

# Local Metropolis Dynamics for Signed Multiline Queues at $q = 1$ and a Strongly Lumped Bottom-Word Chain

## Abstract

Ben Dali–Williams (BDW) give a signed multiline queue (SMLQ) model whose total weight over SMLQs with fixed bottom word  $\mu$  equals the interpolation ASEP polynomial  $F_\mu^*(\mathbf{x}; 1, t)$ , and whose orbit-sum over  $\mu \in S_n(\lambda)$  yields the symmetric interpolation Macdonald polynomial  $P_\lambda^*(\mathbf{x}; 1, t)$  [1]. Working in the “restricted” setting (distinct parts, a unique 0, and no part 1), and assuming a probabilistic regime in which all BDW SMLQ weights are nonnegative, we construct: (i) a concrete, finite-state Markov chain on the *full* SMLQ state space whose stationary distribution is the Gibbs measure  $\Pi(\mathbf{Q}) \propto \text{wt}(\mathbf{Q})$ , built from random-scan Metropolis block updates on bounded windows of adjacent rows; crucially, a bottom block update changes the bottom row word using only BDW local admissibility and local weight ratios; and (ii) an (optional) fiber-refresh composition that produces a *strongly lumpable* SMLQ chain whose projection to bottom words is a genuine Markov chain on  $\text{supp}(\pi)$  with explicit stochastic kernel  $\tilde{K}$ , stationary distribution  $\pi(\mu) \propto F_\mu^*(\mathbf{x}; 1, t)$ , and no transition probabilities defined from global polynomial evaluations.

## Introduction

Fix  $n \geq 2$  and a partition  $\lambda = (\lambda_1 > \cdots > \lambda_n \geq 0)$  with *distinct parts, exactly one part equal to 0, and no part equal to 1*. We call this the *restricted case*. Let  $S_n(\lambda)$  denote the orbit of  $\lambda$  under coordinate permutations.

BDW define *signed multiline queues* (SMLQs) of type  $\mu \in S_n(\lambda)$  and assign each SMLQ  $\mathbf{Q}$  a (generally signed) weight  $\text{wt}(\mathbf{Q})$  [1, Defs. 1.5–1.6, 1.11]. Let  $MLQ^\pm(\mu)$  be the finite set of SMLQs with bottom word  $\mu$ . BDW define the weight-generating function

$$F_\mu^*(\mathbf{x}; q, t) := \sum_{\mathbf{Q} \in MLQ^\pm(\mu)} \text{wt}(\mathbf{Q}) \quad [1, \text{Def. 1.14}],$$

and prove that  $F_\mu^*(\mathbf{x}; q, t)$  equals the interpolation ASEP polynomial  $f_\mu^*(\mathbf{x}; q, t)$  [1, Thm. 1.15]. We specialize to  $q = 1$  and consider the orbit sum

$$P_\lambda^*(\mathbf{x}; 1, t) := \sum_{\mu \in S_n(\lambda)} F_\mu^*(\mathbf{x}; 1, t) \quad (\text{equivalently } Z_\lambda^*(\mathbf{x}; 1, t) \text{ in [1]}),$$

which is the symmetric interpolation Macdonald polynomial at  $q = 1$  [1, Thm. 1.15]. The target bottom-word distribution is

$$\pi(\mu) = \frac{F_\mu^*(\mathbf{x}; 1, t)}{P_\lambda^*(\mathbf{x}; 1, t)}.$$

The goal is to construct a *genuinely nontrivial* Markov chain on the full SMLQ state space with stationary law proportional to  $\text{wt}$ , such that (at stationarity) its bottom-word marginal is  $\pi$ , and then (optionally) to produce a *Markov* chain on bottom words via a rigorous strong lumpability argument. The construction must never define transitions using the polynomial values  $F_\mu^*$  or  $P_\lambda^*$ ; only BDW *local* primitives (row/layer admissibility and local weights) may be used. Our local-update design follows the general AMW “multiline diagram dynamics” philosophy [2], but we do *not* assume any AMW-style collapse/row-exchange identity.

# 1 Concrete, reliable construction on SMLQs

## 1.1 State space, bottom map, and a positivity hypothesis

Let  $L := \lambda_1$  be the largest part. BDW's *enhanced ball systems* have rows labeled

$$1, 1', 2, 2', \dots, L, L'$$

from bottom to top and columns  $1, \dots, n$  on a cylinder [1, Def. 1.5]. Regular rows  $r$  carry labels in  $\{0, 1, \dots, L\}$ ; signed rows  $r'$  carry labels in  $\{0, \pm 1, \dots, \pm L\}$ . In the restricted case, 1 never appears in any row word because  $m_1 = 0$ .

**Definition 1.1** (SMLQ state space and bottom map). Let  $\Omega_\lambda$  be the finite set of BDW signed multiline queues of content  $\lambda$ , i.e. the union

$$\Omega_\lambda = \bigcup_{\mu \in S_n(\lambda)} MLQ^\pm(\mu) \quad [1, \text{Def. 1.6}].$$

For  $Q \in \Omega_\lambda$ , write  $\Phi(Q) \in S_n(\lambda)$  for its *bottom word* (the label configuration on Row 1), and let  $\Omega_\lambda(\mu) = \Phi^{-1}(\mu)$  be the fiber over  $\mu$ .

