

Sorting Algorithms

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Bubble Sort

About

Bubble Sort is the simplest sorting algorithm that works by repeatedly swapping the adjacent elements if they are in the wrong order.

Example

First Pass:

(**5** 1 4 2 8) \rightarrow (**1** **5** 4 2 8), Here, the algorithm compares the first two elements and swaps since $5 > 1$.

(1 **5** 4 2 8) \rightarrow (1 **4** **5** 2 8), Swap since $5 > 4$

(1 4 **5** 2 8) \rightarrow (1 4 **2** **5** 8), Swap since $5 > 2$

(1 4 2 **5** 8) \rightarrow (1 4 2 **5** 8), Now, since these elements are already in order ($8 > 5$), the algorithm does not swap them.

Second Pass:

(**1** 4 2 5 8) \rightarrow (**1** 4 2 5 8)

(1 **4** 2 5 8) \rightarrow (1 **2** **4** 5 8), Swap since $4 > 2$

(1 2 **4** 5 8) \rightarrow (1 2 **4** 5 8)

(1 2 4 **5** 8) \rightarrow (1 2 4 **5** 8)

Now, the array is already sorted, but our algorithm does not know if it is completed. The algorithm needs one **whole** pass without **any** swap to know it is sorted.

Third Pass:

(**1** 2 4 5 8) \rightarrow (**1** 2 4 5 8)

(1 **2** 4 5 8) \rightarrow (1 **2** 4 5 8)

(1 2 **4** 5 8) \rightarrow (1 2 **4** 5 8)

(1 2 4 **5** 8) \rightarrow (1 2 4 **5** 8)

i = 0	j	0	1	2	3	4	5	6	7
	0	5	3	1	9	8	2	4	7
	1	3	5	1	9	8	2	4	7
	2	3	1	5	9	8	2	4	7
	3	3	1	5	9	8	2	4	7
	4	3	1	5	8	9	2	4	7
	5	3	1	5	8	2	9	4	7
	6	3	1	5	8	2	4	9	7
i = 1	j	0	1	2	3	4	5	6	7
	0	3	1	5	8	2	4	7	9
	1	1	3	5	8	2	4	7	
	2	1	3	5	8	2	4	7	
	3	1	3	5	8	2	4	7	
	4	1	3	5	2	8	4	7	
	5	1	3	5	2	4	8	7	
i = 2	j	0	1	2	3	4	5	6	7
	0	1	3	5	2	4	7	8	
	1	1	3	5	2	4	7		
	2	1	3	5	2	4	7		
	3	1	3	2	5	4	7		
	4	1	3	2	4	5	7		
i = 3	j	0	1	2	3	4	5	6	7
	0	1	3	2	4	5	7		
	1	1	3	2	4	5			
	2	1	2	3	4	5			
	3	1	2	3	4	5			
i = 4	j	0	1	2	3	4	5	6	7
	0	1	2	3	4	5			
	1	1	2	3	4				
	2	1	2	3	4				
i = 5	j	0	1	2	3	4	5	6	7
	0	1	2	3	4				
	1	1	2	3					
i = 6	j	0	1	2	3	4	5	6	7
	0	1	2	3					
	1		2						

Implementation -

<https://github.com/lavishabhambri/Weekly-Algo-Newsletter/blob/main/Sorting/Codes/BubbleSort.cpp>

Worst and Average Case Time Complexity - $O(n^2)$. The worst case occurs when the array is reverse sorted.

Best Case Time Complexity - $O(n)$. The best-case occurs when the array is already sorted.

Auxiliary Space - $O(1)$

Boundary Cases - Bubble sort takes minimum time (Order of n) when elements are already sorted.

Sorting In Place - Yes

Stable - Yes

Insertion Sort

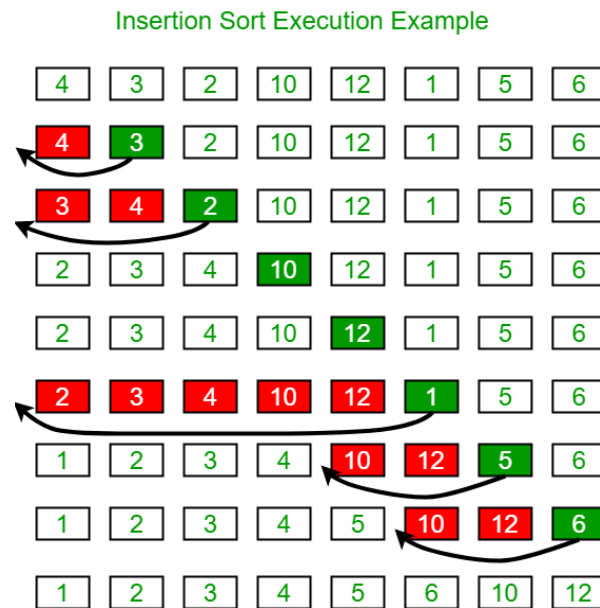
About

Insertion sort is a simple sorting algorithm that works similar to the way you sort playing cards in your hands. The array is virtually split into a sorted and an unsorted part. Values from the unsorted part are picked and placed at the correct position in the sorted part.

Algorithm

To sort an array of size n in ascending order:

- 1: Iterate from $\text{arr}[1]$ to $\text{arr}[n]$ over the array.
 - 2: Compare the current element (key) to its predecessor.
 - 3: If the key element is smaller than its predecessor, compare it to the elements before.
- Move the greater elements one position up to make space for the swapped element.



Implementation -

<https://github.com/lavishabhambri/Weekly-Algo-Newsletter/blob/main/Sorting/Codes/InsertionSort.cpp>

Time Complexity - $O(n^2)$

Auxiliary Space - $O(1)$

Boundary Cases - Insertion sort takes maximum time to sort if elements are sorted in reverse order. And it takes minimum time (Order of n) when elements are already sorted.

Algorithmic Paradigm - Incremental Approach

Sorting In Place - Yes

Stable - Yes