

L-visibility drawings of IC-planar graphs[☆]

Giuseppe Liotta, Fabrizio Montecchiani^{*}

Dipartimento di Ingegneria, Università degli Studi di Perugia, Italy



ARTICLE INFO

Article history:

Received 31 July 2015

Received in revised form 17 November 2015

Accepted 18 November 2015

Available online 28 November 2015

Communicated by Ł. Kowalik

Keywords:

Visibility representations

IC-planar graphs

Computational Geometry

ABSTRACT

An IC-plane graph is a topological graph where every edge is crossed at most once and no two crossed edges share a vertex. We show that every IC-plane graph has a visibility drawing where every vertex is of the form $\{\perp, \lrcorner, \top, \sqcap\}$, and every edge is either a horizontal or vertical segment. As a byproduct of our drawing technique, we prove that every IC-plane graph has a RAC drawing in quadratic area with at most two bends per edge.

© 2015 Elsevier B.V. All rights reserved.

1. Introduction

A *visibility drawing* Γ of a planar graph G maps the vertices of G into non-overlapping horizontal segments (*bars*), and the edges of G into vertical segments (*visibilities*), each connecting the two bars corresponding to its two end-vertices. Visibilities intersect bars only at their extreme points. Γ is a *strong* visibility drawing if there exists a visibility between two bars if and only if there exists an edge in G between the corresponding vertices. Every biconnected planar graph admits a strong visibility drawing (see, e.g., [15]). Conversely, if a visibility may not correspond to an edge of the graph, then Γ is a *weak* visibility drawing. Since every planar graph can be augmented to a biconnected planar graph by adding edges, every planar graph admits a weak visibility drawing.

The problem of extending visibility drawings to non-planar graphs has been first studied by Dean et al. [4]. They introduce *bar k -visibility drawings*, which are visibility

drawings where each bar can see through at most k distinct bars. In other words, each visibility segment can intersect at most k bars, while each bar can be intersected by arbitrary many visibility segments. The graphs that admit a bar 1-visibility drawing are called *1-visible*. Brandenburg et al. and independently Evans et al. prove that *1-planar graphs*, i.e., those graphs that can be drawn with at most one crossing per edge, are 1-visible [2,8]. They focus on a *weak* model, where there is a visibility through at most k bars if there is an edge, while the converse may not be true. In fact, having a strong model would be too restrictive in the case of bar k -visibility drawings. For example, it is easy to see that a cycle of length at least four does not admit a strong bar 1-visibility drawing [2]. In terms of readability, a clear benefit of bar k -visibility drawings is that the crossings form right angles. *Right-angle crossing (RAC) drawings* and their advantages in terms of readability have been extensively studied in the graph drawing literature (see, e.g., [7,11]). However, in a bar k -visibility drawing crossings involve bars and visibilities, i.e., vertices and edges. These crossings are arguably less intuitive than crossings between edges.

Evans et al. introduce a new model of visibility drawings, called *L-visibility drawings* [9]. Their aim is to simultaneously represent two plane *st*-graphs G_r and G_b (whose union might be non-planar). They assume a *strong* model,

[☆] Research supported in part by the MIUR project AMANDA “Algorithms for MAssive and Networked DATa”, prot. 2012C4E3KT_001.

^{*} Corresponding author.

E-mail addresses: giuseppe.liotta@unipg.it (G. Liotta), fabrizio.montecchiani@unipg.it (F. Montecchiani).

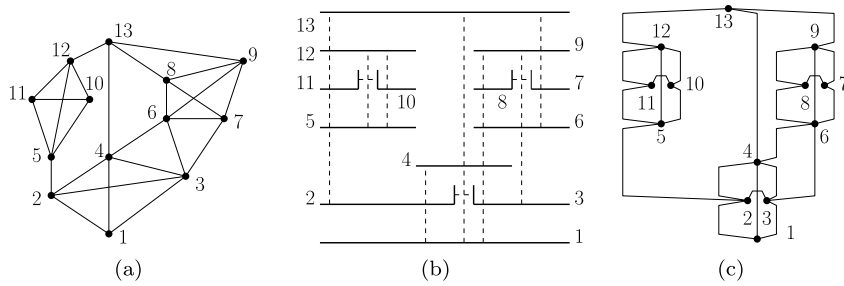


Fig. 1. (a) An IC-plane graph G . (b) A L-visibility drawing of G . (c) A RAC drawing of G with at most two bends per edge.

