Unique Normal Forms in Infinitary Weakly Orthogonal Term Rewriting

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> Universiteit Utrecht Vrije Universiteit Amsterdam

> > TeReSe meeting RWTH Aachen 28 May 2010

Overview

- ► Introduction: weakly orthogonal, UN[∞]
- ▶ Counterexample to UN[∞] for weakly orthogonal TRSs
- ► Counterexample to UN^{∞} for $\lambda^{\infty}\beta\eta$
- Restoring infinitary confluence
- Diamond and triangle properties for developments

Overview

1. Introduction

- 2. Counterexample to $ext{UN}^{\infty}$ for weakly orthogonal iTRSs
- 3. Counterexample to UN $^{\infty}$ in $\lambda^{\infty}\beta\eta$
- 4. Restoring infinitary confluence
- Diamond and triangle properties for developments
- 6. Conclusion

Weakly orthogonal (first-/higher-order) systems:

- left-linear
- all critical pairs are trivial.

Examples.

➤ Successor/Predecessor TRS:

$$P(S(x)) \to x$$
 $S(P(x)) \to x$

with critical pairs:

$$S(x) \leftarrow \underline{S}(\underline{P}(S(x))) \rightarrow S(x)$$
 $P(x) \leftarrow \underline{P}(\underline{S}(P(x))) \rightarrow P(x)$

► Parallel-Or TRS ('almost orthogonal'):

$$por(true, x) \rightarrow true$$

 $por(x, true) \rightarrow true$
 $por(false, false) \rightarrow false$

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- ▶ UN $^{\infty}$: $t_1 \leftarrow t \rightarrow t_1 \land t_2$ normal forms $\Rightarrow t_1 = t_2$.
- \triangleright SN $^{\infty}$: all infinite rewrite sequences are progressive (str. conv.)

- $ightharpoonup SN^{\infty} \Longrightarrow CR^{\infty}$, and $CR^{\infty} \Longrightarrow UN^{\infty}$.
- $ightharpoonup CR^{\infty}$ fails (Kennaway).
 - ▶ But for non-collapsing TRSs: CR[∞] holds
- ► UN[∞] holds (Kennaway/Klop).

- $\qquad \qquad \mathsf{CR}^{\infty} \colon \qquad t_1 \twoheadleftarrow t \twoheadrightarrow t_2 \implies \exists \mathsf{s}. \ t_1 \twoheadrightarrow \mathsf{s} \twoheadleftarrow t_2 \,.$
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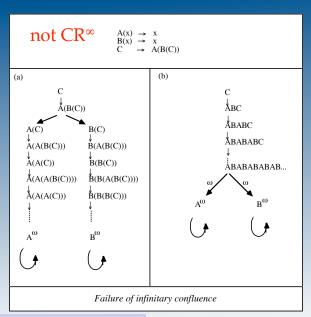
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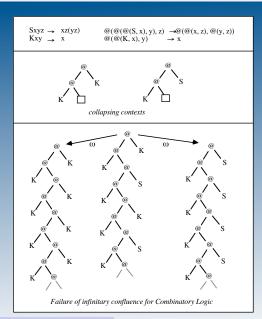
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In the Successor/Predecessor TRS:

$$P(S(x)) \rightarrow x$$
 $S(P(x)) \rightarrow x$

with the normal forms $S^{\omega} = SSS...$ and $P^{\omega} = PPP...$ we consider:

$$\psi = P^1 S^2 P^3 S^4 P^5 S^6 \dots = P SS PPP SSSS PPPPP SSSSSS \dots$$

We find:

$$\psi = \mathbf{PSS} PPP SSSS PPPPP SSSSSS ...$$

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And similar:

$$\psi \rightarrow PPPSSSSSS...$$

 $\twoheadrightarrow SSSSSS... = S^{\omega}$

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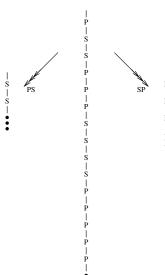
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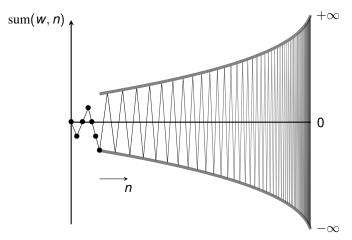
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We find:

$$\psi \twoheadrightarrow \mathsf{PPPPP} \dots = \mathsf{P}^{\omega}$$





Graph for the oscillating PS-word $\psi = P^1 S^2 P^3 \dots$

Some facts about infinite PS-words

The S-norm $\|\mathbf{w}\|_{S}$ and P-norm $\|\mathbf{w}\|_{P}$ of a PS-word \mathbf{w} :

$$\|\mathbf{w}\|_{S} = \sup_{n \in \mathbb{N}} \operatorname{sum}(\mathbf{w}, n) \qquad \|\mathbf{w}\|_{P} = \sup_{n \in \mathbb{N}} (-\operatorname{sum}(\mathbf{w}, n))$$

