

# Programming Exercise 2 - Combinatorial Optimization

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## Algorithm 1 Minimum Mean Weight Cycle

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1: procedure MINMEANWEIGHTCYCLE
2:   if  $G$  has no cycle then
3:     return
4:    $\gamma \leftarrow \max\{c(e) : e \in E(G)\}$ 
5:    $T \leftarrow \emptyset$ 
6:    $c'(e) := c(e) - \gamma$ 
7:   loop:
8:      $J \leftarrow \text{MinWeightEmptyJoin}()$ 
9:     if  $c'(J) = 0$  then
10:      return 0 -  $c'$ -weight-cycle
11:    else
12:       $\gamma' \leftarrow \frac{c'(J)}{|J|}$ 
13:       $c'(e) \leftarrow c'(e) - \gamma', \forall e \in E(G)$ 

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## Algorithm 2 Minimum Weight $\emptyset$ -Join

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1: procedure MINWEIGHTEMPTYJOIN
2:    $T \leftarrow V^-$ 
3:    $G_i \leftarrow G(T)$ , with cost function  $d(e) := |c(e)|$ 
4:    $\overline{G}_i \leftarrow \text{MetricClosure}(G_i)$ 
5:    $M \leftarrow \text{MinWeightPerfectMatching}(G_i), \forall i$ 
6:    $J \leftarrow P_{\{x_1, y_1\}} \triangle \dots \triangle P_{\{x_m, y_m\}}$ , where  $P_{x,y}$  is the min x-y-path
7:   return  $J \triangle E^-$ 

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Algorithm 1 was in exercise 7.4. Algorithm 2 is (almost explicitly) stated in Theorems 51 and 52 of the lecture notes. For Algorithm 2,  $E^-$  is the set of edges with negative weight,  $V^-$  the set of vertices that are incident with an odd number of edges in  $E^-$