where each vertex is represented by a horizontal bar and a vertical bar that share an extreme point, i.e. it is an *L-shape* in the set $\{\mathbf{L}, \mathbf{J}, \mathbf{7}, \mathbf{F}\}$. Two L-shapes are connected by a vertical (horizontal) visibility segment if and only if there exists an edge in G_r (G_b) between the corresponding vertices, no two L-shapes cross one another, and visibilities intersect bars only at their extreme points. A clear advantage of this kind of drawing is that the only possible crossings are between vertical and horizontal visibilities, i.e., between edges of the graph. Furthermore, similar to bar k -visibilities, these crossings form right angles.

In this paper we initiate the study of *weak* L-visibility drawings of non-planar graphs. We focus on the class of graphs that admit a drawing where each edge is crossed at most once, and no two crossed edges share an end-vertex. These graphs are called *IC-planar graphs* (see Fig. 1(a) for an example). Their chromatic number is at most five [12], and they have at most $13n/4 - 6$ edges, which is a tight bound [17]. Recognizing IC-planar graphs is NP-hard [3]. Our main contribution is summarized by the following theorem, proved in Section 3. See Fig. 1(b) for an example of a drawing computed by using Theorem 1.

Theorem 1. Every n -vertex IC-plane graph G admits a L-visibility drawing in $O(n^2)$ area, which can be computed in $O(n)$ time.

We remark that Theorem 1 contributes to the rapidly growing literature devoted to the problem of drawing graphs that are “nearly planar” in some sense, i.e. graphs where only some types of edge crossings are allowed (for example, an edge can be crossed at most a constant number of times); see e.g., [13] for references. In particular, Brandenburg et al. have recently described a cubic-time algorithm that computes IC-planar drawings with right-angle crossings and straight-line edges [3]. However these drawings may require exponential area, which is proved to be worst-case optimal [3]. Brandenburg et al. leave as an open problem to study techniques that compute IC-planar drawings in polynomial area and with good crossing resolution [3]. We also recall that every graph admits a RAC drawing with at most three bends per edge [6], while testing whether a graph has a straight-line RAC drawing is NP-hard [1]. The following corollary follows as a byproduct of Theorem 1 (see also Fig. 1(c)).

Corollary 1. Every n -vertex IC-plane graph G admits a RAC drawing with at most two bends per edge in $O(n^2)$ area, which can be computed in $O(n)$ time.

2. Preliminaries

We assume familiarity with basic graph drawing concepts, see also [5].

Planarity and connectivity. A graph $G = (V, E)$ is *simple*, if it contains neither loops nor multiple edges. We consider simple graphs, if not otherwise specified. A *drawing* Γ of G maps each vertex of V to a point of the plane and each edge of E to a Jordan arc between its two end-points. We only consider *simple drawings*, i.e., drawings such that the arcs representing two edges have at most one point in common, which is either a common end-vertex or a common interior point where the two arcs properly cross. A drawing is *planar* if no two arcs representing two edges cross. A planar drawing divides the plane into topologically connected regions, called *faces*. The unbounded region is called the *outer face*. A *planar embedding* of a graph is an equivalence class of planar drawings that define the same set of faces. A graph with a given planar embedding is a *plane graph*. For a non-planar drawing, we can still talk about embedding considering that the boundary of a face may consist of portions of arcs between vertices and/or crossing points.

A graph is *biconnected* if it remains connected after removing any one vertex. A directed graph (a digraph for short) is biconnected if its underlying undirected graph is biconnected. A *topological numbering* of a digraph is an assignment, X , of numbers to its vertices such that $X(u) < X(v)$ for every edge (u, v) . A graph admits a topological numbering if and only if it is acyclic. An acyclic digraph with a single source s and a single sink t is called an *st-graph*. A *plane st-graph* is an *st-graph* that is planar and embedded such that s and t are on the boundary of the outer face. In any *st-graph*, the presence of the edge (s, t) guarantees that the graph is biconnected. In the following we consider *st-graphs* that contain the edge (s, t) , as otherwise it can be added without violating planarity. Let G be a plane *st-graph*, then for each vertex v of G the incoming edges appear consecutively around v , and so do the outgoing edges. Vertex s only has outgoing edges, while vertex t only has incoming edges. This particular transversal structure is known as a *bipolar orientation* [14,15]. Each face f of G is bounded by two directed paths with a common origin and destination, called the *left path* and *right path* of f .

