



# Mathematical Reasoning & PMI

## Contrapositive and Converse Statements

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- The contrapositive of a statement  $p \Rightarrow q$  is the statement  $\neg q \Rightarrow \neg p$
- The converse of a statement  $p \Rightarrow q$  is the statement  $q \Rightarrow p$

## Compound Statement

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Many mathematical statements are obtained by combining one or more statements using some connecting words like "and", "or" etc. those statement are called a "Compound Statement."

## Sentence

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A sentence is called a mathematically acceptable statement if it is either true or false but not both.

## Negation

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A statement which is formed by changing the true value of a given statement by using the word like 'no', 'not' is called negation of given statement.

## Implications

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- These are statements with word "if then", "only if" and "if and only if". if p then q is the same as following:
- p implies q is denoted by  $p \Rightarrow q$ , then symbol  $\Rightarrow$  stands for implies
  - p is a sufficient condition for q. then symbol  $\Rightarrow$
  - p only if q • q is a necessary condition for p
  - $\neg q$  implies  $\neg p$

## Truth Table for Logical Operations

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Conjunction operation:

p	q	$p \wedge q$
T	T	T
T	F	F
F	T	F
F	F	F

Disjunction operation:

p	q	$p \vee q$
T	T	T
T	F	T
F	T	T
F	F	F

Negation operation:

p	$\neg p$
T	F
T	F
F	T
F	T

Implication operation:

p	q	$p \Rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

Biconditional operation:

p	q	$p \Leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

## Additional Important Points

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- $p \Rightarrow q = \neg p \vee q$
- $\neg (p \Rightarrow q) = \neg (\neg p \vee q) = p \wedge (\neg q)$
- $p \Leftrightarrow q = (p \Rightarrow q) \wedge (q \Rightarrow p)$
- $\neg (p \Leftrightarrow q) = (p \wedge \neg q) \vee (q \wedge \neg p)$
- $(p \Leftrightarrow q) \Leftrightarrow r = p \Leftrightarrow (q \Leftrightarrow r)$

## General Logical Equivalences

It comprises the following laws:

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Idempotent Law

1.  $p \vee p \Leftrightarrow p$
2.  $p \wedge p \Leftrightarrow p$

Associative Law

1.  $(p \vee q) \vee r \Leftrightarrow p \vee (q \vee r)$
2.  $(p \wedge q) \wedge r \Leftrightarrow p \wedge (q \wedge r)$

Commutative Law

1.  $p \vee q \Leftrightarrow q \vee p$
2.  $p \wedge q \Leftrightarrow q \wedge p$

Distributive Law

1.  $p \vee (q \wedge r) \Leftrightarrow (p \vee q) \wedge (p \vee r)$
2.  $p \wedge (q \vee r) \Leftrightarrow (p \wedge q) \vee (p \wedge r)$

Identity Laws

1.  $p \wedge T \Leftrightarrow p$
2.  $p \vee F \Leftrightarrow p$
3.  $p \wedge F \Leftrightarrow F$

Complement Laws

1.  $p \vee \neg p \Leftrightarrow T$
2.  $p \wedge \neg p \Leftrightarrow F$
3.  $\neg T \Leftrightarrow F$
4.  $\neg F \Leftrightarrow T$

Absorption Law

1.  $p \vee (p \wedge q) \Leftrightarrow p$
2.  $p \wedge (p \vee q) \Leftrightarrow p$

Involution Law

1.  $p \sim (\neg p) \Leftrightarrow p$

De-Morgan's Law

1.  $q \sim (p \vee q) \Leftrightarrow \neg p \wedge \neg q$
2.  $q \sim (p \wedge q) \Leftrightarrow \neg p \vee \neg q$

## Tautology and Fallacy

- A tautology asserts that every possible interpretation has only one output, namely true.
- Fallacy implies an assertion of false in every possible interpretation.

**NOTE:** To evaluate tautology and fallacy, we can adapt the concept of the truth table that includes every possible valuation.

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p	q	$p \Rightarrow q$	$q \Rightarrow p$	$(p \Rightarrow q) \vee (q \Rightarrow p)$	$\sim \{(p \Rightarrow q) \vee (q \Rightarrow p)\}$
T	T	T	T	T	F
T	F	F	T	T	F
F	T	T	F	T	F
F	F	T	T	T	F

## Principle of Mathematical Induction

**Base Case:** The given statement is correct for first natural number that is, for  $n=1$ ,  $p(1)$  is true.

**Inductive Step:** If the given statement is true for any natural number like  $n=k$  then it will be correct for  $n = k + 1$  also that is, if  $p(k)$  is true then  $p(k + 1)$  will also be true.

The first principle of mathematical induction says that if both the above steps are proven then  $p(n)$  is true for all natural numbers.

