# Session 3: MPI

### Outline

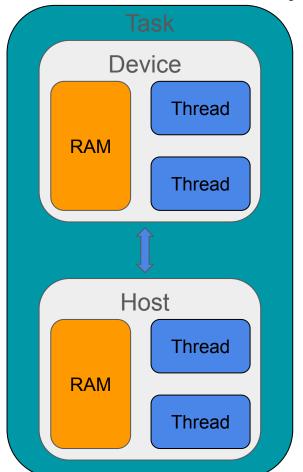
- General description
- Point-to-point communication
- Collective communication
- Closing remarks

The Message Passing Interface (MPI) is the distributed-memory

component to hPIC2's hybrid parallelism

 Unlike with shared-memory, MPI dominates distributed-memory parallelism

- Literally, multiple copies (processes or tasks) of program are launched and connected
- Each task maintains its own memory
  - Each task can also have its own shared-memory instance
- With MPI+Kokkos, we can now run on any supercomputer
  - With GPUs, use 1 GPU per MPI task
  - Without GPUs, generally use one CPU socket per MPI task



. . .

### MPI is a standard, NOT a library in and of itself

- Similar to how C++ is a standard, but with many compilers that actually do the work
- The standard describes function names, call signatures, and functionality
  - The standard describes C and FORTRAN functions and requires the C functions to work with C++
- How these functions are coded is left up to implementations
- Many implementations available from many vendors
  - OpenMPI (not to be confused with OpenMP) is probably the most common today
  - MPICH is an early (but still around) implementation from Argonne
  - MVAPICH2 another popular open-source option
  - o Intel, IBM, Cray all have their own closed-source implementations
- The point: standard-compliant code will work with any up-to-date implementation

### Like Kokkos, MPI must be initialized and finalized

- Include the mpi.h header
- Passing arguments allows MPI to "cut out" its own command-line options
- MPI is C code, so no fancy scope guard object. Just gotta remember to finalize

```
#include <Kokkos Core.hpp>
#include <mpi.h>
int main(int argc, char** argv) {
 MPI Init(&argc, &argv);
 // Can alos use MPI Init(NULL, NULL) with no arguments.
 // If using Kokkos, initialize it after MPI.
 Kokkos::initialize(argc, argv);
    // Do stuff.
 Kokkos::finalize();
 MPI Finalize();
 return 0:
```

### A communicator is an ordered set of tasks, all connected

- A group is an unordered set of tasks
  - So a communicator is a group imbued with an ordering
- Each task's order is called its rank
- One communicator is always available, MPI\_COMM\_WORLD
  - This consists of all tasks that the program starts with
  - This is a macro that can be used in code, we'll see examples later
- A communicator is necessary for almost all MPI functions so that the implementation knows how the available tasks are connected
- New communicators can be created by splitting or reordering old communicators, or from scratch
  - hPIC2 doesn't do this however, we pretty much just stick with the original MPI\_COMM\_WORLD for simplicity

### Communicator size and task's rank found with simple utils

- Since MPI is C code, get function output by pre-allocating the output variable and passing a pointer to it
  - They couldn't just return things, could they...
- When you run this with multiple tasks, each task will fill their local rank variable with a different value
- So each task runs the same code, but with different data

```
int main(int argc, char** argv) {
    MPI_Init(&argc, &argv);

    int comm_size;
    MPI_Comm_size(MPI_COMM_WORLD, &comm_size);

    int rank;
    MPI_Comm_rank(MPI_COMM_WORLD, &rank);

    MPI_Finalize();
    return 0;
}
```

## Data arrays can be sent directly to other ranks

```
void MPI Send(
    void* data,
    int count,
    MPI Datatype datatype,
    int destination,
    int tag,
    MPI Comm communicator
```

Pointer to data. Implicitly cast Number of objects in array

Datatype of array. See next slide

Rank to which to send data
Unique tag for this message

Communicator in which to do send

# MPI has many pre-defined datatypes

- MPI uses these to determine size of sent array (**buffer**) in bytes
- Matching these datatypes is important
- Datatype sizes differ based on architecture
- On some system, a long int is the same size as an int, but on others, it's the same size as a long long int
  - Matching MPI datatype ensures the sizes are consistent

# MPI SHORT

MPI datatype

short int

int

long int

long long int

unsigned char

unsigned int

unsigned short int

unsigned long int

C equivalent

MPI INT

MPI LONG

MPI LONG LONG

MPI UNSIGNED CHAR

MPI UNSIGNED SHORT

MPI UNSIGNED

MPI UNSIGNED LONG

MPI UNSIGNED LONG LONG

MPI FLOAT

MPI DOUBLE

unsigned long long int float

double

long double

MPI LONG DOUBLE MPI BYTE char

### Message must be received on the other end

```
void MPI Recv(
    void* data,
    int count,
    MPI Datatype datatype,
    int source,
    int tag,
    MPI Comm communicator,
    MPI Status *status
```

