

# Languages, automata and computation II

## Tutorial 9 – Weighted automata (2)

Winter semester 2024/2025

### Closure properties of rational series

**Reminder on rational series.** A *linear representation* over  $\Sigma$  is a triple  $A = (x, M, y)$  where the transition matrix  $M : \Sigma \rightarrow \mathbb{Q}^{k \times k}$  maps each letter  $a \in \Sigma$  to a  $k \times k$  rational matrix  $M_a$ ,  $x : \mathbb{Q}^{1 \times k}$  is a row vector, and  $y : \mathbb{Q}^{k \times 1}$  is a column vector. The transition matrix  $M$  is extended homomorphically to a function  $\Sigma^* \rightarrow \mathbb{Q}^{k \times k}$  (where matrices form a ring with the usual notions of matrix sum and product). The semantics of a linear representation is the series  $f : \Sigma^* \rightarrow \mathbb{Q}$  s.t.

$$f(w) = x \cdot M(w) \cdot y, \quad \text{for every } w \in \Sigma^*.$$

Call a function *rational* if it is of the form above.

**Exercise 1.** Show that the set of rational series is closed under the following operations.

1. Multiplication by a constant:  $(c \cdot f)(w) := c \cdot f(w)$ ,  $c \in \mathbb{Q}$ .
2. Addition:  $(f + g)(w) := f(w) + g(w)$ .
3. Pointwise product (Hadamard product):  $(f \cdot g)(w) := f(w) \cdot g(w)$ .
4. Concatenation product (Cauchy product):  $(f * g)(w) := \sum_{uv=w} f(u) \cdot g(v)$ .
5. Iteration (w.r.t. concatenation), when  $f(\varepsilon) = 0$ :  $f^* := f^0 + f^1 + f^2 + \dots$ , where  $f^0(w)$  is 1 if  $w = \varepsilon$  and 0 otherwise, and  $f^{n+1} = f^n * f$  for every  $n \geq 0$ .

The *support* of a series  $f : \Sigma^* \rightarrow \mathbb{Q}$ , denoted  $\text{supp } f$ , is the set of words where  $f$  is nonzero. A *polynomial* in noncommutative variables is a series with finite support.

**Exercise 2.** Show that polynomials are rational series.

**Exercise 3** (Concatenation inverses). Under which condition does  $f$  have an inverse w.r.t. the concatenation product? In the positive case, find an expression for the inverse.

## Regular expressions

**Exercise 4** (Kleene-Schützenberger theorem). Call a function *regular* if it can be generated by the following abstract grammar

$$f, g ::= p \mid c \cdot f \mid f + g \mid f * g \mid f^*,$$

where  $p$  is a polynomial (function with finite support),  $c \in \mathbb{Q}$  is a constant, and iteration  $f^*$  is only applied when defined. Show that a function is regular iff it is rational.

## Supports

A *rational support* is the support of a rational function. Since we do not consider any other kind of support, we just say “support” for “rational support” in the following.

**Exercise 5.** 1. Show that the class of supports includes all regular languages.

2. Are there nonregular supports?

**Exercise 6.** Are the following problems decidable for rational supports:

1. emptiness?
2. universality?
3. equivalence?
4. inclusion?

**Exercise 7.** Are rational supports closed under

1. intersection?
2. union?
3. concatenation?
4. Kleene star?
5. complement?