

GROUP EXERCISE 1

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Group ID: 8

Group Name: Expecto Patronum

1. Group Information:

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2. Questions:

i. Implement the following sorting algorithms using C/C++:

a. Insertion Sort

```
void InsertionSort(long* arr, long size)
{
    long key, i, j;
    for (i = 1; i < size; i++)
    {
        key = arr[i];
        j = i - 1;
        while ((j >= 0) && (arr[j] > key))
            arr[j + 1] = arr[j--];
        arr[j + 1] = key;
    }
}
```

b. Merge Sort

```
void Merge(long* arr, long left, long mid, long right)
{
    long *Temp = new long[right - left + 1];
    long i, j, k;
    for (i = 0, j = left, k = mid + 1; (j <= mid && k <=
right); i++)
    {
        if (arr[j] < arr[k])
            Temp[i] = arr[j++];
        else
```

```

        Temp[i] = arr[k++];
    }
    while (j <= mid)
        Temp[i++] = arr[j++];
    while (k <= right)
        Temp[i++] = arr[k++];
    for (i = 0; i <= right - left; i++)
        arr[i + left] = Temp[i];
    delete Temp;
}

void MergeSort(long* arr, long left, long right)
{
    long mid;
    if (left < right)
    {
        mid = (left + right) / 2;
        MergeSort(arr, left, mid);
        MergeSort(arr, mid + 1, right);
        Merge(arr, left, mid, right);
    }
}

```

c. Quick Sort

```

long Partition(long* arr, long left, long right)
{
    long pivot = (left + right) / 2;
    long i = left;
    long j = right;
    while (i < j)
    {
        while (arr[i] < arr[pivot])
            i++;
        while (arr[j] > arr[pivot])
            j--;
        if (i < j)
        {
            Swap(arr[i], arr[j]);
            i++;
            j--;
        }
    }
    return pivot;
}

void QuickSort(long* arr, long left, long right)
{
    if (left < right)
    {
        long p = Partition(arr, left, right);
        QuickSort(arr, left, p);
        QuickSort(arr, p + 1, right);
    }
}

```

```

    }
}

```

d. Radix Sort

```

void RadixSort(long* arr, long size)
{
    long i, m = arr[0], exp = 1;
    long *Temp = new long[size];
    for (i = 0; i < size; i++)
    {
        if (arr[i] > m)
            m = arr[i];
    }
    while (m / exp > 0)
    {
        long bucket[10] = { 0 };
        for (i = 0; i < size; i++)
            bucket[arr[i] / exp % 10]++;
        for (i = 1; i < 10; i++)
            bucket[i] += bucket[i - 1];
        for (i = size - 1; i >= 0; i--)
            Temp[--bucket[arr[i] / exp % 10]] = arr[i];
        for (i = 0; i < size; i++)
            arr[i] = Temp[i];
        exp *= 10;
    }
    delete Temp;
}

```

e. Counting Sort

```

void countSort(char arr[])
{
    int count[RANGE + 1], i; //RANGE: constant = 1e7, we assume!
    memset(count, 0, sizeof(count));
    for(i = 0; arr[i]; ++i)
        ++count[arr[i]];
    for (i = 1; i <= RANGE; ++i)
        count[i] += count[i-1];
    for (i = 0; arr[i]; ++i)
    {
        output[count[arr[i]]-1] = arr[i];
        --count[arr[i]];
    }
    for (i = 0; arr[i]; ++i)
        arr[i] = output[i];
}

