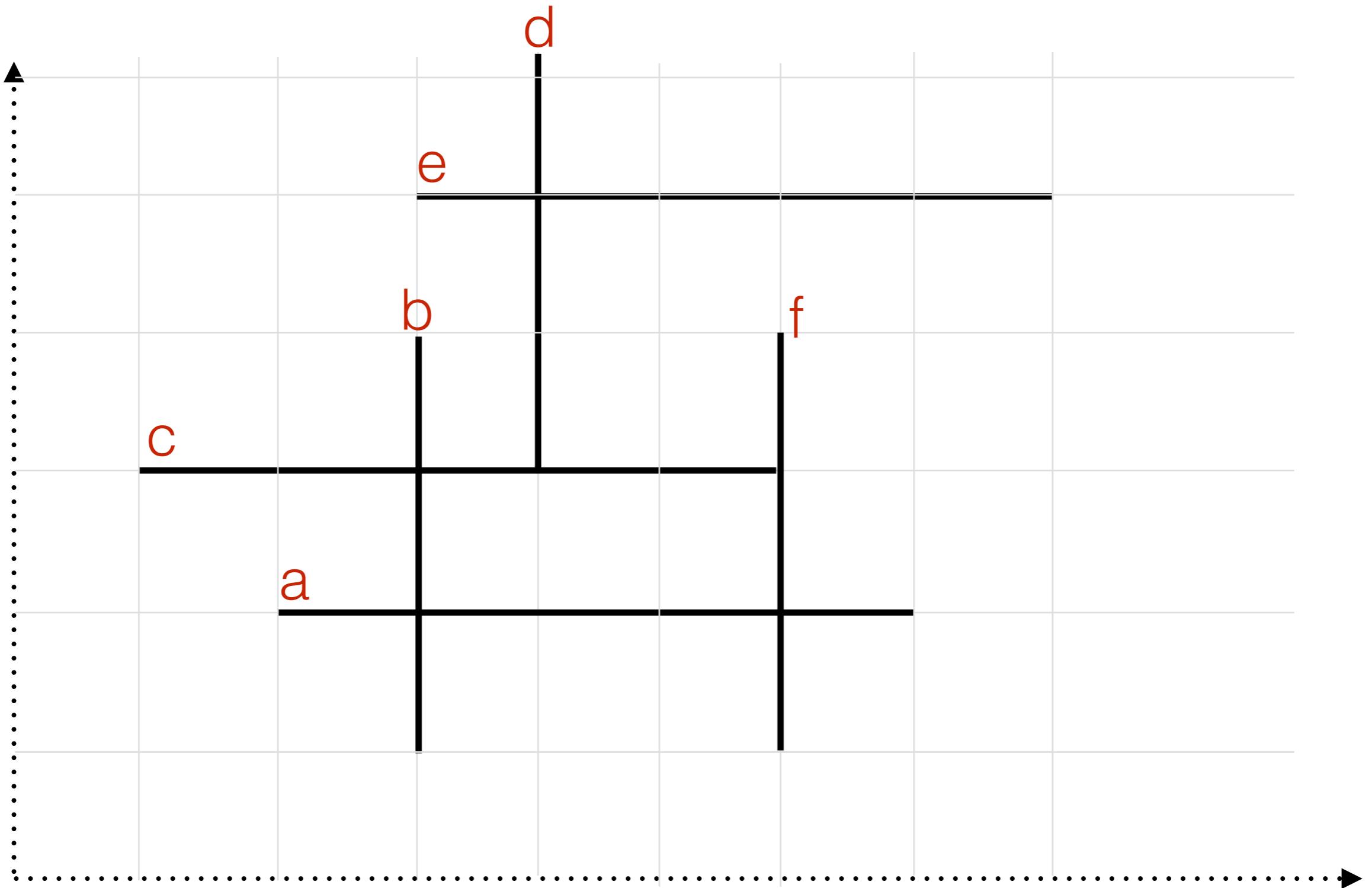
Orthogonal line segment intersection



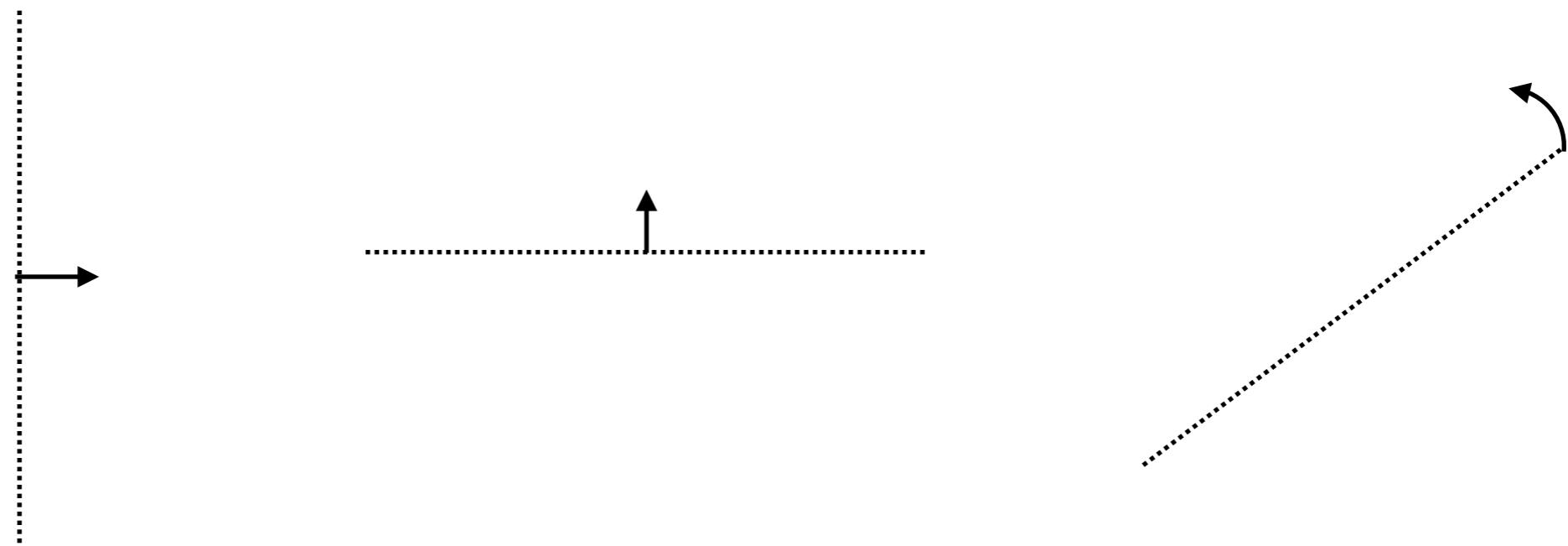
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Orthogonal line segment intersection

- To think
 - How to implement the AS?
 - Analysis
- Let X be the set of x-coordinates of all segments //the events
- Initialize AS = {}
- Sort X and traverse the events in sorted order; let x be the next event in X
