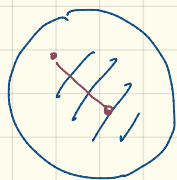
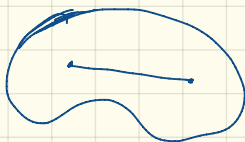


convex set



convex



non-convex

Let $x, y \in \mathbb{R}^n$. The line segments between x and y in the set

$$\{(1-\lambda)x + \lambda y : \lambda \in [0, 1]\}.$$

A set $C \subseteq \mathbb{R}^n$ is convex if $\forall x, y \in C$, C contains the line segment between x and y . ^{of points}

Consider the set

$$H = \{x \in \mathbb{R}^n : a^T x \geq \beta\}$$

Halfspace, a set satisfies a single linear inequality with the coefficients are not all zero. ^{conditions define the set}

$$\text{where } a \in \mathbb{R}^n, a \neq 0 \\ \beta \in \mathbb{R}$$

Take $x, y \in H$,

consider $z = (1-\lambda)x + \lambda y$ for some $\lambda \in [0, 1]$

Then $a^T z = a^T ((1-\lambda)x + \lambda y)$

apply distribute law You can treat the word polyhedron as a fancy name for the

$$\begin{aligned} a^T z &= \underbrace{(1-\lambda)}_{\geq 0} \underbrace{a^T x}_{\geq \beta} + \underbrace{\lambda}_{\geq 0} \underbrace{a^T y}_{\geq \beta} \\ &\geq (1-\lambda)\beta + \lambda\beta = \beta \end{aligned}$$

feasible region of a linear programming problem.

The set of feasible solutions are convex. $a^T z \geq \beta \Rightarrow z \in H$, so H is convex going to be convex.

Fact.

~~The intersection of two convex sets is convex.~~ due to this fact
A well studied problem.

A polyhedron is the intersection of finitely many halfspace. This set is called polyhedron.

$$\begin{aligned} \text{e.g. } \left\{ \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} : x_1 + 2x_2 \geq 5 \right\} \cap \left\{ \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} : -2x_1 + x_2 \geq 0 \right\} \\ = \left\{ \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} : \begin{array}{l} x_1 + 2x_2 \geq 5 \\ -2x_1 + x_2 \geq 0 \end{array} \right\} \end{aligned}$$

Basically. we can see a polyhedron is simply a set of feasible solutions of a linear programming problem.