

$$\begin{array}{lll}
 t & ::= & x \quad \text{Variable} \\
 & | & \lambda x. t \quad \text{Abstraction} \\
 & | & (t, t) \quad \text{Application}
 \end{array}$$

Grammar 1:  $\lambda$ -calculus syntax

## 1 Terms, Types and Kinds

For programming language, three levels: *terms*, *types*, and *kinds*, have proved sufficient .

Recap term-level abstraction and application in the  $\lambda$ -calculus as shown in Grammar 1.

$\Gamma \vdash T :: K$  is read as “type T has kind K in context  $\Gamma$ ”.

*kinding* is a well-formedness relation.

To treat type-level functions, collectively called *type operators* more formally, it is required to:

1. Add a collection of rules of *kinding* which specify how type expressions can be combined to yield new type expressions.
2. Whenever a type T appears in a term  $(\lambda x : T. t)$ , check whether T is well formed.
3. Add a collection of rules for the definitional equivalence relations between types.

