

Learning in game - fictitious play

Introduction

- Most economic theory relies on equilibrium analysis based on Nash equilibrium or its refinements.
- The traditional explanation for when and why equilibrium arises is that it results from analysis and introspection by the players in a situation with all **common knowledge** on
 - the rules of the game
 - the rationality of the players
 - the payoff functions of players are all common knowledge.
- In this lecture, we develop an alternative explanation why equilibrium arises as the **long-run outcome of a process**

Fictitious play

- One of the earliest learning rules, introduced in Brown (1951), is the **fictitious play**.
- The most compelling interpretation of fictitious play is as a “**belief-based**” learning rule
 - players form beliefs about opponent play (from the entire history of past play) and behave rationally with respect to these beliefs.

Fictitious play-Setup

- We focus on a two player strategic form game $G = (\{1,2\}, S, u)$
- The players play this game at times $t = 1, 2, \dots$
- The stage payoff of player i is again given by $u_i(a_i, a_{-i})$ (for the pure strategy profile (a_i, a_{-i}))
- For $t = 1, 2, \dots$ and $i = 1, 2$, define the function $\eta_i^t: A_{-i} \rightarrow \mathbb{N}$
 - $\eta_i^t(a_{-i})$ is the number of times player i has observed the action a_{-i} before time t . Let $\eta_i^0(a_{-i})$ represent a starting point (or fictitious past)
- For example, consider a two player game, with $A_2 = \{U, D\}$.
 - $\eta_1^0(U) = 3$ and $\eta_1^0(D) = 5$
 - player 2 plays U, U, D in the first three periods
 - then, $\eta_1^3(U) = 3 + 2 = 5$ and $\eta_1^3(D) = 5 + 1 = 6$

The basic idea

- The basic idea of fictitious play is that each player assumes that his opponent is using a stationary mixed strategy, and updates his beliefs about this stationary mixed strategies at each step.
- Players choose actions in each period (or stage) to maximize that period's expected payoff given their prediction of the distribution of opponent's actions, which they form according to:

$$\mu_i^t(a_{-i}) = \frac{\eta_i^t(a_{-i})}{\sum_{a'_{-i} \in A_{-i}} \eta_i^t(a'_{-i})}$$

- Player i forecasts player $-i$'s strategy at time t to be the empirical frequency distribution of the past play

Fictitious play model of learning

- Given player i 's belief/forecast about his opponents play, he chooses his action at time t to maximize his payoff, i.e.,

$$a_i^t \in \operatorname{argmax}_{a_i \in A_i} u_i(a_i, \mu_i^t)$$

Remarks:

- Even though fictitious play is “belief based,” it is also myopic, because players are trying to maximize **current payoff** without considering their future payoffs.
- Perhaps more importantly, they are also not learning the “true model” generating the empirical frequencies (that is, how their opponent is actually playing the game).
- In this model, every player **plays a pure best response to opponents' empirical distributions**.
- Not a unique rule due to multiple best responses. Traditional analysis assumes player chooses any of the pure best responses.

Example

- Consider the fictitious play of the following game:

	L	R
U	4, 4	1, 1
D	5, 1	2, 2

- $\eta_i^t(s_{-i})$ is the number of times player i has observed the action s_{-i} before time t
- $\mu_i^t(s_{-i})$ is player i 's forecast on player $-i$'s strategy at time t

- Note that this game is dominant solvable (D is a strictly dominant strategy for the row player), and the unique $NE(D, R)$.
- Assume $\eta_1^0 = (3, 0)$ and $\eta_2^0 = (1, 2.5)$. Then fictitious play proceeds as follows:

