

A machine-independent characterization of timed languages

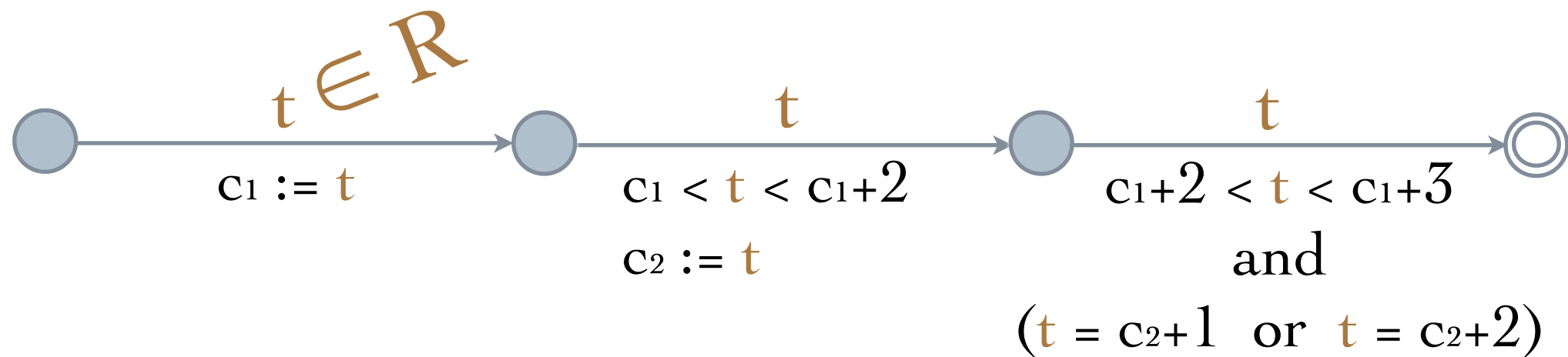
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joint work with Mikołaj Bojańczyk