**Definition 1.2** (BDW weight at  $q = 1$ ). Fix parameters  $(\mathbf{x}, t)$  with  $\mathbf{x} = (x_1, \dots, x_n)$ . For  $Q \in \Omega_\lambda$ , define  $\text{wt}(Q)$  to be the BDW weight specialized at  $q = 1$  [1, Def. 1.11]:

$$\text{wt}(Q) = \text{wt}_{\text{ball}}(Q) \text{wt}_{\text{pair}}(Q).$$

Here  $\text{wt}_{\text{ball}}$  is the product over all signed rows  $r'$  of the shifted ball-weights

$$\text{wt}_{\text{ball}}(r') = \left( \prod_{i: \alpha_i > 0} x_i \right) \left( \prod_{i: \alpha_i < 0} (-t^{n-1}) \right) \quad (\text{since } q^{r-1} = 1 \text{ at } q = 1),$$

where  $\alpha \in \mathbb{Z}^n$  is the signed composition on row  $r'$  [1, (5)–(6)]. The pairing weight  $\text{wt}_{\text{pair}}$  is the product of local factors over all nontrivial pairings in all classic layers and signed layers [1, (4),(7)], with the specialization  $q = 1$ .

The BDW model is signed in general. The problem statement allows us to work in a regime where the weights become probabilistic.

**Assumption 1.3** (Probabilistic regime). We work at  $q = 1$  and assume there is a parameter regime (e.g.  $t \in (0, 1)$ ) together with additional inequalities on  $\mathbf{x}$  such that:

- (i)  $\text{wt}(Q) \geq 0$  for all  $Q \in \Omega_\lambda$ ;
- (ii)  $Z := \sum_{Q \in \Omega_\lambda} \text{wt}(Q) > 0$ .

**Definition 1.4** (Positive support and Gibbs measure). Under Assumption 1.3, define

$$\Omega_\lambda^+ := \{Q \in \Omega_\lambda : \text{wt}(Q) > 0\}, \quad \Pi(Q) := \frac{\text{wt}(Q)}{\sum_{R \in \Omega_\lambda} \text{wt}(R)} \quad (Q \in \Omega_\lambda^+),$$

and  $\Pi(Q) = 0$  for  $Q \notin \Omega_\lambda^+$ . Define the supported bottom words

$$S_\lambda^+ := \{\mu \in S_n(\lambda) : \Omega_\lambda^+(\mu) := \Omega_\lambda^+ \cap \Omega_\lambda(\mu) \neq \emptyset\}.$$

**Remark 1.5** (What is assumed vs. proved about positivity). Assumption 1.3 is an explicit hypothesis: we do not claim here that it holds for all  $(\mathbf{x}, t)$ , only that we restrict to regimes in which the BDW weight is nonnegative and not identically zero. (In particular examples, e.g. [1, Ex. 1.16], some configurations have negative weights unless parameters and/or allowed sign patterns are restricted.)

## 1.2 A local/block base kernel $U_{\text{base}}$ that changes bottom words

We construct a Metropolis–Hastings chain on  $\Omega_\lambda^+$  built from random-scan block updates on small windows of adjacent rows and their incident layers. The key point is that one update block *includes the bottom row as a variable*, hence can change  $\Phi$ .

**Blocks.** We use the row/layer adjacency in BDW’s enhanced ball system: regular row  $r$  is adjacent to signed row  $r'$  (via a signed layer) and to signed row  $(r-1)'$  for  $r \geq 2$  (via a classic layer); signed row  $r'$  is adjacent to regular row  $r$  (signed layer) and regular row  $r+1$  for  $r \leq L-1$  (classic layer). Define the following finite block index set:

$$\mathcal{B} := \{\text{bot}\} \cup \{\text{reg}(r) : 2 \leq r \leq L\} \cup \{\text{sign}(r) : 2 \leq r \leq L-1\} \cup \{\text{top}\}.$$

Each block  $b \in \mathcal{B}$  specifies a *bounded* collection of adjacent rows and the adjacent layers incident to an “interior” row:

- $b = \text{bot}$ : the *bottom window* consisting of Rows  $2, 1', 1$  together with the classic layer between  $2$  and  $1'$  and the signed layer between  $1'$  and  $1$ . In this block, Rows  $1', 1$  and both layers are updated, while Row  $2$  is held fixed as boundary data. This block *can change the bottom word*  $\Phi$ .
- $b = \text{reg}(r)$  for  $2 \leq r \leq L$ : the *regular-row window* consisting of Rows  $r', r, (r-1)'$  together with the signed layer between  $r'$  and  $r$  and the classic layer between  $r$  and  $(r-1)'$ . Here the outer signed rows  $r'$  and  $(r-1)'$  are held fixed and the middle regular Row  $r$  and both incident layers are updated.
- $b = \text{sign}(r)$  for  $2 \leq r \leq L-1$ : the *signed-row window* consisting of Rows  $(r+1), r', r$  together with the classic layer between  $(r+1)$  and  $r'$  and the signed layer between  $r'$  and  $r$ . Here the outer regular rows  $(r+1)$  and  $r$  are held fixed and the middle signed Row  $r'$  and both incident layers are updated.
- $b = \text{top}$ : the *top window* consisting of Rows  $L', L$  together with their signed layer. Here Row  $L$  is held fixed and Row  $L'$  and the signed layer are updated.

