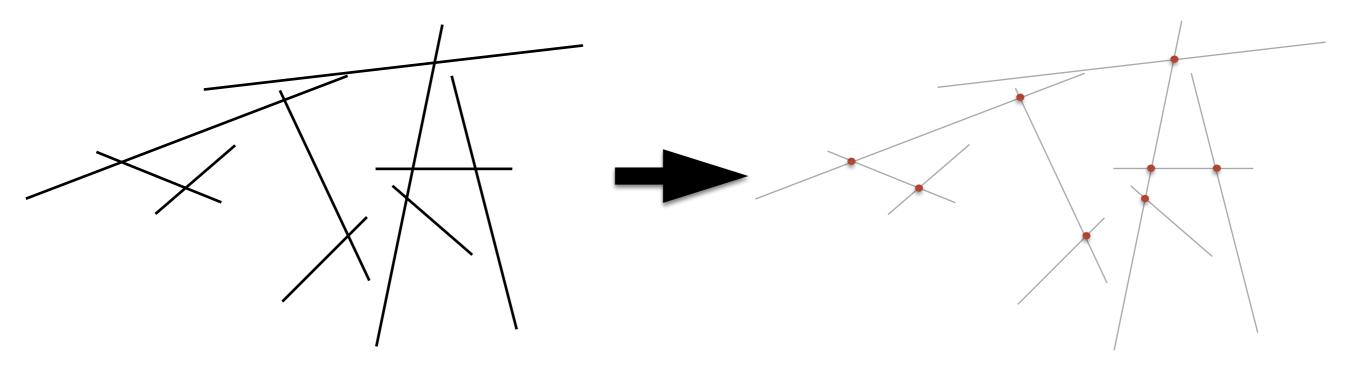
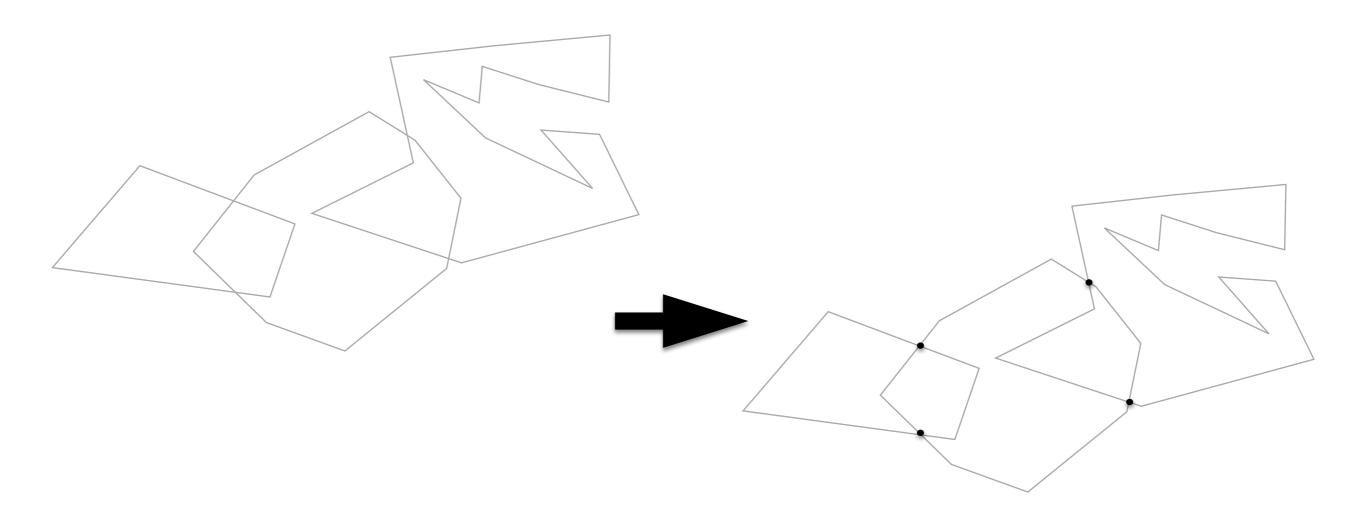


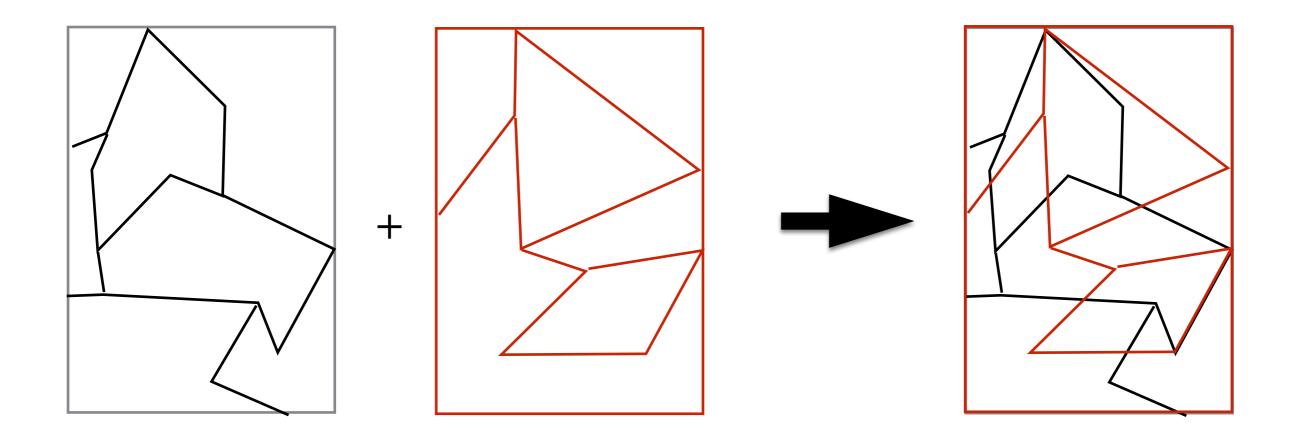
Problem: Given a set of line segments in 2D, find all their pairwise intersections.



