

## Lecture 4: N-Player Games & Nash Theorem

### Last Lecture Recap

- 2 Players zero sum finite game
- 2 Players non-zero sum finite game
- Computing all NE graphically in  $(2 \times 2)$  games  $\Rightarrow$  introduce  $(A, B)$

### Game Elements in This Lecture

**Players**  $\mathcal{I} = \{1, \dots, i, \dots, N\}$  with  $N \geq 2$  (**generalization of the  $N = 2$  case**).

**Strategy** Mixed strategies  $\mathbf{x}_i \in \Delta_i$ , joint profile

$$\mathbf{x} = (\mathbf{x}_1, \dots, \mathbf{x}_N) \in \Delta_X := \prod_{i \in \mathcal{I}} \Delta_i$$

**Cost**  $\bar{J}_i : \Delta_X \rightarrow \mathbb{R}$  for each  $i \in \mathcal{I}$

**Why This Change?** Many real applications involve  $> 2$  players (markets, networks, etc.).

**Central Question:** Does NE exist for  $N$ -player games?

**Key Theorems:**

1. **Brouwer Fixed-Point Theorem:** Continuous  $f : K \rightarrow K$  on compact convex  $K$  has fixed point
2. **Kakutani Fixed-Point Theorem:** Extends to set-valued maps (for BR correspondence)
3. **Nash Theorem:** Every  $N$ -player finite game has at least one mixed-strategy NE

### 4.1 N-Player Finite ("Matrix") Games (Ch. 3.3)

#### 4.1.1 Setup and Nash equilibrium (NE) definition

**Player Set.**  $\mathcal{I} = \{1, \dots, i, \dots, N\}$  with  $N \geq 2$

**Action Set.**

- $\Omega_i$  = set of actions of  $P_i$   $\forall i \in \mathcal{I}$
- $|\Omega_i| = m_i$  (cardinality)
- $M_i$  = index set of actions =  $\{1, \dots, j, \dots, m_i\}$

**Mixed Strategy Space.**

- $\Delta_i$ : set of mixed strategies of  $P_i$ .
- Mixed strategy of  $P_i$ :

$$\mathbf{x}_i = \begin{bmatrix} x_{i1} \\ \vdots \\ x_{ij} \\ \vdots \\ x_{im_i} \end{bmatrix}, \quad x_{ij} = \Pr(P_i \text{ selects action } j).$$

- Definition (probability simplex):

$$\Delta_i := \underbrace{\left\{ \mathbf{x}_i \in \mathbb{R}^{m_i} \mid \sum_{j=1}^{m_i} x_{ij} = 1, x_{ij} \geq 0 \forall j \in M_i \right\}}_{\text{probability simplex in } \mathbb{R}^{m_i}}.$$

- Joint mixed-strategy profile:

$$\mathbf{x} := (\mathbf{x}_1, \dots, \mathbf{x}_N) \in \underbrace{\Delta_1 \times \dots \times \Delta_N}_{\Delta_X \text{ (Cartesian product)}}.$$

- Compact notation for “player  $i$  vs others  $-i$ ”:

$$\mathbf{x} = (\mathbf{x}_i, \underbrace{\mathbf{x}_{-i}}_{\text{all players except } P_i}) := (\mathbf{x}_1, \dots, \mathbf{x}_{i-1}, \mathbf{x}_{i+1}, \dots, \mathbf{x}_N).$$

### Cost Functions.

- **Pure cost:**  $J_i : \Omega_1 \times \dots \times \Omega_N \rightarrow \mathbb{R}$
- **Expected cost:**  $\bar{J}_i(\mathbf{x}_i, \mathbf{x}_{-i}) : \Delta_1 \times \dots \times \Delta_N \rightarrow \mathbb{R}$

$$\begin{aligned} \bar{J}_i(\mathbf{x}) &= \mathbb{E}[J_i(\mathbf{u})] \\ &= \underbrace{\sum_{u_1 \in \Omega_1} \dots \sum_{u_N \in \Omega_N}}_{N\text{-fold sum}} J_i(u_1, \dots, u_N) \quad \underbrace{x(\mathbf{u})}_{\text{prob. of pure profile } \mathbf{u}} \end{aligned}$$

**Note.**  $\bar{J}_i$  is **multilinear** in  $(\mathbf{x}_1, \dots, \mathbf{x}_N)$  and **continuous** on  $\Delta_X$ .  
For fixed  $\mathbf{x}_{-i}$ , the map  $\mathbf{x}_i \mapsto \bar{J}_i(\mathbf{x}_i, \mathbf{x}_{-i})$  is **linear** on  $\Delta_i$ .

