

## COMP 4958: Assignment 1

The purpose of this assignment is to experiment with the RSA algorithm.

The RSA algorithm requires the generation of a pair of keys – a public key and a private key, each in turn is a pair of integers.

The keys are generated by first choosing a pair of large prime numbers  $p$  and  $q$  and calculating their product  $n$ .  $n$  is then one component of the public key and of the private key.

Let  $l$  be the LCM (least common multiple) of  $p - 1$  and  $q - 1$ . Choose a number  $e$  (with  $1 < e < l$ ) that is relatively prime to  $l$ , i.e., the GCD (greatest common divisor) of  $e$  and  $l$  is 1. The pair  $(e, n)$  is then the public key. (According to Wikipedia, a very common value for  $e$  is 65537.)

To generate the private key, we need to find  $d$ , the multiplicative inverse of  $e$  modulo  $l$ . This means we need to solve  $d \times e \equiv 1 \pmod{l}$  for  $d$ . (We'll say more about how to calculate  $d$  below.) The pair  $(d, n)$  is then the private key.

If a message (the plaintext) is represented by an integer  $M$ . The encrypted message (the ciphertext)  $C$  is given by  $M^e \bmod n$ . Given the ciphertext  $C$ , we can recover the plaintext by calculating  $C^d \bmod n$ .

See

[https://en.wikipedia.org/wiki/RSA\\_\(cryptosystem\)#Key\\_generation](https://en.wikipedia.org/wiki/RSA_(cryptosystem)#Key_generation)

to learn more about key generation.

For the mathematical calculations, you'll need to implement the following three functions (where  $a$ ,  $b$ ,  $m$  and  $n$  are all positive integers)

- `pown(a, n, m)` that returns the remainder of  $a^n$  when divided by  $m$ .

For this function, the problem with first finding  $a^n$  and then calculating its remainder is that  $a^n$  can be quite large. To keep the numbers relatively small, we can use the fact that, for  $a$ ,  $b$  and  $m$  all positive integers, `rem(a * b, m)` and `rem(rem(a, m) * rem(b, m), m)` are equal.

For example, `rem(53 * 24, 7)` is 5, `rem(53, 7) * rem(24, 7)` is  $4 * 3 = 12$  which also gives a remainder of 5 when divided by 7.

Secondly, the brute force way of multiplying  $n$  copies of  $a$  together to obtain  $a^n$  involves a lot of multiplications when  $n$  is large. To reduce the number of multiplications, we can do something similar to this example:  $2^{100}$  is  $(2^{50})^2$ ,  $2^{50}$  is  $(2^{25})^2$ ,  $2^{25}$  is  $2 \times (2^{12})^2$ , ...

- `lcm(a, b)` that returns the LCM of  $a$  and  $b$ .

Note that the LCM of  $a$  and  $b$  can be calculated as the quotient of  $a \times b$  by the GCD of  $a$  and  $b$  and the GCD can easily be calculated using the Euclidean algorithm. See

[https://en.wikipedia.org/wiki/Euclidean\\_algorithm#Implementations](https://en.wikipedia.org/wiki/Euclidean_algorithm#Implementations)

- `inverse(a, n)` that returns the multiplicative inverse of  $a$  modulo  $n$ .

The algorithm in “pseudo-code” to find `inverse(a, n)` is as follows:

```

/* https://en.wikipedia.org/wiki/Extended_Euclidean_algorithm#Modular_integers */
function inverse(a, n)
  t := 0;      newt := 1
  r := n;      newr := a

  while newr != 0 do
    quotient := r div newr
    (t, newt) := (newt, t - quotient * newt)
    (r, newr) := (newr, r - quotient * newr)

  if r > 1 then
    return "a is not invertible"
  if t < 0 then
    t := t + n

  return t

```

Use a tail-recursive helper function in your implementation.

In order to encrypt plaintext and decrypt ciphertext, we need to convert between text and integers. The PKCS #1 specification calls for 2 primitives:

- I2OSP (Integer-to-Octet-String Primitive): converts a non-negative integer into an octet string (a sequence of bytes) of a specified length.
- OS2IP (Octet-String-to-Integer Primitive): converts an octet string to a nonnegative integer.

I2OSP performs the conversion by basically writing the integer in base-256 & uses the digits as the octet string. For example,

$$310400273487 = 72 \times 256^4 + 69 \times 256^3 + 76 \times 256^2 + 76 \times 256 + 79$$

Hence the octet string, when written in Elixir bitstring syntax, is `<<72, 69, 76, 76, 79>>`. This corresponds to the string "HELLO". The OS2IP function takes the octet string & converts it back to an integer.

Note that I2OSP also makes sure the octet string has a certain length either by padding or rejecting an octet string that is too long. (It takes 2 arguments; See <https://tools.ietf.org/html/rfc8017#section-4> if you want to know the details.) However, for this assignment, we will omit this length requirement, i.e., the resulting octet string can be any length. Implement two Elixir functions `integer_to_binary` & `binary_to_integer` corresponding to I2OSP & OS2IP. For example

- `integer_to_binary(310400273487)` returns `<<72, 69, 76, 76, 79>>`
- `binary_to_integer(<<72,69,76,76,79>>)` (or `binary_to_integer("HELLO")`) returns 310400273487

You will also need to implement functions to encrypt & decrypt "text". However, the logic for the two operations are essentially the same except that we use different input text & keys. We can implement just one function: `crypt(text, {k, n})` where `{k, n}` is the key. If `k` is the `e` mentioned above, we are using the public key to encrypt `text` (which is a binary, i.e., an octet string). If `k` is `d`, then we are using the private key to decrypt `text`. Note that `crypt` returns a binary.

Put your functions in a module named `A1`. As a test, you will be provided values for `p`, `q`, `e` and a ciphertext. You are asked to decrypt that ciphertext.

Due date & submission details will be announced in class.