

Chapter 4 Arrays

Ground Rules

- Switch off your handphone and pager
- Switch off your laptop computer and keep it
- No talking while lecture is going on
- No gossiping while the lecture is going on
- Raise your hand if you have question to ask
- Be on time for lecture
- Be on time to come back from the recess break to continue the lecture
- Bring your lecture notes to lecture

4.1 Introduction

- Arrays provide means for allocating and accessing memory.
- Instead of declaring variables with different identifiers, we can use the same identifier with indices to represent variables. For instance, we use **int a[3]**; to replace **int x, y, z**;
- E.g.

```
int a[3];      int x, y, z;
```

a[0]:	
a[1]:	
a[2]:	

x:	
----	--

y:	
----	--

z:	
----	--

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4.2 One-Dimensional Arrays

We declare an array in C by giving the name of the array and the type and number of elements. For example,

```
int a[5];
```

a[0]
a[1]
a[2]
a[3]
a[4]

declares a as an array of five int objects.

Likewise,

```
long double x[100];
```

declares x as an array of 100 long double objects:
x[0], x[1], x[2], ..., x[99].

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We can initialize an array when it is declared by listing the initial values for its elements.

For example:

```
int a[5] = {75, 25, 100, -45, 60};
```

a[0]	75
a[1]	25
a[2]	100
a[3]	-45
a[4]	60

```
# define COUNT 999
```

```
# include <stdio.h>
```

```
main()
```

```
{
```

```
    int c[COUNT];
```

```
    int i;
```

```
    for (i=0; i < COUNT; i++) printf (" %d", c[i]);
```

```
    for (i=0; i < COUNT; i++) c[i]=50;
```

```
    getch(); // get a character from keyboard
```

```
    for (i=0; i < COUNT; i++) printf (" %d", c[i]);
```

```
    return 0;
```

```
}
```

c[0]	50
c[1]	50
c[2]	50
:	:
c[997]	50
c[998]	50

3

```
/* demo44.c */
```

```
/* find the average and the difference w.r.t. average of 8 real numbers by the use of array */
```

```
# include <stdio.h>
```

```
int main (void)
```

```
{
```

```
    float x[8], average, sum;
```

```
    int i;
```

```
    printf ("Enter 8 real numbers >");
```

```
    i = 0;
```

```
    sum = 0;
```

```
    while (i < 8)
```

```
    {
```

```
        scanf ("%f", &x[i]);
```

```
        sum += x[i]; /* same as sum += x[i++] */
```

```
        i++;
```

```
    }
```

```
    average = sum/8;
```

```
    printf ("\nThe average value is %.2f \n", average);
```

```
    printf ("Table of difference between x[i] and the average :\n");
```

```
    printf ("\n i    x[i]    x[i]-average");
```

```
    printf ("\n-----\n");
```

```
    i=-1;
```

```
    while (++i < 8)
```

```
    {
```

```
        printf ("%2d.  %.2f    %6.2f\n", i, x[i], x[i]-average);
```

```
    }
```

```
    return 0;
```

```
}
```

Screen Output:

Enter 8 real numbers >10 20 30 40 50 60 70 80

The average value is 45.00

Table of difference between x[i] and the average :

i	x[i]	x[i]-average
0.	10.00	-35.00
1.	20.00	-25.00
2.	30.00	-15.00
3.	40.00	-5.00
4.	50.00	5.00
5.	60.00	15.00
6.	70.00	25.00
7.	80.00	35.00

x[0]

x[1]

x[2]

x[3]

x[4]

x[5]

x[6]

x[7]

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Having learnt array, you are ready to learn some algorithms (算法) used in programming. In this chapter you will learn 3 sorting algorithms, namely, selection sort, bubble sort and heap sort.

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Sorting Problem

Given an array **a**, rearrange its elements so that they are in ascending or descending order.

Why should we study sorting ?

- Needed for arranging data for further analysis
 - Sorting the grades of students in a class
Who will get the SM2 Book Prize ?
 - Sorting the sale of products in a company
Which product should be abandoned ?
- As an initialization step to other algorithms
 - Pre-processing for Binary Search

Selection Sort

- Move the smallest element to **x[0]**.
- Move the second smallest element to **x[1]**.
- Move the third smallest element to **x[2]**.
- :
- Move the **(n-1)**-th smallest element to **x[n-2]**.

