

Game Semantics for Probabilistic μ -Calculi

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Game Semantics for Modal μ -Calculus

Modal μ -Calculus \Leftrightarrow **2-Player Parity Games**

Model M + Formula ϕ \rightsquigarrow Game $G(M, \phi)$

ϕ holds at M iff Player Verifier (\Diamond) wins $G(M, \phi)$

Determinacy: Either Verifier (\Diamond) wins or Refuter (\Box) wins.

Game Semantics for Probabilistic Modal μ -Calculus

Prob. Modal μ -Calculus $\Leftrightarrow 2^{\frac{1}{2}}$ -Player Parity Games

Probabilistic Model M

+

Formula ϕ

\rightsquigarrow

Game $G(M, \phi)$

ϕ holds at M
with probability p

iff

Player Verifier (\Diamond) wins $G(M, \phi)$
with probability p

Determinacy: $\text{Val}(\Diamond) = 1 - \text{Val}(\Box)$

- ▶ $\text{Val}(\Diamond)$ = probability of winning that Player \Diamond can ensure,
- ▶ $\text{Val}(\Box)$ = probability of winning that Player \Box can ensure.

Problem:

Probabilistic Modal μ -Calculus is not sufficiently expressive:

- ▶ e.g., cannot encode probabilistic CTL (**pCTL**).

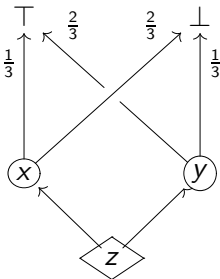
- ▶ Extensions of Prob. Modal μ -Calculus
 - ▶ Capable of encoding pCTL,
- ▶ with corresponding **Game Semantics**

New class of Games

- ▶ Generalizing $2\frac{1}{2}$ -player Parity Games.

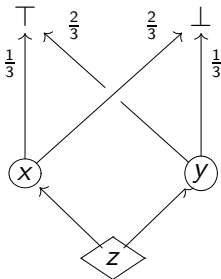
Ordinary $2\frac{1}{2}$ -player games have three kinds of states:

- Verifier (\diamond), Refuter (\square) and Nature (\circ)

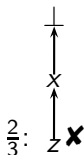


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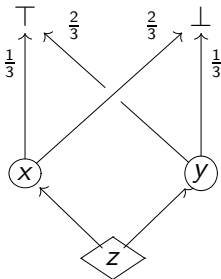


Possible outcomes:

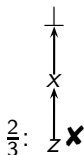


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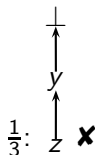
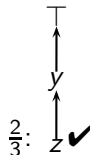
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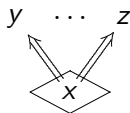
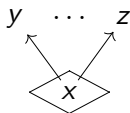
OR:



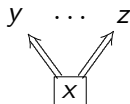
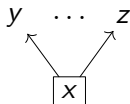
Tree Games

In Tree Games there are **FIVE** kinds of states:

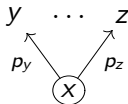
- ▶ Two types of Verifier (\diamond) states,

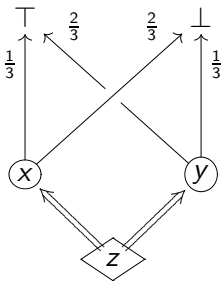


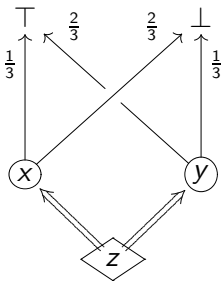
- ▶ Two types of Refuter (\square) states,



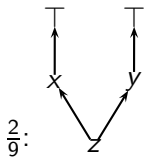
- ▶ Nature (\bigcirc)

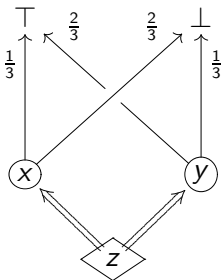




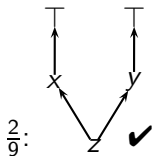


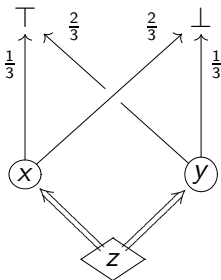
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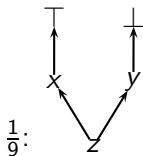
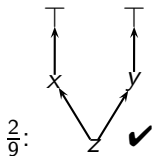