```

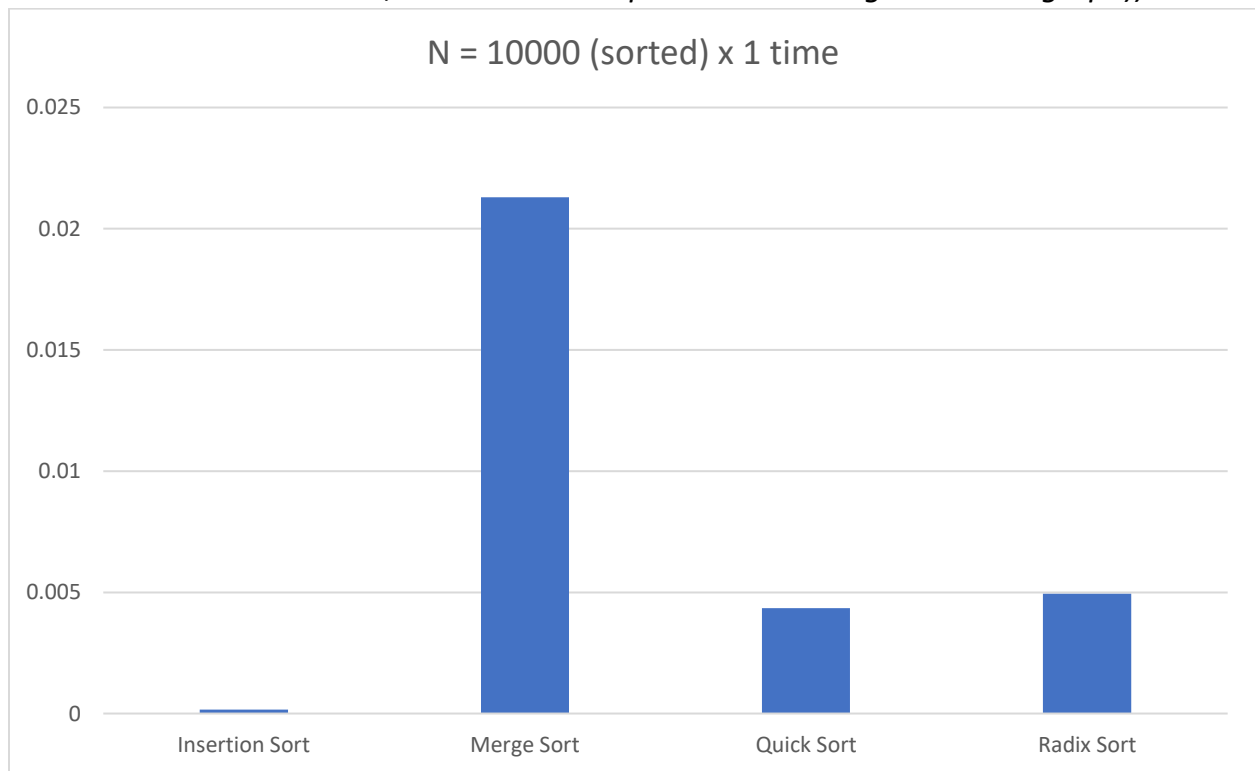
ii. The running time of these Sorting Algorithms:

