A survey of Constraint Handling Techniques in Evolutionary Computation Methods

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Introduction

- Evaluation function:
 - one of the major components of any Evolutionary system.
- Evolutionary computation techniques use efficient evaluation function for feasible individuals but these techniques have not developed any guidelines on how to deal with unfeasible solutions.
- Evolutionary programming and evolution strategy:

(Back et al. 1991)

(Fogel and Stayton 1994)

Reject unfeasible individuals

Genetic Algorithms:

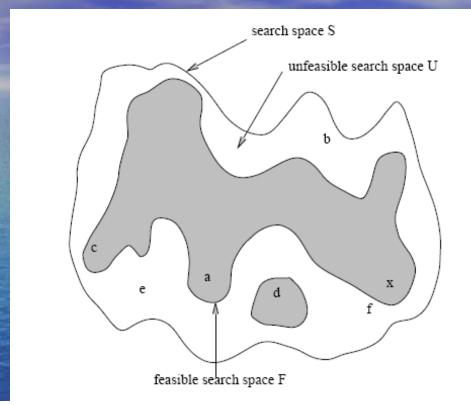
(Goldberg 1989)

Penalize unfeasible individuals

Is there any general rule for designing penalty functions?



Introduction Continue



A search space and its feasible part

We don't make any assumptions about about these subspace:

- They need not be convex
- They not be connected

Introduction Continue.

- We have to deal with various feasible and unfeasible individuals.
- A population may contain some feasible (a, c, d) and unfeasible individuals (b, e, f), while the optimum solution is 'x'
- Problem:

How to deal with unfeasible individuals?

- In general we have to design two two evaluations functions:
 - 1. eval *f*: for feasible domain
 - 2. eval *u*: for unfeasible domain

Introduction Continue

- 1. How should two feasible individuals be compared?
- 2. How should two unfeasible individuals be compared?
- 3. Should we assume that $eval_f(s) \succ eval_u(p)$ for any $s \in \mathcal{F}$ and any $p \in \mathcal{U}$ In particular, which individual is better: individual 'c' or unfeasible individual 'f' (note that the optimum is 'x')
- 4. Should we consider unfeasible individuals harmful and eliminate them from the population? (may be there are useful)
- 5. Should we 'repair' unfeasible solutions by moving them into the closest point of the feasible space ? (should we use a repair procedure for evaluation purpose only)
- 6. Should we chose to penalize unfeasible individuals? $eval_u(p) = eval_f(p) + penalty(p)$

Numerical Optimization and Unfeasible solutions

Non linear programming problem:

The general nonlinear programming problem for continuous variables is to find \overline{X} so as to

optimize
$$f(\overline{X})$$
, $\overline{X} = (x_1, \dots, x_n) \in \mathbb{R}^n$,

where $\overline{X} \in \mathcal{F} \subseteq \mathcal{S}$. The set $\mathcal{S} \subseteq \mathbb{R}^n$ defines the search space and the set $\mathcal{F} \subseteq \mathcal{S}$ defines a *feasible* part of the search space. Usually, the search space \mathcal{S} is defined as an n-dimensional rectangle in \mathbb{R}^n (domains of variables defined as lower and upper bounds):

$$left(i) \le x_i \le right(i), \quad 1 \le i \le n,$$

whereas the feasible set \mathcal{F} is defined by the search space \mathcal{S} and an additional set of constraints:

$$g_j(\overline{X}) \leq 0$$
, for $j = 1, ..., q$, and $h_j(\overline{X}) = 0$, for $j = q + 1, ..., m$.



Condtrain Handling methods

- 1. Methods based on preserving feasibility of solutions:
 - a) Use of specialized operations (The method of Michalewicz & Janikow)
 - b) Searching the boundary of feasible region
- 2. Methods based on penalty functions:
 - a) Static penalty (The method of Homaifar, Lai & Qi)
 - b) Dynamic Penalty (The method of Joines & Houck)
 - c) Annealing penalty (The method of Michalewicz & Attia)
 - d) Death Penalty (The method of Michalewicz)
- 3. Methods based on a search for feasible soluitons:
 - a) Superiority of feasible points (The method of Powell & Skolniks)
 - b) Behavioral memory method (The method of Schoenauer & Xanthakis)
- 4. Muti-Objective optimization methods:
 - a) the method of Paredis
 - b) Cultural Algorithms