**Local admissible fillings.** For a full SMLQ  $\mathbf{Q} \in \Omega_\lambda$ , each block  $b$  comes with a finite set of alternative local configurations inside the rows/layers of  $b$  that are BDW-admissible and consistent with the fixed boundary rows outside the block.

**Definition 1.6** (Block fiber  $\mathcal{S}_b(\mathbf{Q})$ ). Fix  $b \in \mathcal{B}$  and  $\mathbf{Q} \in \Omega_\lambda$ . Let  $\mathcal{S}_b(\mathbf{Q})$  be the set of all  $\mathbf{Q}' \in \Omega_\lambda$  such that:

- $\mathbf{Q}'$  agrees with  $\mathbf{Q}$  on every row word and every layer *not* belonging to the window  $b$ ;
- the restriction of  $\mathbf{Q}'$  to the rows and layers in  $b$  is BDW-admissible (i.e. satisfies the forbidden-configuration rules for classic and signed layers [1, Defs. 1.4, 1.6]) and has the correct row-contents (each row word in the window is a signed/unsigned permutation of the prescribed multiset  $\lambda^{(r)}$  for that row [1, Def. 1.5]).

**Lemma 1.7** (Finite local catalogue:  $\mathcal{S}_b(\mathbf{Q})$  is finite and nonempty). *For every  $\mathbf{Q} \in \Omega_\lambda$  and every block  $b \in \mathcal{B}$ , the set  $\mathcal{S}_b(\mathbf{Q})$  is finite and contains  $\mathbf{Q}$ .*

*Proof.* Nonemptiness holds because  $\mathbf{Q} \in \mathcal{S}_b(\mathbf{Q})$  by construction. Finiteness holds because the rows inside  $b$  have length  $n$  and each such row must be a permutation (signed or unsigned) of a fixed finite multiset  $\lambda^{(r)}$  [1, Def. 1.5]; hence there are finitely many row-word possibilities. For each such choice, the number of possible BDW-admissible pairing patterns on finitely many adjacent layers is finite. Therefore  $\mathcal{S}_b(\mathbf{Q})$  is finite.  $\square$

**The proposal and acceptance rule.** Fix positive block-selection probabilities  $(p_b)_{b \in \mathcal{B}}$  with  $\sum_{b \in \mathcal{B}} p_b = 1$  and  $p_b > 0$  for each  $b$ .

**Definition 1.8** (The base kernel  $U_{\text{base}}$ ). Define  $U_{\text{base}}$  on  $\Omega_\lambda^+$  by the following random-scan Metropolis block update:

**Step 1:** Choose a block  $b \in \mathcal{B}$  with probability  $p_b$ .

**Step 2:** Propose  $Q^{\text{prop}}$  uniformly from  $\mathcal{S}_b(Q)$ .

**Step 3:** If  $\text{wt}(Q^{\text{prop}}) = 0$  (equivalently  $Q^{\text{prop}} \notin \Omega_\lambda^+$ ), *reject* (i.e. set  $Q^+ = Q$ ). Otherwise accept with Metropolis probability

$$\alpha(Q \rightarrow Q^{\text{prop}}) = \min\left\{1, \frac{\text{wt}(Q^{\text{prop}})}{\text{wt}(Q)}\right\},$$

and set  $Q^+ = Q^{\text{prop}}$  if accepted,  $Q^+ = Q$  if rejected.

**Remark 1.9** (Local computability of the acceptance ratio). By BDW's definition,  $\text{wt}(Q)$  factors over signed rows (ball weights) and over layers (pairing weights) [1, Def. 1.11]. A block update changes only the finitely many local factors supported inside the chosen window. Hence the ratio  $\text{wt}(Q^{\text{prop}})/\text{wt}(Q)$  can be computed using only local information inside the chosen window and its fixed boundary rows; no global partition function appears.

### 1.3 Stationarity for $U_{\text{base}}$

**Theorem 1.10** (Reversibility and stationarity of  $U_{\text{base}}$ ). *Under Assumption 1.3, the kernel  $U_{\text{base}}$  is reversible with respect to  $\Pi$  on  $\Omega_\lambda^+$ , and therefore  $\Pi$  is stationary for  $U_{\text{base}}$ .*

*Proof.* Fix  $Q, Q' \in \Omega_\lambda^+$ . If  $Q'$  is not reachable from  $Q$  in one block update, then both one-step probabilities are zero and detailed balance is trivial.