- **(a)** $w woheadrightarrow S^{\omega}$ if and only if $||w||_{S} = \infty$; $w woheadrightarrow P^{\omega}$ if and only if $||w||_{P} = \infty$.
- Every PS-word that reduces to both S^ω and P^ω reduces to every infinite PS-word.
- A PS-word w is root-active if and only if w is the concatenation of infinitely many finite 'zero-sum-words' w₁, w₂, w₃,
- For an infinite PS-word w we have $SN^{\infty}(w)$ if and only if each value sum(w, n) for $n = 0, 1 \dots$ occurs only finitely often.

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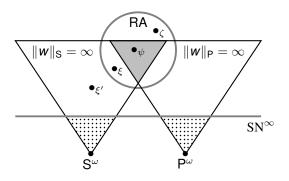
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- Every PS-word that reduces to both S^ω and P^ω reduces to every infinite PS-word.
- **3** A PS-word w is root-active if and only if w is the concatenation of infinitely many finite 'zero-sum-words' w_1, w_2, w_3, \ldots
- **③** For an infinite PS-word \underline{w} we have $SN^{\infty}(\underline{w})$ if and only if each value $sum(\underline{w}, n)$ for $n = 0, 1 \dots$ occurs only finitely often.

Venn diagram



$$\psi = \mathsf{P}^1 \, \mathsf{S}^2 \, \mathsf{P}^3 \, \dots \qquad \qquad \xi = \mathsf{S} \, \mathsf{P} \, \mathsf{S}^2 \, \mathsf{P}^2 \, \mathsf{S}^3 \, \mathsf{P}^3 \dots$$
$$\zeta = (\mathsf{P} \, \mathsf{S})^\omega \qquad \qquad \xi' = \mathsf{S} \, \xi$$

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$$\lambda^{\infty}\beta\eta$$

Terms of $\lambda^{\infty}\beta\eta$: the (potentially) infinite λ -terms in $Ter^{\infty}(\lambda)$

The rewrite rules of $\lambda^{\infty}\beta\eta$ are:

$$(\lambda x.M)N \stackrel{\beta}{\to} M[x:=N]$$
$$\lambda x.Mx \stackrel{\eta}{\to} M \qquad (x \text{ is not free in } M)$$

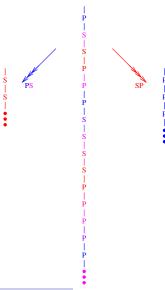
 $\lambda^{\infty}\beta\eta$ is weakly orthogonal, since the critical pairs are trivial:

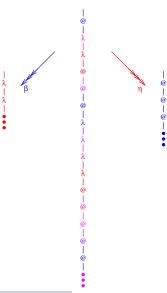
$$Mx \stackrel{\beta}{\leftarrow} (\lambda x. Mx) x \stackrel{\eta}{\rightarrow} Mx \qquad (x \notin fv(M))$$

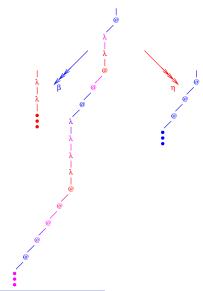
$$\lambda x.M[y:=x] \stackrel{\beta}{\leftarrow} \lambda x.(\lambda y.M)x \stackrel{\eta}{\rightarrow} \lambda y.M \qquad (x \notin fv(M))$$