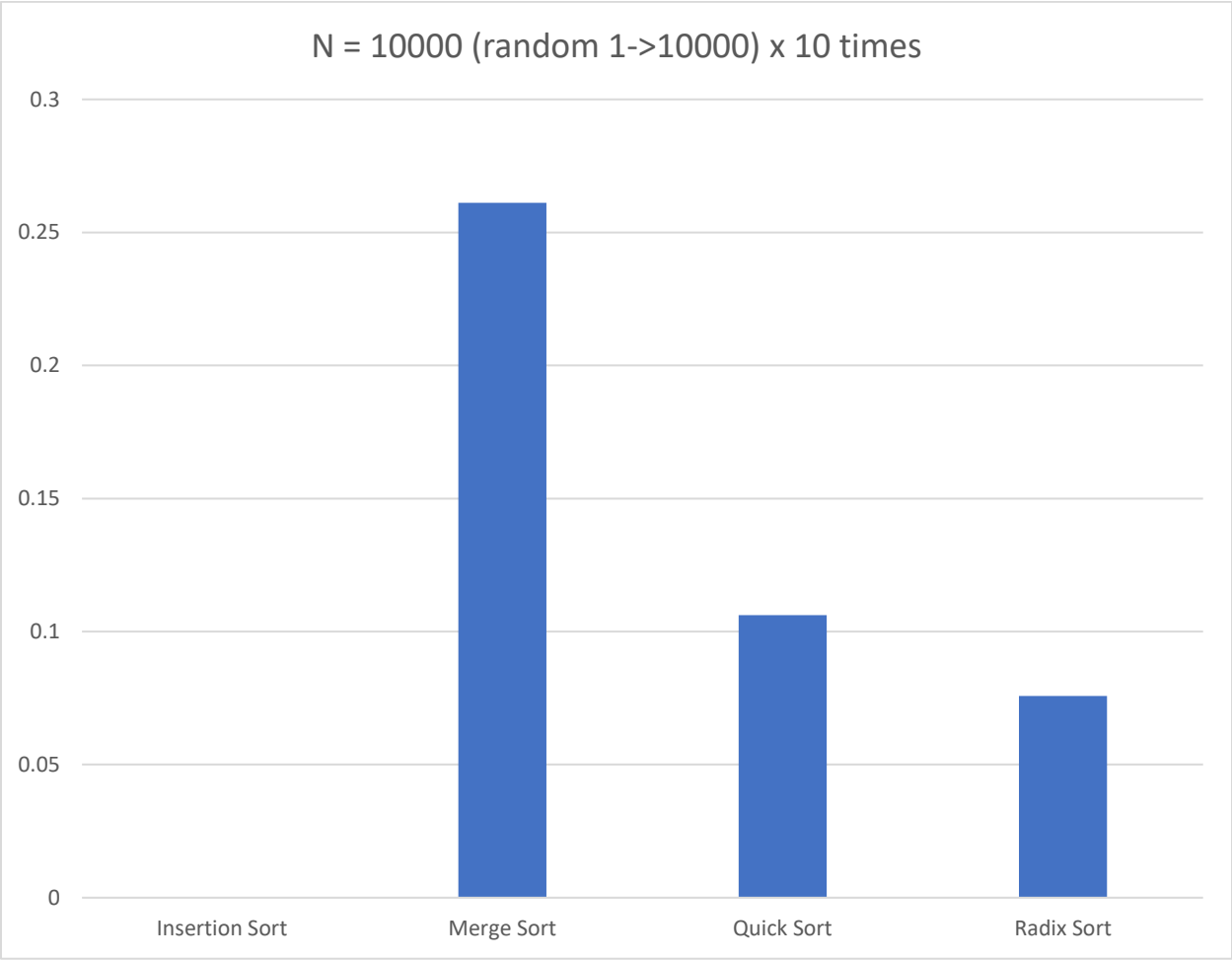
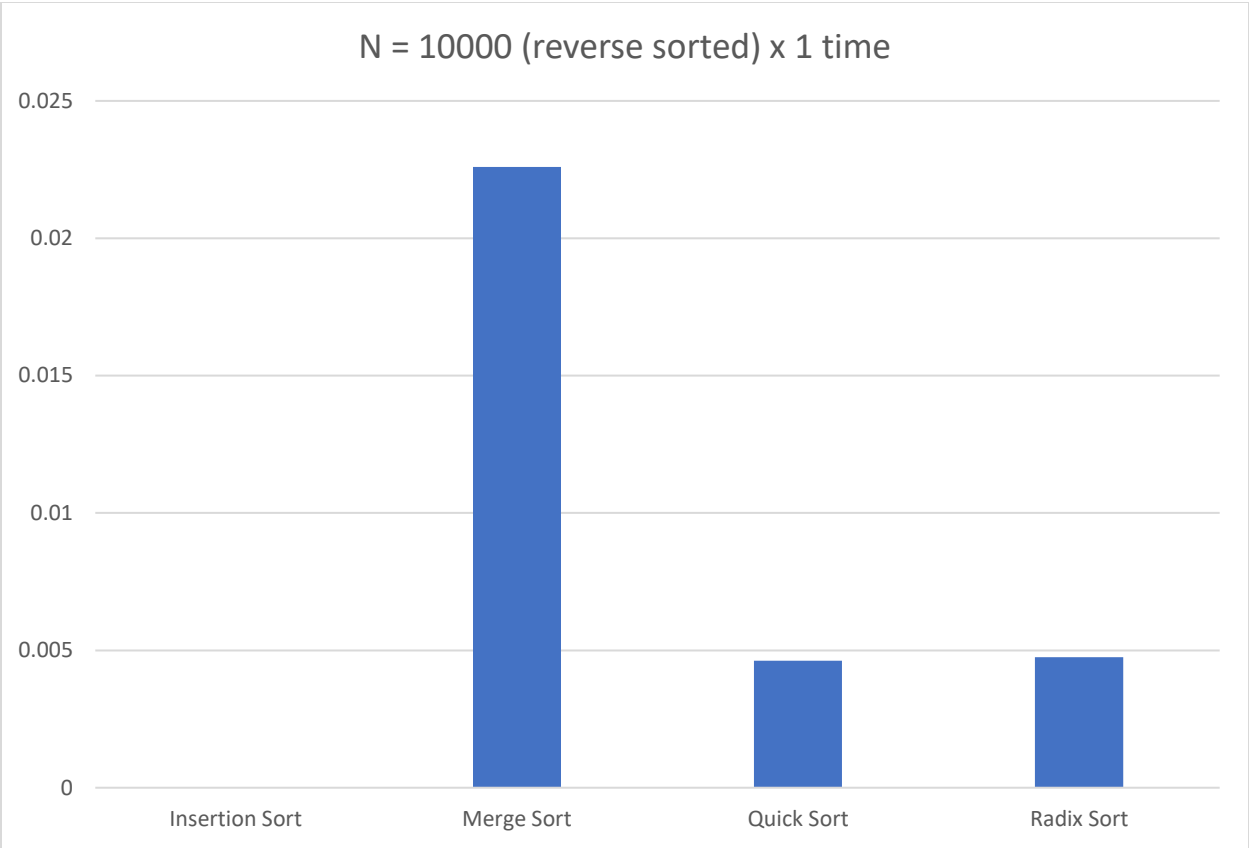
(We took average-times of 10 times we ran the algorithm)

