

Building Compilers for Reconfigurable Switches

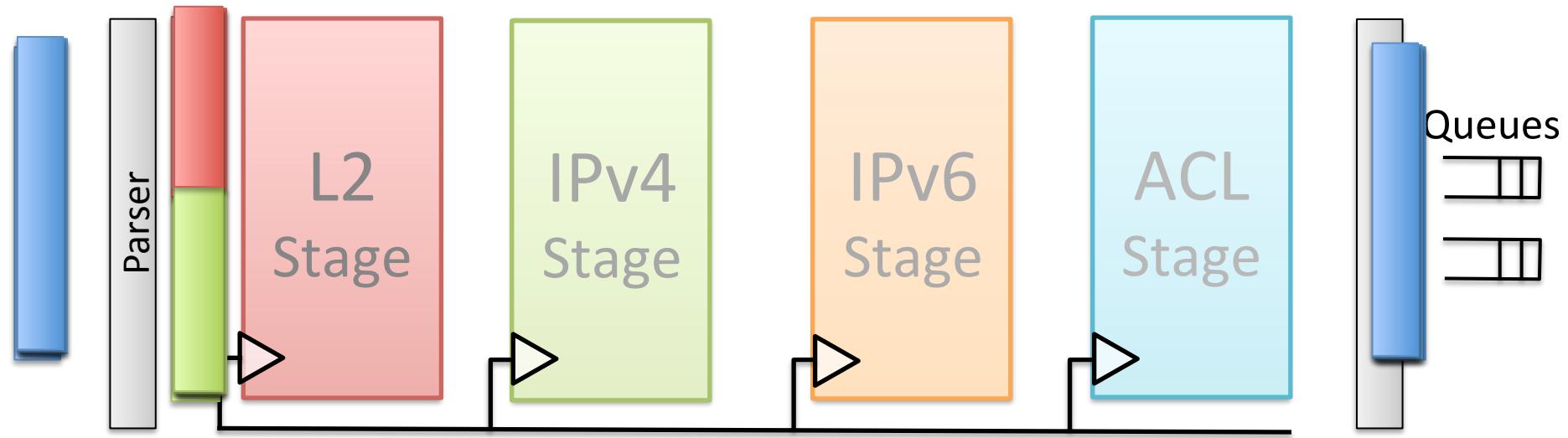
*Lavanya Jose, Lisa Yan,
Nick McKeown, and George Varghese*

Research funded by AT&T, Intel, Open Networking Research Center.

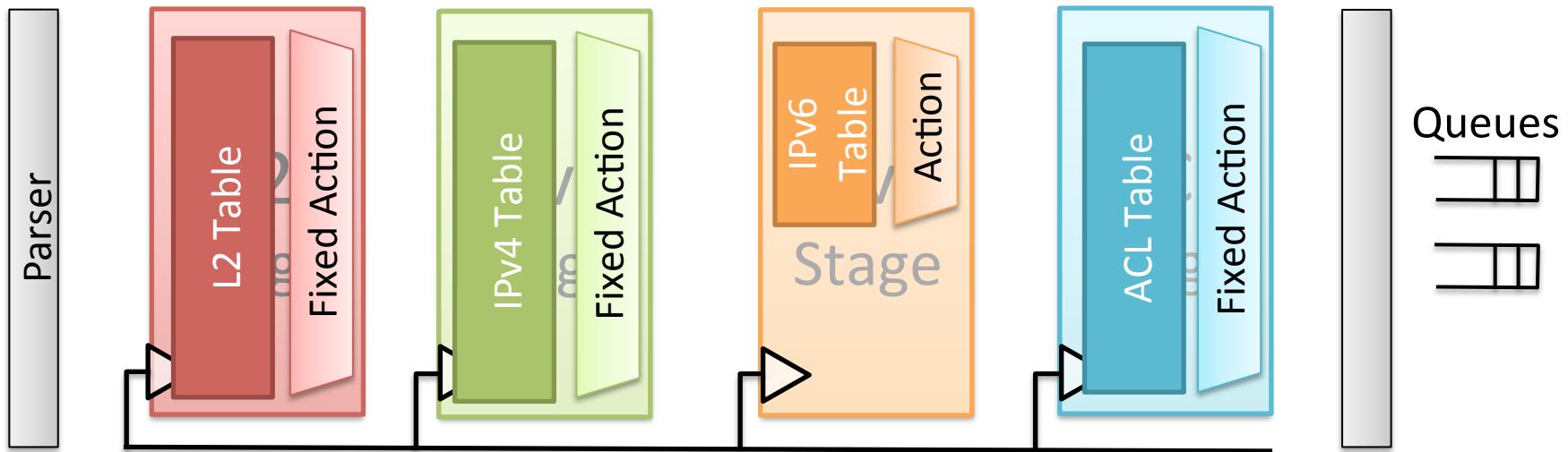
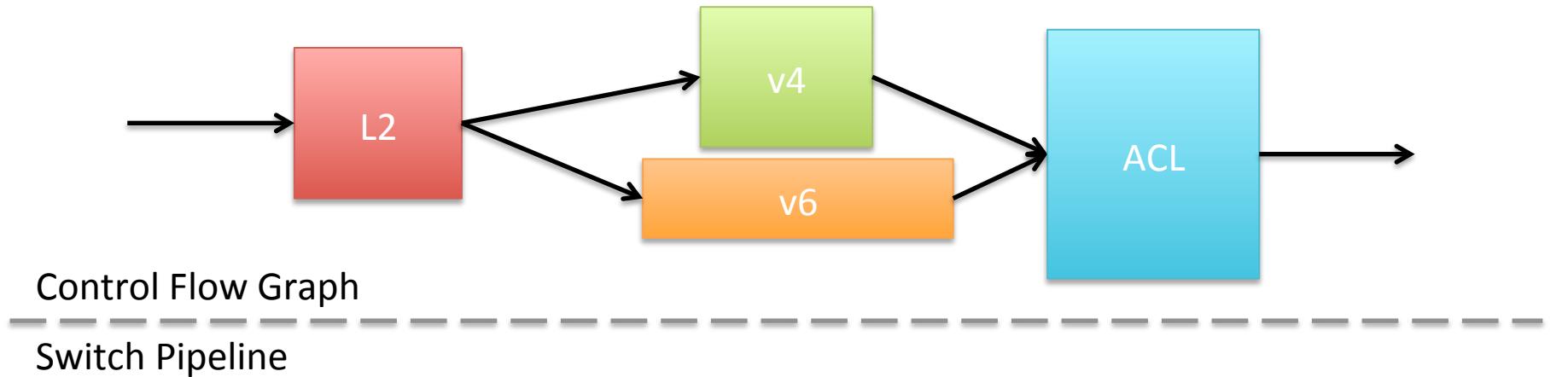
In the next 20 minutes

- Fixed-function switch chips will be replaced by reconfigurable switch chips
- We will program them using languages like P4
- We need a compiler to compile P4 programs to reconfigurable switch chips.