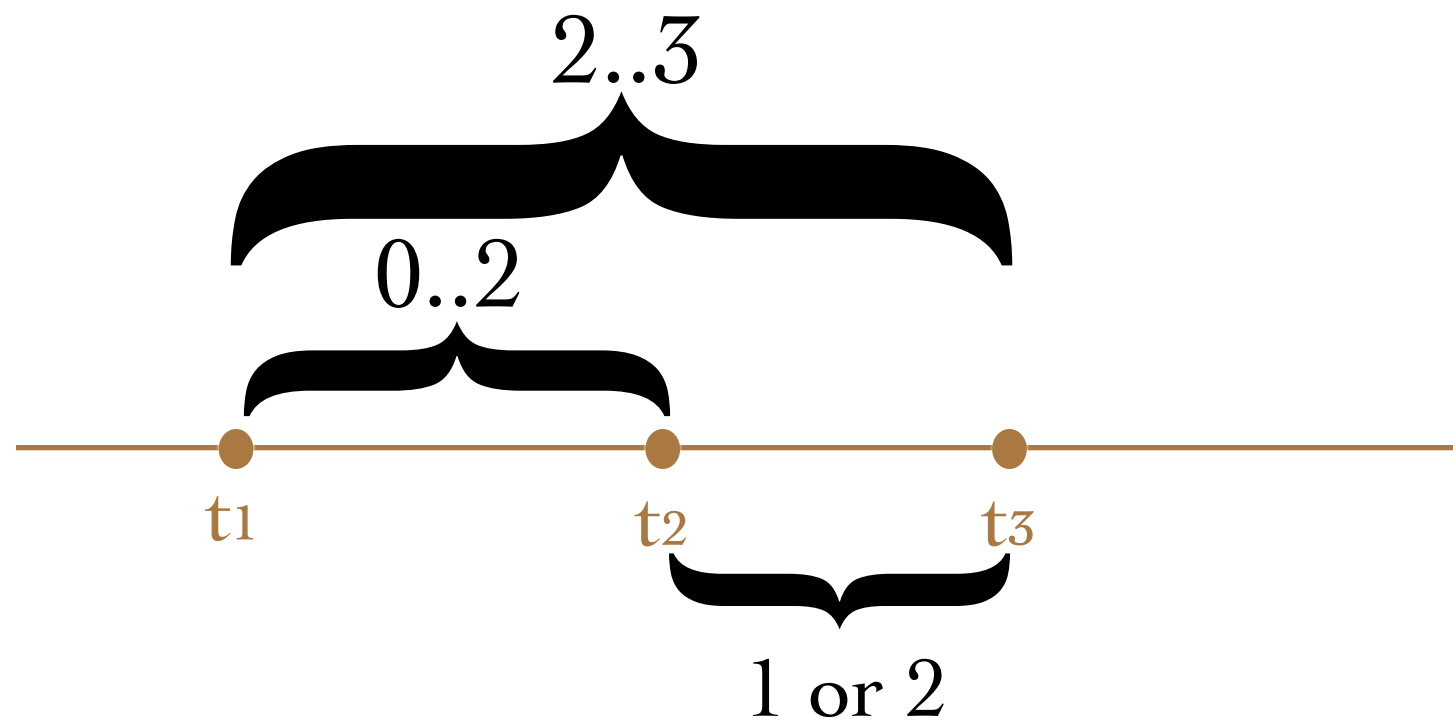
HIGHLIGHTS 2013

deterministic timed automata

with uninitialized clocks



the automaton accepts words $t_1 \ t_2 \ t_3 \in \mathbb{R}^3$ such that



Myhill-Nerode theorem

let L be a language over a finite alphabet A

L is recognized by a DFA

iff

\approx_L has finitely many equivalence classes


$$w \approx_L u \quad \text{iff} \quad \forall v (wv \in L \text{ iff } uv \in L)$$

The same for deterministic timed automata?

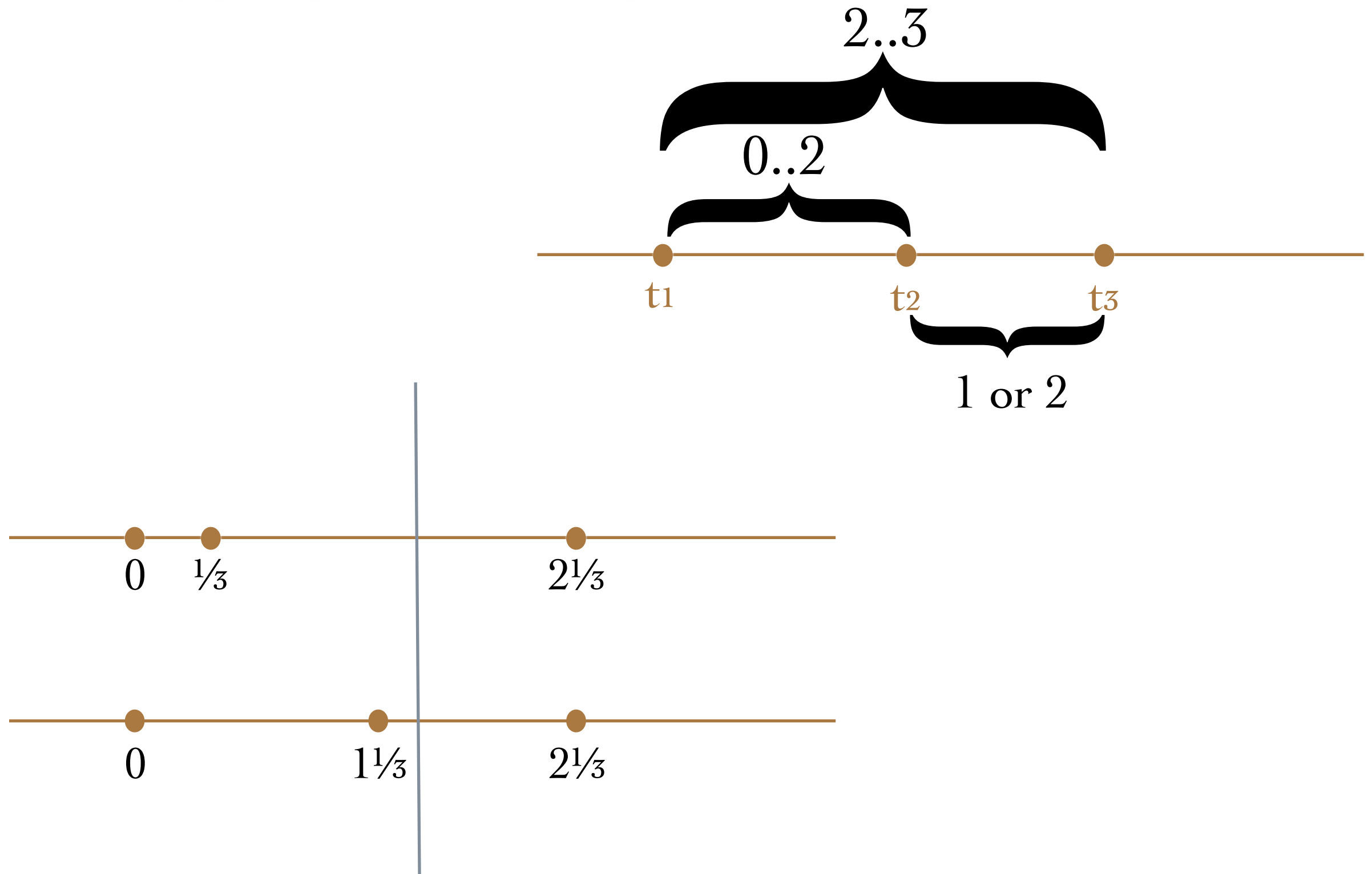
Problems:

