Lecture 5: Outline

- Reminder: functions
- Varieties of functions
- Service-oriented systems
- Example: a master component of a service-oriented system
- Homework: augmented hotel system

Functions (reminder)

- Functions form a special class of relations that satisfy additional requirement: any element of the source set can be related to no more than 1 element of the target
- Functionality requirement mathematically:

$$\forall x, y, z. \ (x \mapsto y) \in R \land (x \mapsto z) \in R \Rightarrow y = z$$

- Any operation applicable to a relation or a set is also applicable to a function. For example, we can talk about the domain and the range of a function or a function as a set of pairs
- If f is a function, then f(x) is the result of the function f for the argument x

Total and partial functions (reminder)

- Functions are called *total* if their domain is the whole source set Syntax: $f \in S \rightarrow T$ (or ascii $f : S \rightarrow T$) where dom(f) = S and $ran(f) \subseteq T$
- Functions are called *partial* if their domain is a subset of the source set Syntax: $f \in S \rightarrow T$ (or ascii f : S +->T) where $dom(f) \subseteq S$ and $ran(f) \subseteq T$

Functional/array assignment (reminder)

 The notation used to describe machine actions (i.e., assignments in the machine events) allows us to directly assign values to indexed elements of arrays:

$$a(i) := E$$

• This is just syntactic sugaring for the following assignment:

$$a := a \Leftrightarrow \{(i \mapsto E)\}$$

- The assignment also works if a is modelled as a partial function. However, if we want to check/read values from such an array, we have to ensure/prove (by using the event guards and/or machine invariants) that the used index belongs to the function domain, i.e., $i \in dom(a)$
- reserved(r) := FALSE

OK!

• nltems(j) := nltems(i) + 1

only if $i \in dom(nltems)$



Varieties of functions

Suppose we have a function f (from the source X to the target Y). Then it is called

```
Total function
                                   if dom(f) = X, ran(f) \subseteq Y
                                   if dom(f) \subseteq X, ran(f) \subseteq Y
Partial function
                    → +->
                                   if dom(f) = X, ran(f) \subseteq Y
Total injection
                    >→ >->
                                    and one-to-one function
                                   if dom(f) \subseteq X, ran(f) \subseteq Y
Partial injection
                    → >+>
                                    and one-to-one function
                                   if dom(f) = X, ran(f) = Y
Total surjection
                    → ->>
Partial surjection
                                   if dom(f) \subseteq X, ran(f) = Y
                    +→ +->>
                                   if dom(f) = X, ran(f) = Y
(Total) Bijection
                    >→> >→
                                    and one-to-one function
```

Varieties of functions: examples

- Injection: a function with 1-to-1 relationship between the source and target sets (e.g., an array without repeating elements)
- Example: VU $id \in PERSON \rightarrow ID$
 - It is a partial injection: not all persons have a VU identification number, however, id is unique for each person
- Advantage: a reverse relation for an injection is also a function!
- Example: $VU_id^{\sim} \in ID \rightarrow PERSON$
 - A total injection this time

Varieties of functions: examples

- Surjection: a function that covers all the target set
- Example:
 married ∈ WIFE → HUSBAND

It is a total surjection

Another example:
 Capital_of ∈ CITY → COUNTRY

A partial surjection this time

Varieties of functions: examples

- Bijection: a total function that is both injection and surjection
- Example:
 married ∈ WIFE → HUSBAND

It is a bijection (in many countries)

- Bijections relate sets with the same power (length)
- Another example:
 VU_account ∈ ID → VU_ACCOUNT

A bijection: both sets are of the same length and one-to-one relationships in both directions

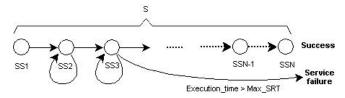
Telecommunicating and service-oriented systems

- Examples: telecommunication networks, service-oriented architectures, web services, cloud-based services
- Usually consists of multiple components that collectively provide a service to the service consumer (the external user or possibly other service provider)
- The components can be further partitioned into those that are responsible for service orchestration (management) or service execution
- Components for service orchestration: master components, service directors, service managers, (sometimes) frontend components
- Components for service execution: worker components, standalone components

Service decomposition

Often, a service request is decomposed and forwarded to different components (sub-service providers)

Service directors / master components are responsible for managing and controlling the whole service flow, forwarding requests to the respective components providing the required sub-services



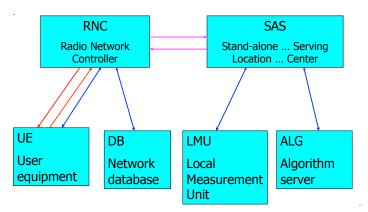
Telecommunicating and service-oriented systems (cont.)