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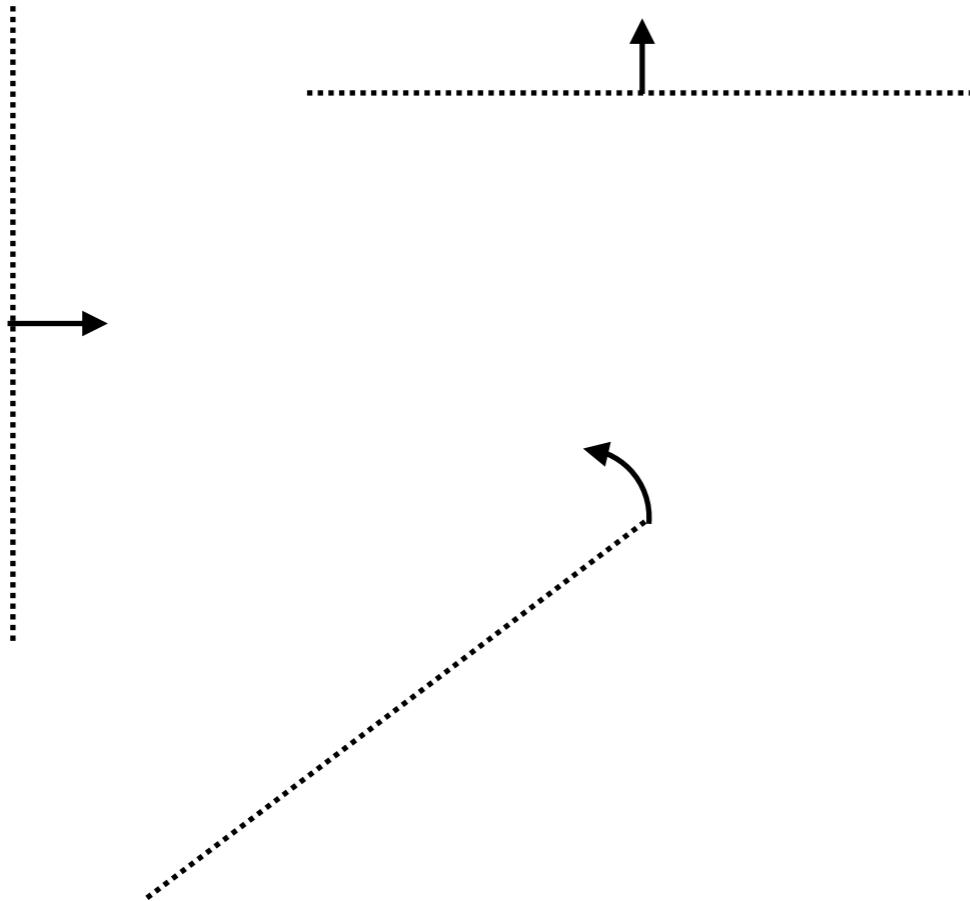


Line sweep



Line sweep algorithms

- Powerful, elegant, frequently used technique
- Line can be horizontal or vertical or radial or



- Traverse events in order and maintain an Active Structure (AS)
 - AS contains objects that are “active” (started but not ended) in other words they are intersected by the current sweep line
 - at some events, insert in AS
 - at some events, delete from AS
 - at some events, query AS