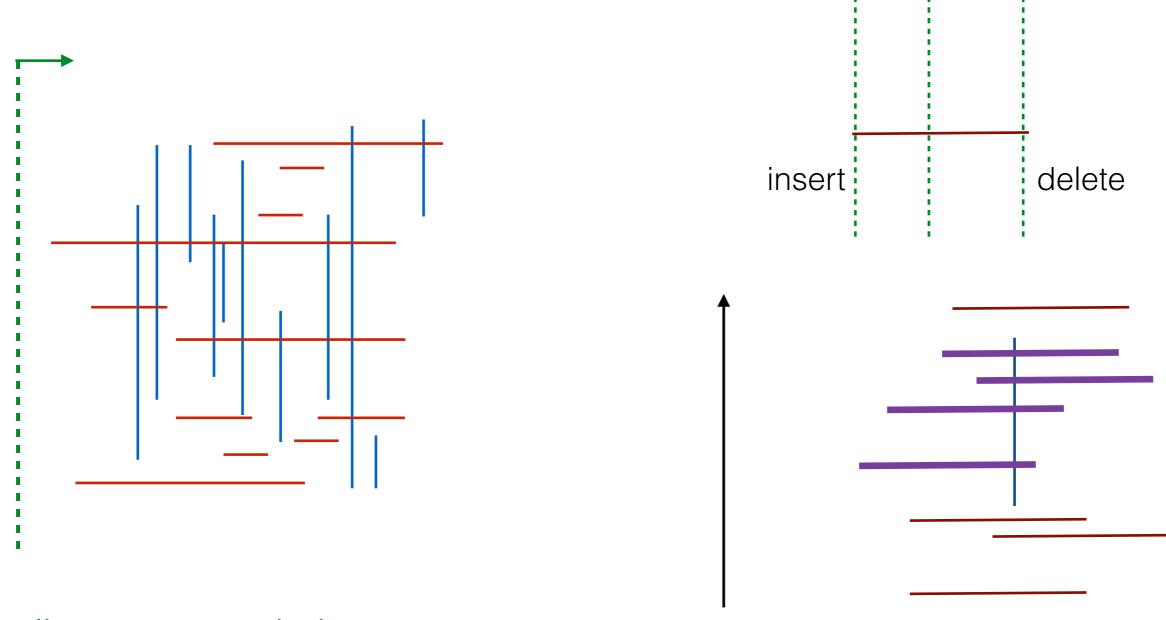
Problem: Given a set of line segments in 2D, find all their pairwise intersections.



Problem: Given a set of line segments in 2D, find all their pairwise intersections.



Last time we looked at a special case: Orthogonal line segment intersection



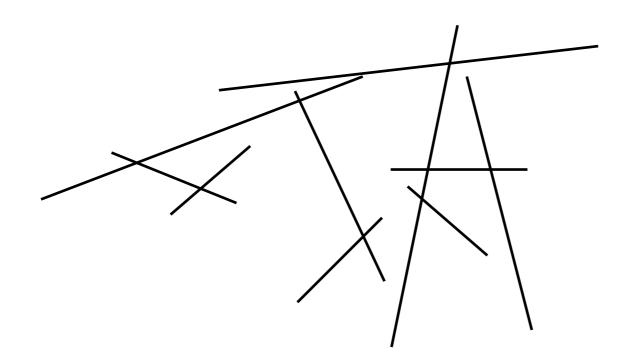
line sweep technique solve the problem behind the line

AS: segments in order of y

Result: The intersections of a set of n orthogonal segments in the plane can be found in $O(n \lg n + k)$ time.

General line segment intersection

- n: size of the input (number of segments)
- k: size of output (number of intersections)

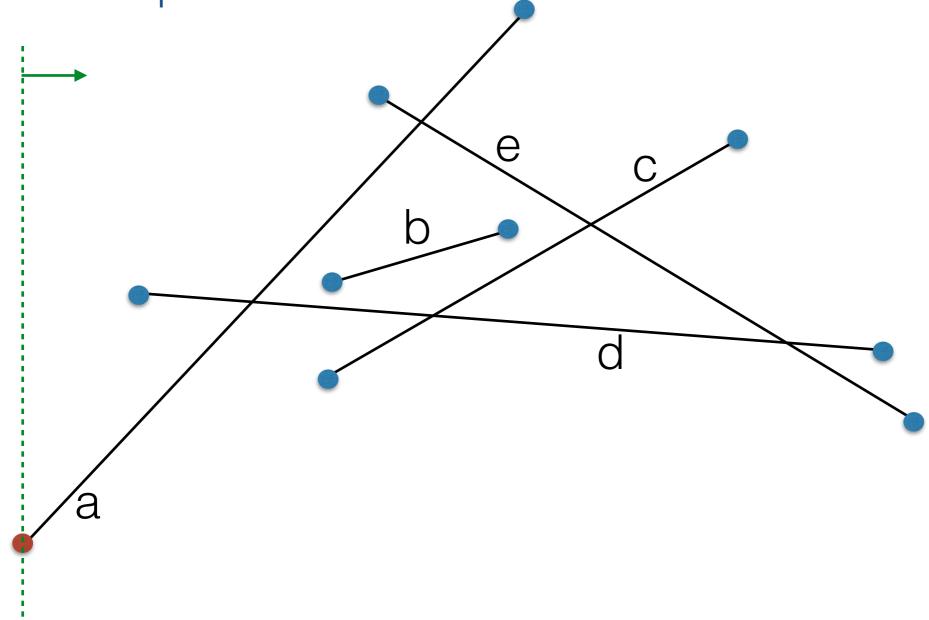


- Extend sweep line idea
- We'll get an overall bound of O(n lg n + k lg n) which improves on the naive O(n²) when k is small
- The algorithm was developed by Jon Bentley and Thomas Ottman in 1979
- Simple (in retrospect), elegant and practical