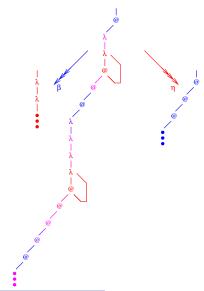


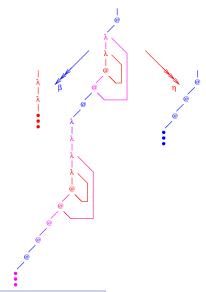


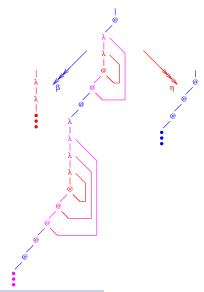


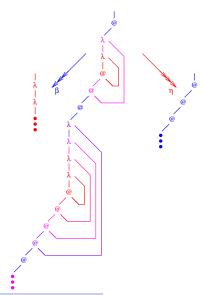


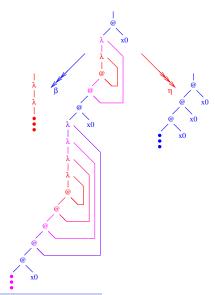


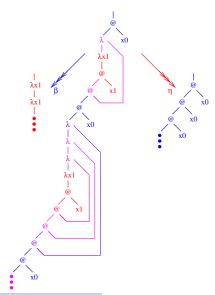


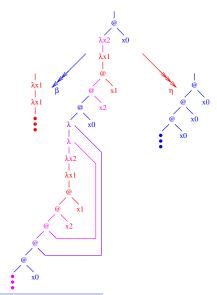


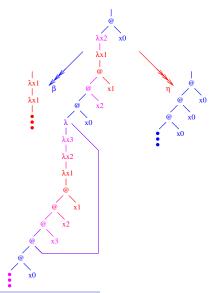




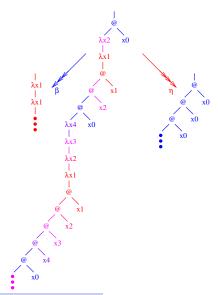








Counterexample: UN $^{\infty}$ fails in $\lambda^{\infty}\beta\eta$



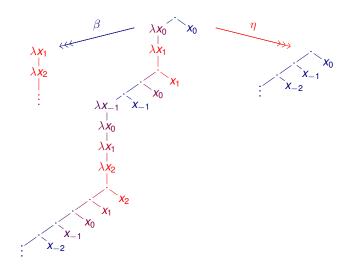
Translating the S-P-example to $\lambda^{\infty}\beta\eta$

- $(_) : \{ P, S \}^{\omega} \to Ter^{\infty}(\lambda)$ defined by:
 - $(w) = (w)_0$;
 - ▶ for all $w \in \{P, S\}^{\omega}$, and $i \in \mathbb{Z}$:

$$(|\mathsf{P}w|)_i = (|w|)_{i-1} \, x_i$$

 $(Sw)_i = \lambda x_{i+1} \cdot (w)_{i+1}$

Counterexample: UN $^{\infty}$ fails in $\lambda^{\infty}\beta\eta$



We saw for $\lambda^{\infty}\beta\eta$:

- $ightharpoonup UN^{\infty}$ fails
- ▶ Consequently: CR[∞] fails

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- ▶ But: UN[∞] holds!

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Contrast with $\lambda^{\infty}\beta$

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- ▶ Consequently: CR[∞] fails

However for $\lambda^{\infty}\beta$ it holds:

- $ightharpoonup CR^{\infty}$ fails
- ► But: UN[∞] holds!

Due to this, $\lambda^{\infty}\beta$ is important for the model theory of λ -calculus: for several models equality is captured by $\lambda^{\infty}\beta$ -convertibility:

- Böhm Trees
- Lévy–Longo Trees
- Berarducci Trees

Overview

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- 3. Counterexample to UN $^{\infty}$ in $\lambda^{\infty}\beta\eta$
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Lemma (Refined Compression Lemma)

Let R be a left-linear iTRS.

Let $\kappa : s \to_{R}^{\alpha} t$ be a rewrite sequence, let d the min. depth of a step,

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For every divergent rewrite sequence $\kappa: s \to_R^{\alpha}$ (length α) there exists a divergent rewrite sequence $\kappa': s \to_{\overline{\rho}}^{\leq \omega}$ (length <

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Orthogonalization (of parallel steps)

Proposition

For parallel steps $\phi: s \longrightarrow t_1$ and $\psi: s \longrightarrow t_2$ in a w-o TRS there exists orthogonal steps ϕ' and ψ' such that $\phi': s \longrightarrow t_1$ and $\psi': s \longrightarrow t_2$ (the pair $\langle \phi', \psi' \rangle$ is an orthogonalization of ϕ and ψ).

Proof

In case of overlaps, we replace the outer redex with the inner one.



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Definition

Let the weakly orthogonal projections ϕ/ψ and ψ/ϕ be the orthogonal projections ϕ'/ψ' and ψ'/ϕ' of the orthogonalization $\langle \phi', \psi' \rangle$ of $\langle \phi, \psi \rangle$.

