

Week 10 Classwork due on Friday Nov 16, 22:00 Hour

Group 5

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Exercise 1

Implement the Diffie-Hellman protocol.

Answers:

The Diffie-Hellman protocol is used to establish a shared secret that can be used for secret communications while exchanging data over a public network using the elliptic curve to generate points and get the secret key using the parameters.

The implementation of the protocol uses the multiplicative group of integers modulo p , where p is prime, and g is a primitive root modulo p . These two values are chosen in this way to ensure that the resulting shared secret can take on any value from 1 to $p-1$. The below is a step by step explanation of the key exchange process between Alice and Bob.

ALICE	BOB
Public Keys available = P, G	Public Keys available = P, G
Private Key Selected = a	Private Key Selected = b
Key generated $x = G^a \text{ Mod } P$	Key generated $y = G^b \text{ Mod } P$
Exchange of generated keys takes place	
Key received = y	key received = x
Generated Secret Key $k_a = y^a \text{ Mod } P$	Generated Secret Key $k_b = x^b \text{ Mod } P$
Algebraically it can be shown that $k_a = k_b$ – therefore both Alice and Bob have a symmetric key to encrypt.	

The implementation in the form of codes are written as follows:

```
p = 23  # P
g = 5   # G
alice = 6  # a
bob = 15  # b

print( "p: " , p )
print( "g: " , g )

# Alice Sends Bob A = g^a mod p
A = pow(g, alice) % p
print( "Alice sending message: " , A )

# Bob Sends Alice B = g^b mod p
B = pow(g, bob) % p
print( "Bob sending message: " , B )

print( "===== " )
# Alice Computes Shared Secret: s = B^a mod p
aliceSecret = (B ** alice) % p
print( "Alice Shared info: " , aliceSecret )

# Bob Computes Shared Secret: s = A^b mod p
bobSecret = (A ** bob) % p
print( "Bob Shared info: " , bobSecret )

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The output is:
('p: ', 23)
('g: ', 5)
('Alice sending message: ', 8)
('Bob sending message: ', 19)
=====
('Alice Shared info: ', 2)
('Bob Shared info: ', 2)
```

Exercise 2

Implement the RSA encryption scheme from scratch. Use the following interface:

- `Gen(minPrime)` generates a public/private keypair (512 bits) where $p, q > \text{minPrime}$.
- `Enc(pubKey, msg)` returns `ctxt` (integer).
- `Dec(privKey, ctxt)` returns `msg` (integer).

Answers:

We will implement the RSA encryption by using the following steps:

`Gen(minPrime)`

- Select (large) random prime numbers p, q ($p \neq q$, but with almost equal size)
- Compute modulus $n = pq$
- Compute $\Phi = (p-1)(q-1)$
- Select public exponente, $1 < e < \Phi$, such that $\text{gcd}(e, \Phi) = 1$
- Compute private exponent $d = e^{-1} \bmod \Phi$
- Return public key (n, e) , and private key (p, q, Φ, d)

`Enc(pubKey=e, msg=m)`

- Return $m^e \bmod n = c$

`Dec(privKey=d, ctxt=c)`

- Return $c^d \bmod n = m$

The codes are written as follows (in Python 2.7):

```
import random

def gcd(a, b):
    while a != 0:
        a, b = b % a, a
    return b

#calculate d
def findModInverse(a, m):
    # Returns the modular inverse of a % m, which is
    # the number x such that a*x % m = 1
    if gcd(a, m) != 1:
        return None
```

```

# Calculate using the Extended Euclidean Algorithm:
u1, u2, u3 = 1, 0, a
v1, v2, v3 = 0, 1, m
while v3 != 0:
    q = u3 // v3
    v1, v2, v3, u1, u2, u3 = (u1 - q * v1), (u2 - q * v2), (u3 - q * v3), v1, v2, v3
return u1 % m

def generatePrime(keysize):
    while True:
        num = random.randrange(2**(keysize-1), 2**(keysize))
        if isPrime(num):
            return num

def isPrime(n):
    if n <= 1: return False
    i = 2
    while i*i <= n:
        if n%i == 0 : return False
        i += 1
    return True

def Gen(minPrime, keySize):

    # Step 1: Create two prime numbers, p and q. Calculate n = p * q.
    p = generatePrime(keySize)
    if p < minPrime:
        return;
    q = generatePrime(keySize)
    if q < minPrime:
        return;
    n = p * q
    print 'p is:', p
    print 'q is:', q

    # Step 2: Create a number e that is relatively prime to (p-1)*(q-1).
    while True:
        # Keep trying random numbers for e until one is valid.

```

```

    e = random.randrange(2 ** (keySize - 1), 2 ** (keySize))
    if gcd(e, (p - 1) * (q - 1)) == 1:
        break
    print 'e is:', e

# Step 3: Calculate d, the mod inverse of e.
d = findModInverse(e, (p - 1) * (q - 1))
print 'd is:', d

publicKey = (n, e)
privateKey = (n, d)
print('Public key:', publicKey)
print('Private key:', privateKey)
return (publicKey, privateKey)

def Enc(pubKey, msg):
    return pow(msg, pubKey[1]) % pubKey[0]

def Dec(privKey, ctxt):
    return pow(ctxt, privKey[1]) % privKey[0]

def main():
    minPrime = 2
    publicKey, privateKey = Gen(minPrime, 10) # set key to 10 bits
    msg = 10000; # set message to be 10000
    ciphertext = Enc(publicKey, msg)
    print 'ciphertext is:', ciphertext
    print 'plaintext is:', Dec(privateKey, ciphertext)

if __name__ == '__main__':
    main()

```

The output is:

```

p is: 919
q is: 601
e is: 623
d is: 251087

```

```
('Public key:', (552319, 623))  
( 'Private key:', (552319, 251087))  
ciphertext is: 164526  
plaintext is: 10000
```

Exercise 3

Implement the RSA signature scheme from scratch. Use the following interface:

- `Gen(minPrime)` generates a public/private keypair (512 bits) where $p, q > \text{minPrime}$.
- `Sign(privKey, msg)` returns a signature (integer).
- `Verify(pubKey, msg, signature)` returns boolean.

Both `Sign()` and `Verify()` take `msg` as integer and use SHA-512.