

Laboratory practice No. 2: Big O Notation

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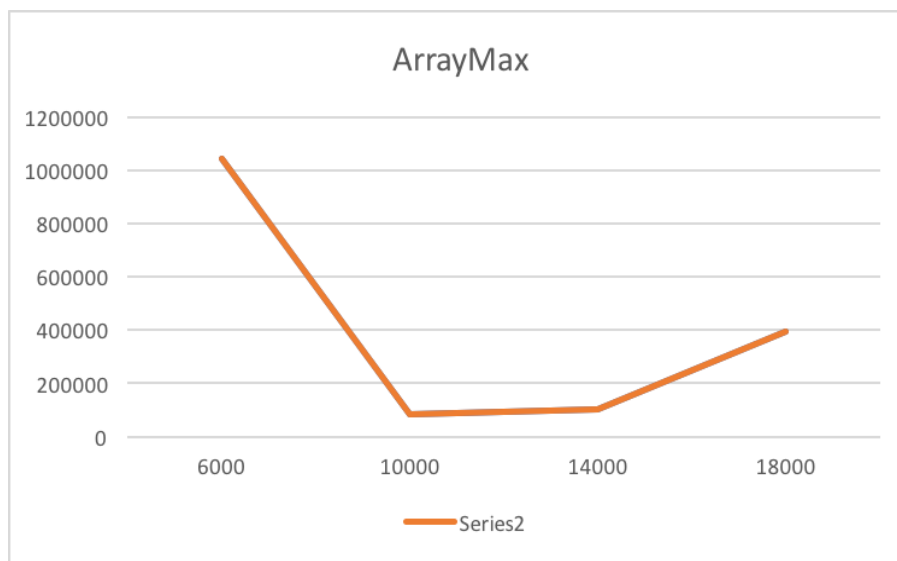
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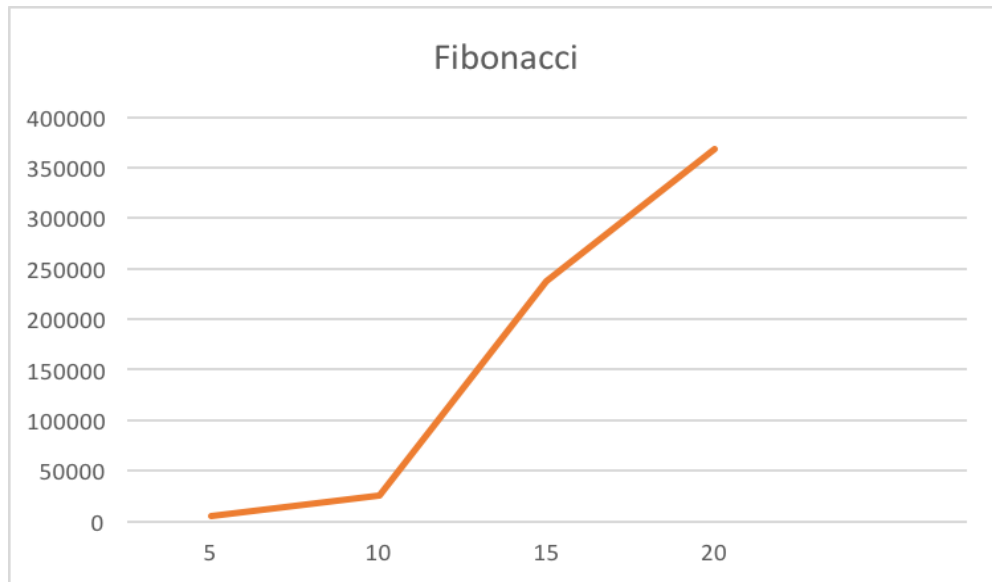
3) Practice for final project defense presentation

1.

	N = 6000	N = 10000	N = 14000	N = 18000
Array sum	More than one minute	More than one minute	More than one minute	More than one minute
Array max	1044868	85404	102393	395846
	N = 5	N = 10	N = 15	N = 20
Fibonacci	5743	26191	237029	368691

2.