Fixed-Function Switch Chips



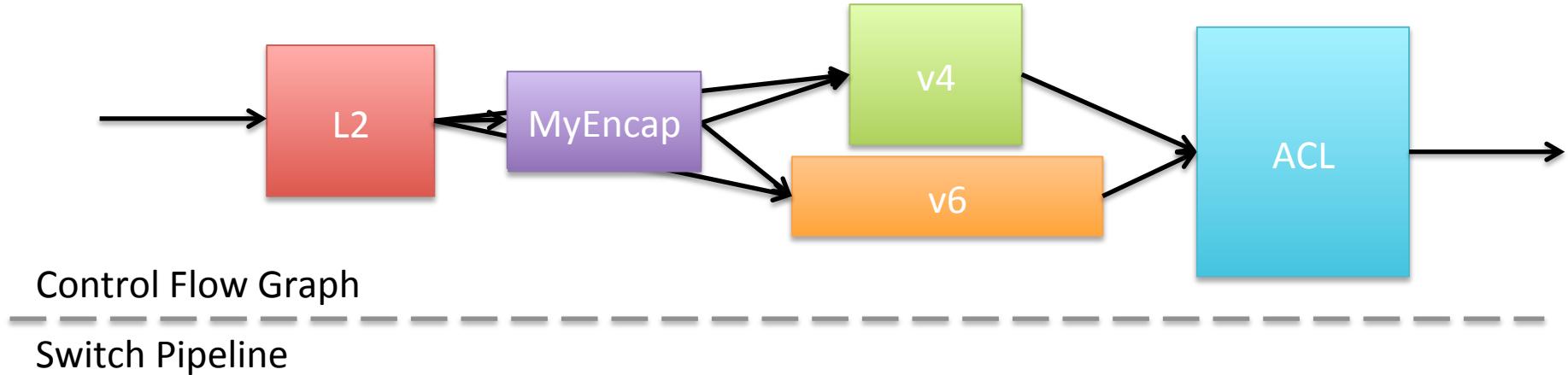
Control Flow Graph



Fixed-Function Switch Chips Are Limited

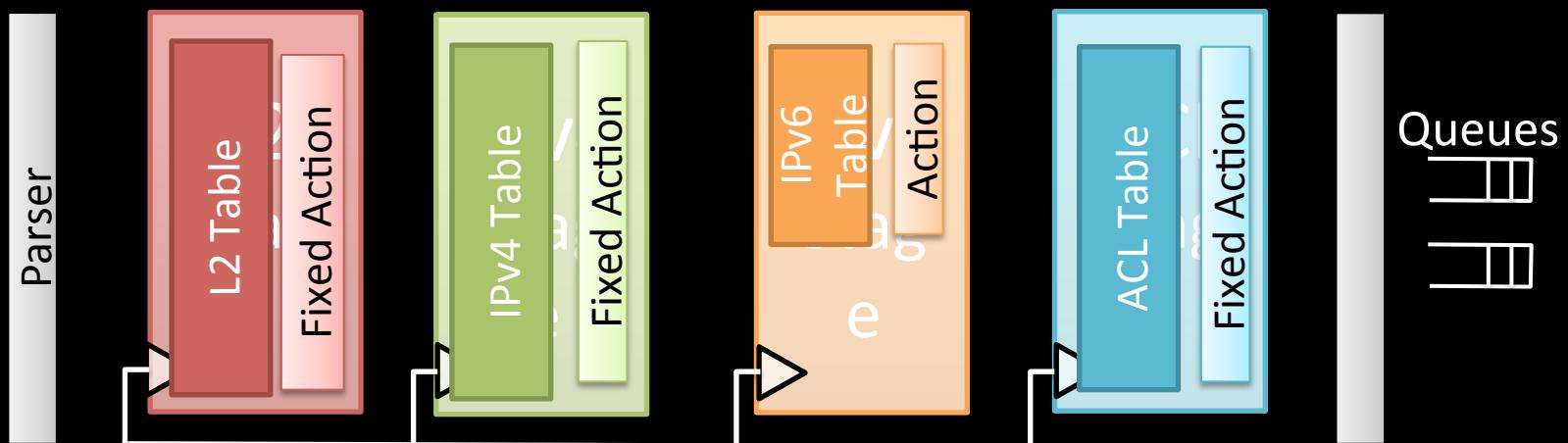
1. Can't add new forwarding functionality
2. Can't add new monitoring functionality

Fixed-Function Switch Chips

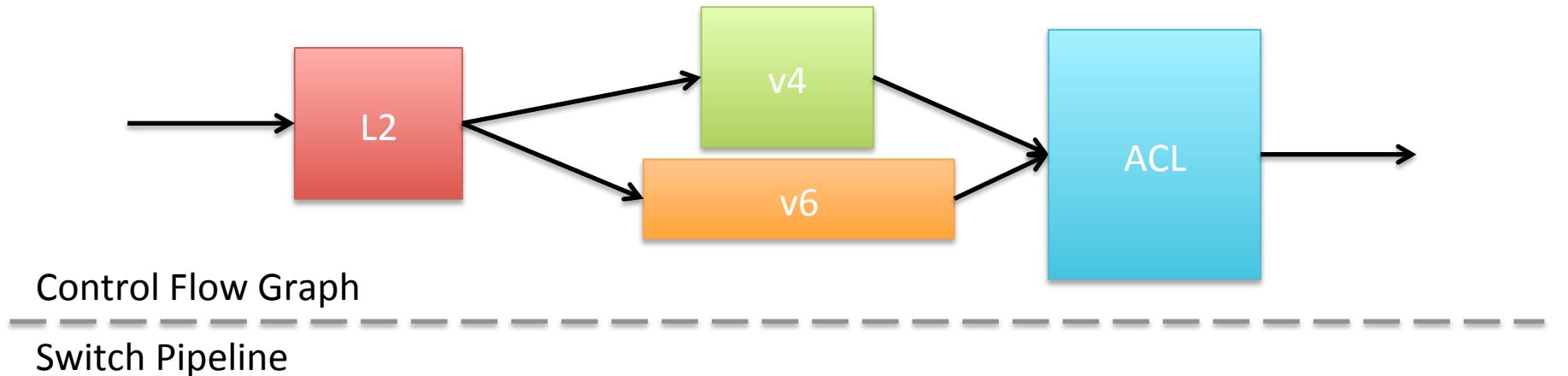


Fixed-Function Switch Chips Are Limited

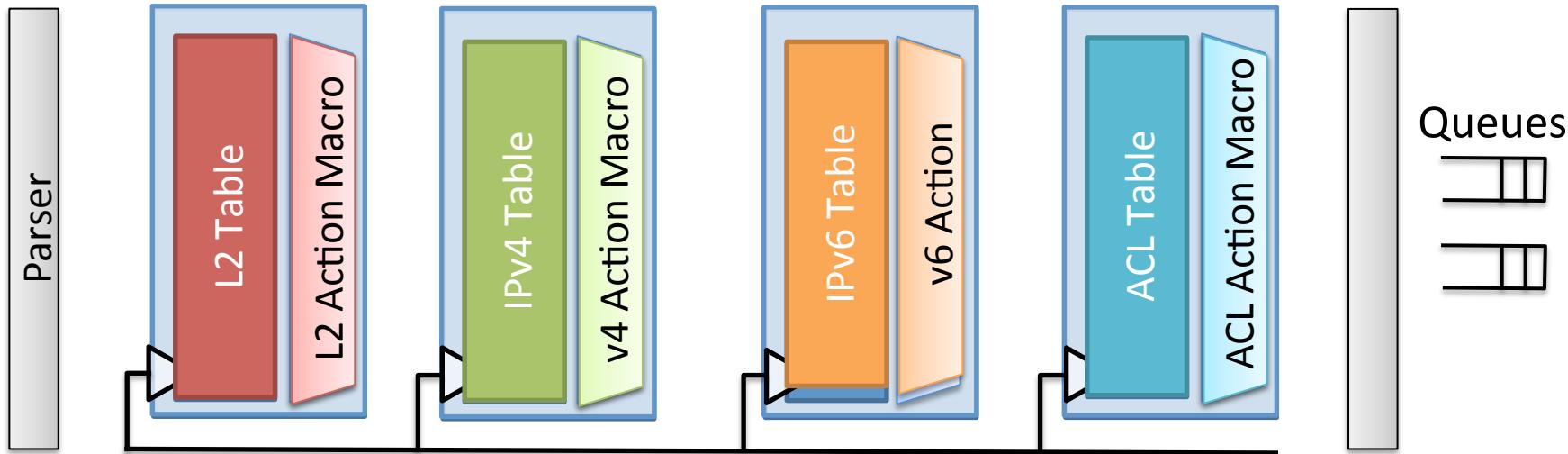
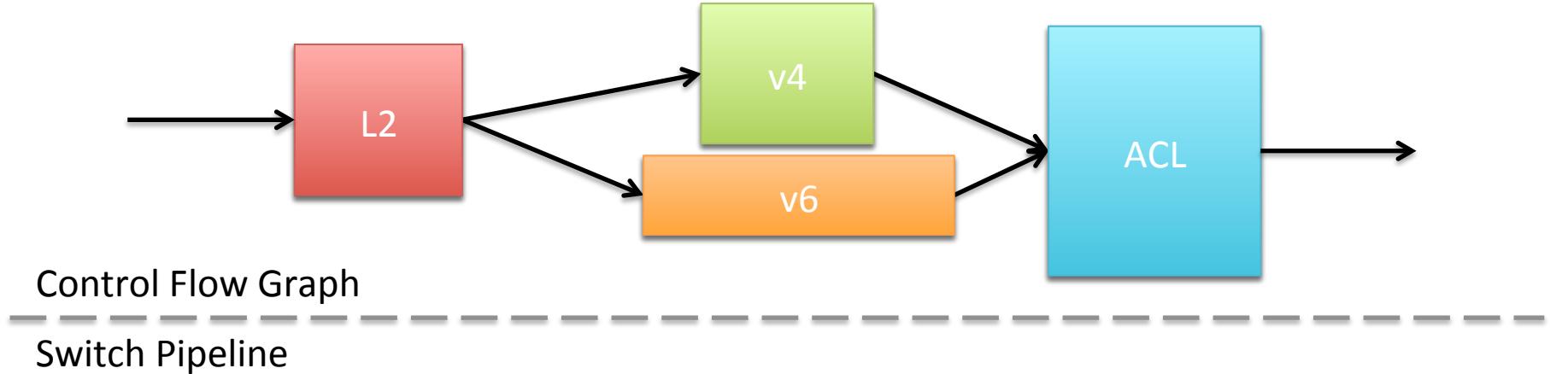
1. Can't add new forwarding functionality
2. Can't add new monitoring functionality
3. Can't move resources between functions



