

Web Data Mining

Lecture 8: PageRank and HITS

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Humla v0.3

Overview

- Web Structure Mining
- PageRank
- HITS

Web Structure Mining (Recall)

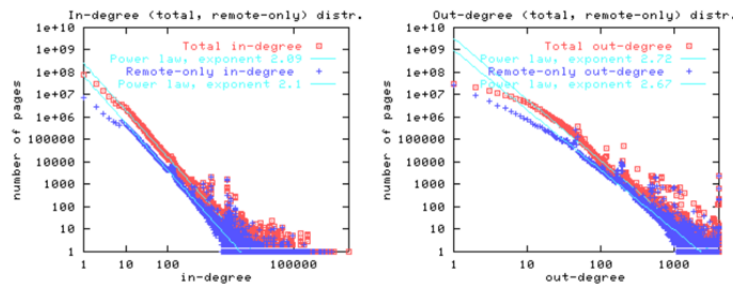
- Main ideas
 - Use graph theory to analyze the node and connection structure of a web site.
 - Help the users to retrieve the relevant documents by analyzing the link structure of the Web.
- Tasks
 - Hyperlink analysis
 - Intra-page vs Inter-page.
 - Analysis of the tree-like structure of page structures
- Applications
 - Document retrieval and ranking
 - Discovery of hubs and authorities
 - Discovery of web communities
 - Citation networks
 - Social network analysis
 - Search engines, SEO, ...

Web Graph

- Terminology
 - Web graph
 - a directed graph representing the web
 - Node
 - web page in the graph
 - Edge
 - hyperlink on the web page
 - In-links (backlinks)
 - links pointing to the node
 - Out-Links
 - links generated from the node
 - In-Degree
 - number of links pointing to the node
 - Out-Degree
 - number of links generated from the node

Web Graph Analysis

- Web graph statistics
 - A. Broder et al., *Graph structure in the web, 2000* - they analyzed the web graph consisting of 200 million pages and 1.5 billion links from AltaVista.
- In (Out) Degree - Power law
 - the probability that a node has in(out)-degree i is proportional to $\frac{1}{i^x}$, where $x=2.1$.

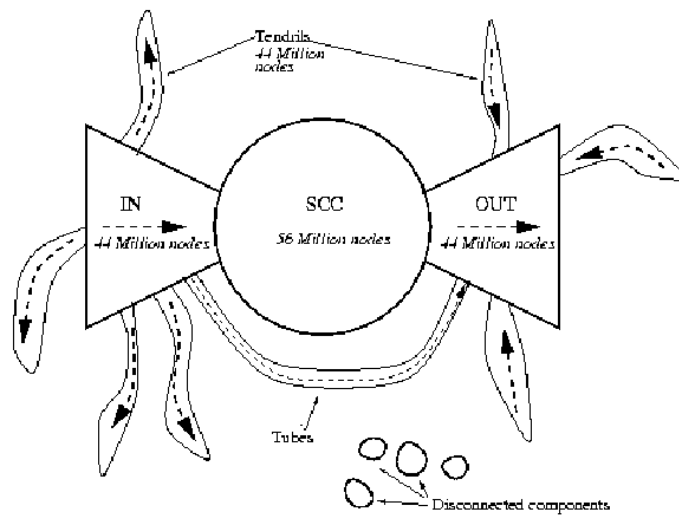


- Web graph
 - Bow Tie Structure
- Applications
 - Important for designs and implementations of main crawlers, search engines, etc.

The Bow-Tie Structure

- Presents the connectivity of the web
 - Web isn't the fully interconnected network
- Components
 - SCC - giant strongly connected component
 - central core, all of whose pages can reach one another along directed links
 - IN
 - pages that can reach the SCC, but cannot be reached from it.
 - e.g. new pages not yet discovered
 - OUT
 - pages that are accessible from the SCC, but do not link back to it.
 - e.g. corporate pages with internal links only
 - Tendrils
 - pages reachable from IN but cannot reach the SCC
 - e.g. single page or document with no out-links
 - pages that can reach the OUT but cannot be reached from the SCC
 - Tubes
 - TENDRILS that fulfill both assumptions
 - e.g. a single page linking only a blog post about a company that links to the pages with internal links
 - Disconnected

