

The weights of leaves in an optimal tree are either $w-1$, w , $w+2$ or w , $w+2$, $w+3$

b) A node of cost w is derived from nodes of cost $w-1$ and $w-3$ so $M_w = M_{w-1} + M_{w-3}$

② $T(n) = 2T(n-1) + T(n-2) \Rightarrow r^2 = 2r + 1$
 $r^2 - 2r - 1 = \frac{2 \pm 2\sqrt{2}}{2} = 1 \pm \sqrt{2}$ $T(n) = a(1+\sqrt{2})^n + b(1-\sqrt{2})^n$

ii) Homogeneous root is $r=4$ $T(n) = c_1 4^n + c_2 2^n$
 $T(0) = c_1 4^0 + c_2 2^0 = 1$
 $T(1) = 4T(0) + 2 = 6$ $T(n) = 2 \cdot 4^n - 2^n$
 $c_1 + c_2 = 1$ $c_1 = 2$
 $4c_1 + 2c_2 = 6$ $c_2 = -1$

iii) $r^2 - 8r + 16 = 0$ $(r-4)^2 = 0$ $r_1 = r_2 = 2$
 $T(n) = c_1 n 2^n + c_2 2^n$ $T(0) = c_2 = 2$
 $T(1) = c_1 2 + 4 = 10$ $c_1 = 3$ $T(n) = 3n 2^n + 2 \cdot 2^n$

③ $T(n-c) = T(\frac{n}{2^c})$ $T(n) = T(\frac{n}{2^c}) + c + \frac{\sqrt{n}}{2^{c/2}}$
 $c + \frac{\sqrt{n}}{2^{c/2}}$ is subproblem size
 $c + \frac{\sqrt{n}}{2^{c/2}} = 1$
 $\sqrt{n} = 2^{c/2}(1-c) \Rightarrow \frac{\sqrt{n}}{1-c} = 2^{c/2}$
 $\left(\frac{\sqrt{n}}{1-c}\right)^{2/c} = 2 \Rightarrow \frac{2}{c} \lg\left(\frac{\sqrt{n}}{1-c}\right) = c$
 Tree has $\frac{2}{c} \lg\left(\frac{\sqrt{n}}{1-c}\right) + 1$ levels

$T(n) = (c + \sqrt{n}) + (c + \frac{\sqrt{n}}{2^{1/2}}) + \dots$
dominated by largest term

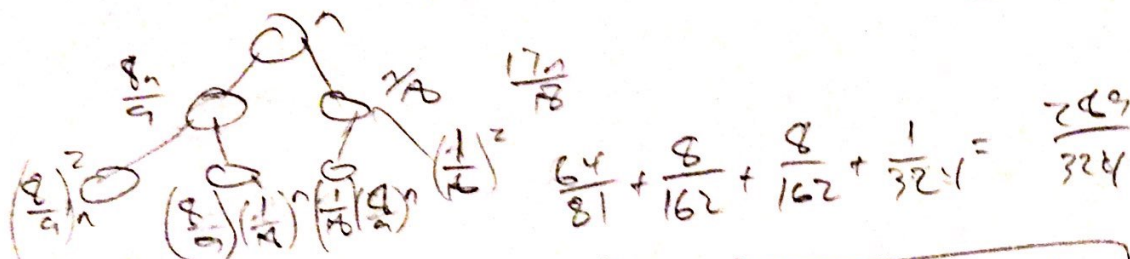
geometric series decreasing

$$T(n) = \Theta(\sqrt{n})$$

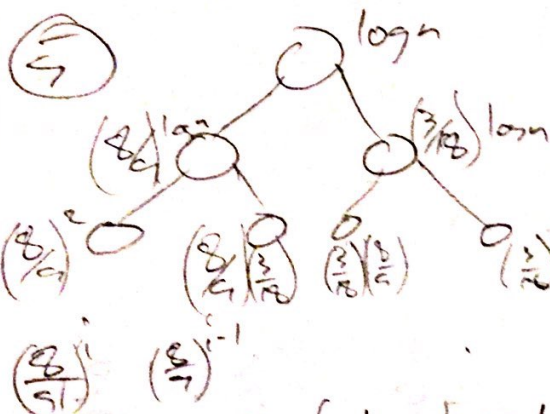
④ $T(n) = n + \frac{17n}{18} + \frac{289}{324}$

geometric series decreasing

$$T(n) = \Theta(n) \quad f(n) = n$$



$$T(n) = n^2 + \left(\frac{17}{18}\right)^2 n^2 + \left(\frac{289}{324}\right)^2 n^2 \quad T(n) = \Theta(n^2) \quad f(n) = n^2$$



$$\log n + \frac{19}{18} \log n + \frac{361}{324} \log n$$

Geometric series increasing

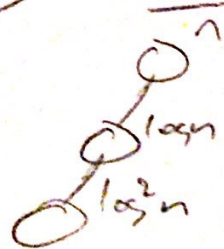
of levels = $\log_{9/8} \log n$

$$T(n) \text{ is dominated by sum at } \log_{9/8} \log n \text{ level} \quad f(n) = \log n$$

$$T(n) \text{ is dominated by sum at } 2 \log_{9/8} n \text{ level} \quad f(n) = n^2$$

⑤ a) $1 \leq i \leq 10 \quad T(i) = 10$

$11 \leq i \leq 1024 \quad T(i) = 4$



Decreasing geometric $T(n) = \Theta(n)$ if $n > 1024$

b) $1 \leq i \leq 10 \quad T(i) = 10$

$11 \leq i \leq 1024 \quad T(i) = \Theta(n)$

$$T(n) = \Theta(n) \text{ if } n > 1024$$

Again decreasing geometric series