$t = 0$

$$\eta_1^0 = \begin{bmatrix} \eta_1^0(a_2 = L) \\ \eta_1^0(a_2 = R) \end{bmatrix} = \begin{bmatrix} 3 \\ 0 \end{bmatrix}$$

$$\mu_1^0 = \begin{bmatrix} \mu_1^0(a_2 = L) \\ \mu_1^0(a_2 = R) \end{bmatrix} = \begin{bmatrix} 3/3 \\ 0 \end{bmatrix}$$

$$\eta_2^0 = \begin{bmatrix} \eta_2^0(a_1 = U) \\ \eta_2^0(a_1 = D) \end{bmatrix} = \begin{bmatrix} 1 \\ 2.5 \end{bmatrix}$$

$$\mu_2^0 = \begin{bmatrix} \mu_2^0(a_1 = U) \\ \mu_2^0(a_1 = D) \end{bmatrix} = \begin{bmatrix} 1/3.5 \\ 2.5/3.5 \end{bmatrix}$$

Example

- Consider the fictitious play of the following game:

	L	R
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- $\mu_i^t(s_{-i})$ is player i 's forecast on player $-i$'s strategy at time t

- Note that this game is dominant solvable (D is a strictly dominant strategy for the row player), and the unique $NE(D, R)$.
- Assume $\eta_1^0 = (3, 0)$ and $\eta_2^0 = (1, 2.5)$. Then fictitious play proceeds as follows:

$t = 0$

$$\eta_1^0 = \begin{bmatrix} \eta_1^0(a_2 = L) \\ \eta_1^0(a_2 = R) \end{bmatrix} = \begin{bmatrix} 3 \\ 0 \end{bmatrix} \quad \mu_1^0 = \begin{bmatrix} \mu_1^0(a_2 = L) \\ \mu_1^0(a_2 = R) \end{bmatrix} = \begin{bmatrix} 3/3 \\ 0 \end{bmatrix}$$

- player 1 thinks player 2 will play L more often, thus $a_1^0 = D$

$$\eta_2^0 = \begin{bmatrix} \eta_2^0(a_1 = U) \\ \eta_2^0(a_1 = D) \end{bmatrix} = \begin{bmatrix} 1 \\ 2.5 \end{bmatrix} \quad \mu_2^0 = \begin{bmatrix} \mu_2^0(a_1 = U) \\ \mu_2^0(a_1 = D) \end{bmatrix} = \begin{bmatrix} 1/3.5 \\ 2.5/3.5 \end{bmatrix}$$

- player 2 thinks player 1 will play D more often, thus $a_2^0 = R$

Example

- Consider the fictitious play of the following game:

	<i>L</i>	<i>R</i>
<i>U</i>	4, 4	1, 1
<i>D</i>	5, 1	2, 2

- $\eta_i^t(s_{-i})$ is the number of times player i has observed the action s_{-i} before time t
- $\mu_i^t(s_{-i})$ is player i 's forecast on player $-i$'s strategy at time t

- Note that this game is dominant solvable (D is a strictly dominant strategy for the row player), and the unique $NE(D, R)$.
- Assume $\eta_1^0 = (3, 0)$ and $\eta_2^0 = (1, 2.5)$. Then fictitious play proceeds as follows:

$t = 0$

$$\eta_1^0 = \begin{bmatrix} \eta_1^0(a_2 = L) \\ \eta_1^0(a_2 = R) \end{bmatrix} = \begin{bmatrix} 3 \\ 0 \end{bmatrix} \quad \mu_1^0 = \begin{bmatrix} \mu_1^0(a_2 = L) \\ \mu_1^0(a_2 = R) \end{bmatrix} = \begin{bmatrix} 3/3 \\ 0 \end{bmatrix}$$

- player 1 thinks player 2 will play L more often, thus $a_1^0 = D$

$$\eta_2^0 = \begin{bmatrix} \eta_2^0(a_1 = U) \\ \eta_2^0(a_1 = D) \end{bmatrix} = \begin{bmatrix} 1 \\ 2.5 \end{bmatrix} \quad \mu_2^0 = \begin{bmatrix} \mu_2^0(a_1 = U) \\ \mu_2^0(a_1 = D) \end{bmatrix} = \begin{bmatrix} 1/3.5 \\ 2.5/3.5 \end{bmatrix}$$