Pointer to pre-alloc'd data. Implicitly cast

Length of received array

Received datatype

Source rank. Can be MPI\_ANY\_SOURCE

Tag matching desired MPI\_Send tag. Can be MPI\_ANY\_TAG

Communicator in which to receive

Object with info about how the receive went. Typically, MPI\_STATUS\_IGNORE to do nothing and save cycles

# Example: Send/Recv short array

- Both ranks allocate a length-3 array
- Rank 0 fills it with data and sends it to rank 1
- Rank 1 receives data and prints it to stdout
- If length is not known by destination, first do a Send/Recv for the length of the array
- Aside: code requires >1 task to run
  - Not good design!
  - Code should be designed to work with any number of tasks

```
#include <mpi.h>
#include <iostream>
int main(int argc, char** argv) {
    MPI Init(&argc, &argv);
    int rank:
    MPI Comm rank(MPI COMM WORLD, &rank);
    double data[3];
    if (rank == 0)
        data[0] = 0.0; data[1] = 1.0; data[2] = 2.0;
        MPI Send(
            data.
            MPI DOUBLE,
            1,
            MPI COMM WORLD
    else if (rank == 1) {
        MPI_Recv(
            data.
            MPI DOUBLE,
            0,
            MPI COMM WORLD,
            MPI STATUS IGNORE
        std::cout << "Received data: " << data[0] << ", "
                  << data[1] << ", " << data[2] << "\n";
    MPI Finalize();
    return 0;
```

### std::vectors can be used as the arrays

- vectors are guaranteed to be contiguous (data is stored in an unbroken line in memory) by the C++ standard
- The "data" method returns a raw pointer to the data stored by a vector
- The "size" method returns length
- Pass this directly to MPI
- vectors are more C++-like and allow for easy dynamic sizing

```
std::vector<double> array(3);
if (rank == 0)
    array[0] = 1.0; array[1] = 2.0; array[2] = 3.0;
    MPI Send(
        array.data(),
        array.size(),
        MPI DOUBLE,
        MPI COMM WORLD
else if (rank == 1) {
    MPI Recv(
        array.data(),
        array.size(),
        MPI DOUBLE,
        MPI COMM WORLD,
        MPI_STATUS_IGNORE
    std::cout << "Received data: " << array[0] << ", "
              << array[1] << ", " << array[2] << std::endl;
```

### Views can sometimes be used too

- MPI can be CUDA-aware, where it'll accept pointers to device memory
- If not, HostMirrors usually work
- Can tell at compile-time with some macro definitions
- Only works if View is contiguous, which is usually true unless the View is actually a SubView
  - Can use the .span\_is\_contiguous()
     method to check at runtime

```
Kokkos::View<double*> array("array", 3);
if (rank == 0) Kokkos::deep_copy(array, 1.0);
#if MPIX_CUDA_AWARE_SUPPORT && defined(KOKKOS_ENABLE_CUDA)
Kokkos::View<double*> mpi_array = array;
#else
Kokkos::View<double*>::HostMirror mpi_array = Kokkos::create_mirror_view(array);
Kokkos::deep_copy(mpi_array, array);
#endif
if (rank == 0) {
    MPI_Send(mpi_array.data(), 3, MPI_DOUBLE, 1, 0, MPI_COMM_WORLD);
}
else if (rank == 1) {
    MPI_Recv(mpi_array.data(), 3, MPI_DOUBLE, 0, 0, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
}
Kokkos::deep_copy(array, mpi_array);
```

### Send/Recv are **blocking** functions

- Send doesn't return until communication is complete
- Very easy to accidentally hang
- In example on right, both ranks
  want to Send an array to the other,
  then receive
  MPI\_Send(array\_A.MPI\_Recv(array\_B.MPI\_Recv(array\_B.MPI\_Send(array\_B.MPI
- But both just wait for the other to Recv forever!

