

# # Segment Tree

## Concepts & Qns... #



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"No more fear of Segment Tree"

video - ①

Segment Tree ???

जो भी है,  
सुना ली है

0	1	2	3	4	5	6

3	1	2	7	5	6	3
---	---	---	---	---	---	---

$n = 7$

index: (1,4) → Sum →

$$O(Q * n)$$

index: (2,5) → Sum →

index: (0,5) → Sum →

large  $n \rightarrow 10^6$

$$Q = 10^4$$

$$O(10^4 * 10^6)$$

An efficient data structure that allows

⇒ Efficient Querying of intervals/range

⇒ Efficient updating of intervals/range

range queries to find

↳ Sum ✓

↳ minimum ✓

↳ maximum ✓

etc....

Segment  
Tree

# How to Build Segment Tree??

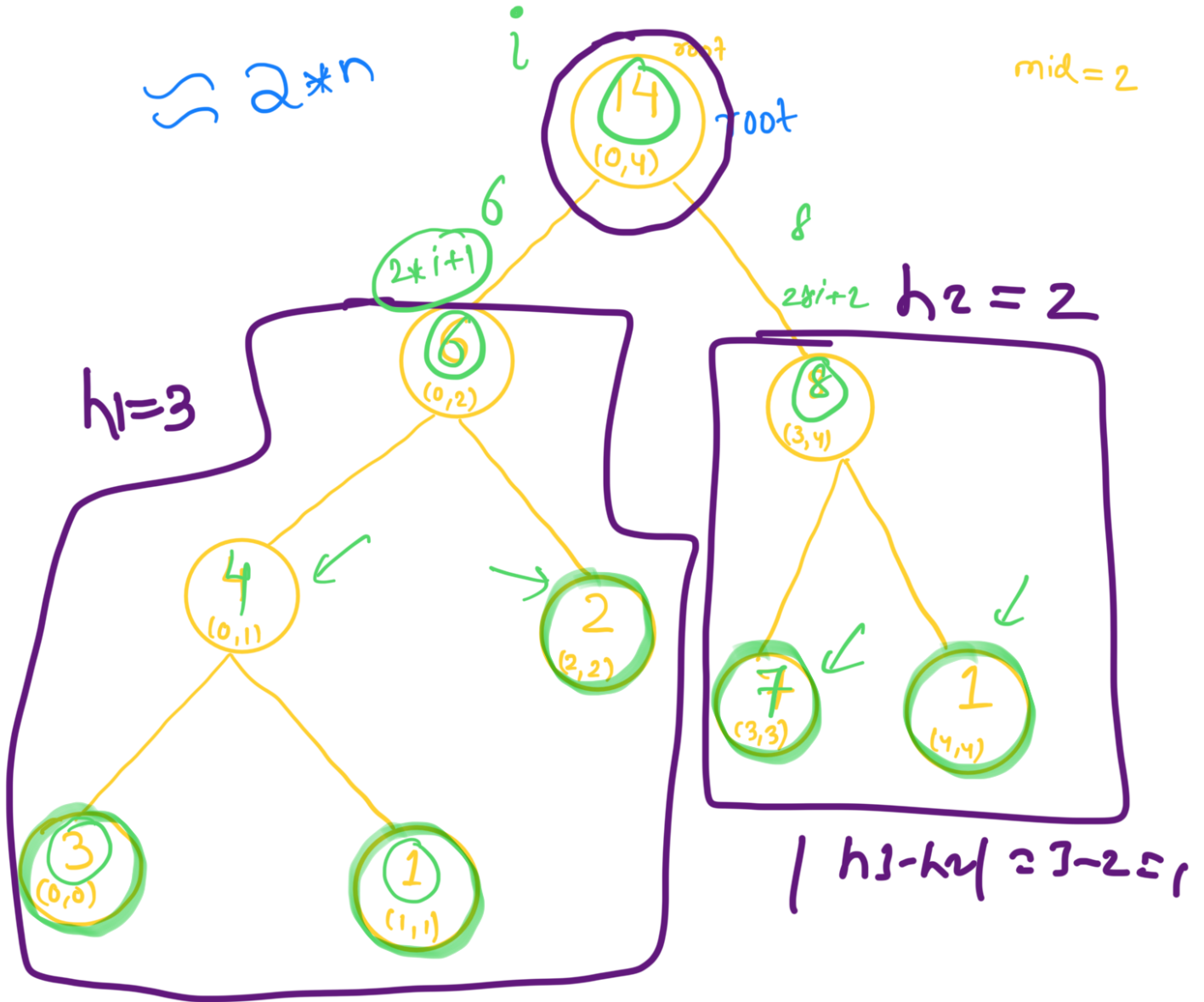
0	1	2	3	4
3	1	2	7	1

Range Sum

$n = 5$

$\approx 2 * n$

mid = 2



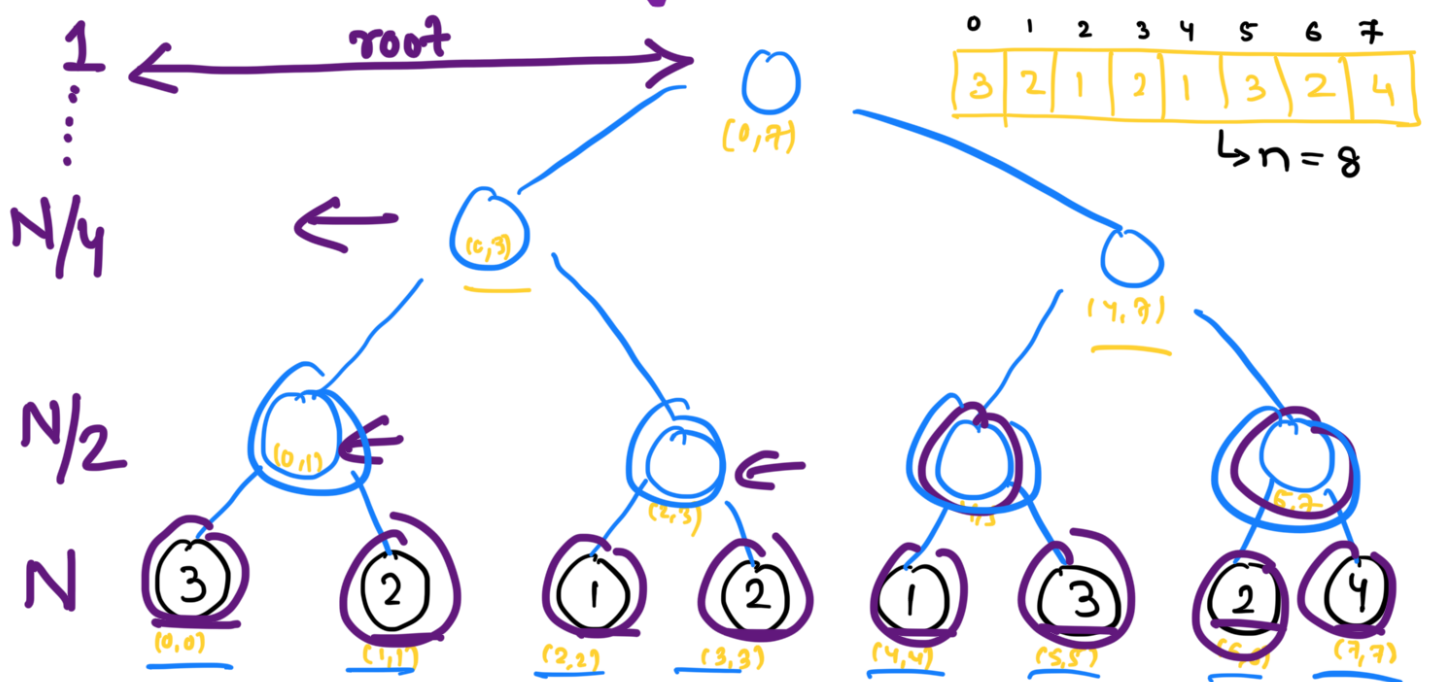
1. Binary Tree
  2. 2 children of all non-leaf nodes.
  3. → Leaf Nodes - Represents a single element in an array.
- Root Node - Represents entire array.