- infinitely many equivalence classes*
- no canonical minimal timed automaton*

deterministic timed automata

with uninitialized clocks

do not minimize



Solution: move to sets with atoms

deterministic orbit-finite automata
in sets with atoms $(\mathbf{R}, <, +1)$

deterministic timed automata
with uninitialized clocks

minimal automata for languages
of deterministic timed automata
with uninitialized clocks

finite automata
up to automorphisms of
 $(\mathbf{R}, <, +1)$

resemble
timed automata
with updates

closed under
minimization

Myhill-Nerode theorem for timed languages

let L be a language over $A \times \mathbf{R}$

such that

- L contains only increasing words
- L is invariant under $\text{Aut}(\mathbf{R}, <, +1)$

L is recognized by a deterministic timed automaton
with uninitialized clocks

iff

- \approx_L has orbit-finite set of equivalence classes
- L is forgetful

deterministic orbit-finite automata
in sets with atoms $(\mathbf{R}, <, +1)$

deterministic timed automata
with uninitialized clocks

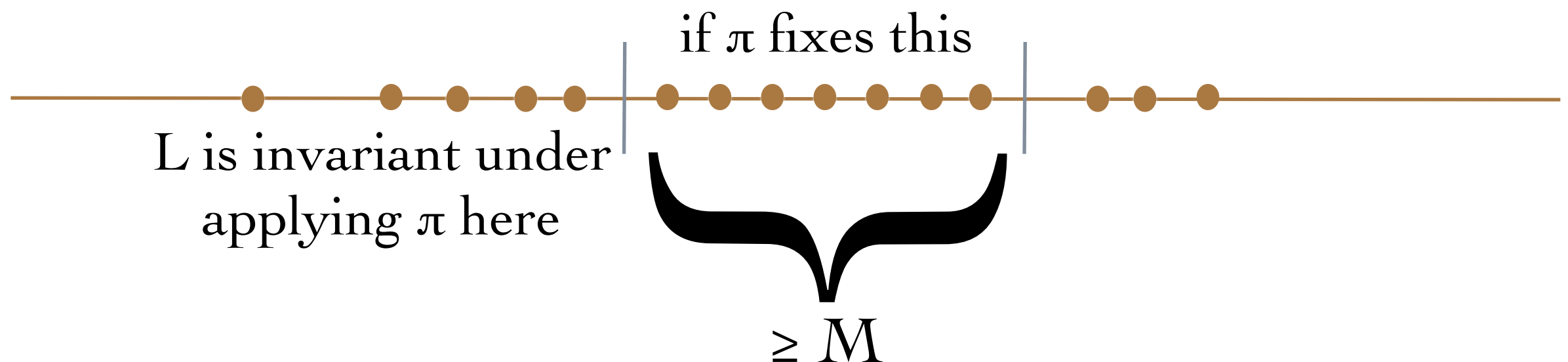
L is forgetful iff

there is $M \in \mathbb{R}$ such that

for every timed word

and $\pi \in \text{Aut}(\mathbb{R}, <, +1)$

for every factorization



summary

- Myhill-Nerode theorem for timed languages
 - superclass of deterministic timed automata
closed under minimization
- } both result due to
sets with atoms

Thank you!