// selection sort, select.c

```
# include <stdio.h>
# include <stdlib.h>
# define n 50
# define seed 17

int x[n];

main ()
{
    int i, j, temp, index, min;
    long checksum;

    printf("\n\nSelection Sort:\n");
    srand(seed);           // fix the seed
    for(i=0;i<n;i++) x[i]= rand() % 300;

    printf("\nBefore:");
    for(i=0;i<n;i++) printf(" %d ",x[i]);

    checksum = 0;
    for(i = 0; i <n; i++){ checksum +=x[i];}
    printf("\nChecksum before sorting: %ld\n", checksum);
```

```
for (i=0; i<n-1; i++)
```

```
{
    index=i;
    min = x[i];
    for (j=i+1; j<n; j++) // n is the size of array
    {
        if (x[j] < min)
        {
            min = x[j];
            index = j;
        }
    }
    if (index !=i)
    {
        temp = x[i];
        x[i] = x[index];
        x[index] = temp;
    }
}
```



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```
printf("\nAfter:");
for(i=0;i<n;i++) printf(" %d ",x[i]);
```

```
checksum = 0;
for(i = 0; i <n; i++ ){ checksum +=x[i];}
printf("\nChecksum after selection sort: %ld \n", checksum);
```

```
return 0;
}
```

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Time Complexity Analysis of Selection Sort in the Worst Case in Term of the Number of Comparison Operations

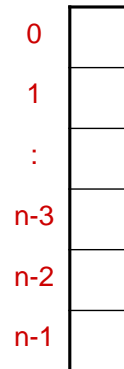
i	Number of Comparisons
0	n-1
1	n-2
2	n-3
:	:
n-3	2
n-2	1

$$\sum_{j=1}^{n-1} j = \frac{(n-1)(n)}{2} = \frac{1}{2}(n^2 - n)$$

$$= O(n^2)$$

Take the highest order and ignore the coefficient

```
for (i=0; i<n-1; i++)
{
    index=i;
    min = x[i];
    for (j=i+1; j<n; j++) // n is the size of array
    {
        if (x[j] < min)
        {
            min = x[j];
            index = j;
        }
    }
    if (index !=i)
    {
        temp = x[i];
        x[i] = x[index];
        x[index] = temp;
    }
}
```



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Bubble Sort

- Repeatedly sweep through array, exchanging adjacent pairs of elements that are out of order.
- Stop when one sweep does not produce any changes.
- Visualize the bubbles coming up a beaker of heated water.

// bubble sort for array. In this version the height of comparison is the same for all loops

include <stdio.h> // This version is not optimal

include <stdlib.h> // bubble1.c

define n 50

define seed 29

int x[n];

main ()

```
{
    int i, exchange, temp;
    long checksum;
```

printf("\n\nBubble Sort:\n\n");

```
srand(seed); // fix the seed
for(i=0;i<n;i++) x[i]= rand() % 300;
```

```
printf("\nBefore:");
for(i=0;i<n;i++) printf(" %d ",x[i]);
checksum = 0;
for(i = 0; i <n; i++) { checksum +=x[i];}
printf("\nChecksum before sorting: %ld\n", checksum);
```

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do

{

exchange=0; // assume no exchange

for (i=0; i<=n-2; i++) // n is the size of array

if (x[i] > x[i+1])

{

temp = x[i];

x[i] = x[i+1];

x[i+1] = temp;

exchange=1;

}

} while (exchange==1);

printf("\nAfter:");

```
for(i=0;i<n;i++) printf(" %d ",x[i]);
```

checksum = 0;

```
for(i = 0; i <n; i++) { checksum +=x[i];}
```

return 0;

}



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// bubble sort for array - the height of next loop of comparison is shorter by 1 step

// The algorithm will stop as soon as the array is sorted.

// bubbl2a.c

include <stdio.h>

include <stdlib.h>

define n 30

define seed 29

int x[n];

main ()

```
{
    int i, height, exchange, temp;
    long checksum;
    int sorted;
```

printf("\n\nBubble Sort:\n\n");

```
srand(seed); // fix the seed
for(i=0;i<n;i++) x[i]= rand() % 300;
```

printf("\nBefore:");

```
for(i=0;i<n;i++) printf(" %d ",x[i]);
```

checksum = 0;