IC-planar graphs. We recall some definitions also given in [3]. A drawing is IC-planar if each edge is crossed at most once, and any two crossed edges do not share an end-vertex. See Fig. 1(a) for an illustration. An *IC-planar*

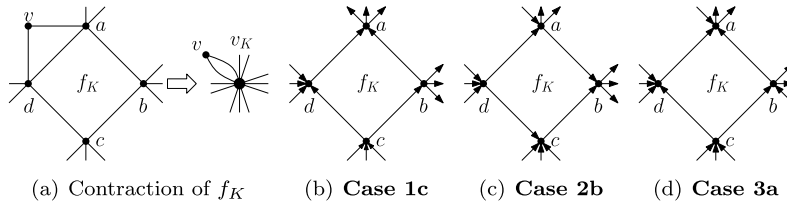


Fig. 3. Illustration for the proof of Lemma 1.

embedding is an embedding derived from an IC-planar drawing. A graph with a given IC-planar embedding is an *IC-plane* graph. Thomassen [16] characterized the possible crossing configurations that occur in a 1-planar drawing, i.e., a drawing where each edge is crossed at most once. This characterization applied to IC-planar drawings gives rise to the following property, where an X-crossing is of the type described in Fig. 2(a), and a B-crossing is of the type described in Fig. 2(b) (the solid edges only).

Property 1. (See [3].) *Every crossing of an IC-plane graph is either an X- or a B-crossing.*

A kite K is a graph isomorphic to K_4 together with an embedding such that all the vertices are on the boundary of the outer face. This implies that two edges of K cross each other, while the other four edges are not crossed and belong to the boundary of the outer face; see Fig. 2(a). Consider a pair of crossing edges of an IC-plane graph G , such that their four end-vertices induce a kite K . The kite K is *empty*, if in G there is no other vertex inside the internal faces of K . The following property is a consequence of the more general Lemma 1 in [3] (in particular of cases c1 and c2 of that lemma).

Property 2. (See [3].) *Let $G = (V, E)$ be an n -vertex IC-plane graph. It is possible to augment G to a biconnected IC-plane graph $G^+ = (V, E^+)$ (with a possibly different embedding), where $E \subseteq E^+$, such that the end-vertices of each pair of crossing edges of G^+ induce an empty kite. This can be done in $O(n)$ time.*

Visibility model. In a L-visibility drawing Γ of a graph G , every vertex is represented by a horizontal and a vertical segment sharing an end-point, i.e., by an L-shape in the set $\{\text{L}, \text{J}, \text{T}, \text{F}\}$. Each edge of G is drawn in Γ as either a horizontal or a vertical visibility segment joining the two L-shapes corresponding to its two end-vertices. Clearly, horizontal visibilities only cross vertical visibilities at right angles. Also, no two L-shapes intersect. If G is an IC-plane graph, then each visibility is crossed at most once and no two crossed visibilities are incident to the same L-shape. In Fig. 1(b), an L-visibility representation Γ of an IC-plane graph G is shown. Finally, we adopt a

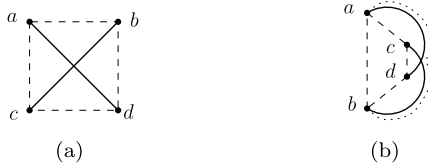


Fig. 2. (a) An X-configuration and (b) a B-configuration.

weak model, where a visibility may not imply the existence of the corresponding edge in the graph. For example, in Fig. 1(b) the L-shapes of vertices 8 and 10 can be joined by a horizontal visibility, although the edge $(8, 10)$ does not exist in G .