(Note: modern compilers are smart enough that this might not actually hang, but leave nothing to chance)

```
std::vector<double> array_A(3);
std::vector<double> array_B(3);
if (rank == 0) {
    array_A[0] = 1.0; array_A[1] = 2.0; array_A[2] = 3.0;
    MPI_Send(array_A.data(), 3, MPI_DOUBLE, 1, 0, MPI_COMM_WORLD);
    MPI_Recv(array_B.data(), 3, MPI_DOUBLE, 1, 0, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
}
else if (rank == 1) {
    array_B[0] = 4.0; array_B[1] = 5.0; array_B[2] = 6.0;
    MPI_Send(array_B.data(), 3, MPI_DOUBLE, 0, 0, MPI_COMM_WORLD);
    MPI_Recv(array_A.data(), 3, MPI_DOUBLE, 0, 0, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
}
```

### Non-blocking communications can be used instead

- Functions preceded with "I" for "immediate return"
- Task continues after call, but data may still be in use, so shouldn't touch it until we are sure data is no longer in use by communicator
- Call comes with request handle that can be used to verify status of communication
- MPI\_Wait will wait until request is finished; safe to modify data afterward
- Non-blocking calls can be used for latency hiding
- Most MPI calls have a non-blocking version

```
std::vector<double> array A(3);
std::vector<double> array B(3);
MPI Request request;
if (rank == 0) {
    array A[0] = 1.0; array A[1] = 2.0; array A[2] = 3.0;
   MPI Isend(array A.data(), 3, MPI DOUBLE, 1, 0, MPI COMM WORLD, &request);
   MPI_Recv(array_B.data(), 3, MPI_DOUBLE, 1, 0, MPI_COMM_WORLD, MPI_STATUS_IGNORE)
   // Don't touch array A until the send is complete!
   MPI Wait(&request, MPI STATUS IGNORE);
   // Now okay to modify array A.
    array A[0] = 7.0;
else if (rank == 1)
    array B[0] = 4.0; array B[1] = 5.0; array B[2] = 6.0;
   MPI Isend(array B.data(), 3, MPI DOUBLE, 0, 0, MPI COMM WORLD, &request);
   MPI Recv(array A.data(), 3, MPI DOUBLE, 0, 0, MPI COMM WORLD, MPI STATUS IGNORE)
   MPI Wait(&request, MPI STATUS IGNORE);
```

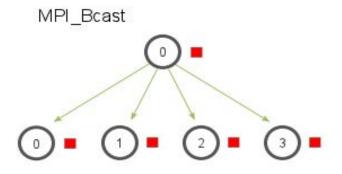
# **Collective** functions must be called by all tasks in a communicator

- Definition: tasks will wait at collective function call until EVERY task has reached it
- Simplest is MPI\_Barrier, which does nothing except effectively wait until all tasks reach it
- If one task is stuck, ALL tasks hang!
- Practically, it makes little sense to do P2P communications unless code is tailored to a certain # of tasks, so we mostly use collective communications

MPI\_Barrier(MPI\_COMM\_WORLD);

### **Bcast broadcasts**

- Copies data from an array on one rank to all other ranks
- In example, all ranks alloc, but only rank
   0 fills data. Then Bcast copies it to all other ranks
- Generally wanna choose rank 0 as the root
  - It's the only rank guaranteed to appear regardless of the number of tasks
  - Scales up to arbitrarily many tasks



https://mpitutorial.com/

```
std::vector<double> array(3);
if (rank == 0) {
    array[0] = 1.0; array[1] = 2.0; array[2] = 3.0;
}
MPI_Bcast(
    array.data(), // Pointer to array
    3, // Length of array
    MPI_DOUBLE, // Array datatype
    0, // Root rank
    MPI_COMM_WORLD // Communicator
);
```

# Scatter splits an array into chunks and sends the chunks

### out to all other tasks

- Rank 0 fills an array with 2 entries per task
- All tasks participate in the Scatter, including the root rank
- The receive buffer from each task is filled with its corresponding chunk from the root task's send buffer

```
MPI Comm size(MPI COMM WORLD, &comm size);
int rank;
MPI Comm rank(MPI COMM WORLD, &rank);
std::vector<double> received(2);
std::vector<double> sent(2*comm size);
if (rank == 0) {
    for (int i = 0; i < sent.size(); i++) {
        sent[i] = i;
MPI Scatter(
                    // send buffer. only needs to be valid on root rank
    sent.data(),
                    // Number of elements to send to each rank
    MPI DOUBLE,
                    // Send buffer datatype
    received.data(), // Buffer into which to put received data
                     // Length of receive buffer
    2,
   MPI DOUBLE,
                    // Receive buffer datatype
    MPI COMM WORLD
std::cout << "Rank " << rank << " received: "
          << received[0] << " " << received[1] << std::endl;</pre>
```

Example output with 4 tasks:

```
Rank 0 received: 0 1
Rank 2 received: 4 5
Rank 1 received: 2 3
Rank 3 received: 6 7
```