Otherwise, there exists a unique block  $b$  such that  $Q'$  differs from  $Q$  only inside  $b$ ; in particular  $Q' \in \mathcal{S}_b(Q)$ . Because the boundary outside  $b$  is unchanged in either direction, we also have  $Q \in \mathcal{S}_b(Q')$  and  $|\mathcal{S}_b(Q)| = |\mathcal{S}_b(Q')|$ . Hence the proposal probabilities satisfy

$$q(Q \rightarrow Q') = \frac{p_b}{|\mathcal{S}_b(Q)|} = \frac{p_b}{|\mathcal{S}_b(Q')|} = q(Q' \rightarrow Q).$$

Metropolis acceptance then gives the standard identity

$$\Pi(Q) q(Q \rightarrow Q') \alpha(Q \rightarrow Q') = \Pi(Q') q(Q' \rightarrow Q) \alpha(Q' \rightarrow Q),$$

because  $\Pi(Q) \propto \text{wt}(Q)$  and  $\alpha(Q \rightarrow Q') = \min\{1, \text{wt}(Q')/\text{wt}(Q)\}$ . Thus detailed balance holds, implying stationarity.  $\square$

**Proposition 1.11** (Aperiodicity via an explicit self-loop). *For every  $Q \in \Omega_\lambda^+$ ,*

$$U_{\text{base}}(Q \rightarrow Q) \geq \sum_{b \in \mathcal{B}} p_b \cdot \frac{1}{|\mathcal{S}_b(Q)|} > 0.$$

*In particular,  $U_{\text{base}}$  is aperiodic on each of its communicating classes.*

*Proof.* For each block  $b$ , Lemma 1.7 implies  $Q \in \mathcal{S}_b(Q)$ , so with probability  $p_b \cdot (1/|\mathcal{S}_b(Q)|)$  the proposal equals the current block configuration. The Metropolis acceptance is then 1. Summing over  $b$  yields the bound.  $\square$

## 1.4 Genuine bottom-word motion from local moves

The following lemma addresses the key critique: the base dynamics must *move bottom words by local BDW primitives*, not only via a global independence move.

**Definition 1.12** (A “diagonal” positive SMLQ with prescribed bottom word). Fix  $\mu \in S_n(\lambda)$ . Define an SMLQ  $Q^{\text{diag}}(\mu)$  by:

- In Row 1 (bottom regular row), place the word  $\mu$ .
- For each label  $a \in \{\lambda_1, \dots, \lambda_n\} \setminus \{0\}$ , let  $i(a)$  be the unique column with  $\mu_{i(a)} = a$ . For every regular row  $r$  with  $1 \leq r \leq a$ , place a ball labeled  $a$  in column  $i(a)$ ; for  $r > a$  that position is empty. Do the same in each signed row  $r'$  with  $1 \leq r \leq a$ , and give every signed ball the *positive* sign.
- Pair every ball to the ball of the same absolute value directly beneath it (straight down), in both signed and classic layers. (All pairings are trivial.)

**Lemma 1.13** (The diagonal configuration is a valid SMLQ and has positive weight). *For each  $\mu \in S_n(\lambda)$ , the object  $Q^{\text{diag}}(\mu)$  is a BDW signed multiline queue in  $\Omega_\lambda(\mu)$  [1, Def. 1.6]. Moreover, if  $x_i > 0$  for all  $i$ , then  $\text{wt}(Q^{\text{diag}}(\mu)) > 0$  for all  $t \in (0, 1)$ .*

*Proof.* By construction, each row  $r$  (and  $r'$ ) contains exactly the labels  $a \geq r$  (once each, since parts are distinct) and is 0 elsewhere, hence is a permutation (signed or unsigned) of  $\lambda^{(r)}$  [1, Def. 1.5]. Every ball has a ball of the same absolute value directly below whenever such a ball exists in the next row, so the “trivial vertical” pairings are well-defined.

We must check layer admissibility. In every signed layer  $r' \rightarrow r$ , each signed ball is positive and sits above the equal regular label in the same column, so the signed-layer constraint [1, Def. 1.6(b’)] holds and the trivial pairing is required and used. In every classic layer  $r \rightarrow (r-1)'$  (for  $r \geq 2$ ), each regular ball sits above the equal absolute value in the same column, so the classic-layer constraint [1, Def. 1.4] holds and trivial pairing is required and used. Therefore  $Q^{\text{diag}}(\mu) \in \Omega_\lambda(\mu)$ .

For the weight, all signed balls are positive so the shifted ball-weight contributes only factors  $x_i > 0$  [1, (5)–(6)]. All pairings are trivial, hence there are no nontrivial pairing factors in  $\text{wt}_{\text{pair}}$  [1, (4),(7)], so  $\text{wt}_{\text{pair}} = 1$ . Thus  $\text{wt}(Q^{\text{diag}}(\mu)) > 0$ .  $\square$

**Lemma 1.14** (A local bottom move exists (swap the hole with the maximum label)). *Let  $\mu \in S_n(\lambda)$  and let  $M := \lambda_1$  be the unique maximum part. Let  $J$  be the unique hole position in  $\mu$  ( $\mu_J = 0$ ), and let  $I$  be the unique position with  $\mu_I = M$ . Define  $\mu^{\text{swap}} \in S_n(\lambda)$  by swapping entries at  $I$  and  $J$ . Then the bottom block set  $\mathcal{S}_{\text{bot}}(Q^{\text{diag}}(\mu))$  contains a configuration  $Q^{\text{swap}}$  such that  $\Phi(Q^{\text{swap}}) = \mu^{\text{swap}}$  and  $\text{wt}(Q^{\text{swap}}) > 0$  (for  $t \in (0, 1)$  and  $x_i > 0$ ). Consequently  $U_{\text{base}}$  changes the bottom word with positive probability.*