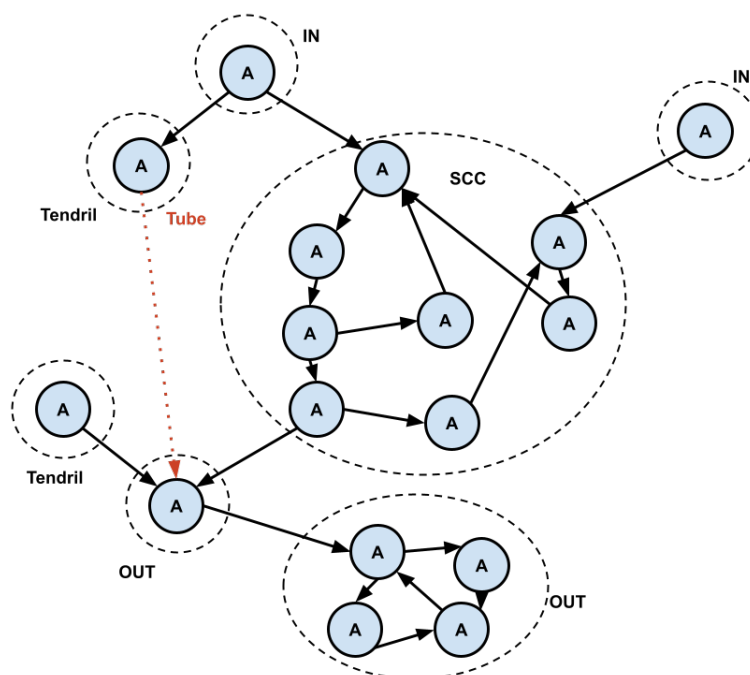
The Bow-Tie Structure (cont.)



- "the chance of being able to surf between two randomly chosen pages is less than one in four"

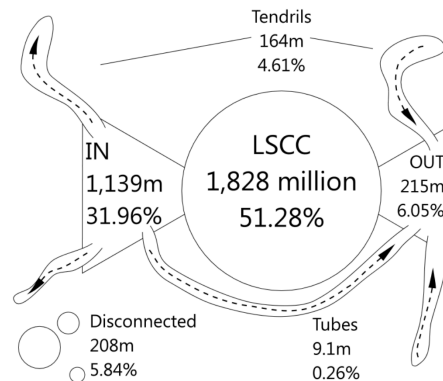
- A. Broder, R. Kumar, F. Maghoul, P. Ragha-van, S. Rajagopalan, S. Stata, A. Tomkins, and J. Wiener. Graph structure in the web. Computer Net-works, 33:309–320, June 2000

The Bow-Tie Structure Example



The Bow-Tie Structure Revisited

- Power law exponent (2000 vs 2012): 2.1 vs 2.24
- Average degree: 7.5 vs 36.8
- SCC: 27.7% vs 51.3%
- IN, OUT: 21%,21% vs 31%,6%
- Pairs of connected pages: 25% vs 48%



- Robert Meusel, Sebastiano Vigna, Oliver Lehmborg, and Christian Bizer. 2014. Graph structure in the web — revisited: a trick of the heavy tail. In Proceedings of the 23rd International Conference on World Wide Web (WWW '14 Companion). ACM, New York, NY, USA, 427-432.

Application: Improving Search Results

- Web Search
 - Can build on top of existing boolean and vector models from Information Retrieval.
 - Vector based model was used in AltaVista.
- Issues of basic IR models
 - Results are too large that the user can explore.
 - All documents are treated equally according to the relevance point of view.
 - Results are returned only using the text based matching approaches.
 - Heavily influenced by many spam techniques
 - e.g. keyword stuffing
- Need for other relevance/popularity scores
 - Web structure is the most well known source of additional information about popularity of web pages.