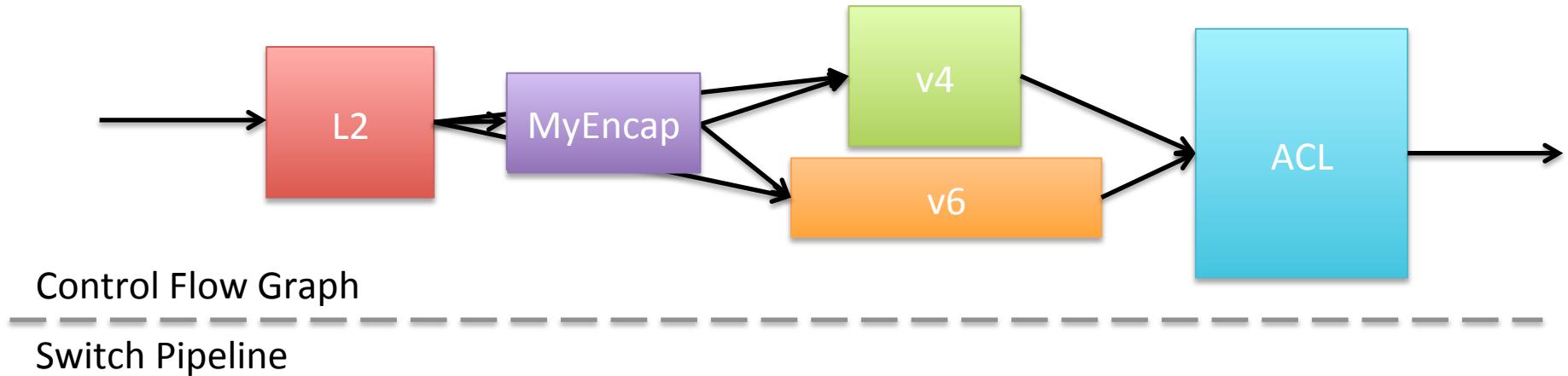
Reconfigurable Switch Chips



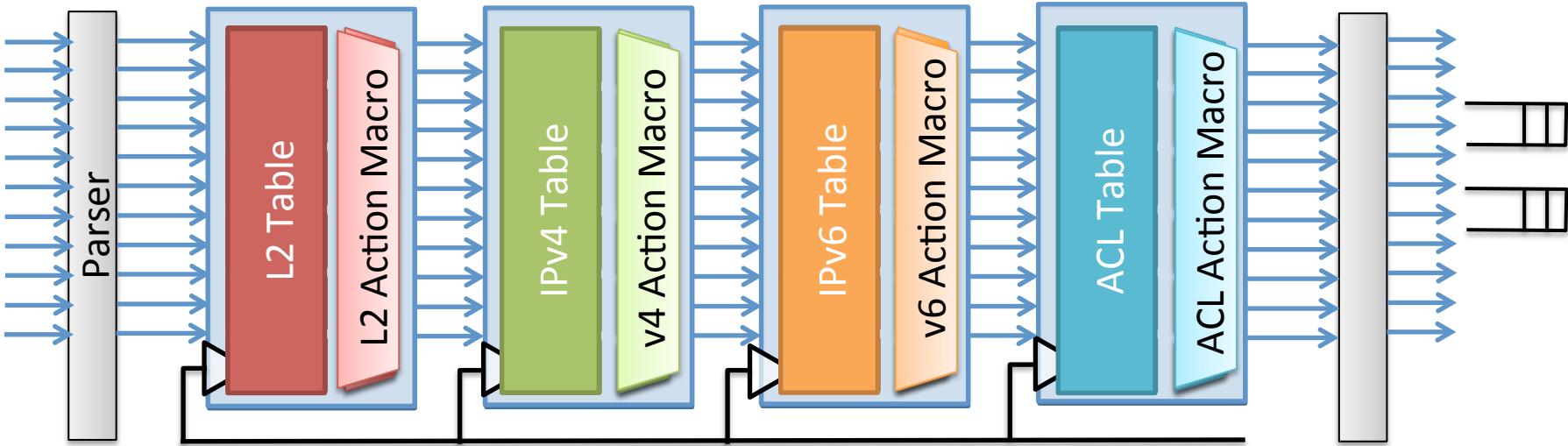
Mapping Control Flow to Reconfigurable Chip.



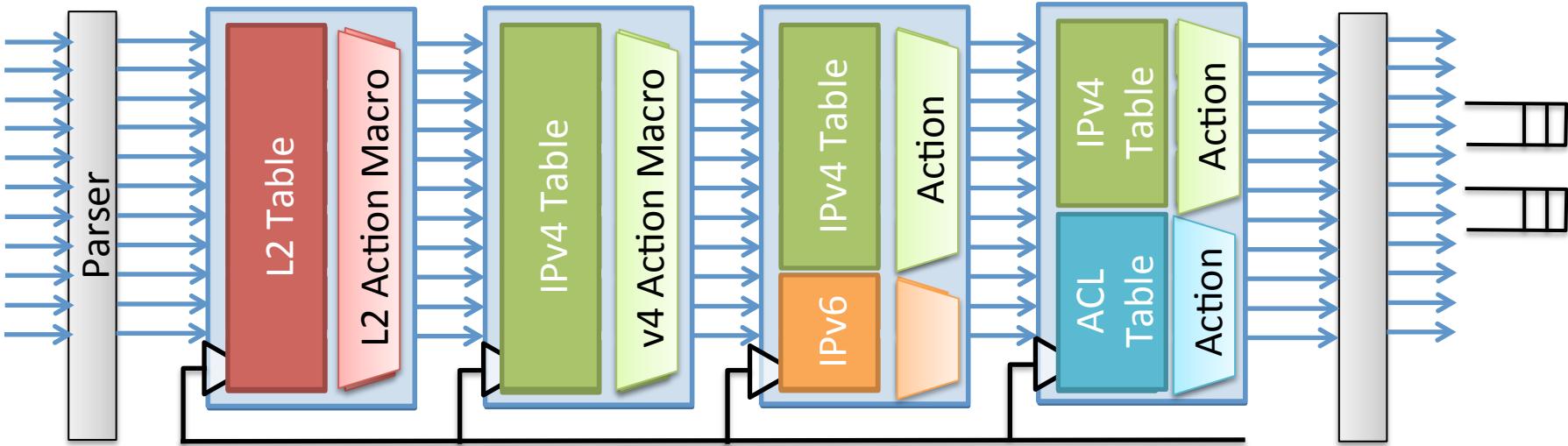
Reconfigurable Switch Chips



Match Action
Memory ALU



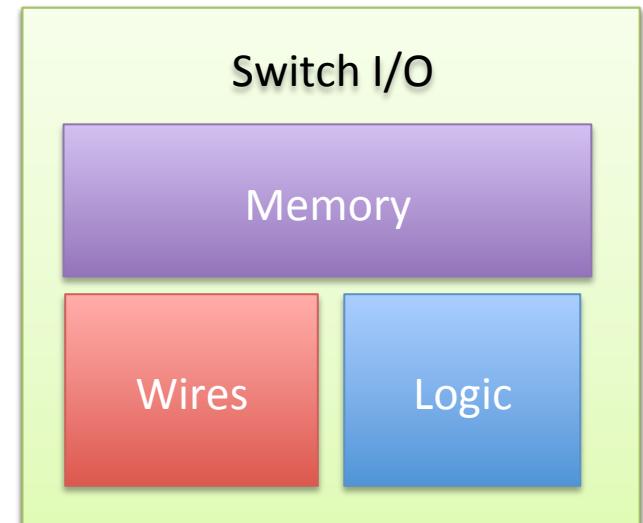
Protocol Independent Switch



Match + Action Processor: pipelined and in-parallel

Reconfigurability: the norm in 5 years

- Reconfigurability adds mostly to logic.
- Logic is getting relatively smaller.
- The cost of reconfigurability is going down.
- Fixed switch chip area today:
 - I/O (40%), Memory (40%),
 - Wires, Logic





Press Fixed Function Broadcom Tomahawk: 3.2 Tbps
Reconfigurable Cavium Xpliant: 3.2 Tbps

Broadcom Delivers Industry's First High-Density 25/100 Gigabit Ethernet Switch for Cloud-Scale Networks



Customers, New StrataXGS® Tomahawk™ Series Delivers 3.2 Tbps SDN Control and Visibility Features

PRODUCTS SUPPORT NEWS & EVENTS SALES SOFTWARE PARTNERS CORPORATE INVESTORS
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News&Events
Press Releases
Cavium Networks in the News