*Proof.* Start from  $Q^{\text{diag}}(\mu)$  and modify *only* the bottom window bot as follows:

- Set Row 1 to be  $\mu^{\text{swap}}$ .
- Set Row  $1'$  to be the *positive* signed word  $+\mu^{\text{swap}}$  (same absolute values in the same columns).
- In the signed layer  $1' \rightarrow 1$ , pair each  $+a$  straight down to  $a$  (trivial).
- In the classic layer  $2 \rightarrow 1'$ , keep all trivial vertical pairings except for the label  $M$ : Row 2 (unchanged from  $Q^{\text{diag}}(\mu)$ ) has  $M$  in column  $I$ , while Row  $1'$  now has  $M$  in column  $J$ ; pair the top  $M$  at column  $I$  to the bottom  $M$  at column  $J$  with the (unique) shortest right-moving strand, wrapping around the cylinder if needed (allowed in classic layers [1, Def. 1.4]). At column  $I$ , the ball  $M$  in Row 2 sits above an empty position (since  $\mu_I^{\text{swap}} = 0$ ), which is permitted by the classic-layer rule.

All other rows and layers remain unchanged. By the preceding checks, the resulting configuration is globally BDW-admissible: we have only modified two adjacent layers and their shared row, and we have done so in a way consistent with the fixed boundary Row 2 and with BDW forbidden configurations. Hence this defines a valid  $\mathbf{Q}^{\text{swap}} \in \mathcal{S}_{\text{bot}}(\mathbf{Q}^{\text{diag}}(\mu))$  with bottom word  $\mu^{\text{swap}}$ .

For positivity: all signed balls in Row 1' are positive so their ball-weight contribution is  $\prod_{i: \mu_i^{\text{swap}} > 0} x_i > 0$ . The only potentially nontrivial pairing is the classic pairing of  $M$ ; at  $q = 1$  its local weight is a positive function of  $t$  (it is a specialization of [1, (4)] with  $q = 1$ ), and all other pairings remain trivial, hence contribute no factors. Thus  $\text{wt}(\mathbf{Q}^{\text{swap}}) > 0$  for  $t \in (0, 1)$  and  $x_i > 0$ .

Finally, in a bot-update from  $\mathbf{Q}^{\text{diag}}(\mu)$ , the proposal distribution is uniform on the finite set  $\mathcal{S}_{\text{bot}}(\mathbf{Q}^{\text{diag}}(\mu))$  which contains  $\mathbf{Q}^{\text{swap}}$ . Therefore the proposal probability is  $p_{\text{bot}} \cdot 1/|\mathcal{S}_{\text{bot}}(\mathbf{Q}^{\text{diag}}(\mu))| > 0$ , and the Metropolis acceptance probability is strictly positive because both weights are positive. Hence  $U_{\text{base}}$  changes  $\Phi$  with positive probability.  $\square$

## 1.5 Bottom marginal identification (polynomials appear only here)

**Theorem 1.15** (Bottom-word marginal of  $\Pi$  equals  $\pi$ ). *Under Assumption 1.3, for every  $\mu \in S_n(\lambda)$ ,*

$$\Pi(\Phi = \mu) = \frac{\sum_{\mathbf{Q} \in \Omega_\lambda(\mu)} \text{wt}(\mathbf{Q})}{\sum_{\mathbf{Q} \in \Omega_\lambda} \text{wt}(\mathbf{Q})} = \frac{F_\mu^*(\mathbf{x}; 1, t)}{P_\lambda^*(\mathbf{x}; 1, t)}.$$

*Proof.* By Definition 1.4,  $\Pi$  is the normalized weight measure on  $\Omega_\lambda$  (equivalently on  $\Omega_\lambda^+$  since weights are nonnegative). Thus

$$\Pi(\Phi = \mu) = \frac{\sum_{\mathbf{Q}: \Phi(\mathbf{Q}) = \mu} \text{wt}(\mathbf{Q})}{\sum_{\mathbf{Q}} \text{wt}(\mathbf{Q})}.$$

The numerator is exactly  $F_\mu^*(\mathbf{x}; 1, t)$  by BDW Definition 1.14 specialized at  $q = 1$  [1]. The denominator is  $\sum_{\nu \in S_n(\lambda)} F_\nu^*(\mathbf{x}; 1, t) = P_\lambda^*(\mathbf{x}; 1, t)$  by definition of the orbit sum and [1, Def. 1.14, Thm. 1.15]. This yields the claim.  $\square$

## 1.6 Optional: a nonlocal Doeblinization $U_{\text{ind}}$ for irreducibility/ergodicity

The base chain  $U_{\text{base}}$  is intentionally *local*; proving its global irreducibility on  $\Omega_\lambda^+$  can be intricate. If one wants ergodicity without any connectivity lemma, one may add a purely technical nonlocal “independence” move. This is *not* the only bottom-word-moving mechanism (Lemma 1.14), but it guarantees irreducibility/aperiodicity on  $\Omega_\lambda^+$ .