Static penalty (The method of Homaifar, Lai & Qi)

- For every constraint we establish a family of intervals that determine appropriate penalty value.
- The methods works as follows:
 - for each constraint, create several (\ell) levels of violation,
 - for each level of violation and for each constraint, create a penalty coefficient R_{ij} $(i = 1, 2, ..., \ell, j = 1, 2, ..., m)$; higher levels of violation require larger values of this coefficient.
 - start with a random population of individuals (i.e., these individuals are feasible or unfeasible),
 - evaluate individuals using the following formula

$$eval(\overline{X}) = f(\overline{X}) + \sum_{j=1}^{m} R_{ij} f_j^2(\overline{X}),$$

where R_{ij} is a penalty coefficient for the *i*-th level of violation and the *j*-th constraint.

Static penalty Continue

- Weakness:
 Number of parameters: m(2l+1)
 m=5, l=4: 45 parameters
- Quality of solutions heavily depends on the values of these parameters:
 - if Rij are moderate: algorithm converge to an unfeasible solution
 - if Rij are too large: algorithm reject unfeasible solution
- Finding an optimal set of parameters for reaching to a feasible solution near optimum is quite difficult.

Dynamic Penalty (The method of Joines & Houck)

In this method individuals are evaluated (at the iteration t) by the following formula:

$$eval(\overline{X}) = f(\overline{X}) + (C \times t)^{\alpha} \sum_{j=1}^{m} f_{j}^{\beta}(\overline{X}),$$

- C,α,β are constant.
- The method is quite similar to Homaifar et al. but it require many fewer parameter.
- This method is independent of the total number of constraints.
- •Penalty component changes with the generation number.
- •The pressure on unfeasible solutions is increased due to the: $(C \times t)^{\alpha}$

Quality of the solution was very sensitive to $(C = 0.5, \alpha = \beta = 2)$ three parameters (reasonable parameters are

Annealing penalty (The method of Michalewicz & Attia)

- In this method linear and nonlinear constraints are processed separately.
- The method works as follow:

- divide all constraints into four subsets: linear equations, linear inequalities, nonlinear equations, and nonlinear inequalities,
- select a random single point as a starting point (the initial population consists of copies of this single individual). This initial point satisfies all linear constraints,
- create a set of active constraints A; include there all nonlinear equations and all violated nonlinear inequalities.
- set the initial temperature τ = τ₀,
- evolve the population using the following formula:

$$eval(\overline{X}, \tau) = f(\overline{X}) + \frac{1}{2\tau} \sum_{j \in A} f_j^2(\overline{X}),$$

(only active constraints are considered),

- if τ < τ_f, stop, otherwise
 - decrease temperature τ,
 - the best solution serves as a starting point of the next iteration,
 - update the set of active constraints A,
 - repeat the previous step of the main part.

- At every iteration the algorithm considers active constraints only.
- The pressure on unfeasible solution is increased due to the decreasing values of temperature ζ.
- This method is quite sensitive to values of its parameters:
 - starting temperature ζ , freezing temperature $\zeta 0$ and cooling scheme ζf to decrease temperature ζ (standard value s $\tau_0 = 1$, $\tau_{i+1} = 0.1 \cdot \tau_i$, with $\tau_f = 0.000001$ are
- This algorithm may converge to a near-optimum solution just in one iteration

Death Penalty (The method of Michalewicz

- This method is popular option in many evolutionary techniques like evolutionary programming.
- Removing unfeasible solutions from population may work well when the feasible search space *F* is convex and it constitutes a reasonable part of the whole search space.
- For some problems where the ratio between the size of F and S is small and an initial population consist of unfeasible individuals only, we must to improve them.
- Experiment shows that when ratio between F and S was between 0% and 0.5% this method performed worse than other methods.
- Unfeasible solutions should provide information and not just be thrown away

Methods based on preserving feasibility of solutions:

- Use of specialized operators
- (The method of Michalewicz & Janikow)
- The idea behind this method is based on specialized operators which transform feasible individuals into feasible individuals.
- This method works only with linear constraint and a feasible starting point.
- A close set of operators maintains feasibility of solutions (the offspring solution vector is always feasible).
- Mutation (which values xi will give from its domain?)
- Crossover: $a\overline{X} + (1-a)\overline{Y}$ (for $0 \le a \le 1$)
- It gave surprisingly good performance on many test functions.
- Weakness of the method lies in its inability to deal with non convex search space (to deal with nonlinear constraints in general)