- player 2 thinks player 1 will play D more often, thus $a_2^0 = R$

Example

- Consider the fictitious play of the following game:

	<i>L</i>	<i>R</i>
<i>U</i>	4, 4	1, 1
<i>D</i>	5, 1	2, 2

- $\eta_i^t(s_{-i})$ is the number of times player i has observed the action s_{-i} before time t
- $\mu_i^t(s_{-i})$ is player i 's forecast on player $-i$'s strategy at time t

- Note that this game is dominant solvable (D is a strictly dominant strategy for the row player), and the unique $NE(D, R)$.
- Assume $\eta_1^0 = (3, 0)$ and $\eta_2^0 = (1, 2.5)$. Then fictitious play proceeds as follows:

$t = 1$

$$\eta_1^1 = \begin{bmatrix} \eta_1^1(a_2 = L) \\ \eta_1^1(a_2 = R) \end{bmatrix} = \begin{bmatrix} 3 \\ 1 \end{bmatrix} \quad \mu_1^1 = \begin{bmatrix} \mu_1^1(a_2 = L) \\ \mu_1^1(a_2 = R) \end{bmatrix} = \begin{bmatrix} 3/4 \\ 1/4 \end{bmatrix}$$

- player 1 thinks player 2 will play L more often, thus $a_1^1 = D$

$$\eta_2^1 = \begin{bmatrix} \eta_2^1(a_1 = U) \\ \eta_2^1(a_1 = D) \end{bmatrix} = \begin{bmatrix} 1 \\ 3.5 \end{bmatrix} \quad \mu_2^1 = \begin{bmatrix} \mu_2^1(a_1 = U) \\ \mu_2^1(a_1 = D) \end{bmatrix} = \begin{bmatrix} 1/4.5 \\ 3.5/4.5 \end{bmatrix}$$

- player 2 thinks player 1 will play D more often, thus $a_2^1 = R$

Example

- Since D is a dominant strategy for the row player, he always plays D , and μ_2^t converges to $(0, 1)$ with probability 1.
- Therefore, player 2 will end up playing R .
- The remarkable feature of the fictitious play is that **players don't have to know anything about their opponent's payoff**. They only form beliefs about how their opponents will play.

Convergence of Fictitious play to pure strategies

- Let $\{a^t\}$ be a sequence of **strategy profiles** generated by fictitious play (FP).
- Let us now study the asymptotic behavior of the sequence $\{s^t\}$, i.e., the convergence properties of the sequence $\{a^t\}$ as $t \rightarrow \infty$
- We first define the notion of convergence to **pure strategies**.

Definition

The sequence $\{a^t\}$ converges to a if there exists T such that $a^t = a$ for all $t \geq T$

Theorem

Let $\{a^t\}$ be a sequence of strategy profiles generated by fictitious play.

- 1) If $\{a^t\}$ converges to \bar{a} , then \bar{a} is a pure strategy **Nash equilibrium**
- 2) Suppose that for some t , $a^t = a^*$, where a^* is a **strict Nash equilibrium**. Then $a^\tau = a^*$ for all $\tau > t$.

Convergence of Fictitious play to pure strategies

- Part 1 is straightforward (Asymptotically stable strategy is Nash equilibrium)
- Consider part 2
- Let $a^t = a^*$. We will show that $a^{t+1} = a^*$.
- Note that for all $a_{-i} \in A_{-i}$

$$\mu_i^{t+1}(a_{-i}) = (1 - \alpha)\mu_i^t(a_{-i}) + \alpha s_{-i}^t(a_{-i}), \text{ with } s_{-i}^t(a_{-i}) = \begin{cases} 1 & \text{if } a_{-i} = a_{-i}^* \\ 0 & \text{otherwise} \end{cases}$$