→ Other nodes - Represents an interval or range of an array.

4. Height =  $\lceil \log_2(n) \rceil$

\* 5. Balanced Binary Tree.

# How many nodes ???



# nodes =  $N + N/2 + N/4 + \dots + 1$

Sum<sub>∞</sub> =  $\frac{a}{1-r} = \frac{N}{1-\frac{1}{2}} = (2 * N)$  ↳ G.P.

$$\begin{aligned} \text{leaf} &= N \\ \text{internal} &= (N-1) \end{aligned} \quad \left. \vphantom{\begin{aligned} \text{leaf} &= N \\ \text{internal} &= (N-1) \end{aligned}} \right\} N + N - 1 = (2N - 1) \\ &\approx 2N$$

Why I am not using Prefix Sum Array???

nums =

0	1	2	3	4
1	3	4	2	2

(1, 4)

PrefixSum =

0	1	2	3	4
1	4	8	10	12

12 - 1

= 11 ✓

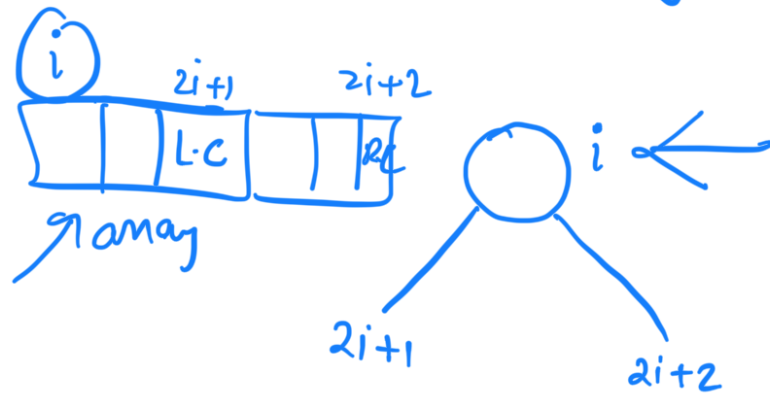
- ① range Sum
- ② range max-len.
- ③ " " " "