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Methods based on a search for feasible solutions: Superiority of feasible points (The method of Powell & Skolniks)

- The key concept behind this method is the assumption of superiority of feasible solutions over unfeasible ones.
- Incorporate a heuristic rule for processing unfeasible solutions:
 - "every feasible solution is better than every unfeasible solution"
- This rule is implemented in the following way: evaluations of feasible solutions are mapped into the interval $(-\infty,1)$

evaluations of unfeasible solutions are mapped into the interval $(1, \infty)$

Evaluation procedure:

$$eval_f(\overline{X}) = f(\overline{X}),$$

 $eval_u(\overline{X}) = f(\overline{X}) + r \sum_{j=1}^m f_j(\overline{X}),$

where r is a constant, and

$$eval(\overline{X}) = \begin{cases} eval_f(\overline{X}), & if \ \overline{X} \in \mathcal{F} \\ eval_u(\overline{X}) + \rho(\overline{X}, t), & if \ \overline{X} \in \mathcal{S} - \mathcal{F}. \end{cases}$$

The function $\rho(\overline{X}, t)$ influences unfeasible solutions only; it is defined as

$$\rho(\overline{X},t) = \max\{0, \max_{\overline{X} \in \mathcal{F}} \{eval_f(\overline{X})\} - \min_{\overline{X} \in \mathcal{S} - \mathcal{F}} \{eval_u(\overline{X})\}.$$

In other words, unfeasible individuals have increased penalties: they may not be better than the worst $(\max_{\overline{X} \in \mathcal{F}} \{eval_f(\overline{X})\})$ feasible individual.

- Experiments show that this method work very well however the topology of the feasible search space might be an important factor.
- For problems with a small ratio |F|/|S| the algorithm is often trapped into an unfeasible solution.
- This method should require at least one feasible individual for start.

Behavioral memory method (The method of Schoenauer & Xanthakis)

- This method based on the idea of handling constraints in a particular order.
 - start with a random population of individuals (i.e., these individuals are feasible or unfeasible),
 - set j = 1 (j is constraint counter),
 - evolve this population to minimize the violation of the j-th constraint, until a given percentage of the population (so-called flip threshold ϕ) is feasible for this constraint. In this case

$$eval(\overline{X}) = g_1(\overline{X}).$$

- set j = j + 1,
- the current population is the starting point for the next phase of the evolution, minimizing the violation of the j-th constraint:

$$eval(\overline{X}) = g_j(\overline{X}).^2$$

During this phase, points that do not satisfy at least one of the 1st, 2nd, ..., (j-1)-th constraints are eliminated from the population. The halting criterion is again the satisfaction of the j-th constraint by the flip threshold percentage ϕ of the population.

• if j < m, repeat the last two steps, otherwise (j = m) optimize the objective function f rejecting unfeasible individuals.

This method require 3 parameter:

sharing factor (to maintain diversity of the population) flip threshold

particular permutation of constraint which determine their order

- In the final step of this algorithm the objective function *f* is optimized.
- But for large feasible spaces the method just provides additional computational overhead.
- For very small feasible search space it is essential to maintain a diversity in the population
- Experiments indicate that the method provides a reasonable performance except when the feasible search "too small"



Test case #1

The problem [4] is to minimize a function:

$$G1(\overline{X}) = 5x_1 + 5x_2 + 5x_3 + 5x_4 - 5\sum_{i=1}^{4} x_i^2 - \sum_{i=5}^{13} x_i,$$

subject to

$$2x_1 + 2x_2 + x_{10} + x_{11} \le 10,$$

$$2x_1 + 2x_3 + x_{10} + x_{12} \le 10,$$

$$2x_2 + 2x_3 + x_{11} + x_{12} \le 10,$$

$$-8x_1 + x_{10} \le 0, \quad -8x_2 + x_{11} \le 0,$$

$$-8x_3 + x_{12} \le 0, \quad -2x_4 - x_5 + x_{10} \le 0,$$

$$-2x_6 - x_7 + x_{11} \le 0, \quad -2x_8 - x_9 + x_{12} \le 0,$$

$$0 \le x_i \le 1, \ i = 1, \dots, 9, \quad 0 \le x_i \le 100,$$

$$i = 10, 11, 12, \quad 0 \le x_{13} \le 1.$$

The problem has 9 linear constraints; the function G1 is quadratic with its global minimum at