Overview

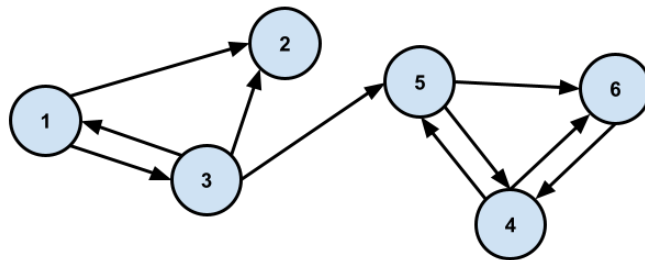
- Web Structure Mining
- PageRank
- HITS

PageRank

- Introduced in April, 1998 at WWW98 by Sergey Brin and Larry Page in a paper titled "The anatomy of a large-scale hypertextual Web search engine."
 - *Uses link structure as an indicator of an individual page's quality.*
 - *The prestige of a page is proportional to the sum of the prestige scores of pages linking to it.*
 - *Prestige is independent of any information need or query.*
- Main formula
 - $\pi^{(k+1)T} = \pi^{(k)T}(\alpha S + (1 - \alpha)E)$
- Characteristics
 - *ability to fight spam, global measure and is query independent, computed off-line, very efficient at the query time.*

PageRank Computation

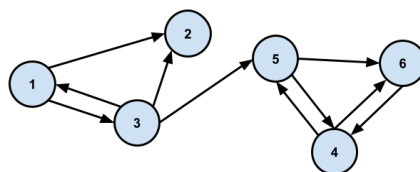
- **Main idea**
 - If a web page is pointed to by other, important pages, then it's also an important page.
 - Think as kind of "fluid" that circulates through networks.
- **PageRank for one page**
 - $r(P_i) = \sum_{P_j \in B_{P_i}} \frac{r(P_j)}{|P_j|}$
 - B_{P_i} - set of pages linking to P_i
 - $|P_j|$ - number of outlinks from P_j
- **Examples:**
 - $r(P_1) = \frac{r(P_3)}{3}, r(P_2) = \frac{r(P_1)}{2} + \frac{r(P_3)}{3}$



Iterative computation of the PageRank

- Next iteration (k+1) uses states from the previous one (k)
 - $r_{k+1}(P_i) = \sum_{P_j \in B_{P_i}} \frac{r_k(P_j)}{|P_j|}$
- PageRank is initialised with a predefined value
 - $\forall i: r_0(P_i) = \frac{1}{n}$

Node	Iteration 0	Iteration 1	Iteration 2	Order (after 2nd iteration)
P1	$\frac{1}{6}$	$\frac{1}{18}$	$\frac{1}{12} \times \frac{1}{3} = \frac{1}{36}$	5.
P2	$\frac{1}{6}$...	$\frac{1}{18}$	4.
P3	$\frac{1}{6}$	$\frac{1}{6} \times \frac{1}{2} = \frac{1}{12}$	$\frac{1}{36}$	5.
...
P6	$\frac{1}{6}$	$\frac{1}{6}$	$\frac{14}{72}$	2.

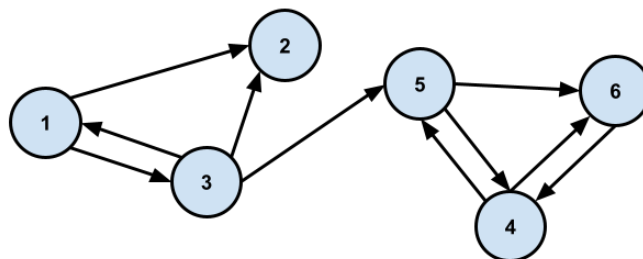


Matrix Representation

- **Mathematically**
 - a system of n linear equations with n unknown variables.
 - Can be represented as a matrix.
- **PageRank vector**
 - $\pi = (r_0(P_1), r_0(P_2), \dots, r_0(P_n))$
- **Use matrix H ($n \times n$)**
 - $H_{ij} = \frac{1}{|P_i|}$ if there is a link from P_i to P_j
 - $H_{ij} = 0$ otherwise
- **Circular definition, where the iterative algorithm is used to solve**
 - $\pi^{(k+1)} = \pi^{(k)} H$
 - The equation is the characteristic equation used for finding the eigensystem of the matrix.
 - π is an eigenvector with the corresponding eigenvalue of 1.
 - 1 is the largest eigenvalue and the PageRank vector P is the principal eigenvector
 - Also called power method
- **Issues:**
 - the Web graph does not meet all conditions
 - There are many pages without any out-links, as well as directed paths leading into a cycle, ...

