
Object-Oriented Database Evolution

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Context

Object-Oriented:

- Object identity
- Inheritance and attribute redefinition

Database:

- Schema:
 - classes (C)
 - inheritance links (\leq)
 - attribute definitions (Σ)
- Instance:
 1. class populations (ν)
 2. references (μ)
- Typing constraints

Motivation

Evolution = specification of a new state of the database from the current one

$$(C, \leq, \Sigma)(\nu, \mu) \mapsto (C', \leq', \Sigma')(\nu', \mu')$$

Evolution language = combination of schema and instance update languages

The database programmer should be provided with a tool allowing to describe the evolutions

Existing approaches do not offer:

- Programming language for evolution of the database (schema + instance)
- Consistency checking for this language

Instance Update, Syntax

Primitives:

- $\text{move } x \ c$
- $\text{new } x \ c$ $\text{delete } x$
- $\text{set } x \ a \ x'$ $\text{cut } x \ a \ x'$

Program over a schema $S =$
set of primitive calls

Instance Update, Semantics

- Program over an instance I = the variables are instantiated with objects in I
- First order formula associated to each instruction of program P

– *new* x_1 c or *move* x_1 c :

$$\forall z (x_1(z) \rightarrow (c(z) \wedge \bigwedge_{c' \in C - \{c\}} \neg c'(z)))$$

– *cut* x_1 a x_2 :

$$\forall z (x_1(z) \rightarrow \forall z' (x_2(z') \rightarrow \neg a(z, z')))$$

$\psi_{S,P}$ = conjunction of the formulas associated to each instruction

Semantics of a Program

Given $\psi_{S,P}$ and an instance I of S

- What is not affected by the program is
 $invariant(I) =$
 $Max\{I' \mid I' \subseteq I \text{ and}$
 $\exists I'' \supseteq I' \text{ s.t. } (I'', v) \models \psi\}$
- the semantics of P over I with the parameters is
 $Min\{I' \mid invariant(I) \subseteq I' \text{ and}$
 $(I', v) \models \psi\}$

Consistency

Inconsistent instance:

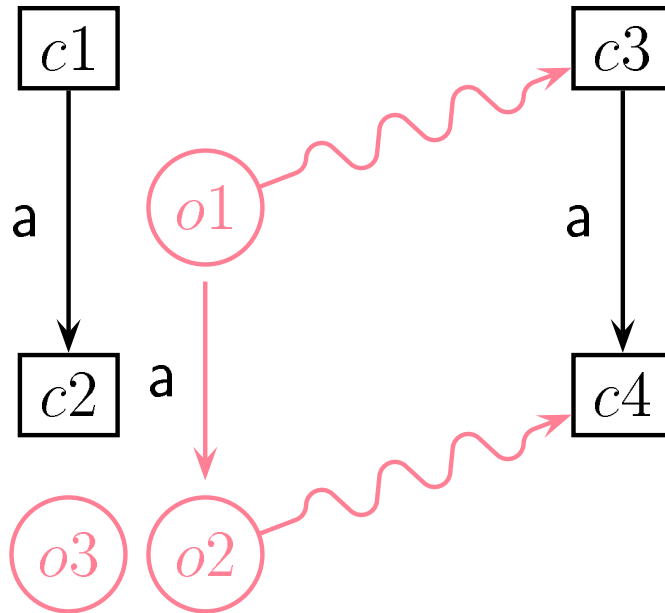
- objects belong to one and only one class
- no ill-typed references
- no undefined values

A program is inconsistent iff there exists an instance and an instantiation s.t. the semantics of the program is inconsistent

Consistency of a program is decidable

Problem

- Function migrating object in $c1$ to $c3$ and object referred through a to $c4$



$\{x_1 : c1; x_2 : c2\}$

$\{move\ x_1\ c3, move\ x_2\ c4\}$

is inconsistent

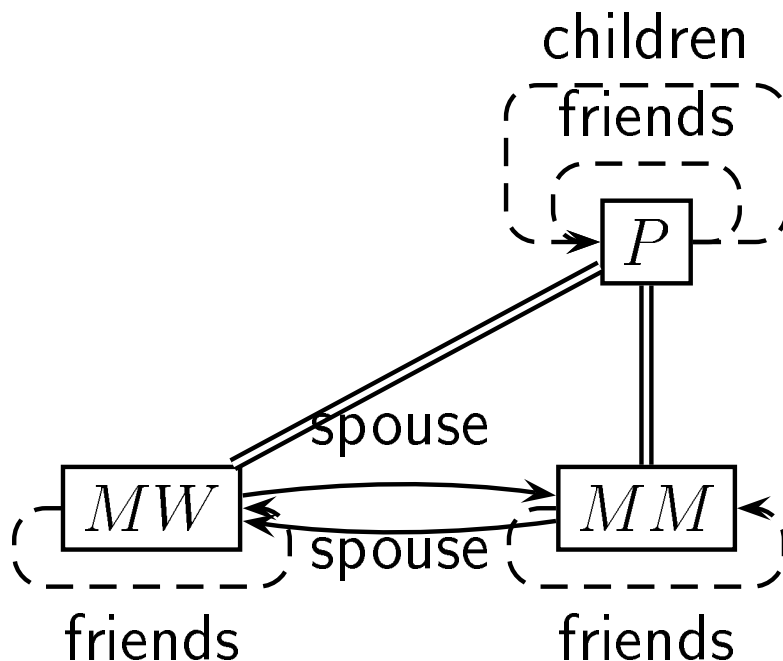
- State that $x_2 = x_1.a$
- Path defined from a typed variable
 $x, l.a, l.a^{-1}, l : c, l \cap l', l \cup l', l - l'$

