

Data Cube Computation: 1) Sort, Aggr, & group tuples 2) Complete higher level aggregates from lower level & cache intermediate 3) Aggregate from child with smaller cardinality 4) Use Apriori to compute iceberg cubes. If a given cell does not meet min-supp neither do its descendants.

Multiseries: Partition into chunks, cell address by "chunkID + offset", compute agg by visiting cells, optimize visit order so ~~partitioning~~ partitioning of agg cells in multiple cuboids can be computed simultaneously. 3-D array of ABC, 2-D & 1-D cuboids must be computed. $\begin{matrix} B \\ \swarrow \downarrow \searrow \\ A \end{matrix}$ One BC chunk needs to be in memory at a

time, for computation of all BC chunks. Simultaneously aggregate for each 2-D plane while one 3-D chunk is in memory. Minimum memory for holding all 2-D planes in chunk memory is whole AB plane $(40 \cdot 100) + (40 \cdot 1000) + (100 \cdot 1000)$ when AB cuboid is being computed. A will have all its chunks in memory & B will have 1 chunk allocated at a time. Most effective when the product of cardinalities is moderate & data not too sparse. When dimensionality high or very sparse not good.

BUC: Starts at apex & drills down. Can use apriori and share data partitioning cost. For each dimension, partition on d, to count # of tuples for each distinct val of d. Each distinct val forms its own partition. Iterate each partition, test for min-supp, if it passes recursively call BUC on d+1 next dim. On return from recursive call, continue on to next partition for d. Traversal tree. Computing AB does not help compute ABC. Process high cardinality dimensions first. Less skew is better.

Star Cubing: Integrates top down & bottom up. Globally bottom up like Multiseries with layers of top down to use Apriori with aggregate values at internal nodes, you can prune based on shared dimensions. Replace nodes not satisfying min-supp with

* Multiseries vs BUC: Multiseries array vs. RDB, Full vs Iceberg, Small Aggr vs Part & sort. Can't ~~handle~~ handle big d.m. Iceberg cannot be incrementally updated. Once pruned it is lost.

Shell Fragments: Compute fragments offline & combine them to dynamically answer online queries. Create inverted index for each fragment which is TID list for each attribute value. Find intersection of TID lists. Frag cube space required is $O(T \cdot \frac{D}{F}) / (2^F - 1)$ F = size of fragments & d.dims.

FP Tree: Find frequent 1-item sets. Sort freq items in order of decreasing support count. Scan each transaction, process items in L order, create branches, increment count. For conditional database find all branches containing that item. This is the conditional pattern base. Make conditional FP-tree column from frequent members of pattern base. Generate frequent items from (conditional) FP-tree, this column doesn't have to link consecutive items.

Apriori: k-itemsets are used to find (k+1) itemsets. First find frequent 1 item sets by scanning the database. If an itemset is not frequent then none of its supersets are. To get (k+1) join k itemsets need first (k-1) in common. Partition to find patterns. Any globally frequent pattern must be frequent in 1 partition.

Vertical data format: Each item in its own row with TID list

Apriori is breadth first, FP is depth first. ~~FP is~~ FP says for frequent k-itemset restrict subsequent search to only items containing it. Mine conditional database recursively find frequent single items, ~~recursively generate FP~~ and partition DB based on each single item pattern, recursively scan for patterns by doing above for each partitioned DB. Mine recursively conditional until empty.

support = $X \cup Y$ Confidence $X \rightarrow Y$ support/support(X) = $P(Y|X)$. A pattern is closed if it is frequent & there exists no super pattern with same support. A pattern is a max pattern if it is frequent & there is no frequent super pattern.
 $lift(B, C) = \frac{c(B \rightarrow C)}{s(C)} = \frac{s(BC)}{s(B) \cdot s(C)}$ if it is 1 B & C are independent $\chi^2 = \sum \frac{O-E}{E}$

all transactions mess with lift & χ^2 . Jaccard & cosine are null invariant
 Geometric techniques: Scatter, Parallel Icon; Chernoff, Stalk figure, Hierarchical, TreeMap
 $\chi^2 = [0, \infty)$ Jaccard $\frac{a}{a+b+c}$ $\cos(i, j) = \frac{a}{a+b+c}$ asymmetric distance

Dimensionality Reduction: Wavelet transform, PCA, attr subset, attribute construction. Naive Bayes, regression, log-linear, cluster, sampling
 Normal distance: $d(i, j) = \frac{r - m}{p}$ Difference in text: cosine sim. $ubmax = \frac{x - min}{max - min}$ (nearest) nearest

$z-score = \frac{x - \mu}{\sigma}$

$(a_1, a_2, \dots, a_{10}) : 1, (a_1, b_2, a_3, b_4, b_5, b_6, b_7, b_8, a_9, b_{10}) : 1$ 2^{10} cuboids 3 closed cells: 2 base & $(a_1, *, a_3, *, *, \dots, a_9, *) : 2$

A cell is closed if no descendant has same count Aggregate cells? $2 \cdot 2^{10} - (2 \text{ base}) - 2^3 \text{ duplicate}$
 If min sup = 2 # of aggregate cells? 9 joining in common
 (gender, *) & (*, course) 400 courses, 2 gender, (E, C) & (M, C) need one space for course 2 for gender

$(a_1, a_2, \dots, a_{10}) : 1, (a_1, b_2, a_3, b_4, b_5, b_6, b_7, b_8, a_9, b_{10}) : 1, (a_1, b_2, a_3, b_4, b_5, b_6, b_7, b_8, a_9, b_{10}) : 1$ 2^3 cuboids combine 3 or less

only 1 cell for those 8, not 3 for each base cell + count once not 3 times, $2^3(3-1)$ dup
 $3 \cdot 2^{10} - 3 - 2^3(3-1)$

for 3 base cell $3 \cdot 2^{10}$, then find dimensions common to all 3^{10} and subtract $2 \cdot 2^{10}$, then those common to any 2 give (j, k, l) subtract $(2^j - 1) - (2^k - 1) - (2^l - 1)$

Sum of largest x values not bounded, holistic. Sum of largest ~~50~~ 50 values. In each partition maintain list of largest 50. Find largest 50 of those

$sum(S, price) \geq 45$ is monotonic: as long as any subset of S has price ≥ 45 $sum(S, price) \geq 45$
 $sum(S, price) \leq 45$ is antimonotonic as long as subset of S if any item doesn't satisfy neither can a-y superset

$avg(S, price) \geq 30$ can be converted to antimonotonic if items w/d in price descending order

conditional pattern base conditional F1 tree frequent patterns
 all paths are suffix frequent paths can merge to get them combinations include suffix
 no suffix

Manhattan $d(i, j) = |m - x_1| + |m - x_2| \dots$ Euclidean $= \sqrt{|m - f_1| + |m - f_2|}$ Supremum is max

KL divergence measures diff between 2 prob distributions no vars $\sum_{x \in X} p(x) \ln \frac{1}{p(x)}$
 using q to approx p

$p = \frac{cov(A, B)}{\sigma_A \sigma_B}$

Need closed patterns to recover frequent patterns with counts