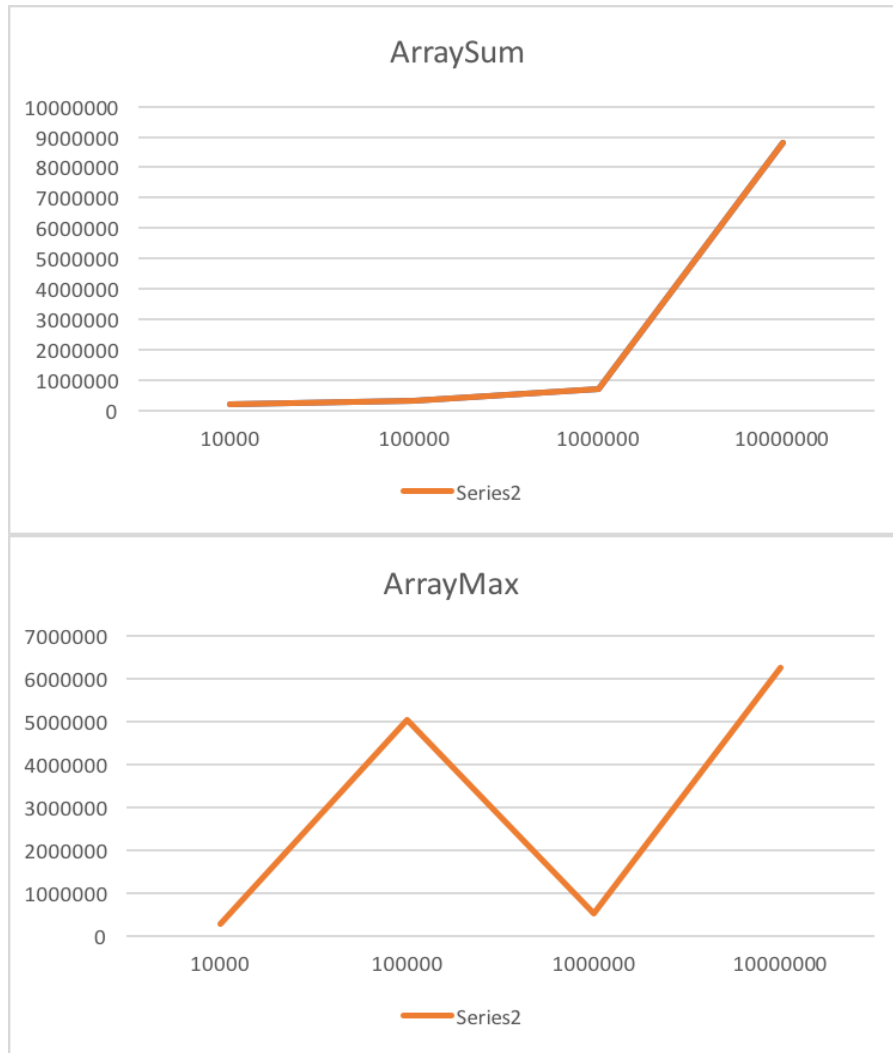


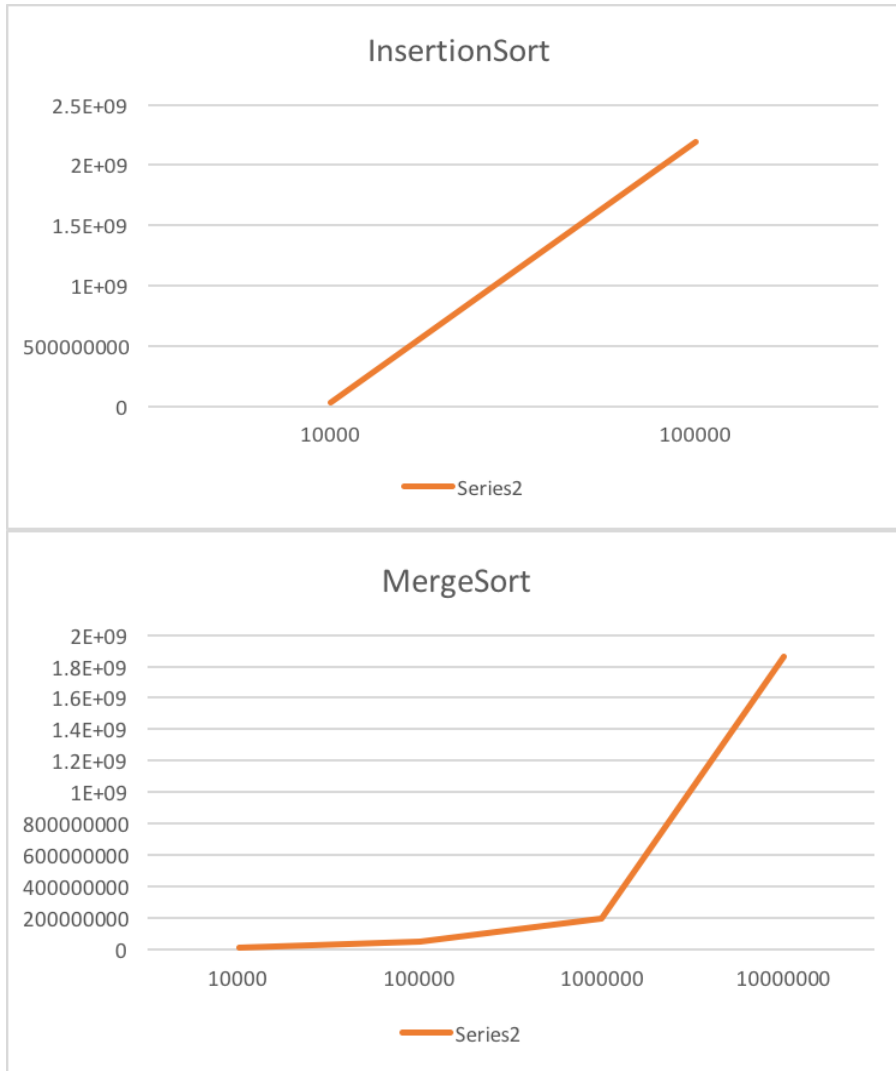
3. Given the experimental data above and the theoretical complexity, we conclude that the graphs do follow the mathematical model of their respective complexity. However, since we used only a handful of data (four values of n) there can be alterations that make the graphs look a little out of shape. But their overall shape does fit into the mathematical model of their complexity.