How to actually build the Segment tree in code:-

0 1 2 3

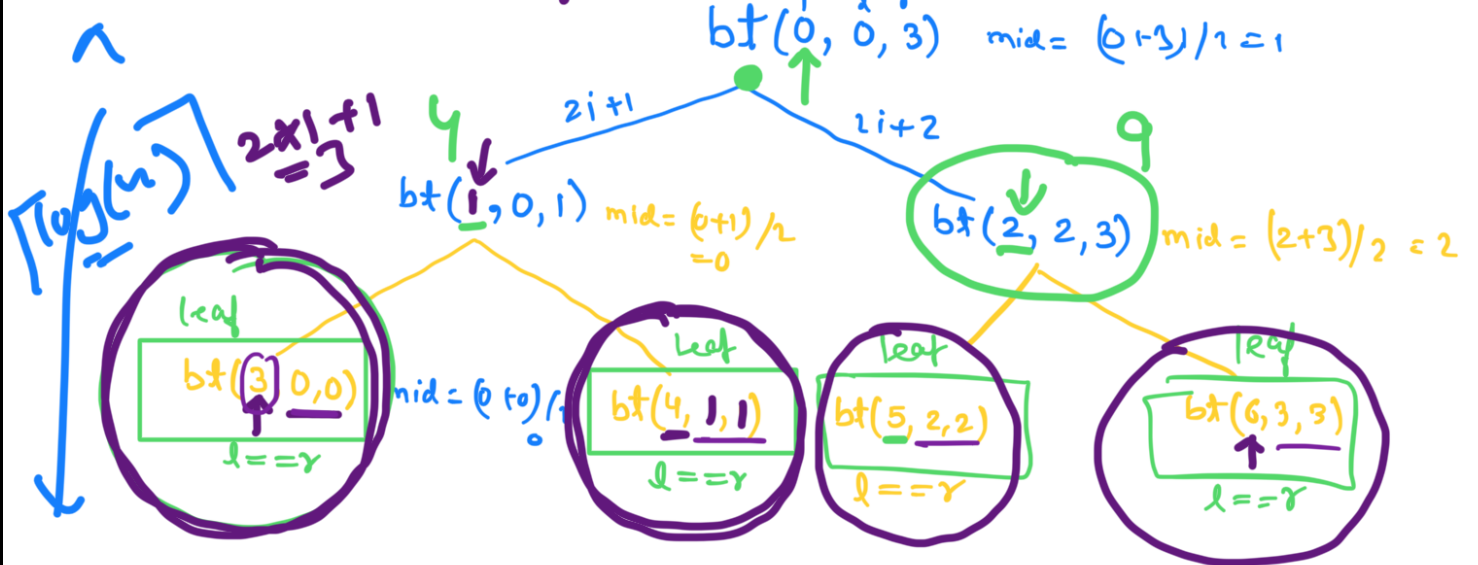
3	1	2	7
---	---	---	---

Tree  $\Rightarrow$  as an array represent



0	1	2	3
3	1	2	4

,  $n = 4$



0	1	2	3	4	5	6	7	8	9	10	11...
13	4	9	3	1	2	7					

$$n$$

$$(n-1)$$

$$n + n - 1$$

$$= 2n - 1$$

$$2 * n$$

$$= 8-1 = 7$$

## Story to Code :-

buildTree (<sup>i</sup>0, <sup>0</sup>l, <sup>n-1</sup>r) ;

↑      ↑      ↑

buildTree ( <sup>↓</sup>i, <sup>↓</sup>l, <sup>↓</sup>r ) {

↑      ↑      ↑

// Base Case

if ( l == r ) {

segTree[i] = nums[r];  
return;

}

int mid = (l+r)/2 ;

buildTree ( 2\*i+1, l, mid );

buildTree ( 2\*i+2, mid+1, r );

} Recursion Leap  
of Faith

segTree[i] = segTree[2\*i+1] + segTree[2\*i+2];

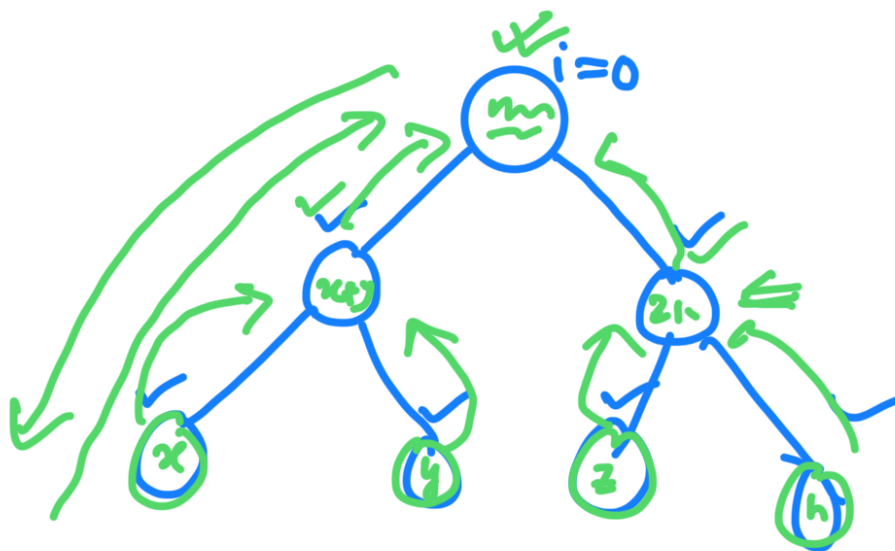
}

What will be the size  
of segtree array ???

$$2 * n.$$

Time Complexity

$$S.C = O(2n)$$



Visiting all nodes twice.



$$\approx \sqrt{2 * n}$$

$$\approx O(n)$$

T.C of building S.T =  $O(n)$ .

↳ what is segment Tree

↳ Build Segment Tree.

↳ Update Segment Tree ✓✓

↳ Query Segment Tree ✓✓

$$(l, r) = \text{sum}$$

(2, 3, 1) - 2

Range Sum Query.

Range Max / Min / Prod /