Instance Update using Paths

- FO2 = First Order
 - with two variables
 - no function symbol
- Satisfiability of path expressions is decidable
- Concrete program = variables replaced by paths
- expresses more functions
- consistency of a concrete program is decidable

Instance Update Example

- Static aspect Real World:
Persons = not-married + married men
+ married women
Attributes = friends, children, spouse
Constraints = “my friends are like me”
- Static aspect Database



Instance Update Program

```
{  
  move  $xx$   $MW$ ,  
  move  $xy$   $MM$ ,  
  set  $xx$  spouse  $xy$ ,  
  set  $xy$  spouse  $xx$ ,  
  cut  $xx$  friends  
    ( $xx.friends - (xx.friends : MW)$ ),  
  cut  $xy$  friends  
    ( $xy.friends - (xy.friends : MM)$ )  
}
```

This program is consistent

It will be used many times. There are people who marry every day using this procedure.

Schema update

Syntax

- primitives:
 - classes: $+c, -c$
 - inheritance links: $+(c, c'), -(c, c')$
 - attribute definitions: $+d, -d$
- program = set of primitive calls

Semantics

- Trivial semantics associated to each instruction
- Semantics of program = combination of semantics of each primitive of the program

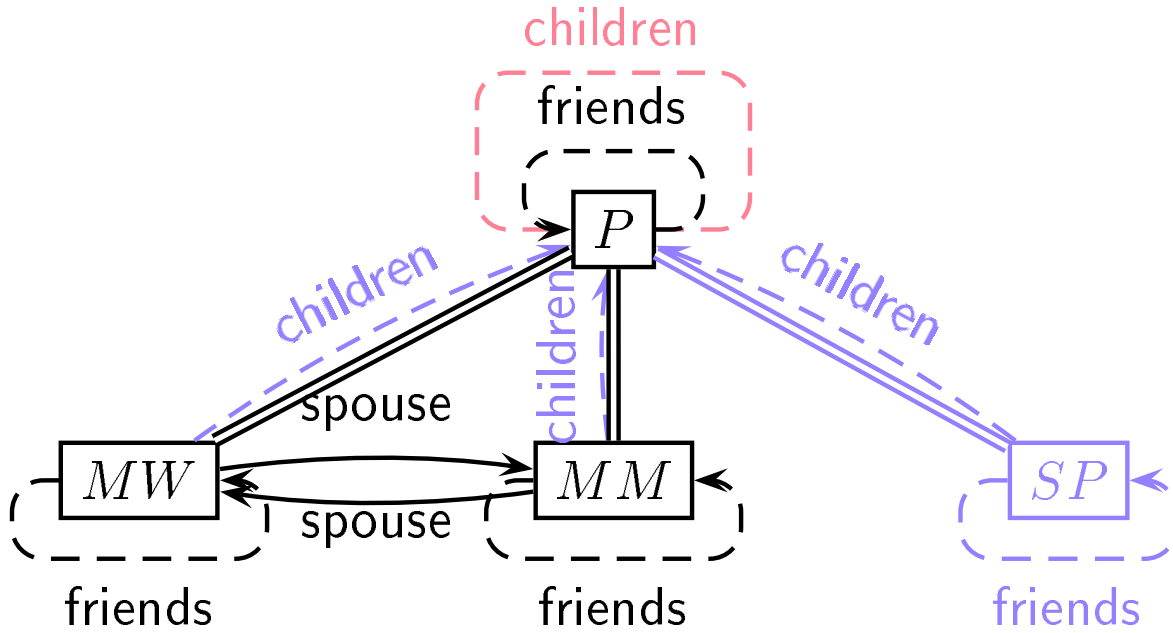
Database Evolution

Schema S , Instance I

- Evolution program = instance update program P_I + schema update program P_S
- Consistency checking of P_I in $S \cup P_S(S)$
- P_I performed over I in schema $S \cup P_S(S)$

consistency is decidable

Evolution Example



$$\{ +_c SP, +_a children : SP \rightarrow P, \\ +_a children : MW \rightarrow P, -_a children P \rightarrow P, \\ +_a children : MM \rightarrow P, \\ +_a friends : SP \rightarrow SP, +_h SP P \}$$

z = all the objects in the database

singleparents =

$$z.children^{-1} - (z.children^{-1} : \\ MW \cup z.children^{-1} : MM)$$

$$P = \{ move\ singleparents\ SP, \\ cut\ singleparents\ friends\ (z - singleparents) \}$$

Conclusion

- Statically typed instance update language featuring object migration
- Evolution mechanism

Future Works

- Evolution usable with different instance update languages by isolating some specific parts of the language that allow consistency checking
- O2 prototype, extension in progress

References

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- [MMW94] A. Mendelzon, T. Milo, and E. Waller. Object migration. In *PODS*, 1994.