- $\mu_i^t(a_{-i})$ is player i 's belief on player $-i$'s strategy at time t
 - player i believes player $-i$ will select action a_{-i} with a probability $\mu_i^t(a_{-i})$
- $s_{-i}^t(a_{-i})$ is the probability that player $-i$ actually select action a_{-i}
- $\alpha = \frac{1}{\left\lceil \sum_{a_{-i}} \eta_i^t(a_{-i}) + 1 \right\rceil}$
- Regard μ_i^{t+1} and s_{-i} are strategies, i.e., probability distribution on the possible actions, i.e., $\mu_i^{t+1}, s_{-i} \in \Delta(A_{-i})$

$$\mu_i^{t+1} = (1 - \alpha)\mu_i^t + \alpha s_{-i}^t$$

Convergence of Fictitious play to pure strategies

$$\mu_i^{t+1} = (1 - \alpha)\mu_i^t + \alpha s_{-i}^t$$

- Therefore, by the linearity of the expected utility, we have for all $a_i \in A_i$,

$$u_i(a_i, \mu_i^{t+1}) = (1 - \alpha)u_i(a_i, \mu_i^t) + \alpha u_i(a_i, s_{-i}^t)$$

- Since a_i^* maximizes both terms

$$a_i^* = a_i^t \in \operatorname{argmax}_{a_i \in A_i} u_i(a_i, \mu_i^t)$$

\because assumption $a^t = a^*$

$$a_i^* = BR(s_{-i}^t) = \operatorname{argmax}_{a_i} u_i(a_i, s_{-i}^t)$$

$$= \operatorname{argmax}_{a_i} u_i(a_i, a_{-i}^*)$$

$$s_{-i}^t(a_{-i}) = \begin{cases} 1 & \text{if } a_{-i} = a_{-i}^* \\ 0 & \text{otherwise} \end{cases}$$

- it follows a_i^* will be played at $t + 1$

$$a_i^* = \operatorname{argmax}_a u_i(a, \mu_i^{t+1})$$

- Thus

$$a_i^{t+1} = a_i^t = a_i^*$$

Convergence of Fictitious play to mixed strategies

- The preceding notion of convergence only applies to pure strategies. We next provide an alternative notion of convergence, i.e., convergence of **empirical distributions or beliefs**.

- Converged in pure strategy profiles

$(A, B) \rightarrow (B, A) \rightarrow \dots \rightarrow (A, B) \rightarrow (A, B) \rightarrow (A, B) \rightarrow (A, B) \rightarrow (A, B) \rightarrow (A, B)$

- Converged in mixed strategy profiles in the time-average sense

$(A, B) \rightarrow (B, A) \rightarrow \dots \rightarrow (A, B) \rightarrow (B, A) \rightarrow (A, B) \rightarrow (B, A) \rightarrow (A, B) \rightarrow (B, B)$

Player 1: (A: 1/2 B: 1/2) Player 2: (A: 1/2 B: 1/2)

Definition

The sequence $\{a^t\}$ converges to $\sigma \in S$ in **the time-average sense** if for all i and for all $a_i \in A_i$, we have

$$\lim_{T \rightarrow \infty} \frac{\sum_{t=0}^{T-1} I\{a_i^t = a_i\}}{T} = \sigma(a_i)$$

- In other words, $\mu_{-i}^T(a_i)$ converges $\sigma_i(a_i)$ as $T \rightarrow \infty$

Convergence in Matching Pennies: An example

- Example illustrates convergence of the fictitious play sequence in the time-average sense.