### What if we want each rank to receive a different size?

- Use Scatterv (v for variable-length)
- Root rank puts all data into one big buffer
- Root also needs an array to tell MPI how many things to send to each task, and where that task's data starts in the big send buffer
- Receiving ranks only need to know how many things to receive

```
Output with 4 tasks:

Rank 0 received:
Rank 2 received: 1 2
Rank 3 received: 3 4 5
Rank 1 received: 0
```

```
int comm size;
MPI Comm size(MPI COMM WORLD, &comm size);
MPI Comm rank(MPI COMM WORLD, &rank);
std::vector<double> received(rank);
std::vector<double> sent;
std::vector<int> counts;
std::vector<int> displacements;
if (rank == 0) {
    int total num to send = comm size * (comm size - 1) / 2;
    sent.resize(total num to send);
    counts.resize(comm size);
    displacements.resize(comm size);
    for (int i=0; i<total num to send; i++) {
        sent[i] = i;
    for (int i=0; i<comm size; i++) {
        counts[i] = i;
        displacements[i] = i * (i - 1) / 2;
MPI Scatterv(
                          // send buffer. only needs to be valid on root rank
    sent.data(),
    counts.data(),
                          // Number of entries to send to each rank.
    displacements.data(), // Offset in send buffer for each rank's data
                          // Send buffer datatype
    MPI DOUBLE,
    received.data(),
                          // Buffer into which to put received data
    rank,
                          // Length of receive buffer
    MPI DOUBLE,
                          // Receive buffer datatype
    MPI COMM WORLD
std::cout << "Rank " << rank << " received: ";</pre>
for (int i=0; i<rank; i++) {
    std::cout << received[i] << " ";
std::cout << std::endl;</pre>
```

### Gather is like the opposite of Scatter

MPI\_Gather

1 2 3 3 https://mpitutorial.com/

- All tasks combine chunks into one big array and send it to a root rank
- Like Scattery, there is a Gathery where number of elements received from each task is variable.

Output with 4 tasks

Received: 0 0 1 2 2 4 3 6

```
int comm size;
MPI Comm size(MPI COMM WORLD, &comm size);
int rank:
MPI Comm rank(MPI COMM WORLD, &rank);
std::vector<double> to send(2);
to send[0] = rank; to send[1] = rank*2.0;
std::vector<double> to recv:
if (rank == 0) {
    to recv.resize(2*comm size);
MPI Gather(
    to send.data(), // Send buffer
                   // Number of elements to send
    2.
    MPI DOUBLE, // Type of elements to send
    to recv.data(), // Receive buffer
                  // Number of elements received per task
    2.
    MPI DOUBLE, // Type of elements received
         // Rank of root
    0.
    MPI COMM WORLD // Communicator
if (rank == 0) {
    std::cout << "Received: ":
    for (int i=0; i<to recv.size(); i++) {
        std::cout << to recv[i] << " ";
    std::cout << std::endl;</pre>
```

# Allgather to get result of gather on every task

- Just a Gather but without a root, every task must be prepared to receive the data
- Yes, there is an Allgatherv!

#### Output with 4 tasks

```
Rank 2 received: 0 0 1 2 2 4 3 6
Rank 3 received: 0 0 1 2 2 4 3 6
Rank 0 received: 0 0 1 2 2 4 3 6
Rank 1 received: 0 0 1 2 2 4 3 6
```

# MPI\_Allgather 1 2 https://mpitutorial.com/

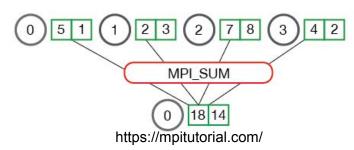
```
int comm size;
MPI Comm size(MPI COMM WORLD, &comm size);
int rank:
MPI Comm rank(MPI COMM WORLD, &rank);
std::vector<double> to send(2);
to send[0] = rank; to send[1] = rank*2.0;
std::vector<double> to recv(2*comm size);
MPI Allgather(
    to send.data(), // Send buffer
             // Number of elements to send
    2.
   MPI DOUBLE, // Type of elements to send
   to recv.data(), // Receive buffer
           // Number of elements received per task
   MPI DOUBLE,
                  // Type of elements received
   MPI COMM WORLD // Communicator
std::cout << "Rank " << rank << " received: ";
for (int i=0; i<to_recv.size(); i++) {
    std::cout << to recv[i] << " ";
std::cout << std::endl;
```

- Essentially sums each entry of an array over the tasks
- But it doesn't just have to be a sum!
  - Can also find the min or max, a product, or logical OR and AND
- Allreduce sends reduction result to all tasks, like Allgather
- There is no Reducev; all tasks must submit array of the same length!