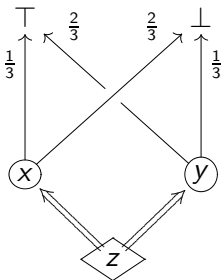
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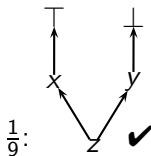
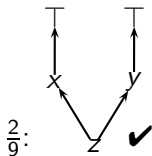


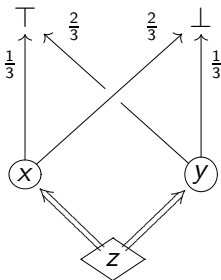
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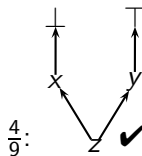
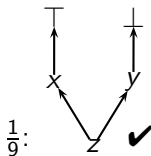
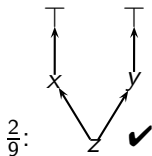


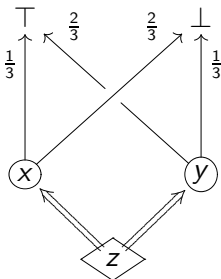
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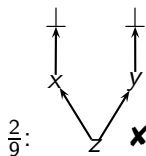
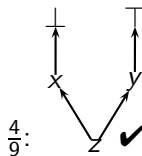
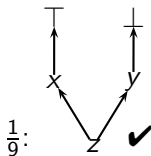
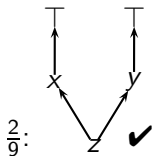


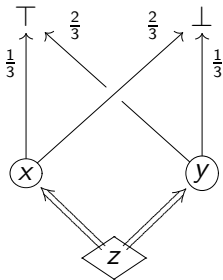
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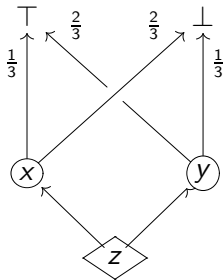


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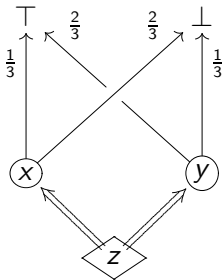




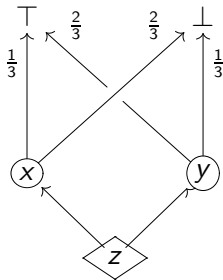
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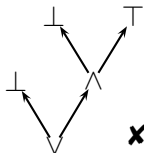
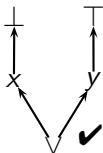
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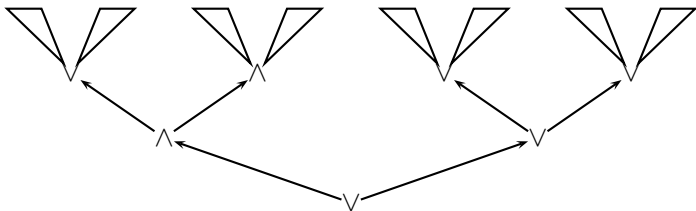
With probabilistic behavior: Parallel-OR \neq Choice-OR

Outcomes of Tree Games are trees: **branching plays**.



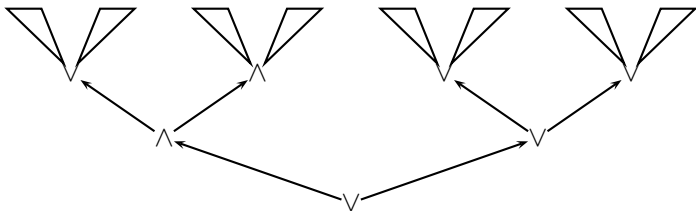
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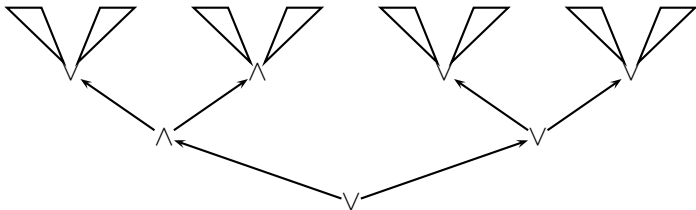
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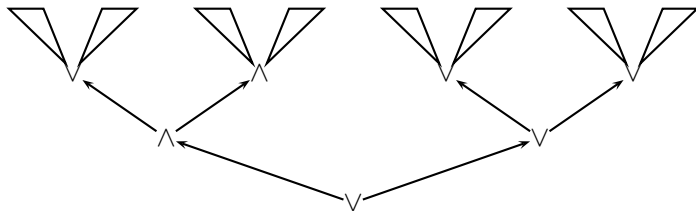


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- ▶ Parity Condition.

