

# LAB EXERCISE #1

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Hochschule Landshut  
**Concepts of Modern Programming Languages**

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DUE: DURING LAB SESSION #1

## **Problem 1: Chapter 4 Mitchell, 4.12 Static Assignment Languages**

If you don't have the book, the question is also in Moodle. Read section 4.4.1 for the distinction between "imperative" and "declarative" (also in Moodle).

## **Problem 2: Pass-by-name has some weird semantics.**

Read about Jensen's device (Mitchell page 96, in Moodle) and implement it in C's Macro processor or another language that offers call by name.

Hint: in Gnu C a compound statement can be used as an expression, c.f.

<https://gcc.gnu.org/onlinedocs/gcc-3.1/gcc/Statement-Exprs.html>

What gets printed?

## **Problem 3: Denotational Semantics**

Read the explanation about denotational semantics in chapter 4.3 in Mitchell (pages 67-76, available in Moodle). Then do exercise 4.8, which is printed below for your convenience:

The Mitchell text describes a denotational semantics for the simple imperative language given by the grammar

$P ::= x := e \mid P_1; P_2 \mid \text{if } e \text{ then } P_1 \text{ else } P_2 \mid \text{while } e \text{ do } P.$

Each program denotes a function from *states* to *states*, in which a *state* is a function

from *variables* to *values*.

1. (a) Calculate the meaning  $C[x := 1; x := x + 1;](s_0)$  in approximately the same detail as that of the examples given in the text, where  $s_0 = \lambda v \in \text{variables}. 0$ , giving every variable the value 0.
2. (b) Denotational semantics is sometimes used to justify ways of reasoning about programs. Write a few sentences, referring to your calculation in part (a), explaining why

$C[x := 1; x := x + 1;](s) = C[x := 2;](s)$  for every state  $s$ .

