

Syntactic NFA Minimization via SAT solving

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GOAL

minimal nfas

Given an nfa M with states Q , nfa N with states S is a *union nfa* of M if

- ▶ each $s \in S$ accepts $\bigcup\{L(N, q) : q \in X\}$ with $X \subseteq Q$
- ▶ $L(N) = L(M)$

nfa $N = (Q, \delta, I, F)$ accepting L is **atomic** if the following equivalent statements hold

- ▶ $\text{rsc}(N^R)$ is a minimal dfa
- ▶ each $q \in Q$ accepts a language from $\text{BLD}(L)$, the closure of $\text{LD}(L)$ under all set-theoretic boolean operations
- ▶ each $q \in Q$ accepts a union of congruence classes of the *Nerode left congruence* \sim_L :

$$\begin{aligned} u \sim_L v &\iff \forall x \in \Sigma^*. u \in x^{-1}L = v \in x^{-1}L \\ &\iff (u^r)^{-1}L^r = (v^r)^{-1}L^r \end{aligned}$$

- ▶ N is a union nfa of the átomaton $\mathbf{dfa}(L^r)^R$

nfa $N = (Q, \delta, I, F)$ accepting L is **subatomic** if the following equivalent statements hold

- ▶ the transition monoid of $\text{rsc}(N^R)$ is isomorphic to $\text{syn}(L^r)$
- ▶ each $q \in Q$ accepts a language from $\text{BLRD}(L)$, the closure of $\text{LD}(L)$ under all boolean operations and right derivatives
- ▶ each $q \in Q$ accepts a union of congruence classes of the *syntactic congruence* \equiv_L :

$$u \equiv_L v \iff \forall x, y \in \Sigma^*. u \in x^{-1}Ly^{-1} = v \in x^{-1}Ly^{-1}$$

- ▶ N is a union nfa of the syntactic nfa **syn**(L)

given $\mathbf{dfa}(L)$ and an rpd nfa $M = (Q, \delta, I, \{q_f\})$ accepting L ,

does a union nfa N of M with k states exist?

We are searching for a union nfa N with states \mathbf{k} :

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► $H \subseteq \mathbf{k} \times Q$:

$$H[i] = X \iff L(N, i) = \bigcup \{L(M, q) : q \in X\}$$

► $G \subseteq \text{LD}(L) \times \mathbf{k}$:

$$G[u^{-1}L] = Y \iff u^{-1}L = \bigcup \{L(N, i) : i \in Y\}$$

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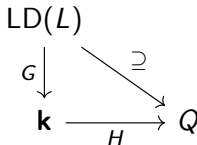
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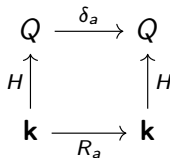
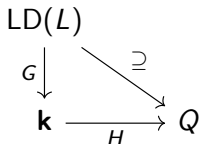
► $G \subseteq \text{LD}(L) \times \mathbf{k}$:

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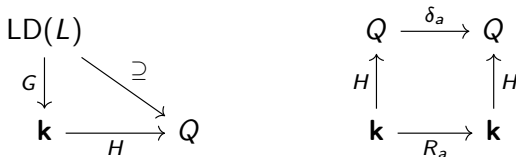
As a commutative diagram:



Require states to accept their assigned language using the transition relations $R_a \subseteq \mathbf{k} \times \mathbf{k}$ ($a \in \Sigma$):

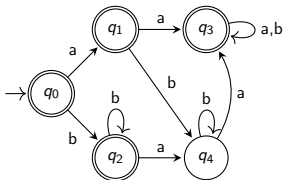


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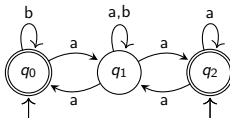


$$N = (\mathbf{k}, (R_a)_{a \in \Sigma}, G[L], H^{-1}[q_f])$$

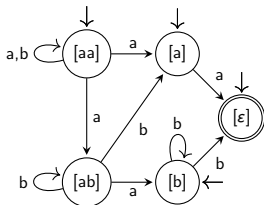
$$L = \{w \in \Sigma^* : |w|_b = 0 \vee |w|_a \neq 1\}$$



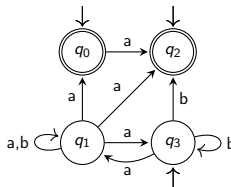
dfa(L)

























N









































dfa(L^r)^R ≅ syn(L)









































N_{atm} = N_{syn}

	$[\epsilon]$	$[a]$	$[b]$	$[aa]$	$[ab]$
$\epsilon^{-1}L$		 			
$a^{-1}L$					 
$b^{-1}L$			 		
$aa^{-1}L$					
$ab^{-1}L$					

	$[\epsilon]$	$[a]$	$[b]$	$[aa]$	$[ab]$
$\epsilon^{-1}L$					
$a^{-1}L$					
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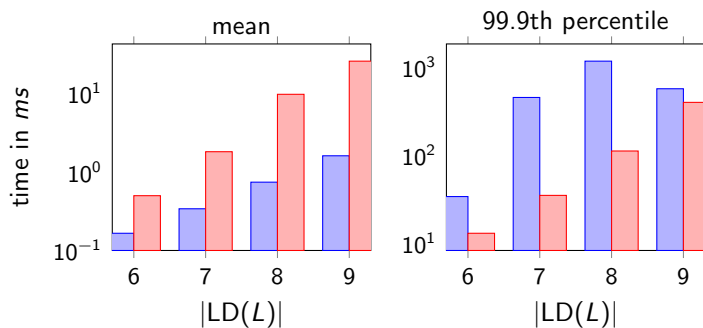
	$[\varepsilon]$	$[a]$	$[b]$	$[aa]$	$[ab]$
$\varepsilon^{-1}L$					
$a^{-1}L$					
$b^{-1}L$					
$aa^{-1}L$					
$ab^{-1}L$					

Two possible options for H :

r	$[\varepsilon], [a]$	$[\varepsilon], [a], [aa]$
g	$[\varepsilon], [b], [aa]$	$[\varepsilon], [b], [aa]$
b	$[a], [aa], [ab]$	$[a], [aa], [ab]$

Implementation

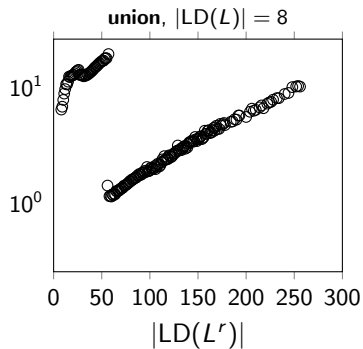
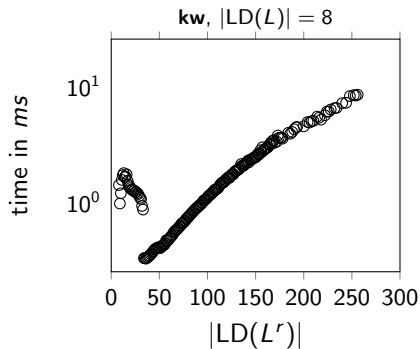
performance results



kw

union

performance results



Thank you for your attention!