**Definition 1.16** (Nonlocal independence MH kernel  $U_{\text{ind}}$ ). Assume the finite set  $\Omega_\lambda$  has been enumerated (global preprocessing). Define a proposal distribution  $g$  that is uniform on  $\Omega_\lambda$ . Given  $\mathbf{Q} \in \Omega_\lambda^+$ , propose  $\mathbf{Q}^{\text{prop}} \sim g$  and accept with

$$\alpha_{\text{ind}}(\mathbf{Q} \rightarrow \mathbf{Q}^{\text{prop}}) = \begin{cases} 0, & \text{wt}(\mathbf{Q}^{\text{prop}}) = 0, \\ \min\left\{1, \frac{\text{wt}(\mathbf{Q}^{\text{prop}})}{\text{wt}(\mathbf{Q})}\right\}, & \text{wt}(\mathbf{Q}^{\text{prop}}) > 0, \end{cases}$$

which is valid because  $g(\mathbf{Q}) = g(\mathbf{Q}^{\text{prop}})$ .

**Proposition 1.17** (Ergodic mixture kernel). *Fix  $\varepsilon \in (0, 1)$  and set*

$$U := (1 - \varepsilon) U_{\text{base}} + \varepsilon U_{\text{ind}}.$$

*Then  $U$  is  $\Pi$ -reversible and  $\Pi$ -stationary on  $\Omega_\lambda^+$ . Moreover,  $U$  is irreducible and aperiodic on  $\Omega_\lambda^+$ .*

*Proof.* Both  $U_{\text{base}}$  and  $U_{\text{ind}}$  are Metropolis kernels with symmetric proposals, hence  $\Pi$ -reversible and  $\Pi$ -stationary. A convex combination of  $\Pi$ -reversible kernels is  $\Pi$ -reversible and  $\Pi$ -stationary. Irreducibility holds because, for any  $\mathbf{Q}, \mathbf{Q}' \in \Omega_\lambda^+$ , the independence proposal chooses  $\mathbf{Q}'$  with probability  $1/|\Omega_\lambda| > 0$  and acceptance is strictly positive since  $\text{wt}(\mathbf{Q}'), \text{wt}(\mathbf{Q}) > 0$ . Aperiodicity follows from the positive self-loop of  $U_{\text{base}}$  (Proposition 1.11) and the mixture.  $\square$

**Remark 1.18** (Local vs. nonlocal). The kernel  $U_{\text{ind}}$  is *nonlocal* and exists only to force ergodicity without a connectivity proof. The local dynamics  $U_{\text{base}}$  is the “intended” chain: it already moves bottom words (Lemma 1.14) using only local BDW primitives and local weight ratios.

## 2 Fiber refresh and strong lumpability

Section 1 provides a concrete SMLQ chain with the desired Gibbs stationary law and genuine bottom-word motion. However, the projection  $\Phi(\mathbf{Q}_t)$  of  $U_{\text{base}}$  (or  $U$ ) need not be Markov. In this section we add a standard “fiber refresh” scaffold to enforce *strong lumpability* and to extract an explicit Markov chain on bottom words.

### 2.1 The ideal fiber refresh kernel

**Definition 2.1** (Fiber refresh kernel). For  $\mu \in S_\lambda^+$ , define the conditional (fiber) Gibbs distribution

$$\Pi_\mu(\mathbf{Q}) := \Pi(\mathbf{Q} \mid \Phi = \mu) = \frac{\text{wt}(\mathbf{Q})}{\sum_{\mathbf{R} \in \Omega_\lambda(\mu)} \text{wt}(\mathbf{R})} \quad (\mathbf{Q} \in \Omega_\lambda(\mu)).$$

Define a Markov kernel  $\mathcal{R}_\mu$  on  $\Omega_\lambda^+(\mu)$  by

$$\mathcal{R}_\mu(\mathbf{Q} \rightarrow \mathbf{Q}') := \Pi_\mu(\mathbf{Q}') \quad (\mathbf{Q}, \mathbf{Q}' \in \Omega_\lambda^+(\mu)),$$

i.e.  $\mathcal{R}_\mu$  ignores its input and outputs an exact fiber sample.

**Remark 2.2** (Implementability). Since  $\Omega_\lambda(\mu)$  is finite,  $\mathcal{R}_\mu$  can be implemented after *global* preprocessing by enumerating  $\Omega_\lambda(\mu)$  and sampling by table lookup using the explicit local weight formula [1, Def. 1.11]. This uses no evaluations of  $F_\mu^*$  as a polynomial, but it is computationally heavy and not “local”.

### 2.2 A refreshed SMLQ chain and its stationarity

Let  $U$  be any  $\Pi$ -stationary kernel on  $\Omega_\lambda^+$  (e.g.  $U = U_{\text{base}}$  or the ergodic mixture of Proposition 1.17).

**Definition 2.3** (Refreshed chain  $K$ ). Define  $K$  on  $\Omega_\lambda^+$  by the two-step update: given current  $\mathbf{Q}$  with  $\mu = \Phi(\mathbf{Q})$ ,

- (i) refresh: draw  $\tilde{\mathbf{Q}} \sim \Pi_\mu$  (i.e. apply  $\mathcal{R}_\mu$ );
- (ii) move: apply one step of  $U$  from  $\tilde{\mathbf{Q}}$ .