• Let X be the set of all x-coords of segments

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- Let X be the set of all x-coords of segments
- Traverse the events in X in order



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The sweep e c

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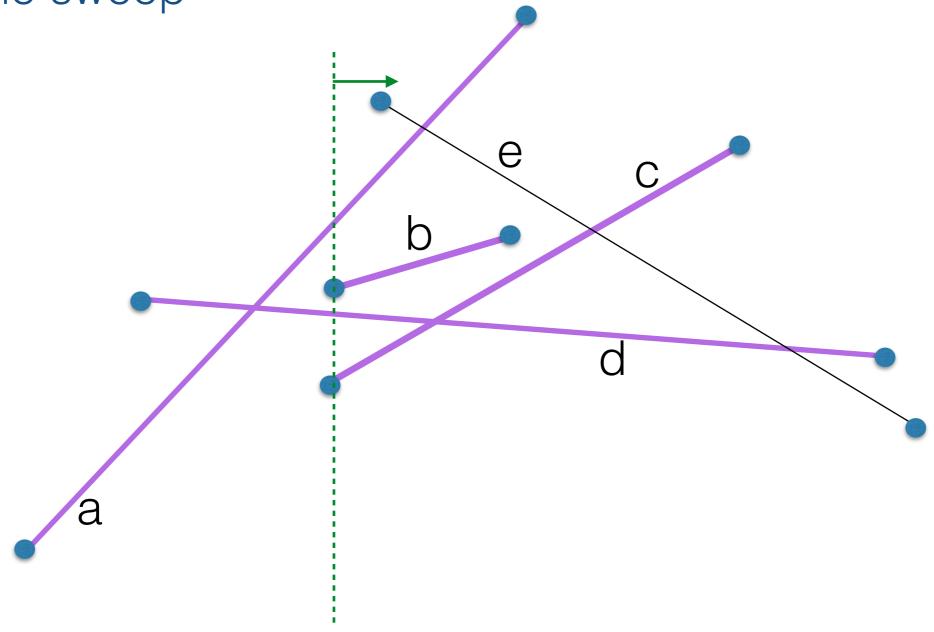
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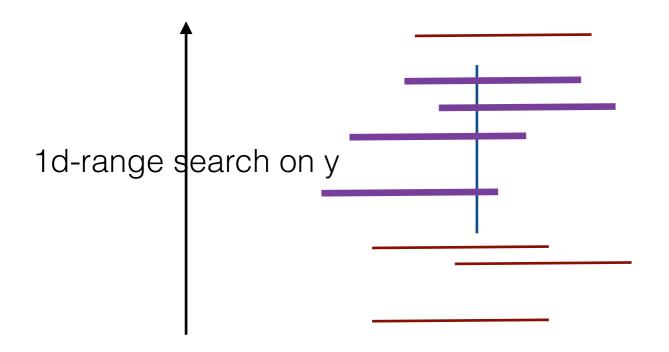
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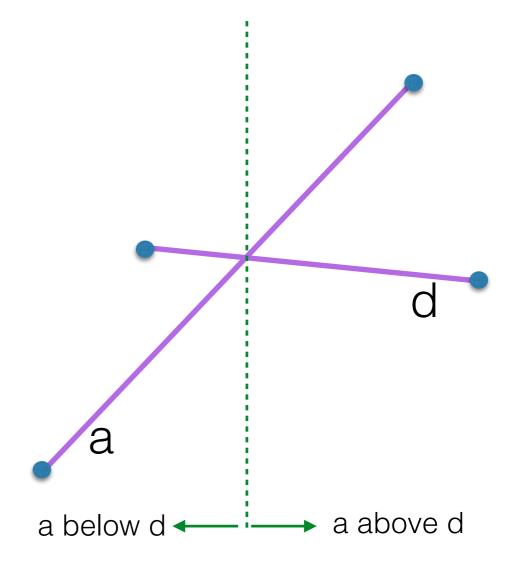
- Let X be the set of all x-coords of segments
- Traverse the events in X in order



- 4 segments are active
- How do we order these segments?
- How do we detect intersections?

How do we order active segments?