$$\overline{X}^* = (1, 1, 1, 1, 1, 1, 1, 1, 1, 3, 3, 3, 1),$$

where $G1(\overline{X}^*) = -15$. Six (out of nine) constraints are active at the global optimum (all except the following three: $-8x_1 + x_{10} \le 0$, $-8x_2 + x_{11} \le 0$, $-8x_3 + x_{12} \le 0$).

Test case #2 The problem [6] is to minimize a function:

$$G2(\overline{X}) = x_1 + x_2 + x_3,$$

where

$$1 - 0.0025(x_4 + x_6) \ge 0,$$

$$1 - 0.0025(x_5 + x_7 - x_4) \ge 0,$$

$$1 - 0.01(x_8 - x_5) \ge 0,$$

$$x_1x_6 - 833.33252x_4 - 100x_1 + 83333.333 \ge 0,$$

$$x_2x_7 - 1250x_5 - x_2x_4 + 1250x_4 \ge 0,$$

$$x_3x_8 - 1250000 - x_3x_5 + 2500x_5 \ge 0,$$

$$100 \le x_1 \le 10000, \quad 1000 \le x_i \le 10000,$$

$$i = 2, 3, \quad 10 < x_i < 1000, \quad i = 4, \dots, 8.$$

The problem has 3 linear and 3 nonlinear constraints; the function G2 is linear and has its global minimum at

$$\overline{X}^* = (579.3167, 1359.943, 5110.071, 182.0174, 295.5985, 217.9799, 286.4162, 395.5979),$$

where $G2(\overline{X}^*) = 7049.330923$. All six constraints are active at the global optimum.

Test case #3

The problem [6] is to minimize a function:

$$G3(\overline{X}) = (x_1 - 10)^2 + 5(x_2 - 12)^2 + x_3^4 + 3(x_4 - 11)^2 + 10x_5^6 + 7x_6^2 + x_7^4 - 4x_6x_7 - 10x_6 - 8x_7,$$

where

$$127 - 2x_1^2 - 3x_2^4 - x_3 - 4x_4^2 - 5x_5 \ge 0,$$

$$282 - 7x_1 - 3x_2 - 10x_3^2 - x_4 + x_5 \ge 0,$$

$$196 - 23x_1 - x_2^2 - 6x_6^2 + 8x_7 \ge 0,$$

$$-4x_1^2 - x_2^2 + 3x_1x_2 - 2x_3^2 - 5x_6 + 11x_7 \ge 0$$

$$-10.0 \le x_i \le 10.0, i = 1, \dots, 7.$$

The problem has 4 nonlinear constraints; the function G3 is nonlinear and has its global minimum at

$$\overline{X}^* = (2.330499, 1.951372, -0.4775414, 4.365726, -0.6244870, 1.038131, 1.594227),$$

where $G3(\overline{X}^*) = 680.6300573$. Two (out of four) constraints are active at the global optimum (the first and the last one).

Test case #4

The problem [6] is to minimize a function:

$$G4(\overline{X}) = e^{x_1x_2x_3x_4x_5},$$

subject to

$$x_1^2 + x_2^2 + x_3^2 + x_4^2 + x_5^2 = 10$$
, $x_2x_3 - 5x_4x_5 = 0$, $x_1^3 + x_2^3 = -1$, $-2.3 \le x_i \le 2.3$, $i = 1, 2$, $-3.2 \le x_i \le 3.2$, $i = 3, 4, 5$.

The problem has 3 nonlinear equations; nonlinear function G4 has its global minimum at

$$\overline{X}^* = (-1.717143, 1.595709, 1.827247, -0.7636413, -0.7636450),$$

where $G4(\overline{X}^*) = 0.0539498478$.