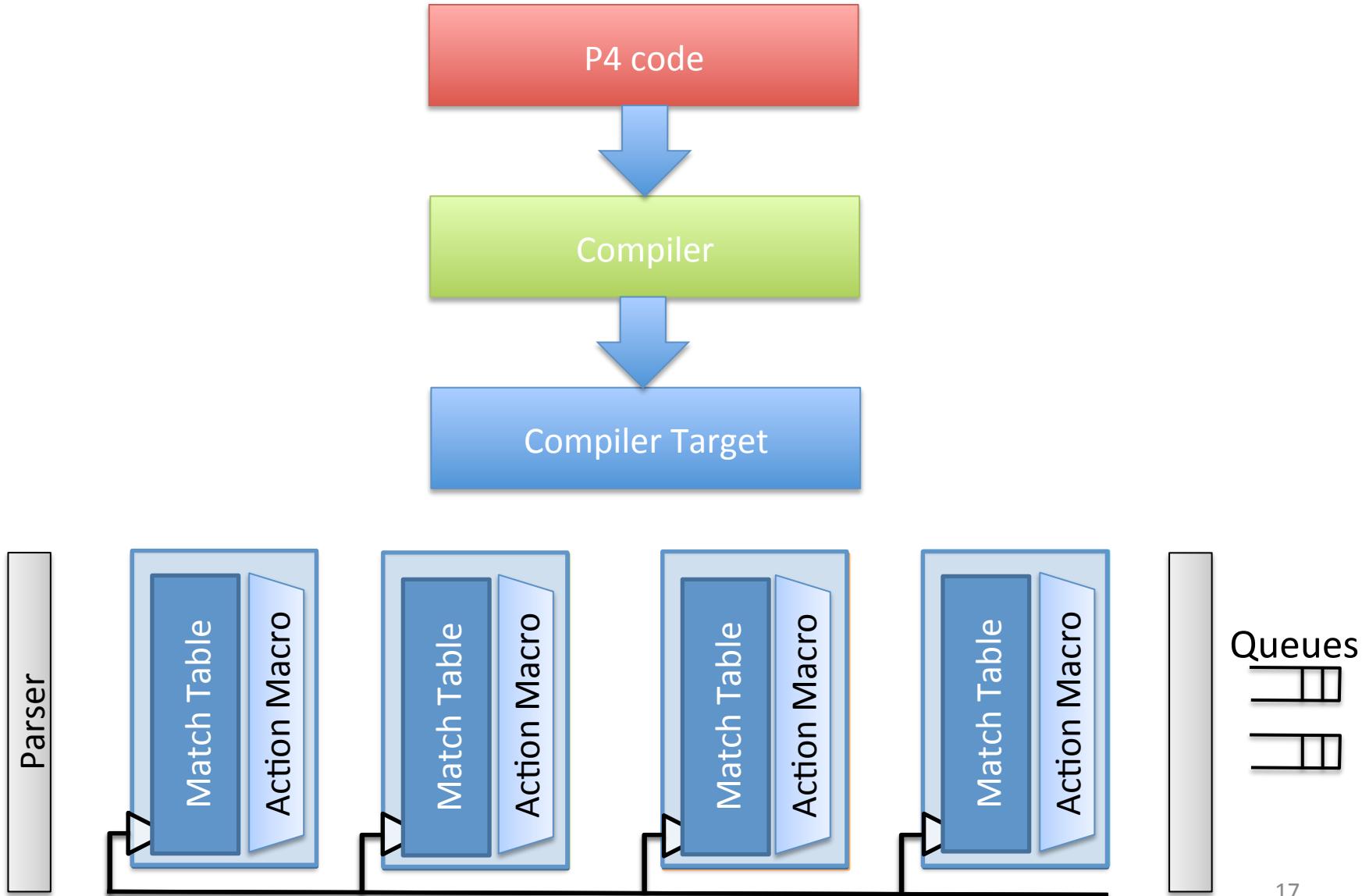
News & Events > Press Releases

Cavium and Xpliant Introduce a Fully Programmable Switch Silicon Family Scaling to 3.2 Terabits per Second
Xpliant™ Packet Architecture (XPA™) Enables Data Plane Flexibility for SDN Without Compromising Raw Speed



Reconfigurable chips are
inevitable.

Configuring Switch Chips



P4 (<http://p4.org/>)

Parser (ANCS'13)

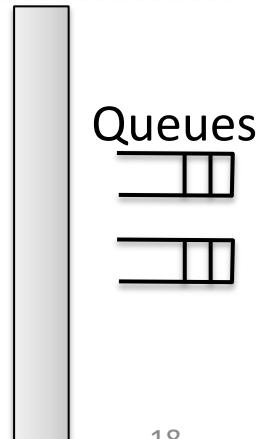
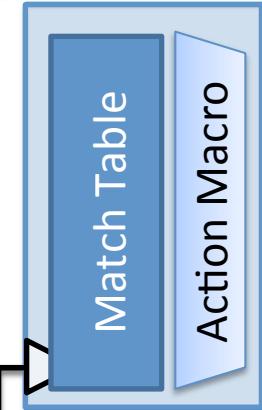
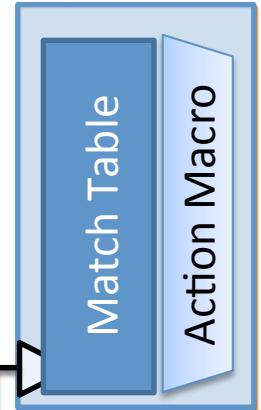
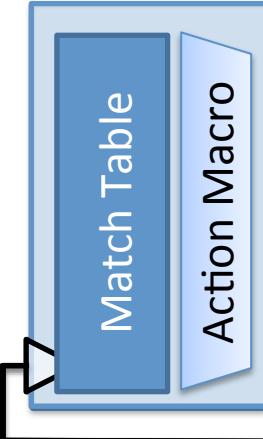
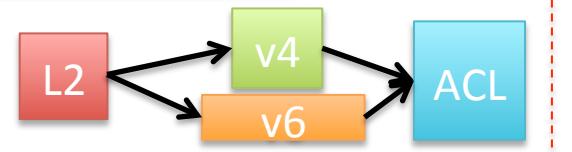
```
parser parse_etherne{  
    extract(etherne);  
    select(latest.etherType)  
{  
    0x800 : parse_ip4;  
    0x86DD : parse_ip6;  
}
```

Match Action Tables

```
table ipv4_lpm {  
    reads {  
        ipv4.dstAddr :  
            lpm;  
    }  
    actions {  
        set_next_hop;  
        drop;  
    }  
}
```

Control Flow Graph

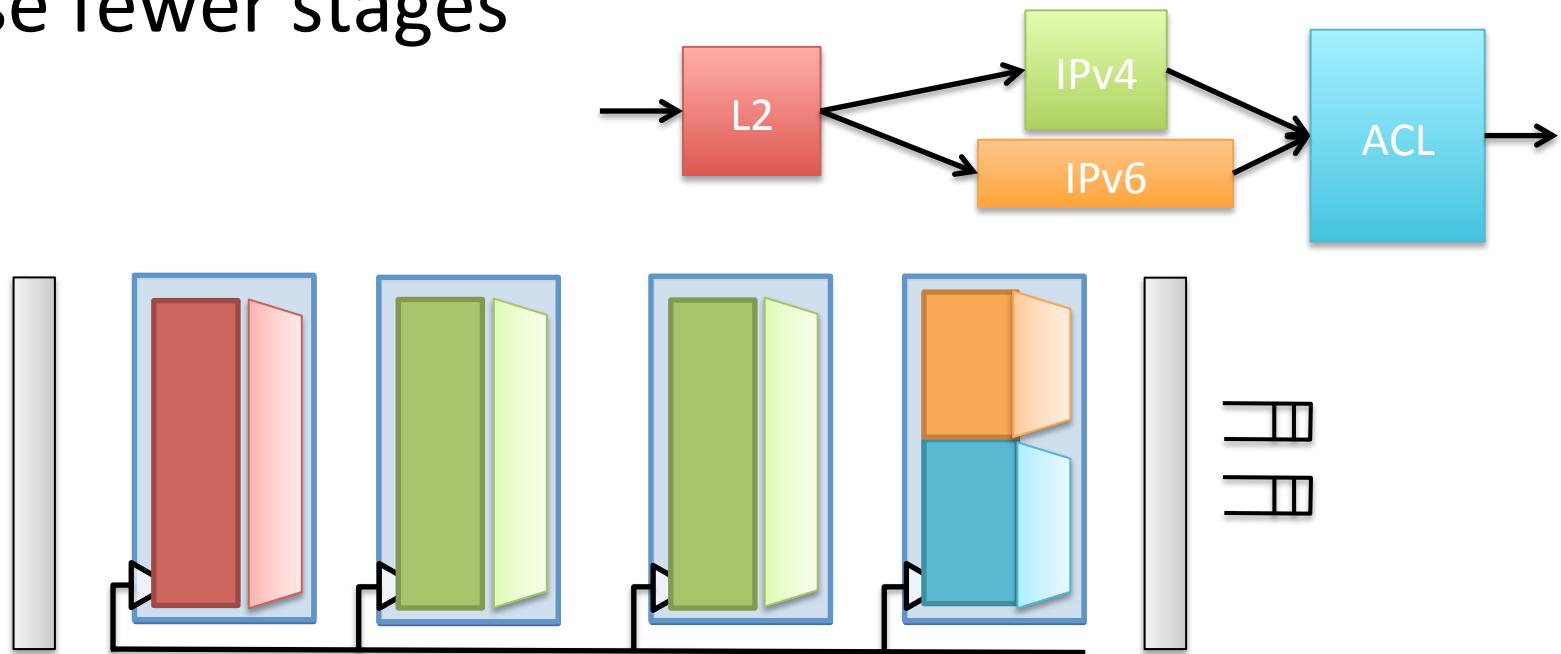
```
control ingress  
{  
    apply(l2_table);  
    if (valid(ipv4)) {  
        apply(ipv4_table);  
    }  
    if (valid(ipv6)) {  
        apply(ipv6_table);  
    }  
    apply (acl);  
}
```



What does reconfigurability
buy us?