	<i>Heads</i>	<i>Tails</i>
<i>Heads</i>	1, −1	−1, 1
<i>Tails</i>	−1, 1	1, −1

Time	η_1^t	η_2^t	Play
0	(0, 0)	(0, 2)	(<i>H</i> , <i>H</i>)
1	(1, 0)	(1, 2)	(<i>H</i> , <i>H</i>)
2	(2, 0)	(2, 2)	(<i>H</i> , <i>T</i>)
3	(2, 1)	(3, 2)	(<i>H</i> , <i>T</i>)
4	(2, 2)	(4, 2)	(<i>T</i> , <i>T</i>)
5	(2, 3)	(4, 3)	(<i>T</i> , <i>T</i>)
6	(<i>T</i> , <i>H</i>)

- In this example, play continues as a deterministic cycle.
- The time average converges to the unique Nash equilibrium,

$$((1/2, 1/2), (1/2, 1/2))$$

More general convergence result

Theorem

Suppose a fictitious play sequence $\{a^t\}$ converges to σ in the time-average sense. Then σ is a Nash equilibrium.

Proof:

- Suppose a^t converges to σ in the time-average sense
- Suppose, to obtain a **contradiction**, that σ is not a Nash equilibrium
- Then there exist some $i, a_i, a'_i \in A_i$ with $\sigma_i(a_i) > 0$ such that

$$u_i(a'_i, \sigma_{-i}) > u_i(a_i, \sigma_{-i})$$

Note that if σ is Nash equilibrium for all $a_i, a'_i \in A_i$ with $\sigma_i(a_i) > 0$ the following is satisfied

$$u_i(a'_i, \sigma_{-i}) \leq u_i(a_i, \sigma_{-i}) = u_i(\sigma_i, \sigma_{-i})$$

because s_i is included to the support for σ , i. e., $\sigma_i(a_i) > 0$

More general convergence result

- Choose $\epsilon > 0$ such that

$$\epsilon < \frac{1}{2} [u_i(a'_i, \sigma_{-i}) - u_i(a_i, \sigma_{-i})] \quad (1)$$

- Choose T sufficiently large that for all $t \geq T$, we have

$$|\mu_i^T(a_{-i}) - \sigma_{-i}(a_{-i})| < \frac{\epsilon}{\max_{a \in A} u_i(a)} \text{ for all } a_{-i} \quad (2)$$

which is possible $\mu_i^t \rightarrow \sigma_{-i}$ by assumption

- Then, for any $t \geq T$, we have

$$\begin{aligned} u_i(a_i, \mu_i^t) &= \sum_{a_{-i}} u_i(a_i, a_{-i}) \mu_i^t(a_{-i}) \\ &\leq \sum_{a_{-i}} u_i(a_i, a_{-i}) \sigma_{-i}(a_{-i}) + \epsilon && \because (2) \\ &< \sum_{a_{-i}} u_i(a'_i, a_{-i}) \sigma_{-i}(a_{-i}) - \epsilon && \because (1) \\ &\leq \sum_{a_{-i}} u_i(a'_i, a_{-i}) \mu_i^t(a_{-i}) = u_i(a'_i, \mu_i^t) && \because (2) \end{aligned}$$

- This shows that after sufficiently large t , a_i is never played, implying that as $t \rightarrow \infty$, $\mu_{-i}^t(a_i) \rightarrow 0$.
- But this contradicts the fact that $\sigma_i(a_i) > 0$, completing the proof.

More general convergence result

- Choose $\epsilon > 0$ such that

$$\epsilon < \frac{1}{2} [u_i(a'_i, \sigma_{-i}) - u_i(a_i, \sigma_{-i})] \quad (1)$$

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which is possible $\mu_i^t \rightarrow \sigma_{-i}$ by assumption

- Then, for any $t \geq T$, we have

$$\begin{aligned} u_i(a_i, \mu_i^t) &= \sum_{a_{-i}} u_i(a_i, a_{-i}) \mu_i^t(a_{-i}) \\ &\leq \sum_{a_{-i}} u_i(a_i, a_{-i}) \sigma_{-i}(a_{-i}) + \epsilon && \because (2) \\ &< \sum_{a_{-i}} u_i(a'_i, a_{-i}) \sigma_{-i}(a_{-i}) - \epsilon && \because (1) \\ &\leq \sum_{a_{-i}} u_i(a'_i, a_{-i}) \mu_i^t(a_{-i}) = u_i(a'_i, \mu_i^t) && \because (2) \end{aligned}$$

- This shows that after sufficiently large t , a_i is never played, implying that as $t \rightarrow \infty$, $\mu_{-i}^t(a_i) \rightarrow 0$.
- But this contradicts the fact that $\sigma_i(a_i) > 0$, completing the proof.