4.

	N = 10.000	N = 100.000	N = 1'000.000	N = 10'000.000
Array sum	190858	322614	682092	8801654
Array max	281835	5019117	516946	6235001
Insertion sort	34135064	2190836413	More than one minute	More than one minute
Merge sort	6331683	45672986	195207975	1862147417

5.





6. Given the experimental data above and the theoretical complexity, the graphs do not meet the mathematical model since we have few data (four values of n) and it is not possible to make out the ideal shape that the graph should have. Additionally, we can tell that the algorithm for InsertionSort is not very efficient considering that we were only able to get data for two values of n . The program took too long to process for remaining two values.
7. In InsertionSort, for bigger n values, the algorithm does not run or takes a very long time to process. Its complexity is $O(n^2)$.
8. Unlike InsertionSort, with bigger values of n , MergeSort is very efficient. Its complexity is $O(n \log n)$, which means that no matter how big the input data gets, the algorithm is able to process it fast and efficiently without much difference in time in between.
9. MergeSort is much more efficient than InsertionSort, considering that its complexity is $O(n \log n)$. We can also see this in the graphs above. The algorithm for MergeSort was able to successfully process all input data, whereas the algorithm for InsertionSort was only able to process two.
10. .

11. Codingbat Array2

- Sum13:
 $T(n) = c_1 + c_2 + c_3 + c_4 + n(c_5 + c_6 + c_7 + c_8) + c_9$
 $T(n) = O(n)$
- Sum67:
 $T(n) = c_1 + c_2 + n(c_3 + c_4 + c_5 + c_6 + c_7 + c_8 + c_9 + c_{10} + c_{11}) + c_{12}$
 $T(n) = O(n)$
- Has22:
 $T(n) = c_0 + n(c_1 + c_2 + c_3 + c_4) + c_5$
 $T(n) = O(n)$
- Luky13:
 $T(n) = c_0 + n(c_1 + c_2 + c_3) + c_4$
 $T(n) = O(n)$
- Sum28:
 $T(n) = c_0 + c_1 + n(c_2 + c_3 + c_4) + c_5$
 $T(n) = O(n)$

Codingbat Array3

- CanBalance:
 $T(n) = c_0 + c_1 + c_2 + n(c_3 + c_4 + c_5 + c_{10} + c_{11} + c_{12}) + n m(c_6 + c_7 + c_8 + c_9) + c_{13}$
 $T(n) = O(nm)$
- SeriesUp:
 $T(n) = c_0 + c_1 + c_2 + c_3 * n + n m(c_4 + c_5 + c_6) + c_7$
 $T(n) = O(nm)$
- Fix45:
 $T(n) = c + c_1 + n(c_2 + c_3 + c_4 + c_5 + c_6 + c_7) + c_8$
 $T(n) = O(n)$
- Fix34:
 $T(n) = c_0 + c_1 + n(c_2 + c_3 + c_4 + c_5) + n m(c_6 + c_7 + c_8) + c_9$
 $T(n) = O(nm)$
- SquareUp:
 $T(n) = c_0 + c_1 + c_2 + c_3 + c_4 + n(c_5 + c_6 + c_{10}) + n m(c_7 + c_8 + c_9) + c_{11}$
 $T(n) = O(nm)$

4) Practice for midterms

1. c
2. b
3. b
4. b
5. d
6. d
7. 1. $T(n) = (n-1) + c$
- 7.2 $O(n)$