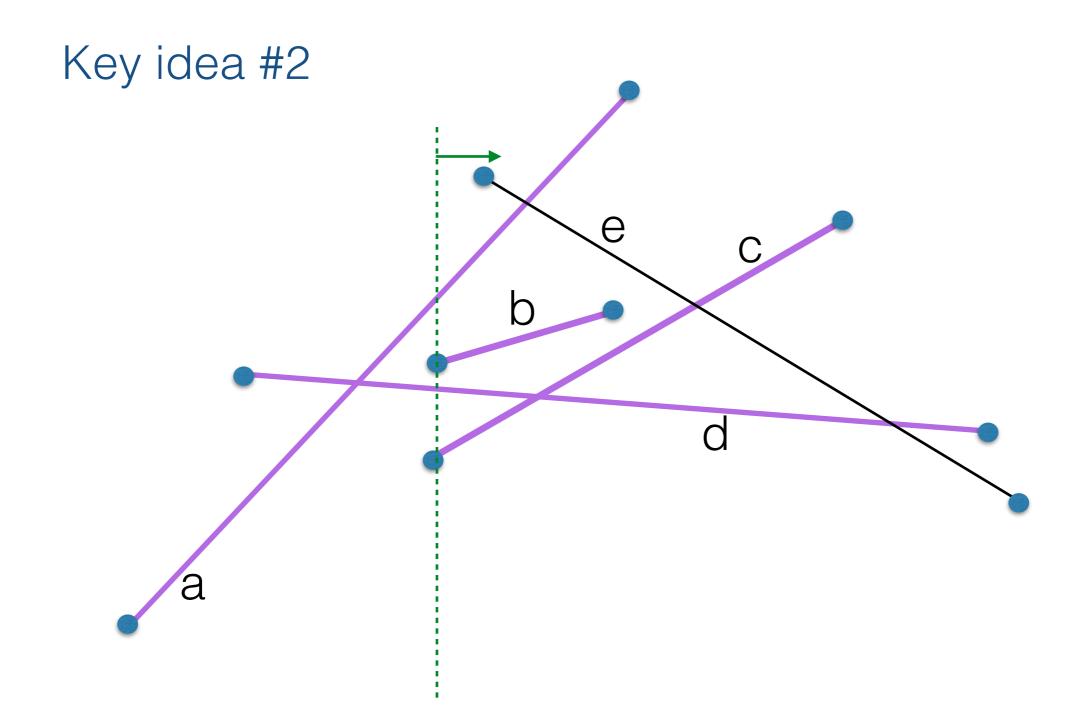


orthogonal segments

general segments

Key idea #1

- Use above-below order
- Order will flip at intersection point



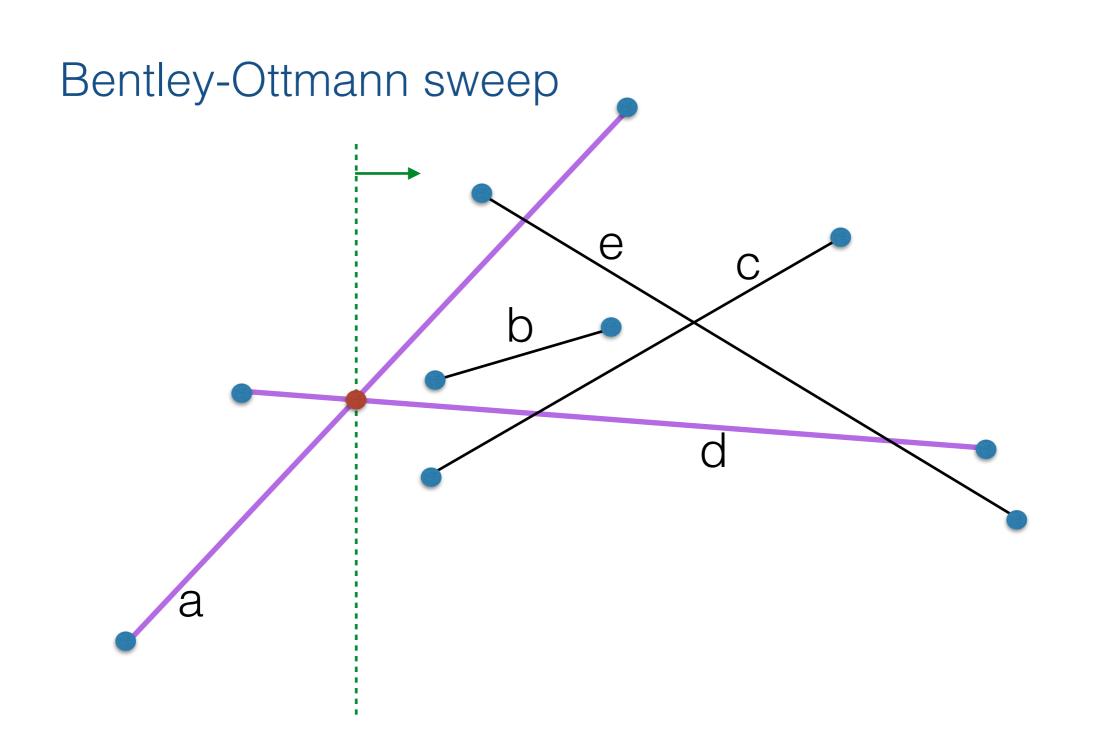
• Segments that intersect are consecutive in above-below order just before they intersect

- Let X be the set of all x-coords of segments
- Initialize AS = {}
- Traverse events in order

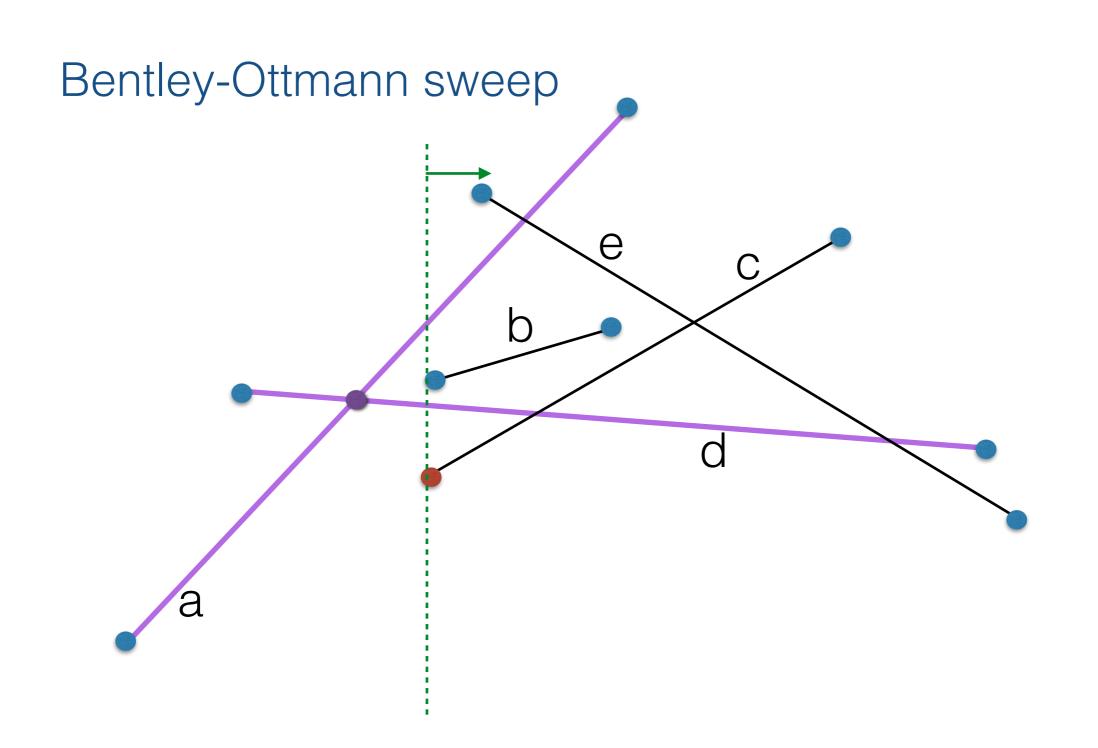
- this event is start of segment a:
 - insert a in AS: a