Benefits of Reconfigurability

- Use resources efficiently
 - Multiple tables per stage
 - Big table in multiple stages
- Use fewer stages



Naïve Mapping: Control Flow Graph

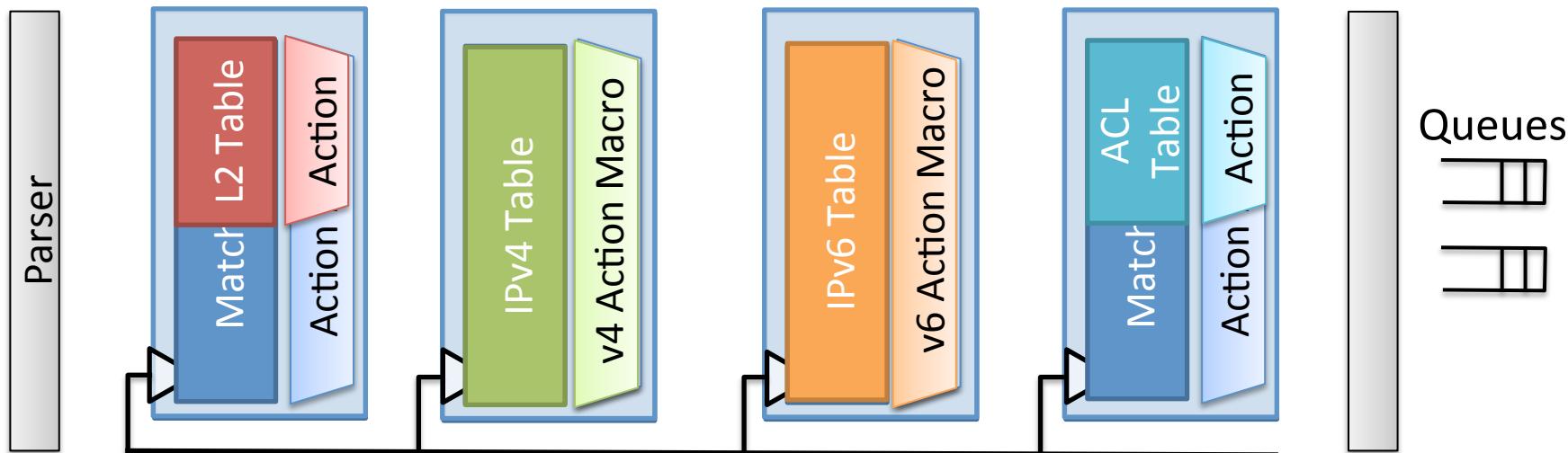
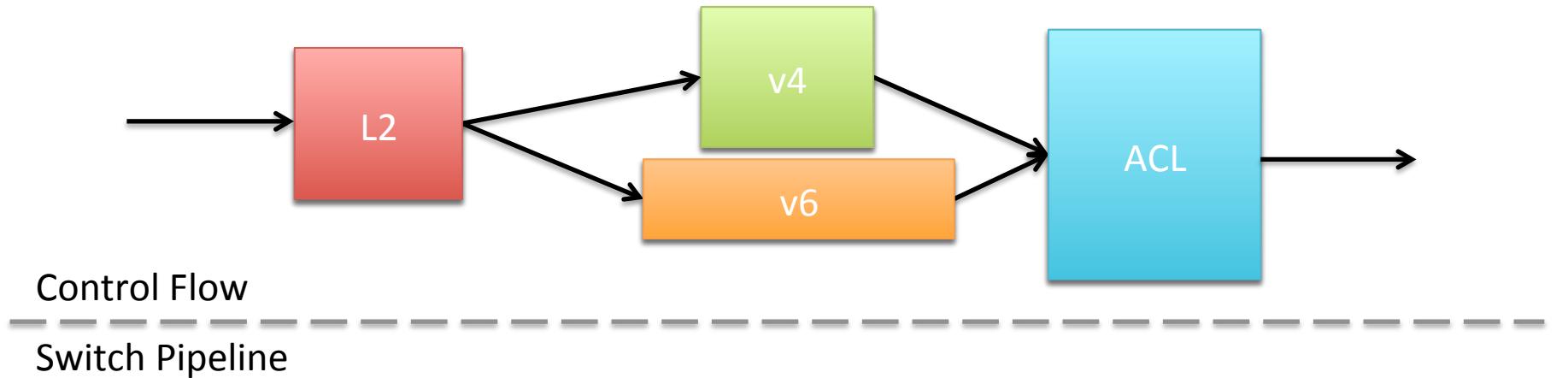
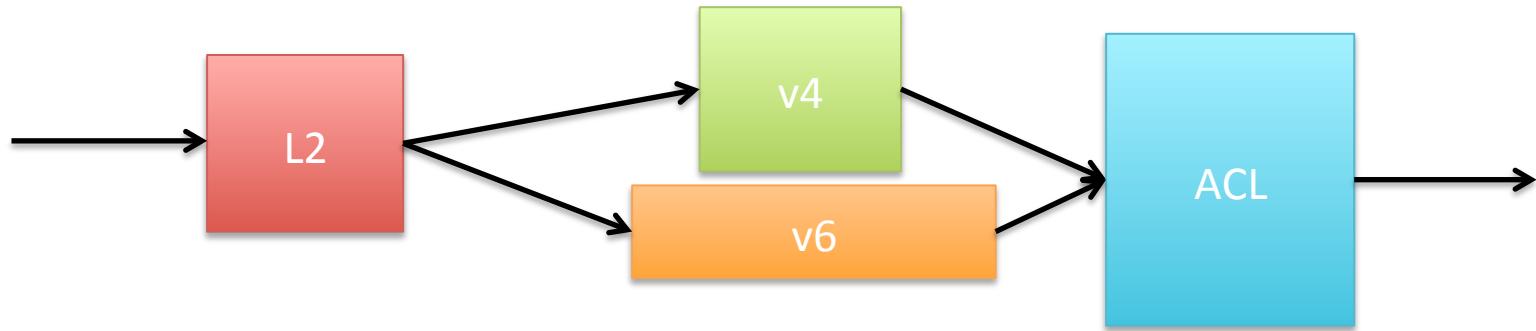
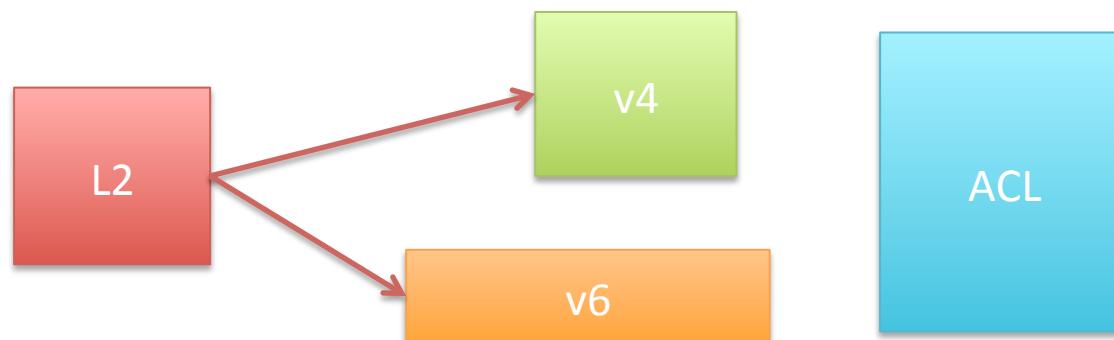


Table Dependency Graph (TDG)