**Def. 3.12 Nash Equilibrium.** Nash Equilibrium (NE) is an N-tuple  $\mathbf{x}^* = (\mathbf{x}_1^*, \dots, \mathbf{x}_N^*) \in \Delta_X$

**Interpretation:** No regret solution for each  $P_i$  (individually optimal)

- 1 **No regret (unilaterally) solution for each  $P_i$**

$$\bar{J}_i(\mathbf{x}_i^*, \mathbf{x}_{-i}^*) \leq \bar{J}_i(\mathbf{x}_i, \mathbf{x}_{-i}^*) \quad \forall \mathbf{x}_i \in \Delta_i, \quad \underbrace{\forall i \in \mathcal{I}}_{N \text{ inequalities to be satisfied}}$$

- 2 **Individually Optimal** Best Response of  $P_i$ :  $\text{BR}_i : \Delta_{-i} \rightrightarrows \Delta_i$  (set-valued map)

$$\mathbf{x}_i^* \in \underbrace{\arg \min_{\mathbf{x}_i \in \Delta_i} \bar{J}_i(\mathbf{x}_i, \mathbf{x}_{-i}^*)}_{\text{usually a set}} =: \text{BR}_i(\mathbf{x}_{-i}^*)$$

- 3 **Fixed-point**

$$\underbrace{\begin{cases} \mathbf{x}_1^* \in \text{BR}_1(\mathbf{x}_{-1}^*) \\ \vdots \\ \mathbf{x}_i^* \in \text{BR}_i(\mathbf{x}_{-i}^*) \\ \vdots \\ \mathbf{x}_N^* \in \text{BR}_N(\mathbf{x}_{-N}^*) \end{cases}}_{\text{NE is the intersection of N individual BR maps}} \Leftrightarrow \underbrace{\mathbf{x}^* \in \text{BR}(\mathbf{x}^*)}_{\text{NE is a fixed point of the overall BR map}}$$

where

$$\text{BR}(\mathbf{x}) := \begin{bmatrix} \text{BR}_1(\mathbf{x}_{-1}) \\ \vdots \\ \text{BR}_N(\mathbf{x}_{-N}) \end{bmatrix} : \Delta_X \rightrightarrows \Delta_X$$

**Note.** Nash equilibrium is a mutual best response where no player can unilaterally improve by deviating. The fixed-point characterization connects NE existence to fixed-point theorems.

#### 4.1.2 Best-response maps and NE

**Def. Best Response Correspondence (pdf note 3.4.1).** For player  $i$  and fixed strategies of others  $\mathbf{x}_{-i} \in \Delta_{-i}$ , the **best response set** is:

$$BR_i(\mathbf{x}_{-i}) = \left\{ \mathbf{x}_i \in \Delta_i \mid \bar{J}_i(\mathbf{x}_i, \mathbf{x}_{-i}) \leq \bar{J}_i(\mathbf{w}_i, \mathbf{x}_{-i}), \forall \mathbf{w}_i \in \Delta_i \right\}$$

This is also called the **optimal response set** or **rational action set**.

**Overall BR map:**  $BR : \Delta_X \rightrightarrows \Delta_X$

$$BR(\mathbf{x}) = BR_1(\mathbf{x}_{-1}) \times \cdots \times BR_N(\mathbf{x}_{-N}) \subset \Delta_X$$

Properties of BR Correspondence For any  $\mathbf{x}_{-i} \in \Delta_{-i}$ :

1.  $BR_i(\mathbf{x}_{-i})$  is **non-empty** (by Weierstrass: continuous function on compact set)
2.  $BR_i(\mathbf{x}_{-i})$  is **convex** (by linearity of  $\bar{J}_i$  in  $\mathbf{x}_i$ )
3.  $BR_i(\mathbf{x}_{-i})$  is **closed** (by continuity)
4.  $BR_i(\mathbf{x}_{-i}) \subseteq \Delta_i$  is a face of  $\Delta_i$  (convex hull of some pure strategy vertices)
5. Can be set-valued: ranges from singleton to entire simplex  $\Delta_i$