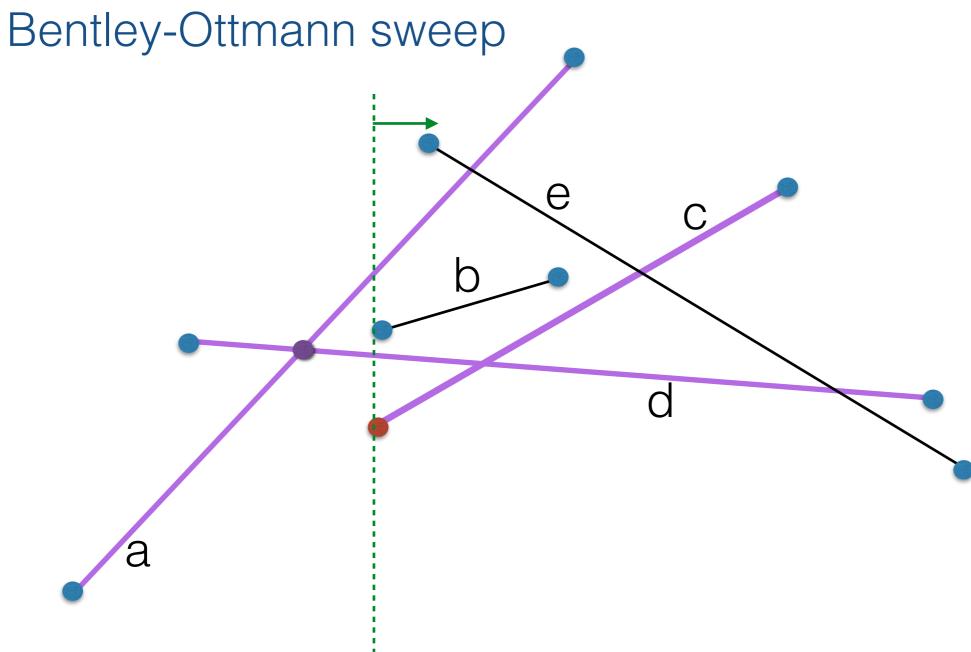
- this event is start of segment d
 - insert d in AS: a < d
 - check if (d,a) intersect to the right of the line; they do; report point and insert it in the list of future events

- this event is start of segment d
 - insert d in AS, d is above a (a < d)
 - check if (d,a) intersect to the right of the line; they do; report point and insert it in the list of future events

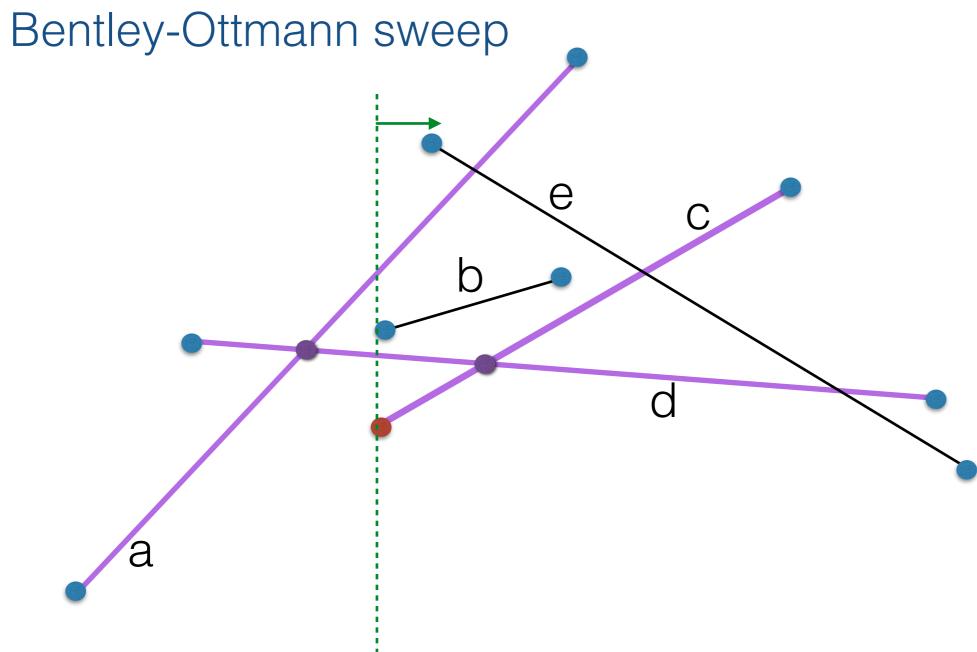


- this event is an intersection point of (a,d):
 - flip a and d is AS: a is now above d (d < a)

