

Emanuel Kieroński, **Jakub Michaliszyn**, Jan Otop

Modal logic

- Many different modal logics (K4, S5, CTL, LTL, ATL, HS, CTL*K)
- Many applications in verification, planning, linguistics
- Many proofs, many papers

Modal logic

- Many different modal logics (K4, S5, CTL, LTL, ATL, HS, CTL*K)
- Many applications in verification, planning, linguistics
- Many proofs, many papers
- Our area of interest: a comprehensive study on the satisfiability problem.

Kripke semantics

- Kripke structure — a *frame* + a *labelling*.

Kripke semantics

- Kripke structure — a *frame* + a *labelling*.
- \mathcal{K} -SAT — local satisfiability problem w.r.t. \mathcal{K} .
- \mathcal{K} -GSAT — global satisfiability problem w.r.t. \mathcal{K} .

Kripke semantics

- Kripke structure — a *frame* + a *labelling*.
- \mathcal{K} -SAT — local satisfiability problem w.r.t. \mathcal{K} .
- \mathcal{K} -GSAT — global satisfiability problem w.r.t. \mathcal{K} .

Our ultimate goal

For all first-order definable classes \mathcal{K} , determine the decidability and complexity of \mathcal{K} -SAT and \mathcal{K} -GSAT.

We are also interested in finite satisfiability.

Negative results

- (E. Hemaspaandra, "The Price of Universality", 1996)
 \mathcal{K} -GSAT is undecidable for some \forall FO-definable \mathcal{K} .
- (E. Hemaspaandra, H. Schnoor, MFCS 2011)
 \mathcal{K} -SAT is undecidable for some \forall FO-definable \mathcal{K} .

Negative results

- (E. Hemaspaandra, "The Price of Universality", 1996)
 \mathcal{K} -GSAT is undecidable for some \forall FO-definable \mathcal{K} .
- (E. Hemaspaandra, H. Schnoor, MFCS 2011)
 \mathcal{K} -SAT is undecidable for some \forall FO-definable \mathcal{K} .
- (E. Kieroński, J. Michaliszyn, J. Otop, FSTTCS 2011)
 \mathcal{K} -GSAT and \mathcal{K}' -SAT are undecidable for some \forall FO³-definable \mathcal{K} and \mathcal{K}' (holds also for finite satisfiability).

$$\neg xRy \vee \neg xRz \vee yRz \vee zRy \vee yRx \vee zRx$$

Positive results

Standard translation

Is φ satisfied w.r.t. the class defined by Φ ? \rightarrow Is $\Phi \wedge ST(\varphi)$ satisfiable?

Positive results

Standard translation

Is φ satisfied w.r.t. the class defined by Φ ? \rightarrow Is $\Phi \wedge ST(\varphi)$ satisfiable?

Decidability in many interesting cases (even multimodal), including

- FO^2 :
 - ▶ with one transitive relation (W. Szwast, L. Tendera, 2012),
 - ▶ with counting quantifiers (I. Pratt-Hartmann, 2005),
 - ▶ with two equivalence relations (E. Kieroński, J. Michaliszyn, I. Pratt-Hartmann, L. Tendera, 2012).
- Guarded Fragment:
 - ▶ with fixed points (E. Grädel, I. Walukiewicz, 1999),
 - ▶ with the transitive closure operator in guards (J. Michaliszyn, 2009).

Positive results

Standard translation

Is φ satisfied w.r.t. the class defined by Φ ? \rightarrow Is $\Phi \wedge ST(\varphi)$ satisfiable?

Decidability in many interesting cases (even multimodal), including

- FO^2 :
 - ▶ with one transitive relation (W. Szwast, L. Tendera, 2012),
 - ▶ with counting quantifiers (I. Pratt-Hartmann, 2005),
 - ▶ with two equivalence relations (E. Kieroński, J. Michaliszyn, I. Pratt-Hartmann, L. Tendera, 2012).
- Guarded Fragment:
 - ▶ with fixed points (E. Grädel, I. Walukiewicz, 1999),
 - ▶ with the transitive closure operator in guards (J. Michaliszyn, 2009).
- High complexity.

Positive results

J. Michaliszyn, J. Otop, LICS 2012

For any \mathcal{K} definable by universal Horn formulas, \mathcal{K} -SAT and \mathcal{K} -GSAT are decidable.

Positive results

J. Michaliszyn, J. Otop, LICS 2012

For any \mathcal{K} definable by universal Horn formulas, \mathcal{K} -SAT and \mathcal{K} -GSAT are decidable.

J. Michaliszyn, E. Kieroński, AIML 2012

Also finite satisfiability of modal logic is decidable w.r.t. the classes definable by universal Horn formulas.

General satisfiability

Type	\mathcal{K}_ϕ -GSAT	\mathcal{K}_ϕ -SAT
S1+	EXPTIME-c	PSPACE-c
S1-	PSPACE-c	NP-c
S2+	NP-c	PSPACE-c
S2-	NP-c	NP-c
S3+	impossible	
S3-	NP-c	NP-c

Except for some trivial formulas like $xRx \wedge (xRx \Rightarrow \perp)$.

Finite satisfiability

Type of Φ	\mathcal{K}_Φ -GFINSAT	\mathcal{K}_Φ -FINSAT
S3+, S3-	FMP, NP-c	
S2+, S2-	NEXPTIME	
S1+ & “merges”	Lack of FMP (always!), PSPACE-c	FMP, PSPACE-c
S1+ & not “merges”	FMP, EXPTIME-c	FMP, PSPACE-c
S1-	FMP, PSPACE-c	FMP, NP-c

Finite vs. General

J. Michaliszyn, J. Otop, P. Witkowski, Gandalf 2012

- There is an undecidable logic that is finitely decidable
- There is a decidable logic that is finitely undecidable

Transitivity

- Transitive modalities are popular in practice:
- F , G of LTL
- B , D , L of HS logic
- K_i , C_G of epistemic logic

J. Michaliszyn, J. Otop, CSL 2013

For any \mathcal{K} of transitive frames definable by universal formulas, \mathcal{K} -SAT and \mathcal{K} -GSAT are decidable. The same holds for the finite satisfiability problem.

So what?

So what?

Our ultimate goal

For all first-order definable classes \mathcal{K} , classify \mathcal{K} -SAT, \mathcal{K} -GSAT (and their finite counterparts) w.r.t. the decidability status and the complexity.

So what?

Our ultimate goal

For all first-order definable classes \mathcal{K} , classify \mathcal{K} -SAT, \mathcal{K} -GSAT (and their finite counterparts) w.r.t. the decidability status and the complexity.

Why?

- Better understanding
- Easy modifications
- Unified theory

So what?

Our ultimate goal

For all first-order definable classes \mathcal{K} , classify \mathcal{K} -SAT, \mathcal{K} -GSAT (and their finite counterparts) w.r.t. the decidability status and the complexity.

Why?

- Better understanding
- Easy modifications
- Unified theory

The “metaproblem”

Input: A first-order formula Φ that defines a class of frames \mathcal{K} .

Question: Is \mathcal{K} -SAT decidable?

Is the metaproblem decidable?

Thank you for your attention!

Summary

- We study the **satisfiability** problem of **modal logic** over **first-order** definable classes of **frames**.
- In some cases the problem is **undecidable**.
- There are wide classes of formulas that lead to **decidable** problems (Horn formulas, transitive formulas, FO^2 , GF).
- Our goal: to classify them all.

Open: Is the “metaproblem” decidable?

Input: A first-order formula Φ that defines a class of frames \mathcal{K} .

Question: Is \mathcal{K} -SAT decidable?