Rewriting Higher-Order Stack Trees

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Context

 Model Checking over some classes of infinite graphs: which logic theories (FO,FO[→*],MSO,...) are decidable?



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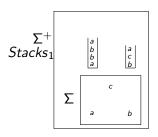
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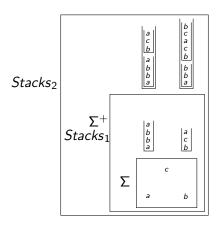
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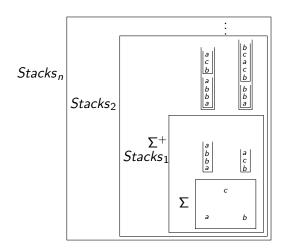
| Configuration graphs of HOPDA [Caucal02] & [Carayol, Wohrle03] | ? | Tree automatic of order n [Colcombet,Loding07] |
|---|--|---|
| Configuration graphs of PDA [Muller,Shupp85] | Ground tree rewriting graphs [Dauchet, Tison90] | Tree automatic [Khoussainov,Nerode94] |
| MSO | $FO[o^*]$ | FO |

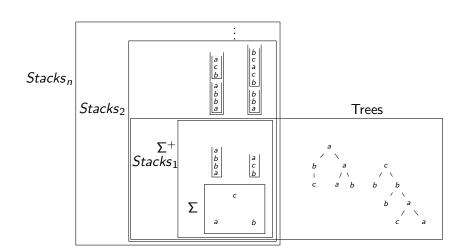


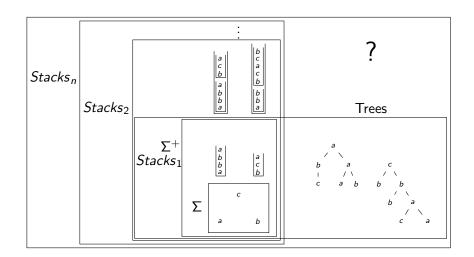






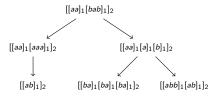






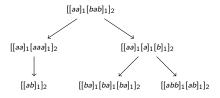
Higher-Order Stack Trees

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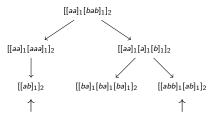


A unary *n*-stack tree is a *n*-stack.

A 1-stack tree is a tree labelled by Σ .

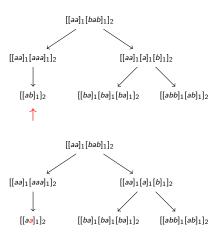
Basic Operations: Level 0

Rewriting operations over Σ : rew_{b,a}



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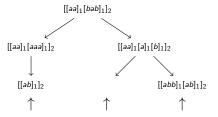


Basic Operations: Level i < n

Copy operations: $copy_i, \overline{copy}_i$

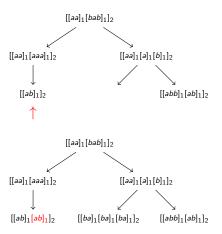
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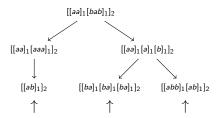


Basic Operations: Level n

Tree copy operations: $\operatorname{copy}_n^i, \overline{\operatorname{copy}}_n^i$

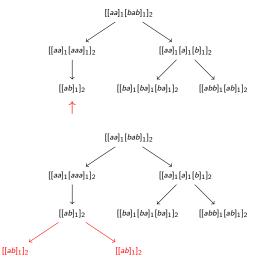
Basic Operations: Level n

Tree copy operations: $copy_3^2$



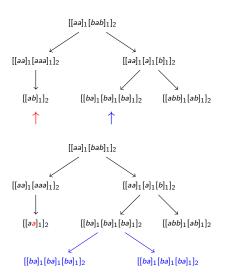
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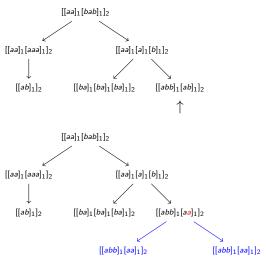
Composition

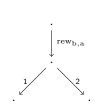
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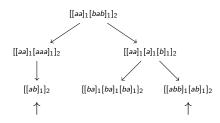


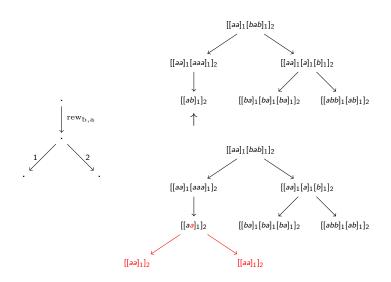
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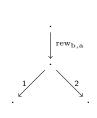
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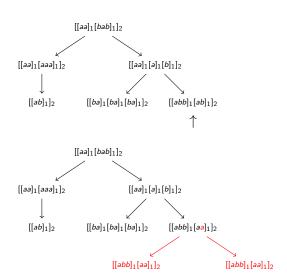


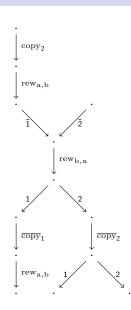




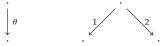






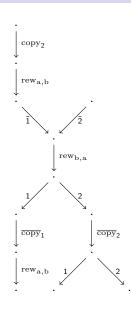


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Only connected operations



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- Only connected operations
- ullet Level 1 o ground tree rewriting rules
- \bullet Unary trees \to finite composition of higher-order pushdown operations

Main Result

Given a set of compound operations R, its rewriting graph \mathcal{G}_R is:

- $V_{G_P} = ST_n$
- $E_{G_R} = \{(t, r, t') \mid r \in R \land r(t, t')\}$

Theorem

Given a finite set of compound operations R, its rewriting graph has a decidable $FO[\stackrel{*}{\to}]$ theory.

Proof ingredients:

- Notion of recognisability over compound operations
- Finite set interpretation of every stack-tree rewriting graph into a graph with a decidable MSO-theory (the level *n* treegraph)



Perspectives

- Languages recognised by rewriting graphs of stack trees. Example $\{u \sqcup u \mid u \in \Sigma^*\}$
- Strictness of the graph hierarchy
- Extension of the model to n-trees labelled by (n-1)-trees