

Languages, automata and computation II

Tutorial 9 – Weighted automata (2)

Winter semester 2024/2025

Closure properties of rational series

Reminder on rational series. A *linear representation* over Σ is a triple $A = (x, M, y)$ where the transition matrix $M : \Sigma \rightarrow \mathbb{Q}^{k \times k}$ maps each letter $a \in \Sigma$ to a $k \times k$ rational matrix M_a , $x : \mathbb{Q}^{1 \times k}$ is a row vector, and $y : \mathbb{Q}^{k \times 1}$ is a column vector. The transition matrix M is extended homomorphically to a function $\Sigma^* \rightarrow \mathbb{Q}^{k \times k}$ (where matrices form a ring with the usual notions of matrix sum and product). The semantics of a linear representation is the series $f : \Sigma^* \rightarrow \mathbb{Q}$ s.t.

$$f(w) = x \cdot M(w) \cdot y, \quad \text{for every } w \in \Sigma^*.$$

Call a function *rational* if it is of the form above.

Exercise 1. Show that the set of rational series is closed under the following operations.

1. Multiplication by a constant: $(c \cdot f)(w) := c \cdot f(w)$, $c \in \mathbb{Q}$.
2. Addition: $(f + g)(w) := f(w) + g(w)$.
3. Pointwise product (Hadamard product): $(f \cdot g)(w) := f(w) \cdot g(w)$.
4. Concatenation product (Cauchy product): $(f * g)(w) := \sum_{uv=w} f(u) \cdot g(v)$.
5. Iteration (w.r.t. concatenation), when $f(\varepsilon) = 0$: $f^* := f^0 + f^1 + f^2 + \dots$, where $f^0(w)$ is 1 if $w = \varepsilon$ and 0 otherwise, and $f^{n+1} = f^n * f$ for every $n \geq 0$.

The *support* of a series $f : \Sigma^* \rightarrow \mathbb{Q}$, denoted $\text{supp } f$, is the set of words where f is nonzero. A *polynomial* in noncommutative variables is a series with finite support.

Exercise 2. Show that polynomials are rational series.

Exercise 3 (Concatenation inverses). Under which condition does f have an inverse w.r.t. the concatenation product? In the positive case, find an expression for the inverse.

Regular expressions

Exercise 4 (Kleene-Schützenberger theorem). Call a function *regular* if it can be generated by the following abstract grammar

$$f, g ::= p \mid c \cdot f \mid f + g \mid f * g \mid f^*,$$

where p is a polynomial (function with finite support), $c \in \mathbb{Q}$ is a constant, and iteration f^* is only applied when defined. Show that a function is regular iff it is rational.

Supports

A *rational support* is the support of a rational function. Since we do not consider any other kind of support, we just say “support” for “rational support” in the following.

Exercise 5. 1. Show that the class of supports includes all regular languages.

2. Are there nonregular supports?

Exercise 6. Are the following problems decidable for rational supports:

1. emptiness?
2. universality?
3. equivalence?
4. inclusion?

Exercise 7. Are rational supports closed under

1. intersection?
2. union?
3. concatenation?
4. Kleene star?
5. complement?