3. Proof of Theorem 1

The proof of Theorem 1 is constructive and is based on a drawing algorithm that takes as input an IC-plane graph G and returns a L-visibility drawing Γ of G . By Property 2, we assume that G is such that each crossing induces an empty kite (see Section 2). In fact, the output of our drawing algorithm maintains the IC-planar embedding obtained by applying Property 2. We begin by removing from G all pairs of crossing edges and orient the resulting graph P to an *st*-graph. The computed orientation is such that, when reinserting a pair of crossing edges in the corresponding planar face of P , one of them is always incident to the origin and the destination of the face. In other words, each face of P that corresponds to an empty kite of G , is oriented so that its left and right paths contain exactly one vertex each. To prove that this is always the case, we first need to introduce additional notation. Let f be a face of a plane graph G . Let v_1, \dots, v_h be the $h \geq 3$ vertices that belong to the boundary of f , and let $N(v)$ be the set of neighbors of a vertex v of G . The *contraction* of f is the operation defined as follows. Add to G a vertex v_f and connect v_f to the vertices in $N(v_1) \cup \dots \cup N(v_h)$. Then remove v_1, \dots, v_h from G . The resulting (multi)graph is still planar. Moreover, the contraction operation can be performed so to preserve the planar embedding of G . Namely, the circular order of the edges incident to v_f is the same circular order encountered walking along the boundary of f . See also Fig. 3(a) for an illustration. The original graph G can be obtained by applying the reverse operation, called the *expansion* of v_f . Namely, vertices v_1, \dots, v_h are reinserted along with their original edges and v_f is removed from the graph.

Lemma 1. *Let G be an n -vertex IC-plane graph such that the end-vertices of each pair of crossing edges induce an empty kite. Let C be the set of crossing edges in G , and consider the plane graph $P = G \setminus C$. Graph P can be oriented to an *st*-graph such that each pair of crossing edges of C has been removed from a face of P whose left and right paths contain exactly one vertex each. This operation can be done in $O(n)$ time.*

Proof. Each pair of crossing edges of G induces an empty kite K , and thus corresponds to a single face f_K in P having exactly four vertices on its boundary. As a first step, we contract each face f_K of P (corresponding to a kite

K in G). See also Fig. 3(a) for an illustration. Notice that, since G is an IC-plane graph, no two faces of P share a vertex, and thus each vertex v_K of P_C corresponds to exactly one face of P . Hence, we can contract the faces following an arbitrary order. The resulting graph P_C is a plane (multi)graph. Indeed, observe that if a face f_K of P shares an edge with a triangular face f , then P_C will contain two parallel edges between v_K and the vertex v of f not in f_K (see also Fig. 3(a)).

As a second step, we orient P_C to an *st*-graph (observe that parallel edges must receive the same orientation). In the third step, we expand one by one all the vertices v_K corresponding to a contracted face f_K . After expanding a vertex v_K , we orient the four reinserted edges of the face f_K maintaining the following invariants: **I1**. The resulting graph has a single source and a single sink; **I2**. The left and right paths of f_K contain exactly one vertex each.

Invariants **I1** and **I2** imply that the graph after expanding v_K is still an *st*-graph, as it has a single source and a single sink by **I1** and is acyclic by **I2**.

Let $\{a, b, c, d\}$ be the four vertices belonging to the boundary of face f_K , encountered in this order clockwise around the boundary of the face. To maintain **I2** we need to orient the edges of f_K such that the origin and the destination of f_K are two non-adjacent vertices, i.e., either a and c or b and d . In order to maintain **I1**, recall that the incoming edges of v_K appear consecutive around it and so do the outgoing edges, unless v_K is the source or the sink of the graph. Thus, if v_K is neither the source nor the sink of the graph, then at most two of the reinserted vertices will have both incoming and outgoing edges incident to v_K , whereas at most three will have only incoming or only outgoing edges incident to v_K . We distinguish the following three cases.