Lemma

Let d_{ϕ} and d_{ψ} be the minimal depth of a step in ϕ and ψ .

- Then the minimal depth of the w-o projections ϕ/ψ and ψ/ϕ is:
 - ▶ in general: $\geq \min(d_{\phi}, d_{\psi})$.
 - if R is non-collapsing: $\geq \min(d_{\phi}, d_{\psi} + 1)$ and $\geq \min(d_{\psi}, d_{\phi} + 1)$, respectively.

Proof.

The orthogonalization does not decrease the height as we always replace the outer by the inner redex. Orthogonal projection cannot lift redexes above the the depth of variables in the right-hand sides.

Infinitary Parallel Moves Lemma PML[∞]

Lemma

Let R be a weakly orthogonal TRS. In general it holds:

$$s \xrightarrow{\geq d_{\kappa}} t_{1} \Rightarrow t_{1}$$

$$\geq d_{\xi} \downarrow \qquad \qquad \geq \min(d_{\kappa}, d_{\xi}) \downarrow \qquad \qquad \downarrow t_{2} \xrightarrow{\sim} \min(d_{\kappa}, d_{\xi}) \qquad \downarrow u$$

If R is non-collapsing, then also:

$$\begin{array}{c}
s \xrightarrow{\qquad \geq d_{\kappa}} & \longrightarrow t_{1} \\
\geq d_{\xi} & \downarrow & \geq \min(d_{\xi}, d_{\kappa} + 1) \downarrow \\
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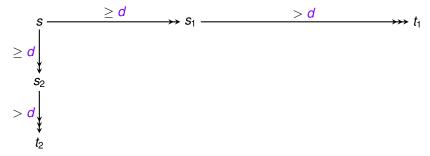
$$\geq d_{\xi} \stackrel{\downarrow}{\downarrow} \underset{t_{2}}{\longrightarrow} \min(d_{\xi}, d_{\kappa} + 1) \stackrel{\downarrow}{\downarrow} \underset{u}{\downarrow}$$

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Weakly orthogonal TRSs without collapsing rules are inf. confluent.

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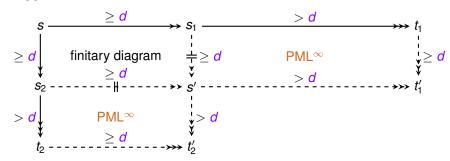
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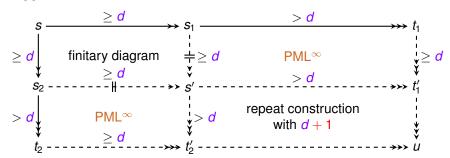
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Diamond and triangle properties for developments

A binary relation \rightarrow on A is said to have:

- ▶ the diamond property if: $\leftarrow \cdot \rightarrow \subseteq \rightarrow \cdot \leftarrow$;
- the triangle property if:

$$\forall a \in A. \ \exists a' \in A. \ a \rightarrow a' \ \land \ (\forall b \in A. \ a \rightarrow b \ \Rightarrow \ b \rightarrow a').$$

Theorem

For every weakly orthogonal TRS without collapsing rules, (infinite) multi-steps have:

- the diamond property;
- 2 the triangle property.

Our proof proceeds by:

- refining an earlier cluster analysis (I-clusters and Y-clusters) from the finite case:
- ▶ a top-down orthogonalization algorithm.

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Overview

- Counterexample to UN[∞]/CR[∞] for weakly orthogonal TRSs
- ▶ By translation: counterexample to UN^{∞}/CR^{∞} for $\lambda^{\infty}\beta\eta$
- ► Restoring CR[∞] for non-collapsing w-o TRSs
- Diamond and triangle properties for developments in non-collapsing w-o TRSs

Summary

	finitary					infinitary			
		PML	CR	UN	NF	PML∞	CR∞	UN∞	NF∞
100000	OTRS	yes	yes	yes	yes	yes	no	yes	yes
	WOTRS	yes	yes	yes	yes	yes	no	no	no
	nc-WOTRS	yes	yes	yes	yes	yes	yes	yes	yes
	1c-WOTRS	yes	yes	yes	yes	yes	no	?	?
6	λeta	yes	yes	yes	yes	no	no	yes	yes
	fe-OCRS	yes	yes	yes	yes	no	no	yes	yes
	$\lambda eta \eta$	yes	yes	yes	yes	no	no	no	no
	WOCRS	yes	yes	yes	yes	no	no	no	no