#### 4.2 Nash Theorem and Idea of Proof (Ch. 3.5)

**Nash Theorem (existence of NE).** Every  $N$ -player finite game has at least one mixed-strategy Nash equilibrium.

##### 4.2.1 Two Fixed-Point Theorems (Appendix A)

**Brouwer's Fixed Point Theorem.**

- Let  $S \subset \mathbb{R}^m$  be a **convex** and **compact** subset.
- Let  $f : S \rightarrow S$  be a **continuous** function.
- Then  $\exists x \in S$  s.t.

$$x = f(x), \quad \text{i.e. } x \text{ is a fixed point of } f$$

**Kakutani's Fixed Point Theorem.**

- Let  $S \subset \mathbb{R}^m$  be a **convex** and **compact** subset.
- Let  $\Phi : S \rightrightarrows S$  be a **set-valued map**, and denote the image of  $x \in S$  by  $\Phi(x) \subset S$ .
- Assume that  $\forall x \in S$ :
  - (1)  $\Phi(x)$  is **non-empty** and **convex**;
  - (2)  $\Phi$  has a **closed graph**

$$\text{Graph}(\Phi) := \{(x, y) \in S \times S \mid y \in \Phi(x)\}$$

- Then  $\exists$  at least one  $x \in S$  s.t.

$$x \in \Phi(x), \quad \text{i.e. } x \text{ is a fixed point of } \Phi$$

### 4.2.2 Applying Kakutani's Theorem to BR

Nash Theorem can be proved by Kakutani's Fixed-Point Theorem.

**Step 1 — Figure out what is what.**

- **Set:**  $S \equiv \Delta_X := \Delta_1 \times \cdots \times \Delta_N$
- **Correspondence:**  $\Phi \equiv \text{BR} : \Delta_X \rightrightarrows \Delta_X$ ,  $\text{BR}(\mathbf{x}) := \text{BR}_1(\mathbf{x}_{-1}) \times \cdots \times \text{BR}_N(\mathbf{x}_{-N})$  where, for each  $i \in \mathcal{I}$ ,

$$\text{BR}_i(\mathbf{x}_{-i}) := \arg \min_{\mathbf{x}_i \in \Delta_i} \bar{J}_i(\mathbf{x}_i, \mathbf{x}_{-i}).$$

**Step 2 — What needs to be shown?**

- (i)  $\Delta_X$  is **convex** and **compact**;
- (ii) For every  $\mathbf{x} \in \Delta_X$ ,  $\text{BR}(\mathbf{x})$  is **non-empty** and **convex**;
- (iii)  $\text{Graph}(\text{BR})$  is **closed**, where

$$\text{Graph}(\text{BR}) := \{ (\mathbf{x}, \mathbf{y}) \in \Delta_X \times \Delta_X \mid \mathbf{y} \in \text{BR}(\mathbf{x}) \}.$$

**Step 3 — Check all the conditions.**

(i)  $\Delta_X$  is **convex** and **compact**.

$$\Delta_i := \{ \mathbf{x}_i \in \mathbb{R}^{m_i} \mid \mathbf{1}^\top \mathbf{x}_i = 1, x_{ij} \geq 0, \forall j \in M_i \}$$

**Convexity of  $\Delta_i$ .** Take  $\mathbf{x}_i, \mathbf{y}_i \in \Delta_i$  and  $\alpha \in [0, 1]$ , and let  $\mathbf{z}_i := \alpha \mathbf{x}_i + (1 - \alpha) \mathbf{y}_i$ .

Each component in  $\mathbf{z}_i$  is non-negative

$$z_{ij} = \alpha x_{ij} + (1 - \alpha) y_{ij} \geq 0, \quad \forall j \in M_i,$$

and the sum of the components in  $\mathbf{z}_i$  is 1.

$$\begin{aligned} \mathbf{1}^\top \mathbf{z}_i &= [1 \quad \cdots \quad 1] \begin{bmatrix} \alpha x_{i1} + (1 - \alpha) y_{i1} \\ \vdots \\ \alpha x_{im_i} + (1 - \alpha) y_{im_i} \end{bmatrix} = \sum_{j=1}^{m_i} (\alpha x_{ij} + (1 - \alpha) y_{ij}) = \alpha \sum_{j=1}^{m_i} x_{ij} + (1 - \alpha) \sum_{j=1}^{m_i} y_{ij} \\ &= \alpha \cdot 1 + (1 - \alpha) \cdot 1 = 1. \end{aligned}$$

$\mathbf{z}_i$  satisfies the defining constraints of  $\Delta_i$ , and hence  $\mathbf{z}_i \in \Delta_i \Rightarrow \Delta_i$  is convex.