- Components can be unavailable/busy or fail during execution
- A master component is responsible for monitoring the availability and failure status of its workers
- A master component can also incorporate fault tolerance mechanisms (e.g., retrying a service request, finding a replacement worker component, etc.)
- A master component can usually handle multiple service requests at once, while worker components are often executing a single service request

Positioning service (Nokia)

Typically, service-oriented systems have a layered architecture

Here, RNC (Radio Network Controller) – the external master (frontend) component, SAS – local service master component, while UE, DB, LMU, and ALG are employed as standalone / worker components



Simple example of a service-oriented system: requirements

- The system (master component) handles the incoming service requests, distributing them to the available worker components
- While being handled, the arrived service requests are stored in the input buffer
- The input buffer has the pre-defined size that cannot be exceeded
- Once handled, all the requests are stored in the output buffer (no size restrictions)
- Each request has the handling status, which can be Waiting, Executing, Completed, or Failed
- If the input buffer is full, a request is immediately transferred to the output buffer with the status Failed

Service-oriented system: requirements (cont.)

- After an arrived request is added to the input buffer, it gets the status Waiting
- Once a request from the input buffer is assigned to one of the available worker components, it gets the status Executing
- Once a request is successfully handled by a worker component, it gets the status Completed and is transferred from the input buffer to the output buffer
- A worker can fail to handle a request. Then the request gets the status Failed and is transferred from the input buffer to the output buffer

Service-oriented system: requirements (cont.)

- Each request is associated with a specific service type
- Each worker component is capable to handle a specific subset of service types. This knowledge is known beforehand and does not change during system execution
- There is a number of active worker components at a particular moment. Any inactive worker can be activated and any active worker (not handling any requests) can be deactivated
- Only an active worker component that is capable to handle the service type of a request can be assigned that request
- A worker component can handle no more than one request at the moment

Service-oriented system: context

```
CONTEXT
 Services ctx
SETS
  REQUEST
  RTYPE
  WORKER
  STATUS
CONSTANTS
  Waiting, Executing, Completed, Failed,
  Reg type, Worker types, buffer size
END
```

Service-oriented system: context (cont.)

```
AXIOMS
partition(STATUS, \{Waiting\}, \{Executing\}, \{Completed\}, \{Failed\})
Req\_type \in REQUEST \twoheadrightarrow RTYPE
Worker\_types \in WORKER \rightarrow \mathbb{P}_1(RTYPE)
buffer\_size \in \mathbb{N}_1
END
```

Revealing a structure of abstract type elements

- Abstract types (like REQUEST or WORKER) are convenient to introduce standard notions of the modelled system
- If we need to reveal the internal structure of abstract type elements, we can introduce constant functions (like Req_type) in the context component to extract the necessary information from abstract type element
- Such functions are similar to a way object attributes (fields) are accessed in OOP
- Another possible function:
 Req_priority ∈ REQUEST → PRIORITY

