First Order Logic of Sparse Random Graphs

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Preliminaries

First order language of graphs

- Variables x_1, \ldots, x_n, \ldots
- Connectives ∧, ∨, equality symbol
 =, and negation symbol ¬.
- Quantifiers ∀, ∃.
- A binary relation symbol R

- Vertices.
- "And", "or", "equals", "not".
- "For all", "there exists".
- Edges *x* ∼ *y*.

$$\forall x_1, x_2 R(x_1, x_2) \implies \exists x_3 \neg (x_3 = x_1) \land \neg (x_3 = x_1) \land R(x_1, x_3)$$

Preliminaries

Binomial model of random graphs G(n, p).

Start with vertex set $[n] = \{1, ..., n\}$ and each edge is added with probability p independently.

If
$$G = ([n], E)$$
:
$$Pr(G) = p^{|E|} \cdot (1 - p)^{\binom{n}{2} - |E|}.$$

The problem

For any F.O. sentence ϕ we can talk about

$$Pr(G(n, p) \models \phi),$$

And if p = p(n) it also makes sense

$$\lim_{n\to\infty} Pr(G(n,p(n)) \models \phi).$$

What do we know about these limits?

The problem

For $p(n) = n^{-\alpha}$ and $\alpha \in [0,1]$ we have a nearly complete answer.

Theorem

(Fagin 19776) If $0 \le p \le 1$ is constant a Zero-One Law is satisfied

Theorem

(Spencer, Shelah 1988)

- If $\alpha \in (0,1) \setminus \mathbb{Q}$ then a Zero-One law holds.
- If $\alpha \in (0,1) \cup \mathbb{Q}$ then a Zero-One lay does not hold. Even more, there is a F.O. sentence ϕ such that

$$\lim_{n\to\infty} Pr(G(n,p(n)) \models \phi)$$

does not exist.

The problem

Theorem

(Lynch, 1992). If p(n) = c/n for some positive real c the limit

$$\lim_{n\to\infty} \Pr(G(n,p(n)) \models \phi)$$

always exists and is analytic in c.

The landscape of G(n, c/n)

We are interested in last theorem.

- The number of cycles of length 3,4..., r are asymptotically distributed like independent Poisson variables.
- Small cycles are a.a.s far away.
- Fixed vertices are a.a.s far away.
- The ball of a given radius centered in fixed vertex is a.a.s a tree. Any tree occurs with a positive probability.

For any quantifier rank k we show (constructively) that there exist some classes of graphs $C_1, ..., C_m$ such that

- (1) A.a.s any two random graphs in the same class have the same rank k type.
- (2) A.a.s. a random graph belongs to some C_i .
- (3) $\lim_{n\to\infty} Pr(G\in C_i)$ exists for any i and is analytic in c.

After this, for any F.O. sentence ϕ we obtain

$$\lim_{n\to\infty} Pr(\phi) = \sum_{i=1}^{m} \lim_{n\to\infty} Pr(\phi \wedge G \in C_i) + \lim_{n\to\infty} Pr(\phi \wedge G \notin C_1 \cup \cdots \cup C_l)$$

$$\stackrel{(2)}{=} \sum_{i=1}^{m} \lim_{n\to\infty} Pr(\phi \mid G \in C_i) Pr(G \in C_i)$$

$$\stackrel{(1)}{=} \sum_{j=1}^{l} \lim_{n\to\infty} Pr(G \in C_{i_j})$$

And finally because of (3) this limit exists and is analytic in c.

The classes C_1, \ldots, C_m are obtained as the equivalence classes of a relation on graphs that **only** depends on the "small" neighborhoods of the small cycles.

$$x \in Core(G, r) \iff d(x, c) \le r$$
, for some cycle c of length $\le 2r + 1$

- The probability of any given "class" of core is given by a nested sequence of products of "Poisson expressions, analytic on c.
- To show that graphs with "equivalent cores share the same rank k type E.F games are used.

Can this be generalized to regular hypergraphs with edge size a y edge probability c/n^{a-1} ?

Yes it seems so. You can even consider arbitrary symmetry groups on the edges and multiple edge sets with different sizes.