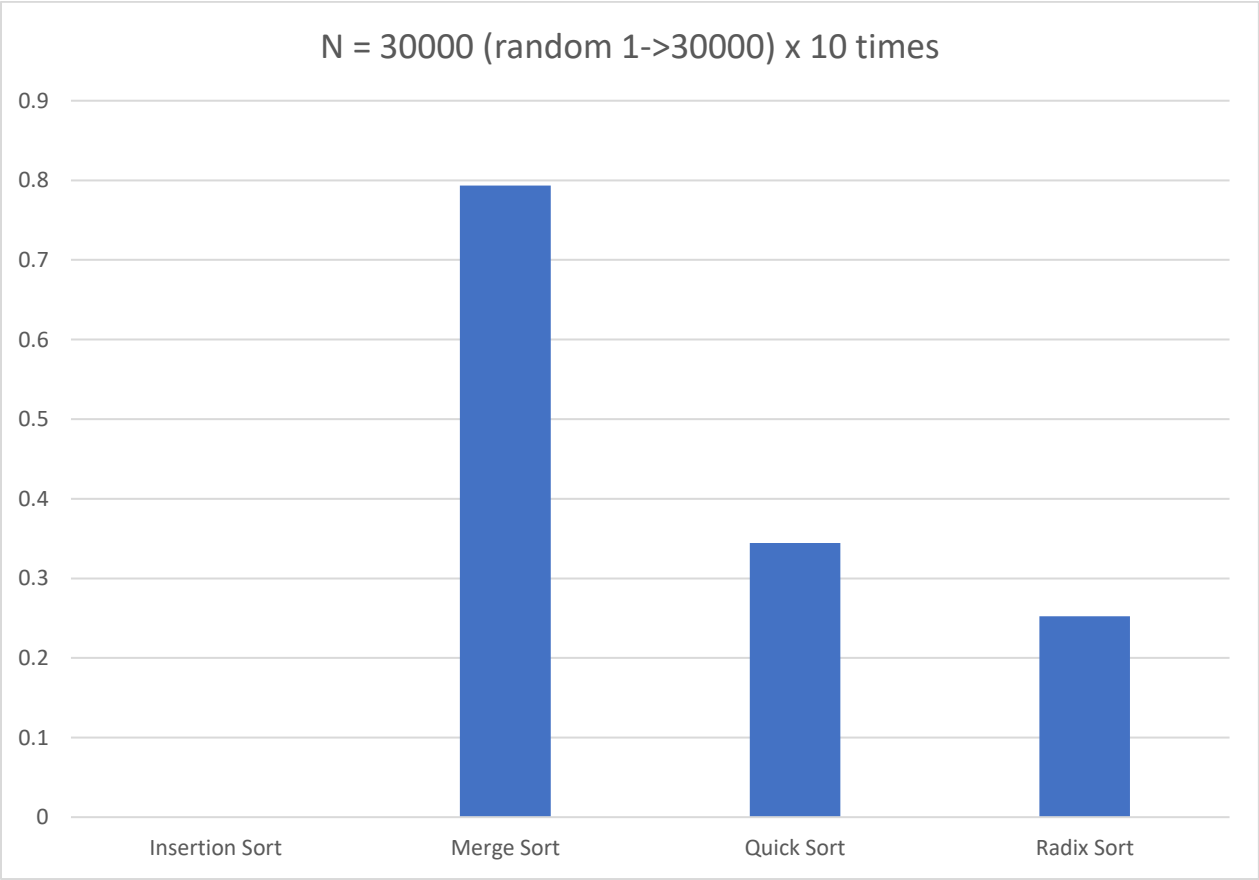
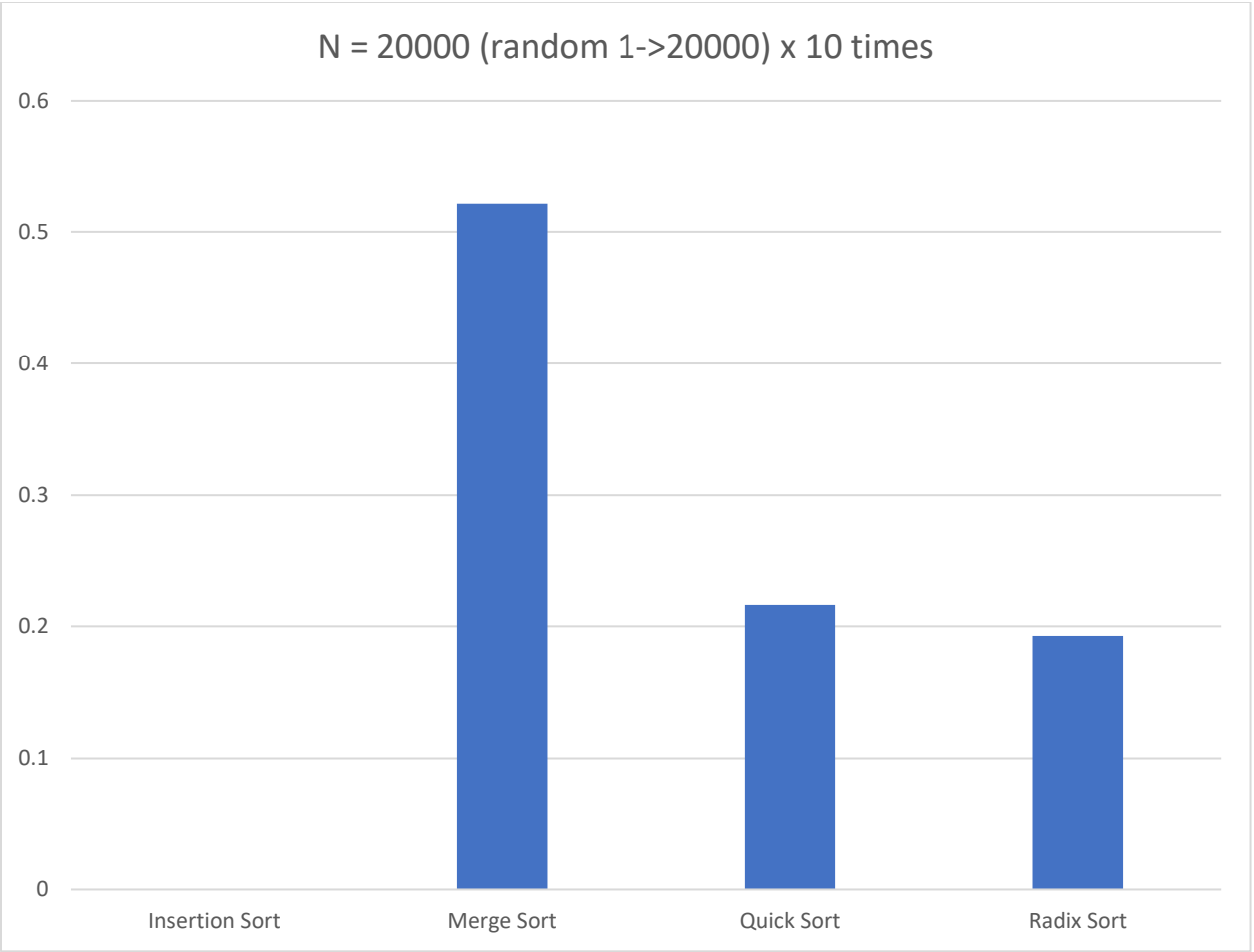
	Insertion Sort	Merge Sort	Quick Sort	Radix Sort
N = 10000 (sorted) x 1 time	0.000161s	0.021294s	0.004345s	0.004947 s
N = 10000 (reverse sorted) x 1 time	0.514267 s	0.022597 s	0.004622 s	0.004753 s
N = 10000 (random 1->10000) x 10 times	2.629274 s	0.261091 s	0.106150 s	0.075821 s
N = 20000 (random 1->20000) x 10 times	10.541194 s	0.521305 s	0.216198 s	0.192793 s
N = 30000 (random 1->30000) x 10 times	24.450407 s	0.793444 s	0.344260 s	0.252208 s
N = 5000000 (random 1->5000000) x 10 times	(too slow)	145.667489 s	23.222543 s	39.818839 s