Output with 4 tasks

Received: 6 12

MPI\_Reduce



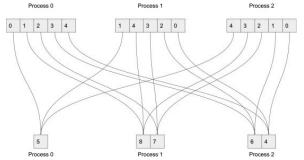
```
int comm size;
MPI Comm size(MPI COMM WORLD, &comm size);
int rank;
MPI Comm rank(MPI COMM WORLD, &rank);
std::vector<double> to reduce(2);
to reduce[0] = rank; to reduce[1] = rank*2.0;
std::vector<double> reduced;
if (rank == 0) {
    reduced.resize(2);
MPI Reduce(
    to reduce.data(), // Send buffer
   reduced.data(), // Receive buffer
    MPI DOUBLE,
                     // Type of elements to reduce
   MPI SUM,
                      // Reduction operator
                      // Rank of root
    MPI COMM WORLD
if (rank == 0) {
    std::cout << "Received: ";</pre>
    for (int i=0; i<reduced.size(); i++) {
        std::cout << reduced[i] << " ";
    std::cout << std::endl;</pre>
```

### Reduce scatter is like Reduce + Scatter

- Performs a reduction over all tasks on an array
- Then split up the result into chunks and send each chunk to a different task
- Useful for setting up potential solve

```
Output with 4 tasks

Rank 0 received:
Rank 3 received: 18 24 30
Rank 1 received: 0
Rank 2 received: 6 12
```



```
int comm size;
MPI Comm size(MPI COMM WORLD, &comm size);
int rank:
MPI Comm rank(MPI COMM WORLD, &rank);
int total_size = comm_size * (comm_size - 1) / 2;
std::vector<double> to reduce(total size);
std::vector<double> reduced(rank);
std::vector<int> counts(comm size);
for (int i=0; i<total_size; i++) {
    to reduce[i] = i * rank;
for (int i=0; i<comm size; i++) {
   counts[i] = i;
MPI Reduce scatter(
    to reduce.data(), // Send buffer
    reduced.data(). // Receive buffer
    counts.data(),
                     // Number of elements to scatter to each task
                     // Type of elements to reduce
   MPI DOUBLE,
   MPI SUM,
   MPI COMM WORLD
std::cout << "Rank " << rank << " received: ";
for (int i=0; i<reduced.size(); i++) {
    std::cout << reduced[i] << " ";
std::cout << std::endl;
```

### To run MPI code, use mpirun or mpiexec

- E.g., rather than just ./hpic2, use mpiexec -np 4 ./hpic2 to run with 4 tasks
- There are a million different command-line options for mpiexec
  - Look up the command-line options for your implementation on its man pages
- Important to note that number of tasks is decided at runtime, NOT when writing code
- We could check to make sure the user is running with a valid number of tasks, but better practice to write code that works no matter how many tasks the user throws at it

### The "MPI Mindset"

- With Kokkos, explicitly define a work load for each thread and launch it
- With MPI, when you write code, you must realize that every task will be executing that code, perhaps with different data
- Code should be designed to work with arbitrarily many tasks
- Communication is the enemy
  - Avoid communication if at all possible. Usually having every task perform some computation is cheaper than having one task do it and broadcast results
  - Reduce size of communications. Instead of sending a whole array, consider sending only the portion needed by other tasks
  - Hide latency. The actual sending/receiving is done by a handler outside of your program. Set up communication as early as possible, do task-local calculations while waiting for data to be received.
- I find the man pages at <a href="https://www.open-mpi.org/doc/v4.1/">https://www.open-mpi.org/doc/v4.1/</a> to be super useful in reminding myself of syntax and functionality

### Exercise: distribute your 1D particle pusher!

We'll reference the exercise from the Kokkos session, so have that finished and handy. We're going to distribute it in the same way that hPIC2's finite-difference mesh pusher is distributed: particles are pushed throughout the entire domain, but the domain is split up for the field solve. Don't worry about the field solve just yet. What it means is that we will reduce the accumulated charge density over the mesh, but split up the reduction result and send chunks to the MPI ranks that will be responsible for the field on that chunk (sound familiar?)

- 1. Do the exercise from Session 2 if you haven't already.
- 2. Add MPI\_Init and MPI\_Finalize.
- 3. Split the particles as evenly as possible over the tasks, ensuring that the total number of particles is as desired even when it is not divisible by the number of tasks.
- 4. Do the same for the nodes.
- 5. Initialize and push only the local number of particles.
- 6. After every push, use MPI\_Reduce\_scatter to reduce the charge density over all tasks and distribute chunks to local arrays having the length that you determined in step 4.