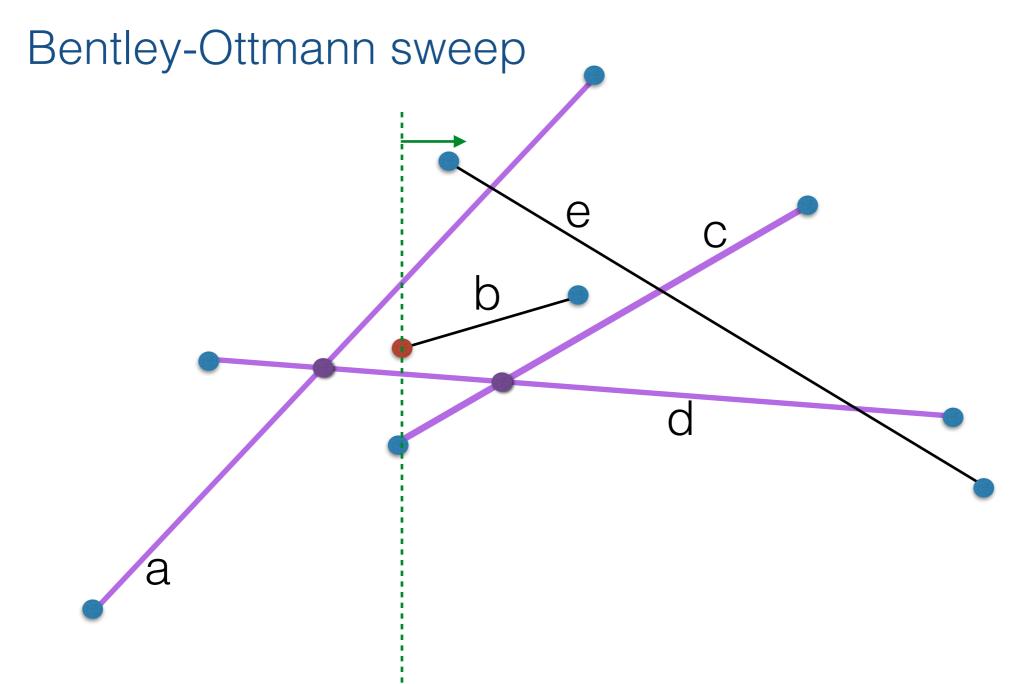




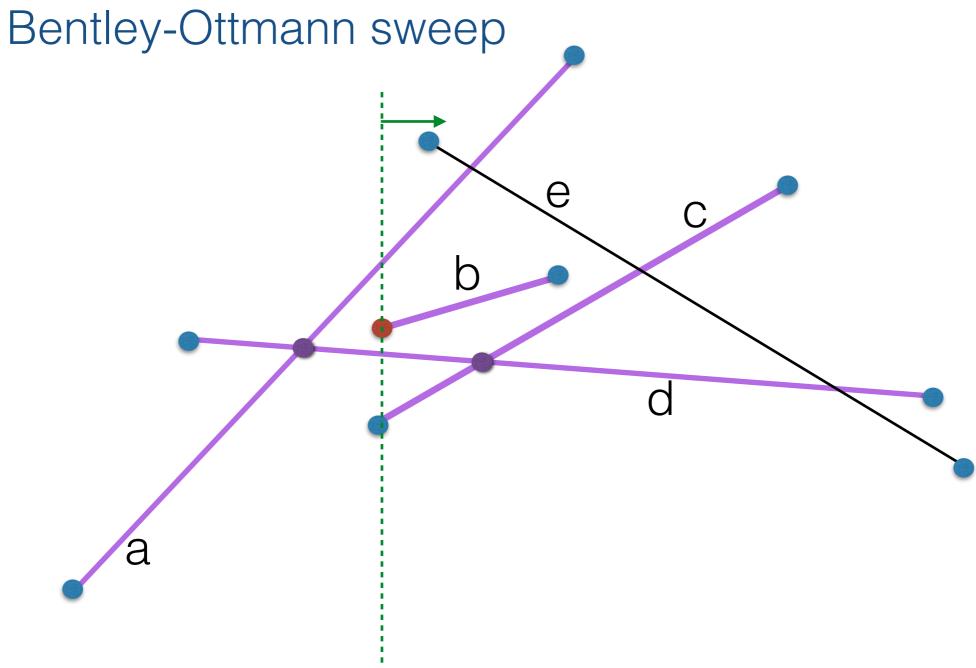
- this event is start of segment c:
 - insert c in AS; c is below d (c < d < a)
 - check c with its above and below neighbors for intersection to the right of the sweep line; this detects the intersection point of c and d; report it and insert it as future event