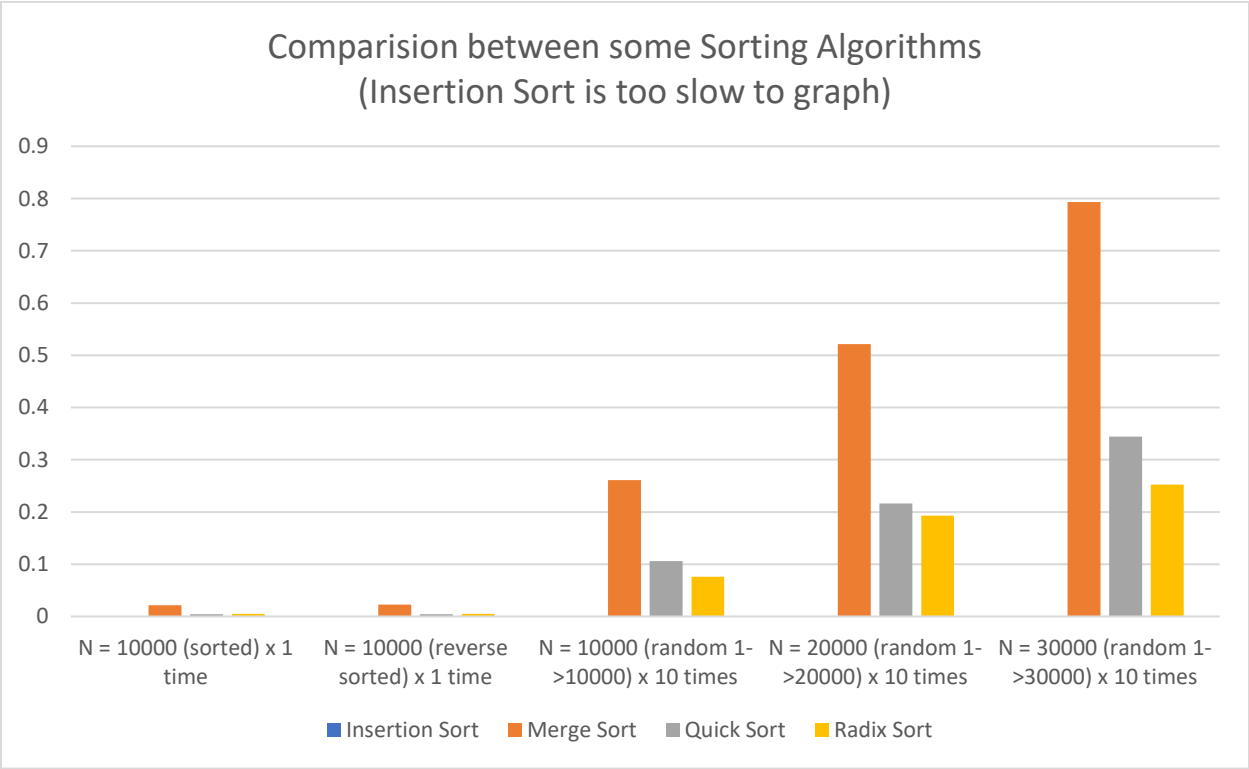
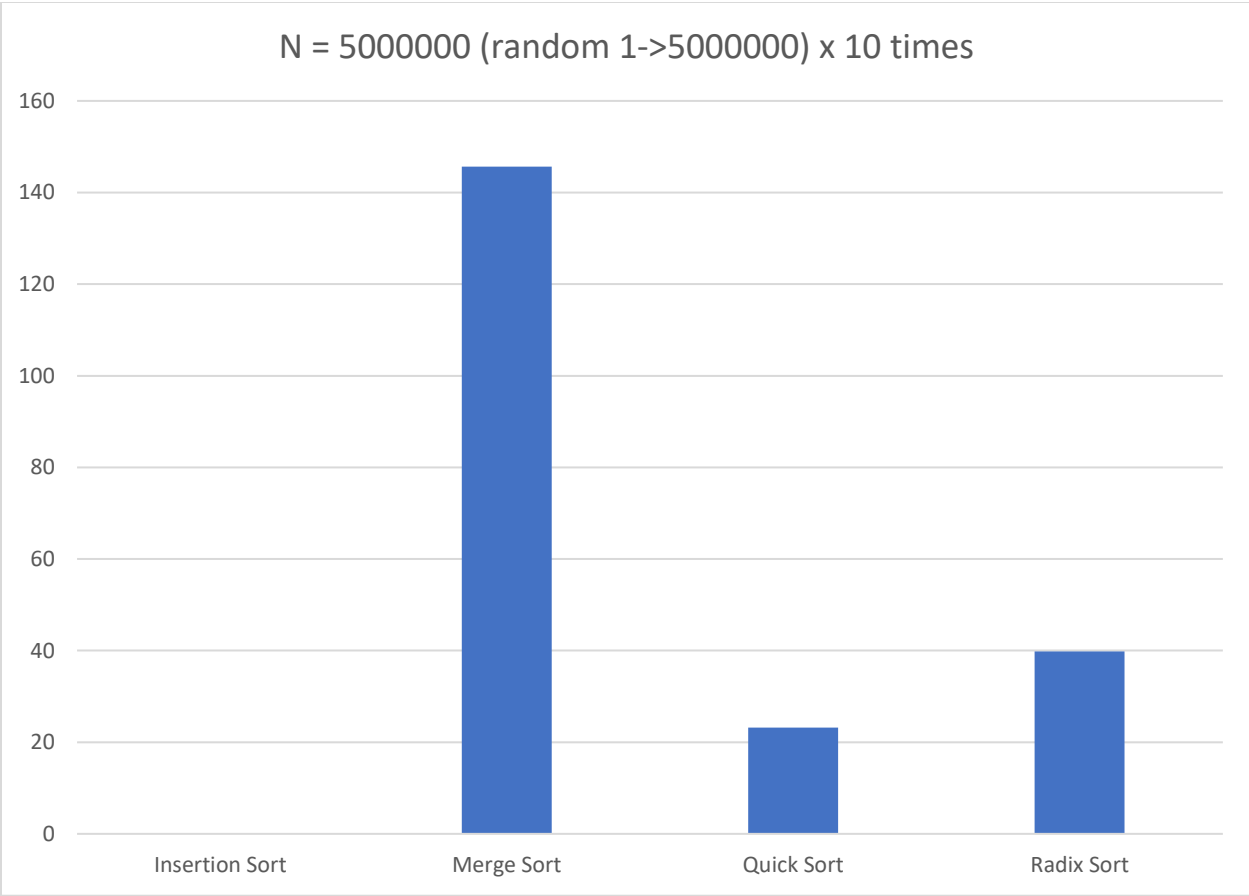
Graph to describe more clearly:

(With $n \geq 30000$, we don't graph the running time of Insertion Sort (because it costs too much time, so we can't compare the other algorithms on graph)).









Comments:

- With small test ($N = 10000$), Insertion Sort runs faster than others, but when we increase the size of array (≥ 10000), it runs too slowly.
- The running time of Quick Sort and Merge Sort seems to be equal, because both takes (average) $O(n \log n)$.
- In some random case, Quick Sort runs better than Merge Sort. I think it depends on test case (we used random array to calculate the running time).
- In most of cases, the Radix Sort runs fastest, because it works on Linear Time.

In conclusion:

- If the size of array is small (less than 10^4), we can use Insertion Sort for simply coding and implementation that without affecting the running time.
- **If the limitation of each element in an array is less than 10^7** , we can use Radix Sort or Counting Sort to sort an array **in Linear Time** (without comparison).
- In the other cases, we know that Heap Sort is the most stable algorithms to work in **$O(n \log n)$** time. But Quick Sort is also the good choice to sort array with size **not bigger than 10^6 elements**. *More than that, we can improve Quick Sort by select pivot randomly (in the test, we chose the mid-element to be a pivot) to avoid the bad partitions..* Merge Sort is also good, but it seems **to be difficult to implement right**.