Matrix Representation (cont.)

	P1	P2	P3	P4	P5	P6
P1	0	$\frac{1}{2}$	$\frac{1}{2}$	0	0	0
P2	0	0	0	0	0	0
P3	$\frac{1}{3}$	$\frac{1}{3}$	0	0	$\frac{1}{3}$	0
P4	0	0	0	0	$\frac{1}{2}$	$\frac{1}{2}$
P5	0	0	0	$\frac{1}{2}$	0	$\frac{1}{2}$
P6	0	0	0	1	0	0



Iterative computation using Matrix

- Using the equation:

$$-\pi^{(k+1)} = \pi^{(k)}H$$

- $\pi^{(0)} = \left(\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}\right)$

	P1	P2	P3	P4	P5	P6
P1	0	$\frac{1}{2}$	$\frac{1}{2}$	0	0	0
P2	0	0	0	0	0	0
P3	$\frac{1}{3}$	$\frac{1}{3}$	0	0	$\frac{1}{3}$	0
P4	0	0	0	0	$\frac{1}{2}$	$\frac{1}{2}$
P5	0	0	0	$\frac{1}{2}$	0	$\frac{1}{2}$
P6	0	0	0	1	0	0

- $\pi^{(1)} = \pi^{(0)}H = \left(\frac{1}{18}, \frac{5}{36}, \frac{1}{12}, \frac{1}{4}, \frac{5}{36}, \frac{1}{6}\right)$

Matrix Representation and Computation

- Complexity
 - Every iteration requires $O(n^2)$
 - Multiplication of PageRank vector of size n and matrix of size $n \times n$
- The matrix is **sparse**
 - Most of the elements are zero
 - Efficient memory representations using LIL (List of List), CSR (Compressed Sparse Row) or CSC (Compressed Sparse Column), ...
 - There are many efficient algorithms for sparse matrix multiplication with complexity $O(nnz)$, where nnz is number of non-zero elements.
- The matrix is close to the stochastic (transition) matrix of probabilities in Markov chain models.
 - Fulfills the "memorylessness" Markov property
 - If one can make predictions for the future without knowing history.
 - Except dangling pages - pages that have no out-links!

Markov Chains

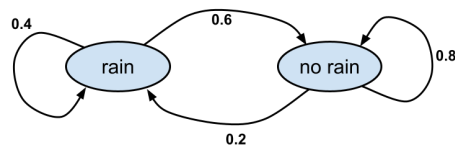
- Markov Chains

- named after Andrey Markov
- mathematical systems that hop from one state to another
- special type of stochastic model
 - the simplest from Markov models
 - the future state depends only on the present state and not on the history

- Example

- Weather
 - raining today
 - 40% rain tomorrow
 - 60% no rain tomorrow
 - not raining today
 - 20% rain tomorrow
 - 80% no rain tomorrow

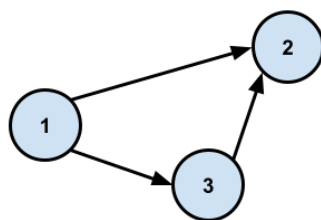
$$- P = \begin{bmatrix} 0.4 & 0.6 \\ 0.2 & 0.8 \end{bmatrix}$$



Issues of the Matrix Representation

- Rank sinks

- pages that have no out-links
- it does not distribute the PageRank to others
- continuously decrease the overall PageRank in the graph



