Branching-Time Model Checking Gap-Order Constraint Systems

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Gap Clauses/Constraints

Def: *positive* Gap Constraints

$$\bigwedge_{0 \leq i \leq n} (x_i - y_i \geq k_i)$$

where x_i, y_i are integer variables or constants and $k_i \in \mathbb{N}$.

positive GC are not negation-closed!

Write

$$\begin{aligned} \mathit{Var} &= \{x, y, \dots\} \text{ for the } \textit{variables} \\ \mathit{Const} &\subset \mathbb{Z} \text{ for the } \textit{constants} \text{ and} \\ \mathit{Val} &\text{ for the set of } \textit{valuations } \nu : \mathit{Var} \to \mathbb{Z}. \end{aligned}$$

Gap Constraints

1. can characterise subsets $S \subseteq \mathit{Val}$ (of satisfied valuations)
2. can determine how valuations evolve: For instance,

$$x - x' \geq 0$$

means the value of x does not increase.

Gap-Order Constraint Systems

Definition (CGS)

are given by finite sets

Var of variables ranging over $\mathbb{Z}$,

Const of integer constants, and

Δ of *positive* transitional gap contraints.

Step semantics:

$$\nu \longrightarrow \nu' \text{ iff } \nu \oplus \nu' \models \mathcal{C} \text{ for some } \mathcal{C} \in \Delta.$$

Example: lex. Countdown of (y, x)

$\mathcal{C}_1 = (x > x' \geq 0) \wedge (y' = y)$
$\mathcal{C}_2 = (x \leq x') \wedge (y > y' \geq 0)$

Overapproximating Counter Machines

Zero-tests

$(c_1 = 0)$

Finite Control

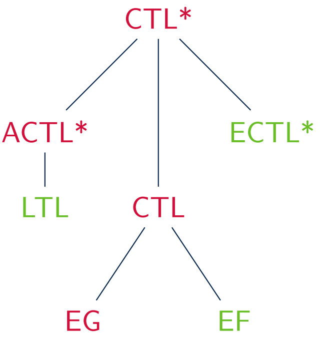
$(state = 0) \wedge (state' = 1)$ // $s_0 \longrightarrow s_1$

Increments/Decrements **are imprecise!**

$(c_1' - c_1 \geq 0)$ //c_1++

$(c_1 - c_1' \geq 0) \wedge (c_1' - 0 \geq 0)$ //c_1- -

Model Checking GCS



[Čer94] Model checking CTL undecidable, but LTL is decidable for IRA.

[BP12] LTL and ECTL* are *PSPACE*-complete; ACTL* is undecidable for GCS.

We EF is decidable and EG undecidable for GCS.

EF over Gap Clauses

Syntax

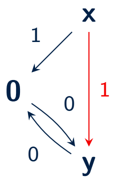
$$\psi ::= \mathcal{C} \mid \neg\psi \mid \psi \vee \psi \mid X\psi \mid EF\psi \mid \cancel{EG\psi} \quad \cancel{(\psi U \psi)}$$

- EF model checking GCS is decidable.
- Proof by finding finite representation for $Sat(\mathcal{C})$ that is closed under negation, union, Pre and Pre*.

Monotonicity Graphs

Gap Constraints as finite labeled graphs over *Var* $\cup$ *Const*

$$\begin{aligned}\mathcal{C} = &(x - 0 \geq 1)\\ &\wedge(y - 0 \geq 0)\\ &\wedge(0 - y \geq 0)\end{aligned} \quad \sim$$



- *Degree* of MG: inverse of minimal negative value
- *Closure* of MG has same denotation
- represent $S \subseteq \mathit{Val}$ by finite sets of (arbitrary) MG. Example: $\{M_{\mathcal{C}}\}$ represents $S = \mathit{Sat}(\mathcal{C}) = \{\nu \mid \nu(x) > \nu(y) = 0\}$.

$$\psi ::= \mathcal{C} \mid \psi \vee \varphi \mid \neg\psi \mid X\psi \mid EF\psi$$

Negation

$$\neg Rep(S) = \neg\{M_0, M_1, \ldots, M_k\}$$

$$\sim$$

$$\neg\mathcal{C}_{M_0} \wedge \neg\mathcal{C}_{M_1} \wedge \cdots \wedge \neg\mathcal{C}_{M_k}$$

... is a Gap-Formula in DNF.

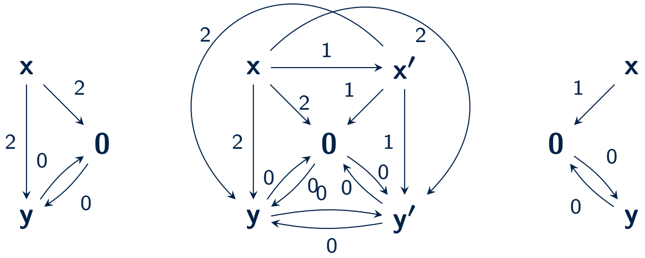
- $\rightsquigarrow$ propagate negations to clauses
- $\rightsquigarrow$ negate clauses $\leftarrow$ increases degree

$$x - y \not\geq k \iff y - x \geq -(k-1)$$

- $\rightsquigarrow$ bring to DNF
- $\rightsquigarrow$ interpret as set of MG

Computing *Pre*

$$S \;= \{\nu \mid \nu(x) > \nu(y) = 0\}$$
$$\mathcal{C}_1 = (x - x' \geq 1) \wedge (y' - y \geq 0) \wedge (y - y' \geq 0) \wedge (x' - 0 \geq 0)$$



$Pre(\mathcal{C}_1, S)$ S

NB: Degree does not increase

Computing Pre^*

Definition ($\sqsubseteq$)

$M \sqsubseteq M'$ if $M(x, y) \leq M'(x, y)$ for all $x, y \in Var \cup Const$.

1. $M \sqsubseteq M'$ implies $[\![M]\!] \supseteq [\![M']\!]$
2. $\sqsubseteq$ is a well-order over MG^n

Compute $Pre^*(M)$:

iteratively unfold the finite! backwards coverability tree and take the union of all nodes. . .

EF Model Checking GCS is decidable

Theorem

For given GCS and EF formula φ, the set $Sat(\varphi)$ is effectively Gap-definable.

Works even with

- arbitrary gap-formulae as atoms and
- positive (trans.) gap-constraints on X/EF operators.

WIP: Equivalence Checking

1. Bisimulation
   - $GCS \approx FS$ is decidable using char. formulae in EF
   - Strong Bisimulation $GCS \sim GCS$ is undecidable
2. Trace inclusion/equivalence
   - $GCS \subseteq GCS$ is in EXPSPACE
   - Universality is EXPSPACE-hard
3. Simulation Preorder
   - $GCS \preceq FS$ and vv. are decidable (wqo)
   - $GCS \preceq GCS$? WIP.

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