Service-oriented system: machine

```
MACHINE
  Services mch
SFFS
  Services ctx
VARIABLES
  input, output, rstatus, active, assigned
INVARIANT
  input \in \mathbb{P}(REQUEST)
  output \in \mathbb{P}(REQUEST)
  rstatus \in REQUEST \rightarrow STATUS
  active \in \mathbb{P}(WORKER)
  dom(rstatus) = input \cup output
  input \cap output = \emptyset
  card(input) \leq buffer \ size
. . .
```

```
assigned \in active \rightarrow REQUEST
   ran(assigned) \subseteq input
   \forall r \cdot r \in ran(assigned) \Rightarrow rstatus(r) = Executing
  \forall r \cdot r \in output \Rightarrow rstatus(r) \in \{Completed, Failed\}
   \forall r \cdot r \in input \Rightarrow rstatus(r) \in \{Waiting, Executing\}
INITIALISATION
   input, output := \varnothing, \varnothing
   rstatus, assigned := \emptyset, \emptyset
   active :\in \mathbb{P}(WORKERS)
```

```
EVENTS
  request \ arrival = ANY r
    WHERE r \notin dom(rstatus) \land card(input) \le buffer size
    THEN
       input := input \cup \{r\}
       rstatus(r) := Waiting
    END
  request rejection = ANY r
    WHERE r \notin dom(rstatus) \land card(input) = buffer size
    THEN
       output := output \cup \{r\}
       rstatus(r) := Failed
    FND
```

```
request assignment = ANY r, w
  WHERE r \in input \land w \in active \land rstatus(r) = Waiting
            Reg type(r) \in Worker types(w)
  THEN
    assigned(w) := r
    rstatus(r) := Executing
  END
request completed = ANY r
  WHERE r \in ran(assigned)
  THEN
    output := output \cup \{r\}
    input := input \setminus \{r\}
    assigned := assigned \triangleright \{r\}
    rstatus(r) := Completed
  END
```

```
request failed = ANY r
  WHERE r \in ran(assigned)
  THEN
     output := output \cup \{r\}
    input := input \setminus \{r\}
    assigned := assigned \triangleright \{r\}
    rstatus(r) := Failed
  END
worker \ activate = ANY w
  WHERE w \notin active
  THEN active := active \cup \{w\} END
worker deactivate = ANY w
  WHERE w \in active \land w \notin dom(assigned)
  THEN active := active \setminus \{w\} END
```

Invariants: summary

Different kinds/forms of invariants:

- 1 Typing invariants: $x \in S, y \subseteq S, r \in S \leftrightarrow T, f \in S \leftrightarrow T, ...$
- ② Interrelationships between different variables (and constants) : $input \cap output = \varnothing, f \subseteq r, dom(f) = y$ $card(input) \leq buffer_size$

The variable values can be used to directly define the typing of other variables:

 $items \in POW(ITEMS)$, $price \in items \rightarrow NAT$

Invariants: summary (cont.)

3 Conditional (local) invariants of the form $cond_1 \wedge ... \wedge cond_k \Rightarrow property$

$$served = TRUE \Rightarrow payed = TRUE$$

 $choice \neq None \land served = FALSE \Rightarrow nItems(choice) > 0$

In such a way, invariants can be "narrowed" down to specific points of system execution

Using quantifiers in invariants on more complex data structures:

```
\forall r \cdot r \in ran(assigned) \Rightarrow rstatus(r) = Executing
\forall r \cdot r \in output \Rightarrow rstatus(r) \in \{Completed, Failed\}
```

Homework: an extended hotel booking system (slightly changed requirements from the first homework)

- The hotel booking system handles room reservation by customers;
- The system must have operations (events) for room reservation, cancellation, customer check-in (with a reservation), customer check-in (without a reservation), and customer check-out;
- Each reservation is stored by the system until it is cancelled or the reservation customer checks-in (with a reservation);
- A reservation stores the information about the reserved room, the reserved dates, and the customer;
- For any reservation, its dates cannot overlap with any other reservation dates for the same room;

Homework: an extended hotel booking system (requirements, cont.)

- A reservation can be cancelled;
- After cancellation, the stored reservation is removed from the system;
- After a customer's check-in (with a reservation), the reserved room gets the status "occupied" and the stored reservation is removed from the system;
- The system keeps the information about the occupied rooms, the customers, and the dates they are staying;
- After a customer's check-in (with a reservation), the information from the room reservation is associated with (copied to) that of the occupied room;

Homework: an extended hotel booking system (requirements, cont.)

- For any occupied room, its dates cannot overlap with any reservation dates for the same room;
- A customer can check-in (without a reservation), if there is an unoccupied and unreserved room for the specified dates;
- Once a customer checks-out, the information about the occupied room (the customer and dates) is removed from the system.

Homework: a hotel booking system (cont.)

Hints:

- Again, carefully define the necessary data structures in the model context
- Hotel reservations can be introduced as an abstract set, while the corresponding constant functions can be defined to extract necessary information (like the customer or dates) from them
- A similar or the same approach can be applied to model the info about the occupied rooms
- Dates can be also modelled as elements (subsets) of a pre-defined abstract set. Alternatively, a subset of natural numbers can be used
- The requirement 5 and 11 must become model invariants