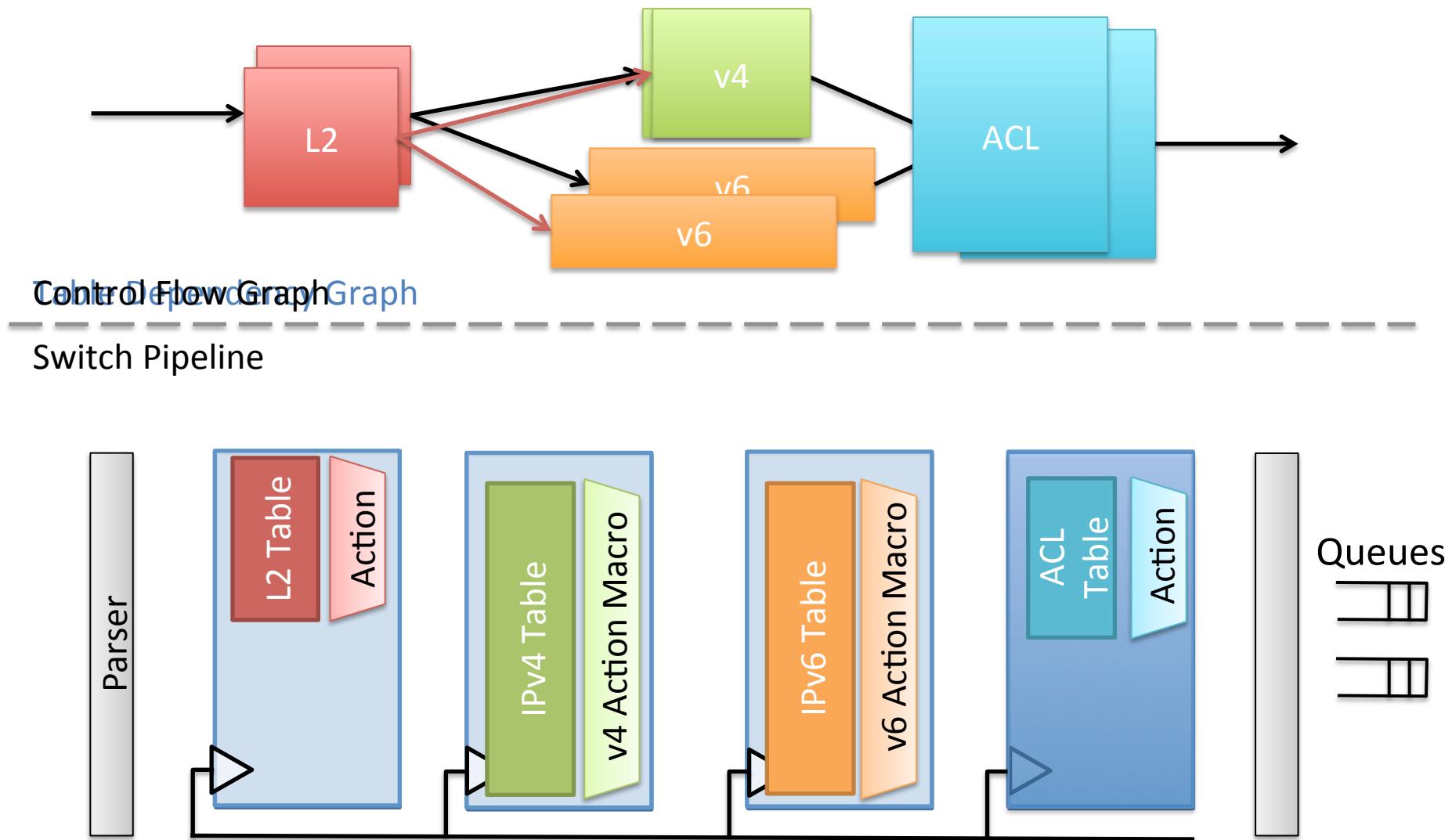


Control Flow Graph

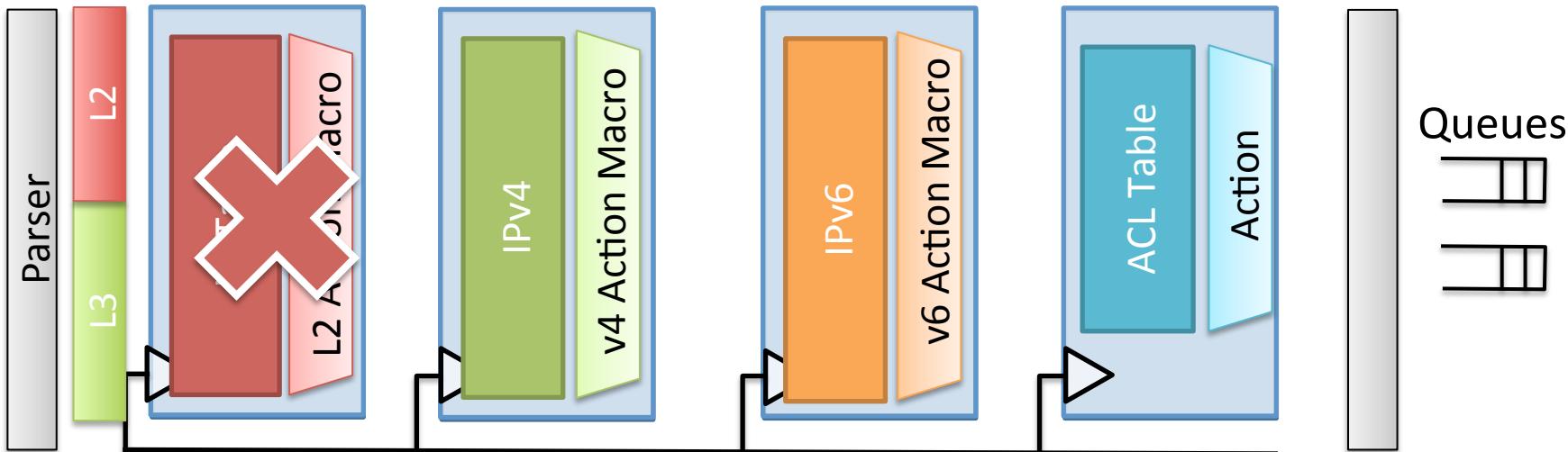
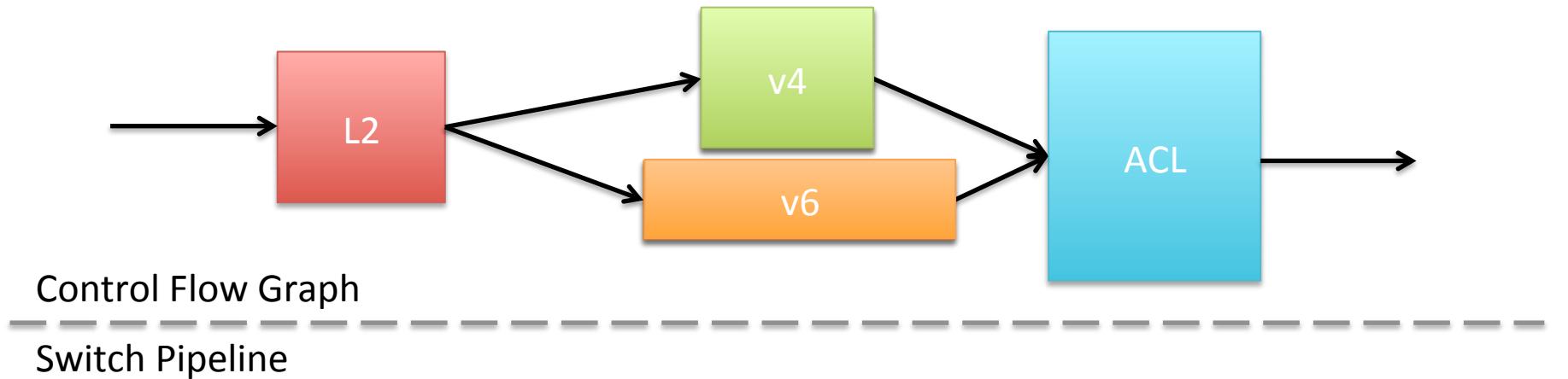
Table Dependency Graph



Efficient Mapping: TDG



Resource constraints



More resource constraints

Table parallelism

Action Memory

Memory Type

Action ALU input

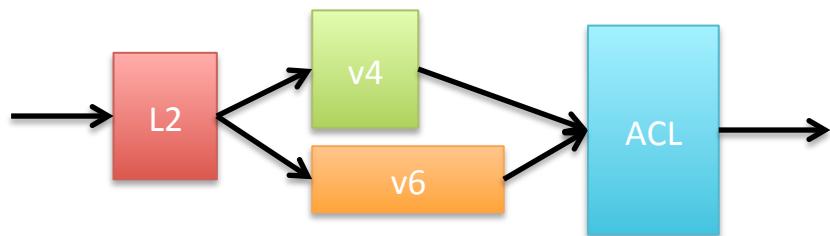
Header widths

The Compiler Problem

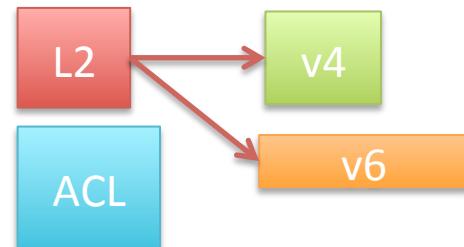
Map match action tables in a TDG to a switch pipeline while respecting dependency and resource constraints.

Step 1: P4 Program

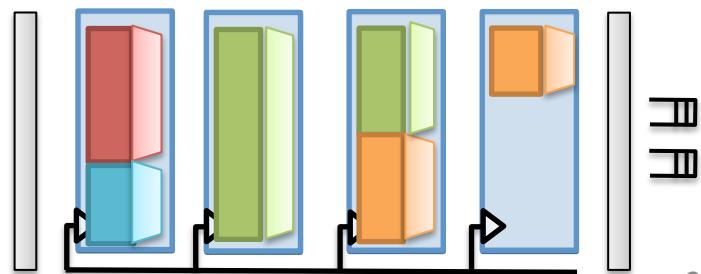
Step 2: Control Flow Graph



Step 3: Table Dependency Graph

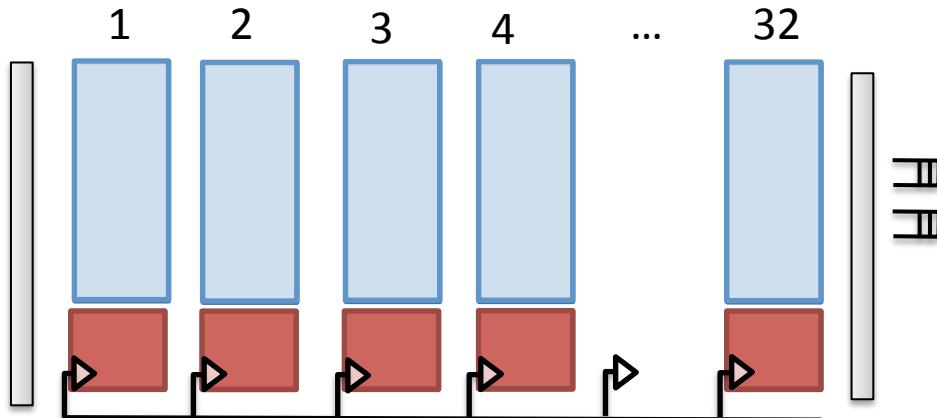


Step 4: Table Configuration



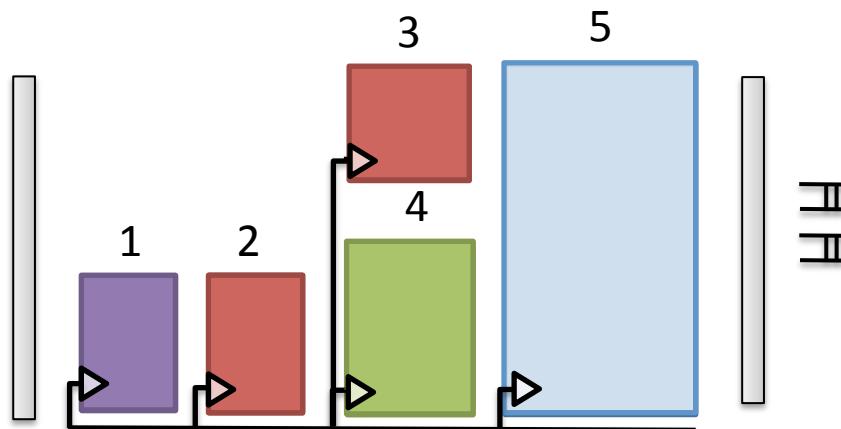
Is that it?