- Example

$$- \pi^{(0)} = (1/3, 1/3, 1/3)$$

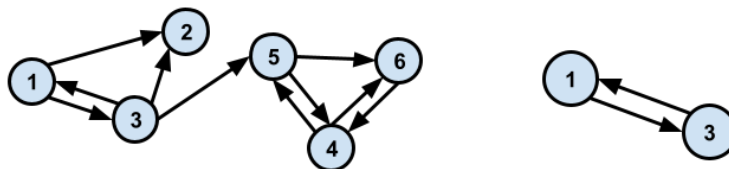
$$- (1/3, 1/3, 1/3) \times \begin{bmatrix} 0 & 1/2 & 1/2 \\ 0 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix} = (0, 1/2, 1/6) = \pi^{(1)}$$

$$- (0, 1/2, 1/6) \times \begin{bmatrix} 0 & 1/2 & 1/2 \\ 0 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix} = (0, 1/6, 0) = \pi^{(2)}$$

$$- (0, 0, 0) = \pi^{(3)}$$

Issues of the Matrix Representation (cont.)

- **Link farms**
 - group of pages that link to every other page in the group
 - a link farm is a clique
 - they support each other
- **Cycles**
 - cause oscillation of the PageRank between them



- **Example**
 - $\pi^{(0)} = (0, 1)$
 - $(0, 1) \times \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} = (1, 0) = \pi^{(1)}$
 - $(1, 0) \times \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} = (0, 1) = \pi^{(2)}$
 - ...

Alternative PageRank Definition

- **Random walk/random surfer**
 - Someone who is randomly browsing a network.
 - Choosing a page at random, picking each page with equal probability.
 - Follow links for a sequence of k steps.
 - In each step, they pick a random out-going link from their current page, and follow it to where it leads.
- Randomly following links is called a random walk
- **Claim**
 - The probability of being at a page X after k steps of this random walk is precisely the PageRank of X after k applications of the Basic PageRank Update Rule.
- **Issue**
 - Rank sink and cycles
- **Solution**
 - Teleportation to a random node

Transition Probability Matrix

- Stochasticity adjustment of matrix H to matrix S

- Update of the dangling node row
→ setting all the zeros to $1/n$
- Random teleport/jump

- $S = H + a(\frac{1}{n}e^T)$

- a is a vector of length n
→ $a_i = 1$ if there is no outlink from P_i
→ $a_i = 0$ otherwise
- $e^T = (1, 1, 1, 1, 1, 1)$

$$\begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 1/3 & 1/3 & 0 & 0 & 1/3 & 0 \\ 0 & 0 & 0 & 0 & 1/2 & 1/2 \\ 0 & 0 & 0 & 1/2 & 0 & 1/2 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{bmatrix} \rightarrow \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 & 0 \\ 1/6 & 1/6 & 1/6 & 1/6 & 1/6 & 1/6 \\ 1/3 & 1/3 & 0 & 0 & 1/3 & 0 \\ 0 & 0 & 0 & 0 & 1/2 & 1/2 \\ 0 & 0 & 0 & 1/2 & 0 & 1/2 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{bmatrix} = S$$

Transition Probability Matrix (cont.)

- S matrix is stochastic

- sum of values in row is equal to 1
- non-negative and square

- Transition matrix for a finite Markov chain

- Probability of using the link for the random walk

- Issue

- It is not irreducible
→ Web graph is strongly connected
→ for each pair of nodes, there is a path from one to another one
- It is not aperiodic
→ periodic - all paths leading from one node back to that node
- Convergence issue!

$$S = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 & 0 \\ 1/6 & 1/6 & 1/6 & 1/6 & 1/6 & 1/6 \\ 1/3 & 1/3 & 0 & 0 & 1/3 & 0 \\ 0 & 0 & 0 & 0 & 1/2 & 1/2 \\ 0 & 0 & 0 & 1/2 & 0 & 1/2 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{bmatrix}$$

Google Matrix

- Solution for irreducible and aperiodic situation
 - Primitivity adjustment
 - We add a link from each page to every page and give each link a small transition probability controlled by a parameter **damping factor d** (e.g. 0.85)
- Updated model
 - Random surfer has two options
 - With probability d , he randomly chooses an out-link to follow.
 - With probability $1-d$, he jumps to a random page without a link.
 - Surfer may get bored, or interrupted
- Google matrix
 - Becomes strongly connected
 - link from each page to every page
 - Becomes aperiodic
 - random surfer does not have to traverse a fixed cycle
- $G = d \times S + (1 - d) \frac{E}{n}$
 - d is damping factor
 - E is $e \times e^T$ - is a $n \times n$ square matrix of all 1

Google Matrix (cont.)