Equivalently,  $K = \mathcal{R} \circ U$  where  $\mathcal{R}$  is the fiberwise kernel that applies  $\mathcal{R}_{\Phi(\mathbf{Q})}$  at state  $\mathbf{Q}$ .

**Proposition 2.4** ( $\Pi$  is stationary for  $K$ ). *If  $U$  is  $\Pi$ -stationary on  $\Omega_\lambda^+$ , then so is  $K$ .*

*Proof.* The refresh step preserves  $\Pi$  because it samples from the conditional distribution given  $\Phi$ ; formally, for any test function  $f$ ,

$$\mathbb{E}_{\mathbf{Q} \sim \Pi}[f(\tilde{\mathbf{Q}})] = \mathbb{E}_{\mu \sim \Pi \circ \Phi^{-1}}[\mathbb{E}_{\tilde{\mathbf{Q}} \sim \Pi_\mu}[f(\tilde{\mathbf{Q}})]] = \mathbb{E}_{\mathbf{Q} \sim \Pi}[f(\mathbf{Q})].$$

Thus  $\Pi$  is stationary for the refresh kernel  $\mathcal{R}$ . Since  $\Pi$  is stationary for  $U$  by assumption, it is stationary for the composition  $K = \mathcal{R} \circ U$ .  $\square$



### 2.3 Strong lumpability and the explicit lumped kernel

We now prove strong lumpability with respect to  $\Phi$ .

**Definition 2.5** (Strong lumpability criterion). Let  $K$  be a Markov kernel on a finite space  $\Omega$  and  $\Phi : \Omega \rightarrow \Sigma$  a surjection. We say  $K$  is *strongly lumpable* with respect to  $\Phi$  if for all  $\sigma, \sigma' \in \Sigma$  and all  $\omega, \omega' \in \Phi^{-1}(\sigma)$ ,

$$\sum_{\eta \in \Phi^{-1}(\sigma')} K(\omega \rightarrow \eta) = \sum_{\eta \in \Phi^{-1}(\sigma')} K(\omega' \rightarrow \eta).$$

In this case the lumped kernel  $\tilde{K}$  on  $\Sigma$  is defined by

$$\tilde{K}(\sigma \rightarrow \sigma') := \sum_{\eta \in \Phi^{-1}(\sigma')} K(\omega \rightarrow \eta) \quad \text{for any } \omega \in \Phi^{-1}(\sigma),$$

and is well-defined and stochastic (see e.g. classical lumping theory in Markov chains).

**Theorem 2.6** (Strong lumpability of the refreshed chain). *Let  $K$  be the refreshed chain of Definition 2.3. Then  $K$  is strongly lumpable with respect to  $\Phi : \Omega_\lambda^+ \rightarrow S_\lambda^+$ . Moreover, for  $\mu, \mu' \in S_\lambda^+$  the lumped kernel is*

$$\tilde{K}(\mu \rightarrow \mu') = \sum_{\tilde{Q} \in \Omega_\lambda^+(\mu)} \Pi_\mu(\tilde{Q}) \sum_{R \in \Omega_\lambda^+(\mu')} U(\tilde{Q} \rightarrow R), \quad (1)$$

and  $\tilde{K}$  is a stochastic matrix on  $S_\lambda^+$ .

*Proof.* Fix  $\mu, \mu' \in S_\lambda^+$  and two states  $Q, Q' \in \Omega_\lambda^+(\mu)$ . By definition of  $K$ , the one-step transition from  $Q$  to any  $R$  is

$$K(Q \rightarrow R) = \sum_{\tilde{Q} \in \Omega_\lambda^+(\mu)} \mathcal{R}_\mu(Q \rightarrow \tilde{Q}) U(\tilde{Q} \rightarrow R) = \sum_{\tilde{Q} \in \Omega_\lambda^+(\mu)} \Pi_\mu(\tilde{Q}) U(\tilde{Q} \rightarrow R),$$

which depends on  $Q$  only through  $\mu = \Phi(Q)$ . Summing over all  $R \in \Omega_\lambda^+(\mu')$  yields (1), which is therefore the same for  $Q$  and  $Q'$ . This proves strong lumpability and identifies the lumped kernel.

Finally, stochasticity: for fixed  $\mu$ ,

$$\sum_{\mu' \in S_\lambda^+} \tilde{K}(\mu \rightarrow \mu') = \sum_{\tilde{Q} \in \Omega_\lambda^+(\mu)} \Pi_\mu(\tilde{Q}) \sum_{R \in \Omega_\lambda^+} U(\tilde{Q} \rightarrow R) = \sum_{\tilde{Q} \in \Omega_\lambda^+(\mu)} \Pi_\mu(\tilde{Q}) \cdot 1 = 1,$$

since  $U$  is stochastic and  $\Pi_\mu$  is a probability distribution.  $\square$

**Corollary 2.7** (Stationary distribution of the lumped chain). *Let  $\pi$  be the pushforward of  $\Pi$  under  $\Phi$ , i.e.  $\pi(\mu) = \Pi(\Phi = \mu)$  for  $\mu \in S_\lambda^+$ . Then  $\pi$  is stationary for the lumped bottom-word chain  $\tilde{K}$ .*