Two Switches We Studied



RMT
32 Stages
(SIGCOMM 2013)

FlexPipe
5 Stages
(Intel FM6000)



Additional switch features

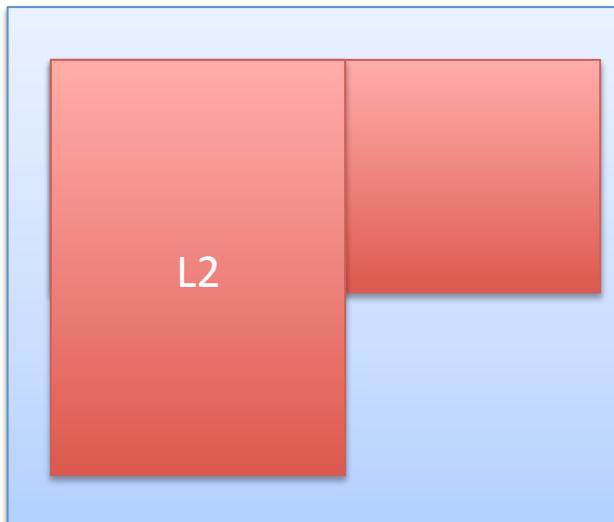
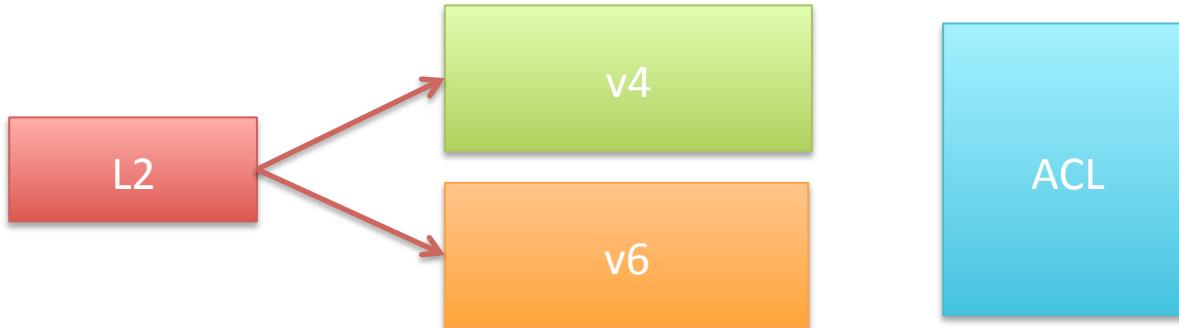


Table shaping in RMT

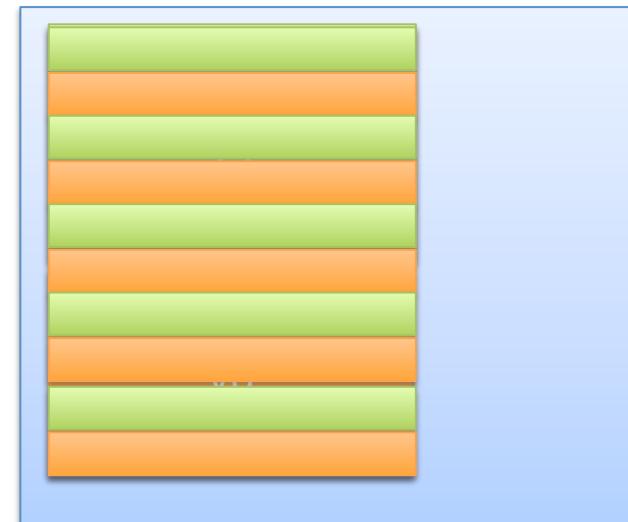


Table sharing in FlexPipe

Action Memory

Table parallelism

Memory Type

The Compiler Problem

Header widths

Action ALU input

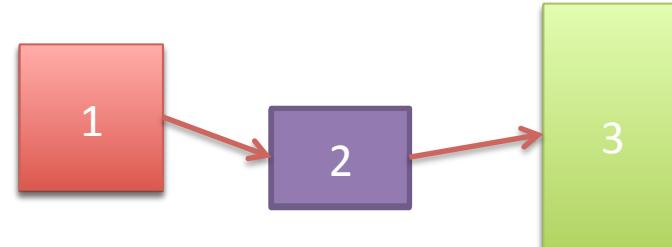
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Table shaping