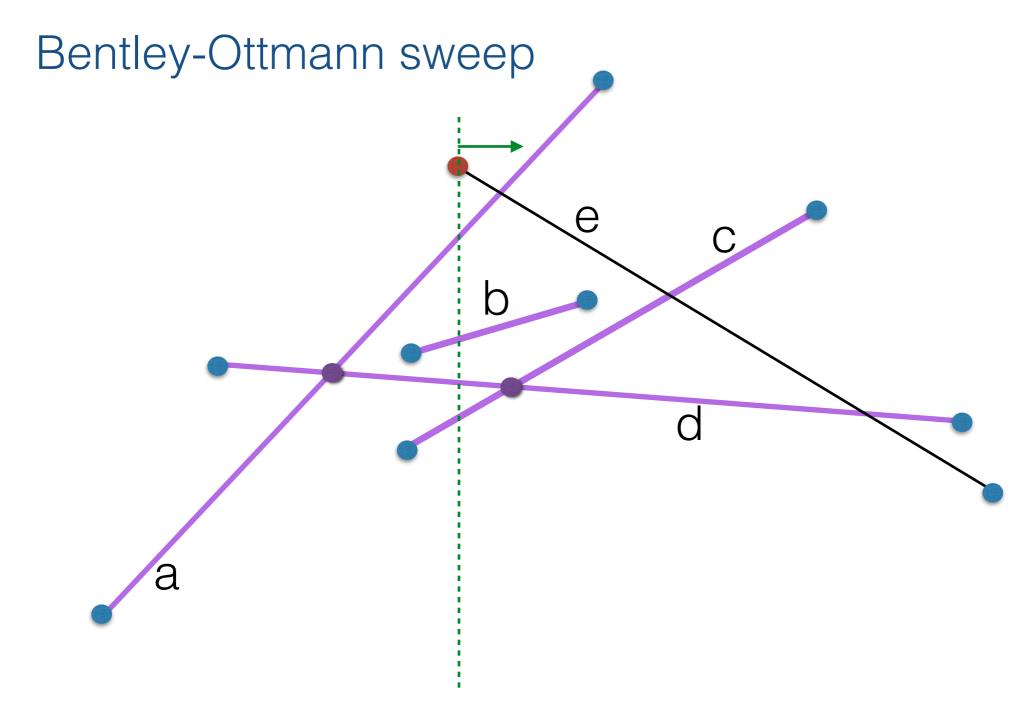
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• this event is start of segment b:



- this event is start of segment b:
 - insert b in AS; c < d < b < a
 - check b with its above and below neighbors for intersection to the right of the sweep line; (d,b) don't intersect; (b, a) don't intersect



• this event is start of segment e:

- this event is start of segment e:
 - insert e in AS: c < d < b < a < e
 - check e with its above and below neighbors for intersection to the right of the sweep line; this detects intersection point of (a,e); report it and insert it as future event

- this event is start of segment e:
 - insert e in AS: c < d < b < a < e
 - check e with its above and below neighbors for intersection to the right of the sweep line; this detects intersection point of (a,e); report it and insert it as future event

- this event is intersection of (a,e):
 - flip a and e: c < d < b < e < a
 - check new neighbors (e,b) for intersection to the right of the sweep line; (e,b) don't intersect

- this event is intersection of (c,d):
 - flip c and d: d < c < b < e < a
 - check new neighbors (c,b) for intersection to the right of the sweep line; (c,b) don't intersect

- this event is end of segment b:
 - delete b from AS: d < c < e < a
 - check new neighbors (c,e) for intersection to the right of the sweep line; this detects the intersection point of (c,e); report it and insert it as future event

- this event is end of segment b:
 - delete b from AS: d < c < e < a
 - check new neighbors (c,e) for intersection to the right of the sweep line; this detects the intersection point of (c,e); report it and insert it as future event

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 - check new neighbors (c,e) for intersection to the right of the sweep line; this detects the intersection point of (c,e); report it and insert it as future event

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 - delete a from AS: d < c < e
 - no new neighbors

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 - delete a from AS: d < c < e
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 - flip c,e in AS: d < e < c
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 - flip c,e in AS: d < e < c
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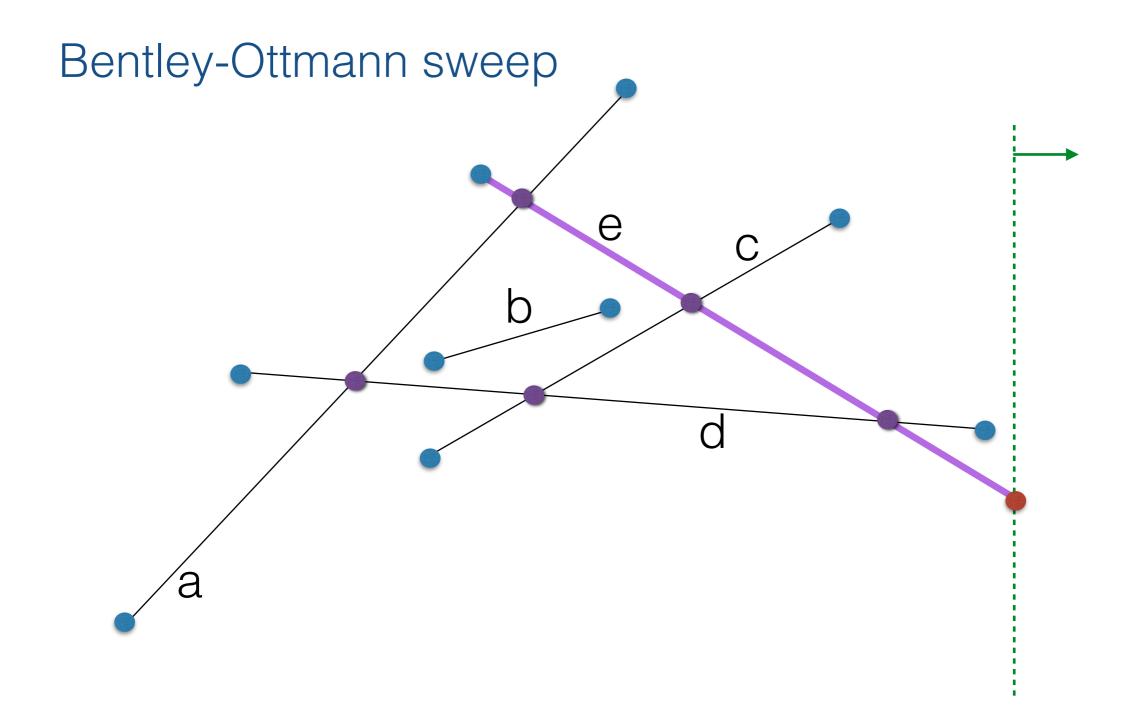
- this event is end of segment c:
 - delete c in AS: d < e
 - no new neighbors