Case 1. No vertex of f_K has both incoming and outgoing edges. See Fig. 3(b) for an illustration. Then we consider the following subcases. **Case 1a.** There are three vertices having only outgoing edges, say a , b and c , which implies that all the edges of d are incoming. In this case, we orient the edges of f_K so that b is the destination and d is the origin of the face, which ensures **I2**. All vertices have now both incoming and outgoing edges, and thus **I1** is also maintained. **Case 1b.** There are three vertices having only incoming edges, say a , b and c , which implies that all the edges of d are outgoing. In this case, the orientation that ensures **I1** and **I2** is the one where d is the destination and b is the origin of the face. **Case 1c.** Two (consecutive) vertices only have incoming edges, say c and d , and two (consecutive) vertices only have outgoing edges, a and b . Then to maintain **I1** and **I2** we let a to be the destination and c the origin of the face. This particular case is shown in Fig. 3(b).

Case 2. Only one vertex of f_K , say a , has both incoming and outgoing edges. Moreover, assume that the incoming edges of a are between the edge (a, d) and the outgoing edges of a , as in Fig. 3(c), since the case in which the incoming edges of a are between the outgoing edges of a and the edge (a, b) is symmetric. We have two subcases. **Case 2a.** The other three vertices only have incoming (resp., outgoing) edges. In this case, we orient the edges of f_K so

that c (resp., a) is the origin and a (resp., c) is the destination. This choice ensures both **I1** and **I2**. **Case 2b.** Two vertices only have incoming edges and the other vertex only have outgoing edges. Due to the bipolar orientation of v_K , the two vertices having only incoming edges are c and d . Then we choose d as origin of the face and b as destination, which ensures both **I1** and **I2**. This particular case is shown in Fig. 3(c). The same orientation works also if two vertices only have outgoing edges (b and c) and the other vertex only have incoming edges (d).

Case 3. Two vertices of f_K have both incoming and outgoing edges. We have two subcases. **Case 3a.** Suppose first that these two vertices, say a and b , are adjacent in P . Moreover, assume that the incoming edges of a are between the edge (a, d) and the outgoing edges of a , as in Fig. 3(d), since the other case is symmetric. This implies that the incoming edges of b are between the outgoing edges of b and the edge (b, c) . Moreover, c and d only have incoming edges. Then, if we let a to be the destination of the face and c the origin, **I1** and **I2** are maintained.

Case 3b. Suppose now that the two vertices, say a and c , are not adjacent in P . Moreover, assume that the incoming edges of a are between the edge (a, d) and the outgoing edges of a , as in Fig. 3(d), since the case in which the incoming edges of a are between the outgoing edges of a and the edge (a, b) is symmetric. This implies that the incoming edges of c are between the outgoing edges of c and the edge (c, d) . Moreover, b and d only have outgoing and incoming edges, respectively. Then, if we orient the edges of f_K so that b is the destination of the face and d the origin, then **I1** and **I2** are ensured.

Finally, suppose that v_K is the source (resp., sink) of the graph. Then all the vertices of f_K have either no incident oriented edges, or all outgoing (resp., incoming) edges. Also, at least one of them has at least one outgoing (resp., incoming) edge, say a . We orient the edges of f_K so that a is the destination (resp., origin) and c the origin (resp., destination) of f_K .

The described algorithm works in $O(n)$ time. Namely, the graph P_C can be constructed and oriented to an *st*-graph in $O(n)$ time (see, e.g., [10]). Furthermore, orienting the edges of an expanded face requires first to analyze the orientation of the edges incident to each vertex of the face, and then to orient the edges of the face. Since the faces are vertex-disjoint, this costs at most $\sum_{v \in P} \deg(v) = 2m_P$, where m_P is the number of edges of P , which is $O(n)$. \square

Let f_K be a face of the plane *st*-graph P which corresponds to an (empty) kite K in G . By Lemma 1, f_K is such that its left and right path have both length two. This implies that one of the two crossing edges of K is incident to the origin and to the destination of f_K , whereas the other one is incident to the two vertices belonging one to the left and one to the right path of f_K . Let P^+ be the biconnected plane graph obtained from P by reinserting, for each pair of crossing edges of G , the edge incident to the origin and to the destination of the corresponding face in P . Furthermore, let Γ' be a strong visibility drawing of P^+ in $O(n^2)$ area, which can be computed in $O(n)$ time (see, e.g., [15]).