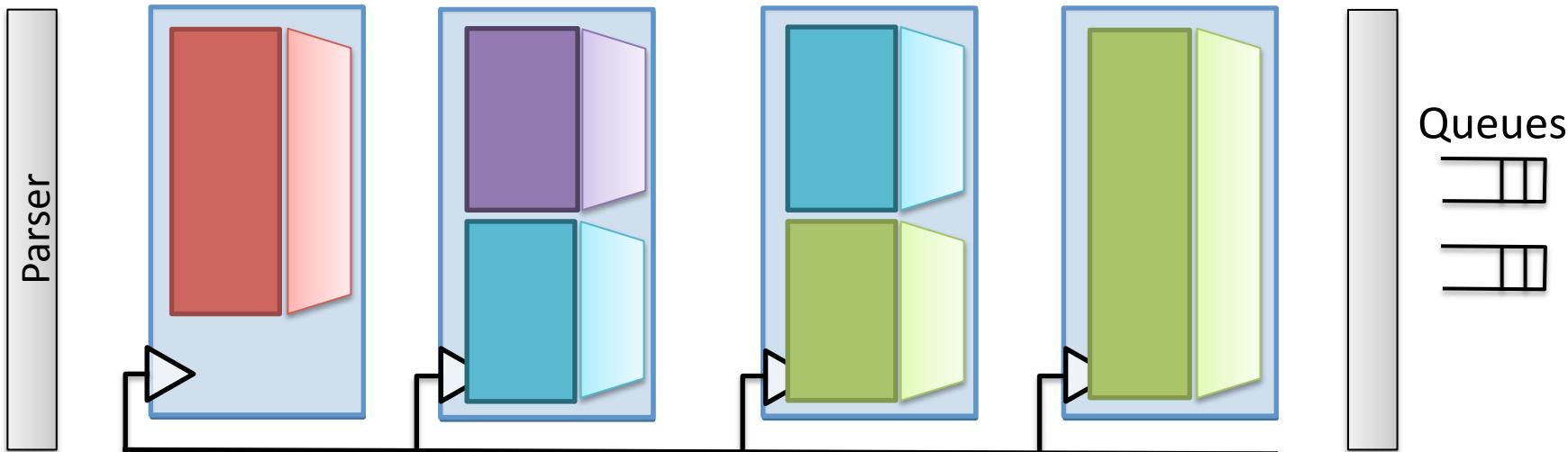
Table sharing

First approach: Greedy

- Prioritize one constraint
- Sort tables
- Map tables one at a time

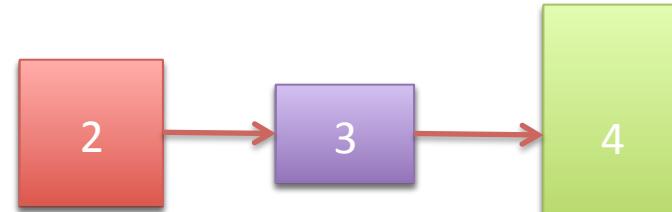


Sort by
dependencies

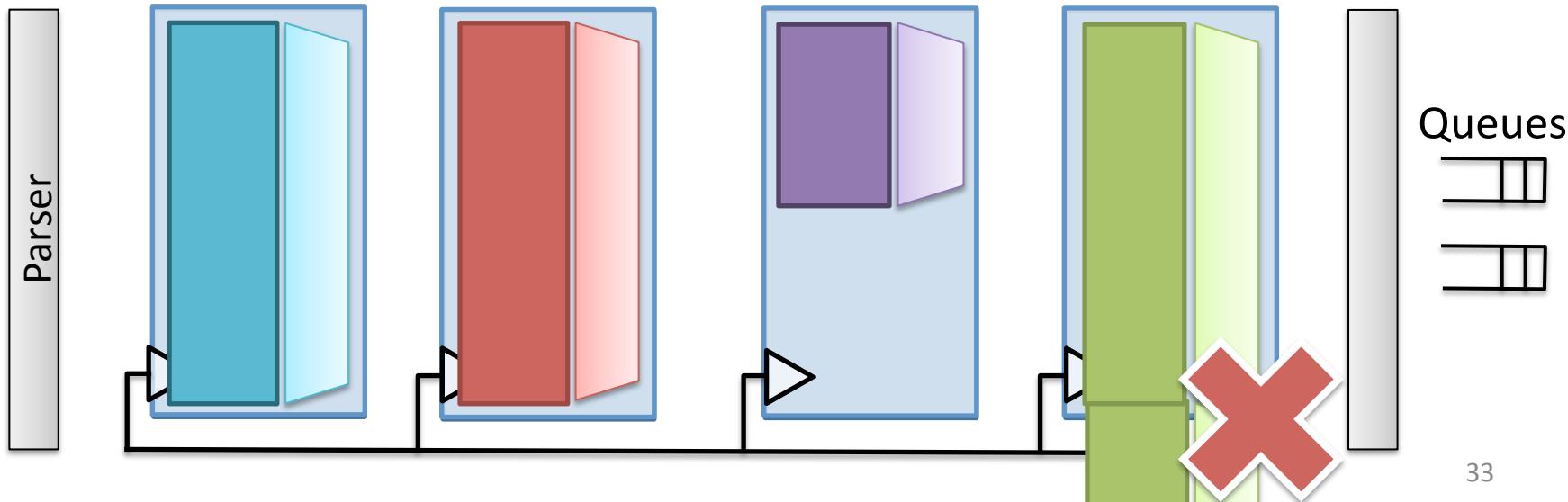


First approach: Greedy

- Prioritize one constraint
- Sort tables
- Map tables one at a time



Sort by
match width



Too many constraints for Greedy

- Any greedy must sort tables based on a metric that is a *fixed* function of constraints.
- As the number of constraints gets larger, it's harder for a fixed function to represent the interplay between all constraints.
- Can we do better than greedy?

Second approach: Integer Linear Programming (ILP)

Find an optimal mapping.

Pros:

- Takes in all constraints
- Different objectives
- Solvers exist (CPLEX)

Cons:

- Blackbox solver
- Encoding is an art
- Slow

ILP Setup

min # stages

subject to:

$$\begin{array}{ll} \text{table sizes} & \geq \text{table sizes} \\ \text{assigned} & \text{specified} \\ \\ \text{memories} & \leq \text{memories in} \\ \text{assigned} & \text{physical stage} \end{array}$$

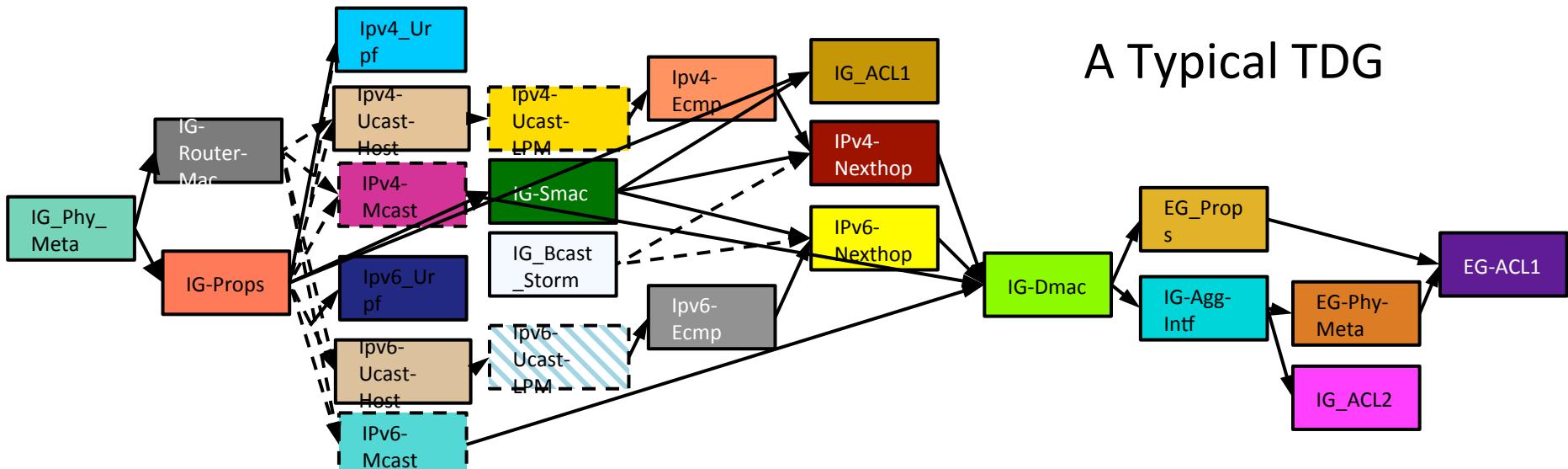
dependency constraints

Experiment Setup

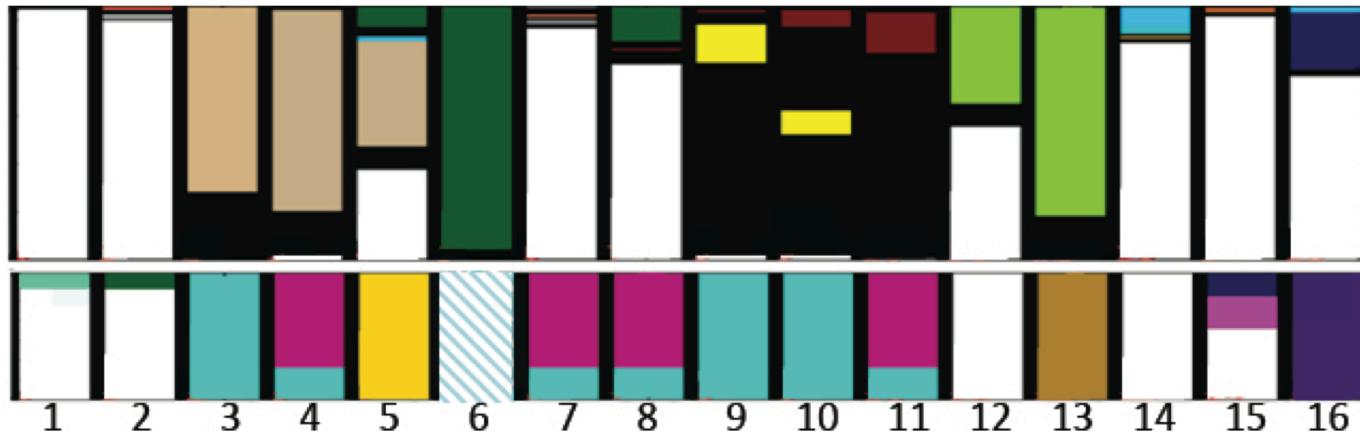
- 4 datacenter use cases from Intel, Barefoot
- Differ in tables, table sizes, and dependencies

Example Use Case

A Typical TDG



Configuration for RMT



Metrics: Greedy vs ILP

1. Ability to fit program in chip
2. Optimality
3. Runtime

Setup: Greedy vs ILP

1. Ability to fit: FlexPipe
 - Variants of use cases in 5-stage pipeline.
2. Optimality: RMT
 - Minimum stage, pipeline latency, power
3. Runtime: both switches

Results: Greedy vs ILP

1. Can Greedy fit my program?
 - Yes, if resources aplenty (RMT, 32 stages)
 - No, if resources constrained (FlexPipe, 5 stages),
Can't fit 25% of programs.
2. How close to optimal is Greedy?
 - 30% more time for packet to get through RMT pipeline.
3. Hmm.. looks like I need ILP. How slow is it?
 - 100x slower than Greedy
 - Reasonable if programs don't change often.

If we have time,
we should run ILP.

Use ILP to suggest best Greedy for program type.

Critical constraints

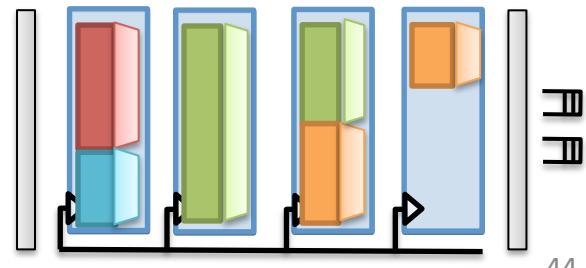
- Dependency critical: $16 \rightarrow 13$ stages
- Additional resource constraints less important

Critical resources

- TCAM memories critical: $16 \rightarrow 14$ stages
 - Results for one of our datacenter L2/L3 use cases

Conclusion

- **Challenge:** Parallelism and constraints in reconfigurable chips makes compiling difficult.
- **TDG:** highlights parallelism in program.
- **ILP:** better if enough time, fitting is critical, or objectives are complicated.
- **Best Greedy:** ILP can choose via notion of *critical* constraints and *critical* resources.



Thank you!

Research funded by AT&T, Intel, Open Networking Research Center.

ILP Run time

- Number of constraints? Not obvious. E.g., RMT
 - Min. stage: few secs.
 - Min. power: few secs.
 - Min. pipeline latency 10x slower
- Number of variables? How fine-grained is the resource assignment? E.g., FlexPipe
 - One match entry at a time: many days..
 - 100-500 match entries at a time: < 1 hr