

# Mitjos lab2 实验报告

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## 一、实验内容

**Exercise 1.** In the file `kern/pmap.c`, you must implement code for the following functions (probably in the order given).

```
boot_alloc()
mem_init() (only up to the call to check_page_free_list(1))
page_init()
page_alloc()
page_free()
```

`check_page_free_list()` and `check_page_alloc()` test your physical page allocator. You should boot JOS and see whether `check_page_alloc()` reports success. Fix your code so that it passes. You may find it helpful to add your own `assert()`s to verify that your assumptions are correct.

```
// LAB 2: Your code here.
result = nextfree;
nextfree = ROUNDUP((char *) (nextfree+n), PGSIZE);
return result;
```

修改 `boot_alloc` 函数，`boot_alloc(uint32_t n)` 主要是申请 `n` 个字节的地址空间，返回申请空间的首地址。由于未初始化的全局变量和静态变量会被自动初始化为 0，系统第一次调用 `boot_alloc()` 这个函数的时候，`nextfree` 会指向第一个空闲页的首地址。接下来，根据输入的 `n`，来分配地址。如果 `n=0`，则返回 `nextfree`，否则分配 `n` 字节的地址，返回分配地址的首地址。在整个过程中，需要 4K 对齐。

```
pages = boot_alloc(npages * sizeof(struct PageInfo));
memset(pages, 0, npages*sizeof(struct PageInfo));
```

修改 `mem_init()` 函数，为 `pages` 申请 `npages` 的页面，存放这些结构体，并且用 `memset` 来初始化

```
size_t i;
for (i = 0; i < npages; i++) {
    if (i == 0)
    {
        pages[i].pp_ref = 1;
        pages[i].pp_link = NULL;
    }
    else if (i >= 1 && i < npages_basemem)
    {
        pages[i].pp_ref = 0;
        pages[i].pp_link = page_free_list;
        page_free_list = &pages[i];
    }
    else if (i >= IOPHYSMEM/PGSIZE && i < EXTPHYSMEM/PGSIZE )
    {
        pages[i].pp_ref = 1;
        pages[i].pp_link = NULL;
    }
    else if ( i >= EXTPHYSMEM / PGSIZE && i < ( (int)(boot_alloc(0))
- KERNBASE)/PGSIZE)
    {
        pages[i].pp_ref = 1;
        pages[i].pp_link = NULL;
    }
}
```

```

        pages[i].pp_ref = 1;
        pages[i].pp_link = NULL;
    }
    else
    {
        pages[i].pp_ref = 0;
        pages[i].pp_link = page_free_list;
        page_free_list = &pages[i];
    }
}

```

修改 page\_init 函数，系统初始化是分配物理内存

```

if(page_free_list == NULL)
    return NULL;

struct PageInfo* page = page_free_list;
page_free_list = page->pp_link;
page->pp_link = 0;
if(alloc_flags & ALLOC_ZERO)
    memset(page2kva(page), 0, PGSIZE);
return page;

if(pp->pp_link != 0 || pp->pp_ref != 0)
    panic("page_free is not right");
pp->pp_link = page_free_list;
page_free_list = pp;
return;

```

修改 page\_alloc 函数和 page\_free 函数，申请和释放页面

**Exercise 4.** In the file kern/pmap.c, you must implement code for the following functions.

```

pgdir_walk()
boot_map_region()
page_lookup()
page_remove()
page_insert()

```

check\_page(), called from mem\_init(), tests your page table management routines. You should make sure it reports success before proceeding.

```

int pdeIndex = (unsigned int)va >>22;
if(pgdir[pdeIndex] == 0 && create == 0)
    return NULL;
if(pgdir[pdeIndex] == 0){
    struct PageInfo* page = page_alloc(1);
    if(page == NULL)
        return NULL;
    page->pp_ref++;
    pte_t pgAddress = page2pa(page);
    pgAddress |= PTE_U;
    pgAddress |= PTE_P;
    pgAddress |= PTE_W;
    pgdir[pdeIndex] = pgAddress;
}
pte_t pgAdd = pgdir[pdeIndex];
pgAdd = pgAdd>>12<<12;
int pteIndex =(pte_t)va >>12 & 0x3ff;
pte_t * pte =(pte_t*) pgAdd + pteIndex;
return KADDR( (pte_t) pte );

```

pgdir\_walk(): 用于返回 va 对应的二级页表的地址(PTE)。给定一个虚拟地址 va 和 pgdir(page director table 的首地址), 返回 va 所对应的 pte(page table entry)。当 va 对应的二级页表存在时, 只直接给出 PTE 的地址, 当 va 对应的二级页表还没有被创建时, 需要手动申请页面并创建, 最后返回 PTE 的地址的时, 返回 PTE 地址对应的虚拟地址。

```

while(size)
{
    pte_t* pte = pgdir_walk(pgdir, (void* )va, 1);
    if(pte == NULL)
        return;
    *pte= pa |perm|PTE_P;
    size -= PGSIZE;
    pa += PGSIZE;
    va += PGSIZE;
}

```

boot\_map\_region: [va, va+size)映射到[pa, pa+size]

```

pte_t* pte = pgdir_walk(pgdir, va, 0);
if(pte == NULL)
    return NULL;
pte_t pa = *pte>>12<<12;
if(pte_store != 0)
    *pte_store = pte ;
return pa2page(pa);

```

page\_lookup: 返回虚拟地址 va 对应的物理地址的页面 page

```

pte_t* pte;
struct PageInfo* page = page_lookup(pgdir, va, &pte);
if(page == 0)
    return;
*pte = 0;
page->pp_ref--;
if(page->pp_ref == 0)
    page_free(page);
tlb_invalidate(pgdir, va);

```

page\_remove: 将 va 和其对应的页面取消映射

```

pte_t* pte = pgdir_walk(pgdir, va, 1);
if(pte == NULL)
    return -E_NO_MEM;
if( (pte[0] & ~0xfff) == page2pa(pp))
    pp->pp_ref--;
else if(*pte != 0)
    page_remove(pgdir, va);

*pte = (page2pa(pp) & ~0xfff) | perm | PTE_P;
pp->pp_ref++;
return 0;

```

page\_insert():把 va 映射到指定的物理页表 page

**Exercise 5.** Fill in the missing code in `mem_init()` after the call to `check_page()`.

Your code should now pass the `check_kern_pgdir()` and `check_page_installed_pgdir()` checks.

```

int perm = PTE_U | PTE_P;
int i=0;
n = ROUNDUP(npages*sizeof(struct PageInfo), PGSIZE);
for(i=0; i<n; i= i+PGSIZE)
    page_insert(kern_pgdir, pa2page(PADDR(pages) + i), (void *)
(UPAGES +i), perm);

perm = 0;
perm = PTE_P | PTE_W;
boot_map_region(kern_pgdir, KSTACKTOP-KSTKSIZE, ROUNDUP(KSTKSIZE,
PGSIZE), PADDR(bootstack), perm);

```

```
int size = ~0;
size = size - KERNBASE + 1;
size = ROUNDUP(size, PGSIZE);
perm = 0;
perm = PTE_P | PTE_W;
boot_map_region(kern_pgdir, KERNBASE, size, 0, perm );
```

计算出 `pages` 结构体的大小，通过 `page_insert()` 进行映射。利用 `boot_map_region()` 实现地址之间的静态映射。

## 二、实验结果

```
make[1]: Leaving directory '/home/lidongwen/Documents/git/os/lab'
running JOS: (1.0s)
  Physical page allocator: OK
  Page management: OK
  Kernel page directory: OK
  Page management 2: OK
Score: 70/70
```