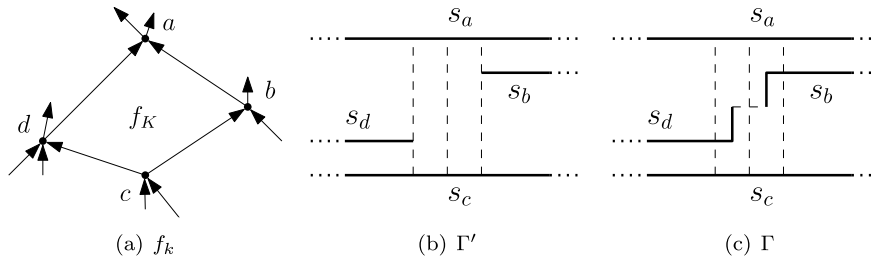


Fig. 4. Illustration for the proof of Lemma 2.

Lemma 2. Γ' can be transformed into a L-visibility drawing Γ of G that requires $O(n^2)$ area. This operation can be done in $O(n)$ time.

Proof. Consider a face f_K of P corresponding to an empty kite K in G . Let $\{a, b, c, d\}$ be the four vertices of f_K , encountered in this order clockwise around the boundary of the face. See also Fig. 4(a). Without loss of generality, let a be the destination of the face. Then Lemma 1 implies that c is the origin of f_K , and thus the edge reinserted in P^+ for this face is (a, c) . In other words, f_K is split in two faces in P^+ , and these two faces share the edge (a, c) . Consider the subdrawing of Γ' induced by the four vertices $\{a, b, c, d\}$. An illustration is also shown in Fig. 4(b). Let s_v be the bar representing a vertex v in Γ' . Either s_d and s_b are drawn at the same y -coordinate, or, one is above the other. Also, the two bars do not overlap as there is no edge between b and d in P^+ (and Γ' is a strong visibility drawing). Between the two bars there is actually (at least) one unit gap needed to draw the visibility from s_c to s_a . Moreover, s_c is below both s_d and s_b , while s_a is above both of them. In any case, we first extend s_d (resp., s_b) by 0.25 units to the right (resp., left). Next, if s_b and s_d have the same y -coordinate, then it suffice to draw two vertical bars s'_d and s'_b , such that the bottommost end-point of s'_d (resp., s'_b) coincides with the rightmost end-point of s_d (resp., leftmost end-point of s_b), and such that the other end-point is 0.5 units above it. If one is above the other, say s_b is above s_d , and the difference in terms of y -coordinates between the two bars is $k \geq 1$ units, then we draw two vertical bars s'_d and s'_b , such that the end-point of s'_d (resp., s'_b) coincides with the rightmost end-point of s_d (resp., leftmost end-point of s_b), and such that the other end-point is $k/2$ units above it (resp., below it). In both cases, the two resulting L-shapes see each other through a horizontal visibility segment. See also Fig. 4(c). Since every vertex is adjacent to at most one crossed edge we have that the final drawing Γ is a L-visibility drawing of G . Since Γ' contains $O(n)$ segments which have to be transformed to L-shapes, Γ is computed in $O(n)$ time. Finally, in order to restore integer coordinates, we scale by a factor 4 the grid of Γ , which thus takes $O(n^2)$ area. \square

Lemmas 1 and 2 imply Theorem 1. To prove Corollary 1, consider a visibility drawing Γ of an n -vertex IC-plane graph G . By Theorem 1, Γ can be computed in $O(n)$ time and fits on a grid of $O(n^2)$ size. Let ℓ be an L-shape of Γ . The representative point r of ℓ is defined as follows. If both the horizontal and the vertical segments of ℓ have

non-zero length, then r is the point where they touch. Otherwise, r is the midpoint of the segment of ℓ having non-zero length. Replace each vertical visibility segment with a polyline as follows. Let s be a visibility segment connecting the two L-shapes ℓ_1 and ℓ_2 . Let r_1 and r_2 be the representative points of ℓ_1 and ℓ_2 , respectively. Also, let p_1 and p_2 be the points that s shares with ℓ_1 and ℓ_2 , respectively. Suppose that r_1 is below r_2 (and thus p_1 is below p_2). Replace s with the polyline starting at r_1 , bending 0.25 grid units above p_1 , bending again 0.25 units below p_2 , and ending at r_2 . With a symmetric operation we can also replace each horizontal visibility segment. Finally, replace each L-shape with its representative point. The resulting drawing is an IC-plane drawing of G where edges are polylines with (at most) two bends that cross at right-angles. Finally, scaling by a factor 4 the grid of the drawing we restore integer coordinates. Fig. 1(c) shows a RAC drawing computed from the L-visibility drawing in Fig. 1(b).

