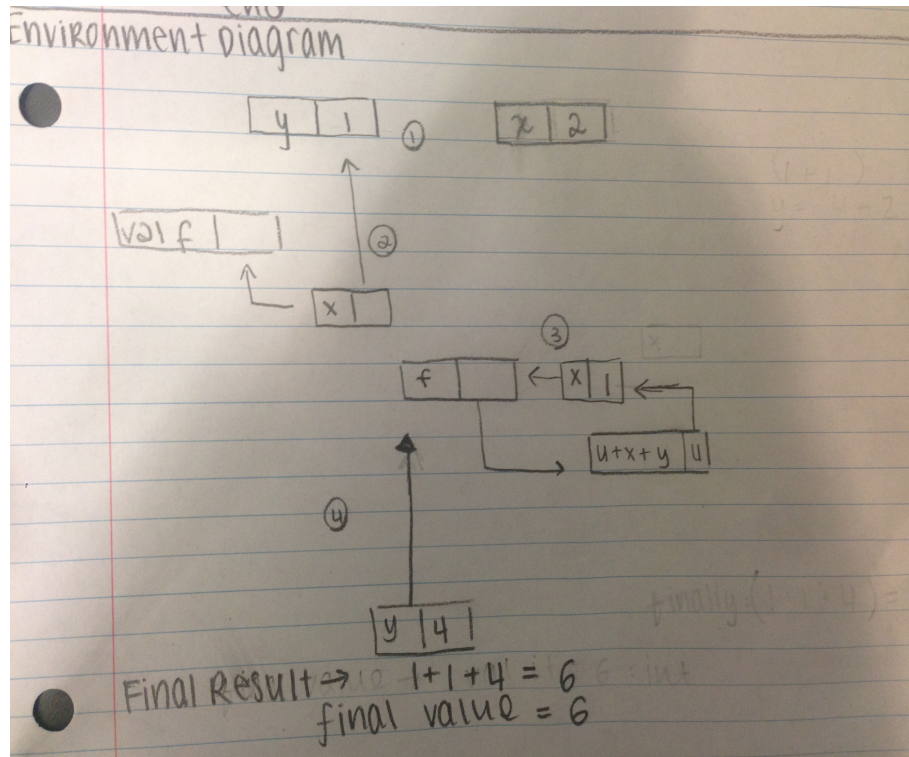


# COMP302 - Assignment 4

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Sourced from Stackoverflow and Wikipedia, Slides, Bridgette Pienka's book

## 1 Question 3



## 2 Question 4

**Specification** In this proof, we will consider:

```
(* reduce: 'a list * 'a * ('a * 'a -> 'a) -> 'a *)  
fun reduce (nil, base, op) = base  
| reduce (h::t, base, op) = op (h, reduce(t, base, op))
```

```

(* reduce tr: 'a list * 'a * ('a * 'a -> 'a) -> 'a *)
fun reducetr (nil, base, op) = base
| reduce tr (h::t, base, op) = reducetr(t, op(h, base), op)

```

**Proof** We will prove that the two below programs return the same result if applied to the same arguments. That is we

```

reduce (h::t, base, op) == reduce tr (h::t, base, op)

```

**Conditions** We will complete the proof using the lemma and the below conditions:

- (a)  $op(base, m) = m$
- (b)  $op(n, m) = op(m, n)$
- (c)  $op(n, op(m, k)) = op(op(n, m), k)$

**Lemma**  $op(h, reduce(l, n, op)) = reduce\ tr(l, op(h, n), op)$

**Initial Inductive Hypothesis** To prove the two give the same result with the same arguments, we need to consider the properties of a,b,c.

**Induction Step** Take  $k:: t = k:: h:: t$

```

op(k, reduce(h::k, base, op)) = reduce tr (t, op, h, n, op)
reduce tr (h::t, op(k, base, op)) -> reduce tr (t, op(h, op(k, base), op))

```

**Induction Expansion** By the induction hypothesis, we can further expand while considering the conditions:

```

op (h, reduce (t, op (k, base), op))
reduce (h::t, op (k, base), op)
op (h, op (t, op (k, base)))
op (k, op (h, op, (t, base)))
op (h, op (t, base)) x reduce (h::t, base, op)
op (k, reduce (h::t, base, op))

```

**Final Step**  $reduce(h::t, base, op) = reduce\ tr(t, op(h, base), op)$   
 $op(h, reduce(t, base, op)) = reduce\ tr(t, op(h, base), op)$   
 by the lemma this is true, completing the proof