Example: The Anti-Coordination game

- The theorem gives sufficient conditions for the empirical distribution of the players' action to convergence to a mixed-strategy equilibrium
- However, it does not make any claims about the distribution of the particular outcomes (payoffs that each player can have)
- Consider the following *Anti-Coordination game*

	<i>A</i>	<i>B</i>
<i>A</i>	0, 0	1, 1
<i>B</i>	1, 1	0, 0

- What are the Nash equilibriums?

$$(A, A), (B, B), \left(A: \frac{1}{2}, B: \frac{1}{2} \right)$$

Example: The Anti-Coordination game

Round	1's action	2's action	1's belief	2's belief
0			(1, 0.5)	(1, 0.5)
1	<i>B</i>	<i>B</i>	(1, 1.5)	(1, 1.5)
2	<i>A</i>	<i>A</i>	(2, 1.5)	(2, 1.5)
3	<i>B</i>	<i>B</i>	(2, 2.5)	(2, 2.5)
4	<i>A</i>	<i>A</i>	(3, 2.5)	(3, 2.5)
5	⋮	⋮	⋮	⋮

- The strategy of each player converges to the mixed strategy (0.5, 0.5), which is the mixed strategy Nash equilibrium
- However, the payoff received by each player is 0, since the players never hit the outcomes with positive payoff
- Thus, although the empirical distribution of the strategies converges to the mixed strategy Nash equilibrium, **the players may not receive the expected payoff of the Nash equilibrium.**

Example: Shapley's Almost-Rock-Paper-Scissors game

- The empirical distributions of players actions need not converge at all.
- Consider the following rock-paper-scissors game proposed by Shapley

	Rock	Paper	Scissors
Rock	0, 0	0, 1	1, 0
Paper	1, 0	0, 0	0, 1
Scissors	0, 1	1, 0	0, 0

- The unique Nash equilibrium of this game is for each player to play the mixed strategy is $(1/3, 1/3, 1/3)$

Example: Shapley's Almost-Rock-Paper-Scissors game

Round	1's action	2's action	1's belief	2's belief
0			$(0, 0, 0.5)$	$(0, 0.5, 0)$
1	R	S	$(0, 0, 1.5)$	$(1, 0.5, 0)$
2	R	P	$(0, 1, 1.5)$	$(2, 0.5, 0)$
3	R	P	$(0, 2, 1.5)$	$(3, 0.5, 0)$
4	S	P	$(0, 3, 1.5)$	$(3, 0.5, 1)$
5	S	P	$(0, 3, 2.5)$	$(3, 1.5, 1)$
\vdots	\vdots	\vdots	\vdots	\vdots

- The empirical play of this game never converges to any fixed distribution

When empirical distribution converges?

Theorem

Each of the following is a sufficient condition for the empirical frequencies of play to converge in fictitious play:

- The game is zero sum:
- The game is solvable by iterated elimination of strictly dominated strategies;
- The game is potential game
- The game is $2 \times n$ and has generic payoffs

Summary

- Fictitious play is very sensitive to the players' initial beliefs
- Fictitious play is somewhat paradoxical in that each agent assumes a stationary policy of the opponent, yet, no agent plays a stationary policy except when the process happens to converge to one
- It is simple to state and gives rise to nontrivial properties
- Because players only are thinking about their opponent's actions, they are not playing attention to whether they are actually been doing well.

Extension:

- How to define fictitious play if one has continuous action space?