**Compactness of  $\Delta_i$ .**

- Bounded:  $x_{ij} \geq 0$  and  $\sum_j x_{ij} = 1$  imply  $0 \leq x_{ij} \leq 1$ , so  $\Delta_i \subset [0, 1]^{m_i}$ .
- Closed: Let  $\{\mathbf{x}_i^{(n)}\} \subset \Delta_i$  with  $\mathbf{x}_i^{(n)} \rightarrow \mathbf{x}_i$ . Then  $x_{ij}^{(n)} \geq 0$  for all  $n$  implies  $x_{ij} = \lim_n x_{ij}^{(n)} \geq 0$ , and

$$\sum_j x_{ij} = \sum_j \lim_n x_{ij}^{(n)} = \lim_n \sum_j x_{ij}^{(n)} = \lim_n 1 = 1.$$

Thus  $\mathbf{x}_i \in \Delta_i$ , so  $\Delta_i$  is closed.

Therefore  $\Delta_i$  is closed and bounded, hence compact.

**From  $\Delta_i$  to  $\Delta_X$ .**

$$\Delta_X = \Delta_1 \times \cdots \times \Delta_N \subset \mathbb{R}^{m_1 + \cdots + m_N}.$$

- Cartesian product of convex sets is convex  $\Rightarrow \Delta_X$  is convex.
- Cartesian product of compact sets in finite-dimensional Euclidean space is compact  $\Rightarrow \Delta_X$  is compact.

Thus condition (i) holds.

(ii) **BR(x) is non-empty and convex.** Fix  $i \in \mathcal{I}$  and  $\mathbf{x}_{-i} \in \Delta_{-i}$ .

**Non-emptiness of  $\text{BR}_i(\mathbf{x}_{-i})$ .**

By Weierstrass theorem, a continuous function on a compact set attains its minimum. Since  $\bar{J}_i(\cdot, \mathbf{x}_{-i})$  is continuous and  $\Delta_i$  is compact, the minimizer set  $\text{BR}_i(\mathbf{x}_{-i})$  is non-empty.

**Convexity of  $\text{BR}_i(\mathbf{x}_{-i})$ .**

Suppose  $\mathbf{y}, \mathbf{z} \in \text{BR}_i(\mathbf{x}_{-i})$  and let  $\alpha \in [0, 1]$ . Then for any  $\mathbf{w} \in \Delta_i$ :

$$\bar{J}_i(\alpha \mathbf{y} + (1 - \alpha) \mathbf{z}, \mathbf{x}_{-i}) = \alpha \bar{J}_i(\mathbf{y}, \mathbf{x}_{-i}) + (1 - \alpha) \bar{J}_i(\mathbf{z}, \mathbf{x}_{-i}) \leq \bar{J}_i(\mathbf{w}, \mathbf{x}_{-i})$$

by linearity of  $\bar{J}_i$  in  $\mathbf{x}_i$ . Hence  $\alpha \mathbf{y} + (1 - \alpha) \mathbf{z} \in \text{BR}_i(\mathbf{x}_{-i})$ .

Since  $\text{BR}(\mathbf{x}) = \prod_i \text{BR}_i(\mathbf{x}_{-i})$  is a product of non-empty convex sets, it is non-empty and convex.

(iii) **Graph(BR) is closed (next lecture).**

To show that  $\text{Graph}(\text{BR})$  has a graph set that is closed, i.e. it contains all its limit points.

We prove by contradiction: assume that  $\text{Graph}(\text{BR})$  is not closed, i.e., it does not contain all its limit points.

#### Real analysis & topology detour.

- **Compact set (in  $\mathbb{R}^m$ ).**  $S$  is compact  $\iff S$  is closed and bounded (Heine–Borel).
- **Convex set.**  $S$  is convex if  $\forall x, y \in S$  and  $\forall \alpha \in [0, 1]$ ,

$$\alpha x + (1 - \alpha)y \in S.$$

- **Closed set.**  $S$  is closed iff for every sequence  $\{x^{(n)}\} \subset S$  with  $x^{(n)} \rightarrow x$ , we have  $x \in S$  (i.e.,  $S$  contains all its limit points).
- **Graph of a correspondence.** For  $\Phi : S \rightrightarrows S$ ,

$$\text{Graph}(\Phi) := \{(x, y) \in S \times S \mid y \in \Phi(x)\}.$$