- $G = d \times S + (1 - d) \frac{E}{n}$
 - $d = 0.9$

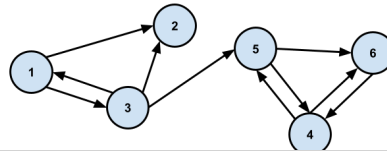
$$S = \begin{bmatrix} 0 & 1/2 & 1/2 & 0 & 0 & 0 \\ 1/6 & 1/6 & 1/6 & 1/6 & 1/6 & 1/6 \\ 1/3 & 1/3 & 0 & 0 & 1/3 & 0 \\ 0 & 0 & 0 & 0 & 1/2 & 1/2 \\ 0 & 0 & 0 & 1/2 & 0 & 1/2 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{bmatrix}$$

- $G = 0.9 \times S + 0.1 \frac{E}{6}$

$$G = \begin{bmatrix} 1/60 & 7/15 & 7/15 & 1/60 & 1/60 & 1/60 \\ 1/6 & 1/6 & 1/6 & 1/6 & 1/6 & 1/6 \\ 19/60 & 19/60 & 1/60 & 1/60 & 19/60 & 1/60 \\ 1/60 & 1/60 & 1/60 & 1/60 & 7/15 & 7/15 \\ 1/60 & 1/60 & 1/60 & 7/15 & 1/60 & 7/15 \\ 1/60 & 1/60 & 1/60 & 11/12 & 1/60 & 1/60 \end{bmatrix}$$

PageRank Computation

- Power Iteration Method
 - $\pi^{(k+1)} = \pi^{(k)} G$
- Google matrix
 - Stochastic, Irreducible, Aperiodic, Primitive
 - No-zero elements
 - It is not sparse any more!
- Computation
 - Complexity $O(n^2)$
- Example
 - 50 iterations
 - $\pi = (0.03721, 0.05396, 0.04151, 0.3751, 0.206, 0.2862)$
 - Order: 4, 6, 5, 2, 3, 1

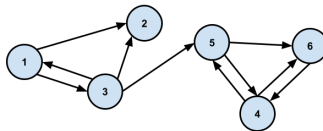


PageRank Computation (cont)

- Convert to operations with sparse matrix
 - $\pi^{(k+1)} = \pi^{(k)} G$
 - $\pi^{(k+1)} = d\pi^{(k)} H + (d\pi^{(k)} a + 1 - d) \frac{e^T}{n}$
- The most computational intensive operation
 - multiplication of vector and matrix uses sparse matrix H
- Convergence criteria
 - 1-norm
 - the iteration ends after the 1-norm of the residual vector is less than a pre-specified threshold δ
 - 1-norm for a vector is simply the sum of all the components
 - page-order
 - no significant change of the page order between iterations
 - usually around 50

PageRank Example

- Example 1
 - iterations: 50
 - damping factor: 1.0
 - following links
 - $\pi = (7.18e - 10, 1.24e - 09, 8.36e - 10, 0.44, 0.22, 0.33)$
- Example 2
 - iterations: 50
 - damping factor: 0.0
 - random choosing
 - $\pi = \left(\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}\right)$
- Google
 - damping factor ≈ 0.85



PageRank Modifications

- Intelligent surfer
 - Modification of probabilities in transition matrix
 - Analysis of users behavior
 - Using click logs, ...
 - Similarities of pages
 - Using cosine similarity
 - Anchor text, or the surrounding information
- Personalization
 - Modification of the teleportation
 - $G = d \times S + (1 - d) \frac{E}{n}$
 - E is $e \times e^T$ - is a $n \times n$ square matrix of all 1
 - Change $e \times e^T$ to $e \times v^T$, where v^T provides information about preferences for specific pages

Overview

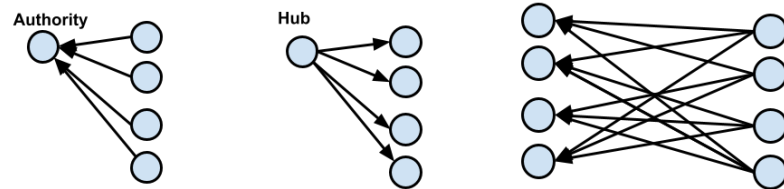
- Web Structure Mining
- PageRank
- **HITS**