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 - delete c in AS: d < e
 - no new neighbors

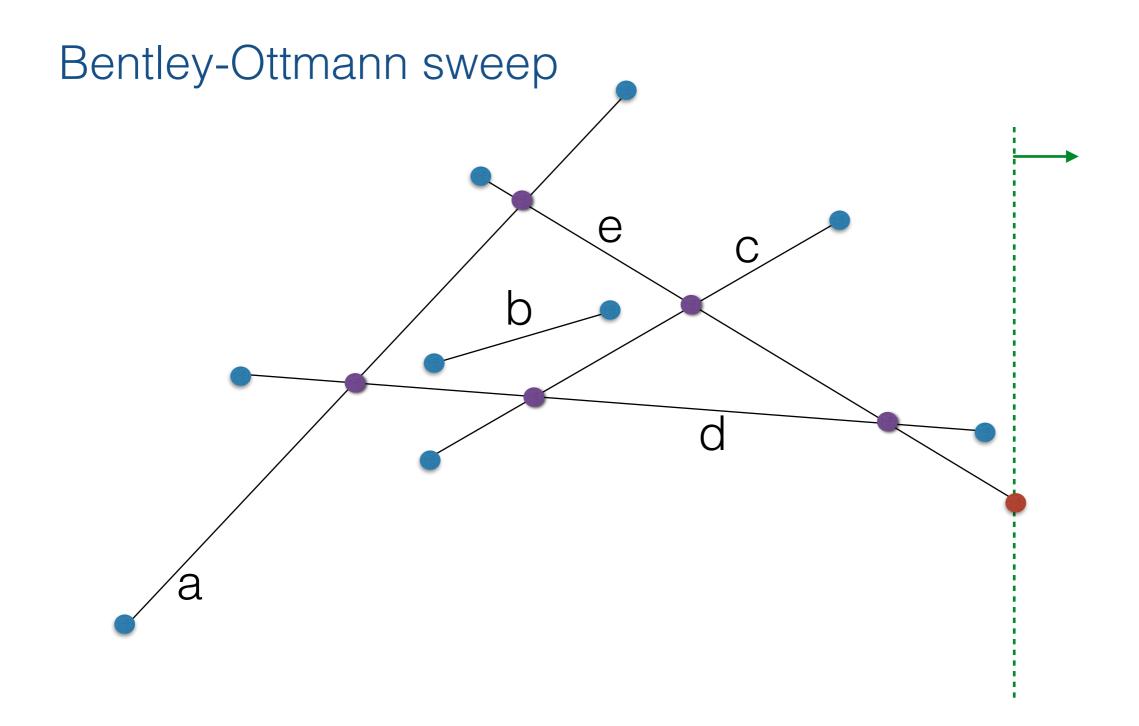
- this event is the intersection of (d,e):
 - flip d,e in AS: e < d
 - no new neighbors

- this event is the end of d:
 - delete d in AS: e
 - no new neighbors

- this event is the end of d:
 - delete d in AS: e
 - no new neighbors



- this event is the end of e:
 - delete e in AS:
 - no new neighbors



- this event is the end of e:
 - delete e in AS:
 - no new neighbors

- Simplifying assumptions
 - no vertical segments
 - no two segments intersect at their endpoints
 - no three (or more) segments have a common intersection
 - all endpoints (of segments) and all intersection points have different xcoordinates
 - no segments overlap

- These assumptions are not realistic for real data..
- But, they don't provide insight into the plane sweep technique, so we omit them
- Real data challenges
 - dealing with degenerate cases
 - dealing with finite precision arithmetic and precision problems

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- Active structure AS:
 - For any position of the sweep line SL, AS contains all active segments (ie segments that start before SL and end after SL)
 - AS is sorted by their y-coordinates of their intersection with SL

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Event list EL:

- For any position of SL, EL contains segment endpoints to the right of SL, and also the intersections to the right of SL of active segments that were/ are neighbors in SL
- EL is sorted by x-coordinate

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Active structure AS:

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Event list EL:

- For any position of SL, EL contains segment endpoints to the right of SL, and also the intersections to the right of SL of active segments that were/ are neighbors in SL
- EL is sorted by x-coordinate
- For any position of the sweep line SL, all pairs of intersecting dead segments have been reported.

Algorithm Bentley-Ottmann (S)

//S is a set of n line segments in the plane

- initialize AS= {}
- sort 2n endpoints of all segments in S by x-coord and store them in EventList
- while EventList not empty
 - let e be the next event from EventList; delete it from EL

//sweep line moves to x=e.x

• if e is left endpoint of a segment I

//I becomes active

- insert I in AS in the right place
- check if I intersects with I->prev and I->succ in AS to the right of the sweep line; if they do, insert their intersection point in the EventList

//optional: since I.prev and I.succ are not neighbors anymore, we check if they intersect and if they do, delete that intersection point from the EventList

- if e is the right endpoint of a segment
 - ...
- if e is the intersection of two segments
 - ...
- end.

Questions

- AS
 - What operations do we do on AS?
 - What data structure should we use for AS?
- EL
 - Note that we know a priori the 2n events corresponding to start and endpoints of segments, but EL is not static; the events corresponding to intersection points are generated on the fly
 - What operations do we do on EL?
 - What data structure should we use for EL?

Running time

- AS
 - What is the size of AS?
 - O(n)
 - How many operations?
 - O(n+k)
 - Overall time?
 - O((n+k)lg n)
- EventList
 - What is the size of EventList?
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 - How many operations?
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Bentley-Ottmann sweep runs in O((n+k) lg n) time.