4. Conclusions and open problems

We have proved that every IC-plane graph G has a L-visibility drawing which can be computed in linear time. As a corollary, our result implies that G has a RAC drawing in quadratic area and at most two bends per edge which can also be computed in linear time. We conclude the paper with two open problems: (i) Does every 1-planar graph admit a visibility drawing where the shape associated with each vertex is either an L-shape, a T-shape, or a +-shape? (ii) Does every IC-plane graph admit a RAC drawing with at most one bend per edge in polynomial area?

References

- [1] E.N. Argyriou, M.A. Bekos, A. Symvonis, The straight-line RAC drawing problem is NP-hard, *J. Graph Algorithms Appl.* 16 (2) (2012) 569–597.
- [2] F.J. Brandenburg, 1-visibility representations of 1-planar graphs, *J. Graph Algorithms Appl.* 18 (3) (2014) 421–438.
- [3] F.J. Brandenburg, W. Didimo, W.S. Evans, P. Kindermann, G. Liotta, F. Montecchiani, Recognizing and drawing IC-planar graphs, in: *GD 2015*, in: LNCS, Springer, 2015, in press.
- [4] A.M. Dean, W.S. Evans, E. Gethner, J.D. Laison, M.A. Safari, W.T. Trotter, Bar k -visibility graphs, *J. Graph Algorithms Appl.* 11 (1) (2007) 45–59.
- [5] G. Di Battista, P. Eades, R. Tamassia, I.G. Tollis, *Graph Drawing*, Prentice Hall, 1999.
- [6] W. Didimo, P. Eades, G. Liotta, Drawing graphs with right angle crossings, *Theor. Comput. Sci.* 412 (39) (2011) 5156–5166.
- [7] W. Didimo, G. Liotta, The crossing angle resolution in graph drawing, in: J. Pach (Ed.), *Thirty Essays on Geometric Graph Theory*, Springer, 2012.

- [8] W.S. Evans, M. Kaufmann, W. Lenhart, T. Mchedlidze, S.K. Wismath, Bar 1-visibility graphs vs. other nearly planar graphs, *J. Graph Algorithms Appl.* 18 (5) (2014) 721–739.
- [9] W.S. Evans, G. Liotta, F. Montecchiani, Simultaneous visibility representations of plane *st*-graphs using L-shapes, in: *WG 2015*, in: LNCS, Springer, 2015, in press.
- [10] S. Even, R.E. Tarjan, Computing an *st*-numbering, *Theor. Comput. Sci.* 2 (3) (1976) 339–344.
- [11] W. Huang, S.-H. Hong, P. Eades, Effects of crossing angles, in: *PacificVis 2008*, IEEE, 2008, pp. 41–46.
- [12] D. Král, L. Stacho, Coloring plane graphs with independent crossings, *J. Graph Theory* 64 (3) (2010) 184–205.
- [13] G. Liotta, Graph drawing beyond planarity: some results and open problems, in: *ICTCS'14*, 2014, pp. 3–8.
- [14] P. Rosenstiehl, R.E. Tarjan, Rectilinear planar layouts and bipolar orientations of planar graphs, *Discrete Comput. Geom.* 1 (1986) 343–353.
- [15] R. Tamassia, I.G. Tollis, A unified approach to visibility representations of planar graphs, *Discrete Comput. Geom.* 1 (1) (1986) 321–341.
- [16] C. Thomassen, Rectilinear drawings of graphs, *J. Graph Theory* 12 (3) (1988) 335–341.
- [17] X. Zhang, G. Liu, The structure of plane graphs with independent crossings and its applications to coloring problems, *Cent. Eur. J. Math.* 11 (2) (2013) 308–321.