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They can be regarded as infinitary Boolean formulas.

- ▶ Infinite branches are either winning for \diamond or for \square .
 - ▶ Parity Condition.
- ▶ “Value” of the branching play is given by a Parity Game.

We call these games

Two-Player Stochastic **meta-** Parity Games

- ▶ The moves of \diamond , \square and \bigcirc determine a play
- ▶ The play is a tree (Branching play),
- ▶ itself interpreted as a Parity Game to determine winner.

Two-Player Stochastic meta-Parity games

- ▶ are games of **imperfect information**.
 - ▶ Each thread executes **ignoring** other parallel threads.
- ▶ Do not seem to be straightforwardly reducible to other known classes of games.

Outcomes are not finite or infinite **paths**,

- ▶ they are finite or infinite trees: **branching plays**.
- ▶ As usual, space BP of possible branching plays is **uncountable**
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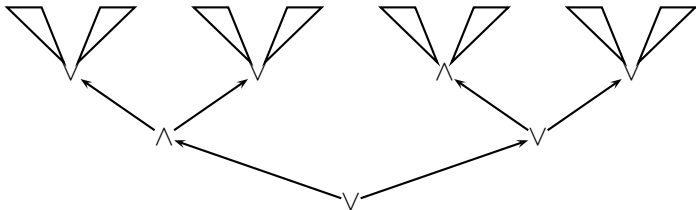
Δ_2^1 in the Projective Hierarchy.

DANGER : It is consistent with ZFC that some Δ_2^1 set is not measurable!!!

The set $\mathcal{W} \subseteq \text{BP}$ is defined as

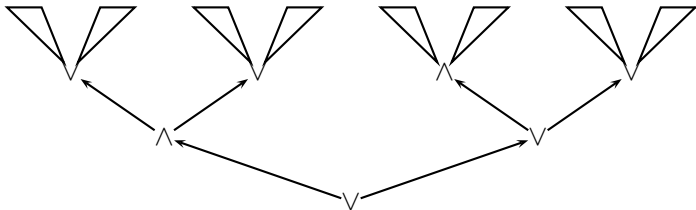
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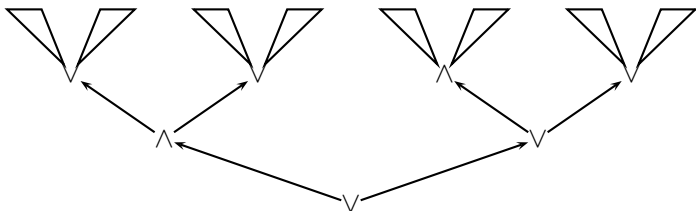
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Inductive Characterization: $\mathcal{W} = \bigcup_{\alpha < \omega_1} \mathcal{W}^\alpha$

- ▶ \mathcal{W} is the limit of a transfinite increasing chain.
- ▶ ω_1 is the smallest uncountable ordinal. $|\omega_1| = \aleph_1$.

Measure theory works well with **countable** chains:

$$X_0 \subseteq X_1 \subseteq \dots X_n \dots$$

$$\text{if } X = \bigcup_{\alpha < \omega} X^\alpha \quad \text{then} \quad \mu(X) = \bigsqcup_{\alpha < \omega} \mu(X^\alpha)$$

But not much can be said about transfinite chains:

$$\text{if } X = \bigcup_{\alpha < \omega_1} X^\alpha \quad \text{then} \quad \mu(X) \geq \bigsqcup_{\alpha < \omega_1} \mu(X^\alpha)$$

Set Theory: Martin's Axiom at the cardinal \aleph_1 (MA_{\aleph_1})

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Remark: $\text{ZFC} + \text{MA}_{\aleph_1} \vdash \aleph_1 < 2^{\aleph_0}$,

Theorem ($\text{ZFC} + \text{MA}_{\aleph_1}$) **[Thesis]**: Determinacy.

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Problem left open: is ZFC sufficient?

- Our theorems can not be refuted in ZFC, if ZFC is consistent.

THANK YOU