Test Case #5

The problem [6] is to minimize a function:

$$G5(\overline{X}) = x_1^2 + x_2^2 + x_1x_2 - 14x_1 - 16x_2 + (x_3 - 10)^2 + 4(x_4 - 5)^2 + (x_5 - 3)^2 + 2(x_6 - 1)^2 + 5x_7^2 + 7(x_8 - 11)^2 + 2(x_9 - 10)^2 + (x_{10} - 7)^2 + 45,$$

where

$$\begin{aligned} &105 - 4x_1 - 5x_2 + 3x_7 - 9x_8 \ge 0, \\ &-10x_1 + 8x_2 + 17x_7 - 2x_8 \ge 0, \\ &8x_1 - 2x_2 - 5x_9 + 2x_{10} + 12 \ge 0, \\ &-3(x_1 - 2)^2 - 4(x_2 - 3)^2 - 2x_3^2 + 7x_4 + 120 \ge 0, \\ &-5x_1^2 - 8x_2 - (x_3 - 6)^2 + 2x_4 + 40 \ge 0, \end{aligned}$$

$$-x_1^2 - 2(x_2 - 2)^2 + 2x_1x_2 - 14x_5 + 6x_6 \ge 0,$$

$$-0.5(x_1 - 8)^2 - 2(x_2 - 4)^2 - 3x_5^2 + x_6 + 30 \ge 0,$$

$$3x_1 - 6x_2 - 12(x_9 - 8)^2 + 7x_{10} \ge 0,$$

$$-10.0 \le x_i \le 10.0, \quad i = 1, \dots, 10.$$

The problem has 3 linear and 5 nonlinear constraints; the function G5 is quadratic and has its global minimum at

$$\overline{X}^* = (2.171996, 2.363683, 8.773926, 5.095984, 0.9906548, 1.430574, 1.321644, 9.828726, 8.280092, 8.375927).$$

where $G5(\overline{X}^*) = 24.3062091$. Six (out of eight) constraints are active at the global optimum (all except the last two).

Summary of test case

TC	n	Type of f	ρ	LI	NE	NI	a
#1	13	quadratic	0.0111%	9	0	0	6
#2	8	linear	0.0010%	3	0	3	6
#3	7	polynomial	0.5121%	0	0	4	2
#4	5	nonlinear	0.0000%	0	3	0	3
#5	10	quadratic	0.0003%	3	0	5	6

In the table below:

Method #1 is: Homaifar's method

Method #2 is: joines 's method

Method #3 is: Xanthakis's mehtod

Method #4 is: Attia's method

Method #5 is: powell's method

Method #6 is: Death penalty



Experimental, Result

TC	Exact		Method	Method	Method
	opt.		#1	#2	#3
		b	-15.002	-15.000	-15.000
#1	-15.000	m	-15.002	-15.000	-15.000
		w	-15.001	-14.999	-14.998
		c	0, 0, 4	0, 0, 0	0, 0, 0
		b	2282.723	3117.242	7485.667
#2	7049.331	m	2449.798	4213.497	8271.292
		w	2756.679	6056.211	8752.412
		c	0, 3, 0	0, 3, 0	0, 0, 0
		b	680.771	680.787	680.836
#3	680.630	m	681.262	681.111	681.175
		w	689.660	682.798	685.640
		c	0, 0, 1	0, 0, 0	0, 0, 0
		b	0.084	0.059	
#4	0.054	m	0.955	0.812	*
		w	1.000	2.542	
		c	0, 0, 0	0, 0, 0	
		b	24.690	25.486	
#5	24.306	m	29.258	26.905	_
		w	36.060	42.358	
		c	0, 1, 1	0, 0, 0	

Method	Method	Method	Method
#4	#5	#6	#6(f)
-15.000	-15.000		-15.000
-15.000	-15.000	_	-14.999
-15.000	-14.999		-13.616
0, 0, 0	0, 0, 0		0, 0, 0
7377.976	2101.367		7872.948
8206.151	2101.411	_	8559.423
9652.901	2101.551		8668.648
0, 0, 0	1, 2, 0		0, 0, 0
680.642	680.805	680.934	680.847
680.718	682.682	681.771	681.826
680.955	685.738	689.442	689.417
0, 0, 0	0, 0, 0	0, 0, 0	0, 0, 0
0.054	0.067		
0.064	0.091	*	*
0.557	0.512		
0, 0, 0	0, 0, 0		
18.917	17.388		25.653
24.418	22.932		27.116
44.302	48.866		32.477
0, 1, 0	1, 0, 0		0, 0, 0