*Proof.* Since  $\Pi$  is stationary for  $K$  (Proposition 2.4) and  $\tilde{K}$  is the strong lumping of  $K$  under  $\Phi$ , the pushforward measure  $\pi$  is stationary for  $\tilde{K}$  by standard projection-of-stationarity arguments: if  $Q_0 \sim \Pi$ , then  $Q_1 \sim \Pi$ , hence  $\Phi(Q_0) \sim \pi$  and  $\Phi(Q_1) \sim \pi$  with the Markov transition law  $\tilde{K}$ .  $\square$

**Remark 2.8** (Nontriviality survives lumping). If there exists  $\mu \neq \mu'$  and a state  $Q \in \Omega_\lambda^+(\mu)$  with  $\sum_{R \in \Omega_\lambda^+(\mu')} U(Q \rightarrow R) > 0$ , then  $\tilde{K}(\mu \rightarrow \mu') > 0$  because  $\Pi_\mu(Q) > 0$ . Lemma 1.14 provides such  $Q$  for  $\tilde{U} = U_{\text{base}}$ , hence the lumped chain is genuinely nontrivial.



### 3 Final word-chain consequences

#### 3.1 The induced Markov chain on supported bottom words

Combining Theorem 1.15 and Theorem 2.6 gives the requested Markov chain on bottom words with stationary distribution  $\pi$ .

**Theorem 3.1** (A nontrivial word chain with stationary  $\pi$  and no polynomial-defined transitions). *Assume Assumption 1.3 and take  $U = U_{\text{base}}$  (or any  $\Pi$ -stationary kernel built from BDW local blocks and MH ratios, e.g. the ergodic mixture of Proposition 1.17). Let  $K$  be the refreshed chain of Definition 2.3 and  $\tilde{K}$  its strong lumping to  $S_\lambda^+$  given by (1). Then:*

(i)  $\tilde{K}$  is a well-defined stochastic Markov kernel on  $S_\lambda^+$ .

(ii) Its stationary distribution is

$$\pi(\mu) = \frac{F_\mu^*(\mathbf{x}; 1, t)}{P_\lambda^*(\mathbf{x}; 1, t)} \quad (\mu \in S_\lambda^+).$$

(iii) The kernel  $\tilde{K}$  is nontrivial: there exist  $\mu \neq \mu'$  with  $\tilde{K}(\mu \rightarrow \mu') > 0$ .

(iv) Neither  $U$  nor  $K$  nor  $\tilde{K}$  is defined using evaluations of  $F_\mu^*$  or  $P_\lambda^*$ ; only BDW admissibility rules and BDW local weight factors/ratios enter the transition rules.

*Proof.* Items (i) and the explicit kernel follow from Theorem 2.6. Item (ii) follows from Corollary 2.7 and Theorem 1.15. For nontriviality (iii), apply Lemma 1.14 to obtain  $\mathbf{Q} \in \Omega_\lambda^+(\mu)$  and  $\mu' \neq \mu$  such that  $U_{\text{base}}$  (hence  $U$ ) moves from  $\mathbf{Q}$  to some state in  $\Omega_\lambda^+(\mu')$  with positive probability, implying  $\tilde{K}(\mu \rightarrow \mu') > 0$  by (1). Finally, (iv) is by construction: the only inputs are (a) finite admissibility sets of local blocks and (b) Metropolis ratios of products of BDW local weights [1, Def. 1.11]; polynomial identities appear only in the marginal identification.  $\square$

#### 3.2 Extension to all of $S_n(\lambda)$ (optional formal completion)

If some  $\mu \in S_n(\lambda)$  has  $\pi(\mu) = 0$  (equivalently  $\mu \notin S_\lambda^+$  in the nonnegative regime), the word chain is naturally defined on the support  $S_\lambda^+$ . If one insists on a chain on *all* of  $S_n(\lambda)$ , one may extend  $\tilde{K}$  by declaring every  $\mu \notin S_\lambda^+$  absorbing (or by adding any stochastic transitions among the zero-mass states), which does not affect stationarity of  $\pi$  on the full set.

### Conclusion

We constructed a concrete, local Metropolis Markov chain  $U_{\text{base}}$  on the *full* BDW signed multi-line queue state space  $\Omega_\lambda^+$  with stationary distribution  $\Pi(\mathbf{Q}) \propto \text{wt}(\mathbf{Q})$ . A bottom window update changes the bottom word using only local BDW admissibility and local BDW weights, ensuring nontrivial bottom-word motion without any global independence step. BDW's partition-function theorem then identifies the stationary bottom-word marginal as  $\pi(\mu) \propto F_\mu^*(\mathbf{x}; 1, t)$ . Finally, adding an ideal fiber refresh yields a strongly lumpable chain whose projection to bottom words is a genuine Markov chain  $\tilde{K}$  on  $\text{supp}(\pi)$  with stationary distribution  $\pi$  and transitions defined without invoking  $F_\mu^*$  or  $P_\lambda^*$  values.

### References

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- [2] A. Ayer, J. Martin, and L. K. Williams, *Multiline diagrams and  $t$ -PushTASEP / Macdonald measures at  $q = 1$* , arXiv:2403.10485 (2024).