HITS

- **HITS**
 - *Hypertext Induced Topic Search*
 - *Presented by Jon Kleinberg in January, 1998 at the Ninth Annual ACM-SIAM Symposium on Discrete Algorithms.*
 - *utilizes the web structure as important aspect*
 - *a query is used to select a subgraph from the Web*
- **Main characteristics**
 - *Search query dependent*
 - *Two rankings*
 - *authority ranking and hub ranking*
- **Approach**
 - *For a search query, HITS first expands the list of relevant pages returned by a search engine and then produces rankings of the expanded set of pages.*

Hubs and Authorities

- **Hub**
 - *page with many outlinks*
 - *page is a source of many important links to authority pages relevant for the topic*
- **Authority**
 - *a page with many inlinks*
 - *if people trust the page, they link to it and it becomes the authority*
- **The goal**
 - *Find best hubs and authorities*
 - *Good authorities are linked by good hubs*
 - *Good hubs link to good authorities*



HITS Algorithm

- **Collecting pages**
 - *HITS sends a query to a search engine and collects top t highest ranked pages that are relevant to the query (e.g. $t=200$)*
 - *Called root set W*
 - *Grows W by including pages that link to any page in W or are linked by any page from W . At most k per page. (e.g. $k=50$)*
 - *Called base set S (size 1000-5000)*
- **Graph**
 - *HITS works with the graph $G(V,E)$ composed from all pages in the base set S .*
 - *L is the adjacency matrix of the graph G .*
- **Scores**
 - *Authority score*
 - $a(i)^k = \sum_{(j,i) \in E} h(j)^{(k-1)}$
 - *Hub score*
 - $h(i)^k = \sum_{(i,j) \in E} a(j)^{(k-1)}$

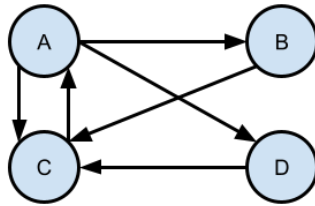
HITS Algorithm (cont.)

- Matrix representation
 - Similar to PageRank
 - $\rightarrow a = L^T h$
 - $\rightarrow h = La$
- Iterative computation
 - using the power iteration method
 - $\rightarrow a_k = L^T L a_{k-1}$
 - $\rightarrow h_k = L L^T h_{k-1}$
 - $\rightarrow a_0 = h_0 = (1, 1, 1, \dots)$
 - normalization
 - $\rightarrow \sum_{i=1}^n a_i = 1$
 - $\rightarrow \sum_{i=1}^n h_i = 1$
 - ends after the 1-norms of the residual vectors are less than some thresholds (e.g. 5 iteration)
- Return top ranked pages as authorities and hubs.

HITS Algorithm (cont.)

- Convergence issues
 - HITS will always converge
 - can provide different hub and authority vectors
 - \rightarrow depending on the initialization
 - \rightarrow caused by the problem that $L^T L$ (respectively $L L^T$) is reducible
- Modification
 - When pages are relevant to the query, but they can be separated in the graph G
 - \rightarrow e.g. words with different meaning
 - Compute HITS on smaller communities
- Characteristics
 - ability to rank pages according to the query topic
 - \rightarrow more relevant hubs and authorities
 - query time execution
 - \rightarrow time consuming operation
 - does not have the anti-spam capability
 - \rightarrow a simple page with many links can easily become a hub
 - topic drift
 - \rightarrow expanded pages are not relevant

HITS Example



$$L = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}, L^T = \begin{bmatrix} 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 1 \\ 1 & 0 & 0 & 0 \end{bmatrix}, h_0 = \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}, a_0 = \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

$$a_1 = L^T h_0 = \begin{bmatrix} 1 \\ 1 \\ 3 \\ 1 \end{bmatrix}, h_1 = L a_0 = \begin{bmatrix} 3 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$