```
for(i = 0; i <n; i++) { checksum +=x[i];}
```

```
printf("\nChecksum before sorting: %ld\n", checksum);
```

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exchange=1;

for (height=n-2 ; height>=0 && exchange==1 ; height--)

{

exchange=0; // assume no exchange

for (i=0; i<=height; i++) // rise the number up to height

if (x[i] > x[i+1])

{

temp = x[i];

x[i] = x[i+1];

x[i+1] = temp;

exchange=1;

}

}

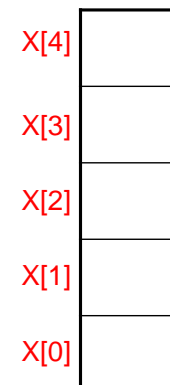
printf("\nAfter:");

```
for(i=0;i<n;i++) printf(" %d ",x[i]);
```

checksum = 0;

```
for(i = 0; i <n; i++) { checksum +=x[i];}
```

```
printf("\nChecksum after sorting bubble1: %ld\n", checksum);
```

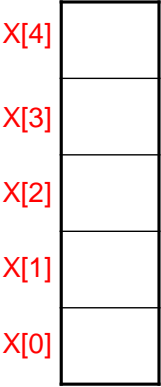


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```
sorted = 1;
for(i = 0; i < n-1; i++)
{
    if (x[i] > x[i+1])
    {
        sorted = 0;
        break;
    }
}

if (sorted == 1)
    printf ("\n sorted!");
else
    printf ("Not sorted");

return 0;
}
```



Time Complexity Analysis of Bubble Sort (with decreasing height) in the Worst Case in Term of the Number of Comparison Operations

The first bubble will go to the top (position **n-1**) and this takes n-1 comparisons.
The second bubble will go to position **n-2** and this takes n-2 comparisons.
The third bubble will go to position **n-3** and this takes n-3 comparisons.

:

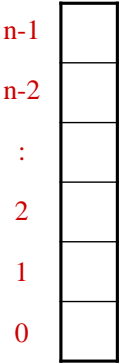
The second last bubble will go to position **2** and this takes 2 comparisons.
The last bubble will go to position **1** and this takes 1 comparison.

$$\sum_{j=1}^{n-1} j = \frac{(n-1)(n)}{2} = \frac{1}{2}(n^2 - n)$$
$$= O(n^2)$$

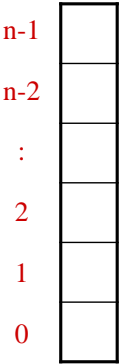
Decreasing height in subsequent iteration

$$\sum_{j=1}^{n-1} (n-1) = (n-1)(n-1)$$
$$= n^2 - 2n + 1$$
$$= O(n^2)$$

Always bubble to the top in all iterations



What is the best time complexity of Bubble Sort to sort n data?



Structured Sorting Algorithm

Heap Sort

Why do we want a structure in sorting algorithm?

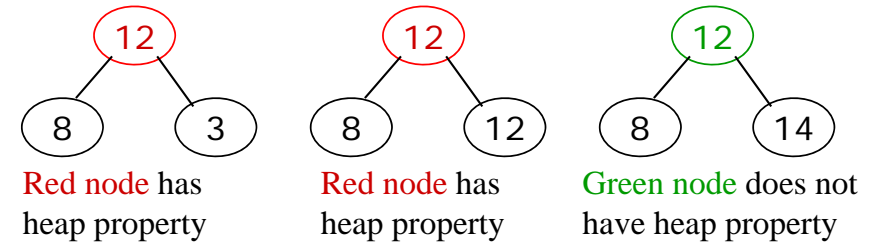
Although bubble sort is slow $O(n^2)$ in the worst case, it can be very fast $O(n)$ for some pre-orderings of data. For example, when the input data are not too messed up initially, bubble sort can stop in a few iterations. But this best performance cannot be guaranteed. We want guarantee in the time complexity for any pre-orderings of data because guarantee is critical in many real-time applications such as in the aircraft control system, aircraft navigation system etc.

Unless there is a structure to constrain or condition the algorithm, there is no way to give any guarantee in the performance. (Likewise in life.)

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The heap property

- A node has the heap property if the value in the node (father/parent) is as large as or larger than the values in its sons (children).

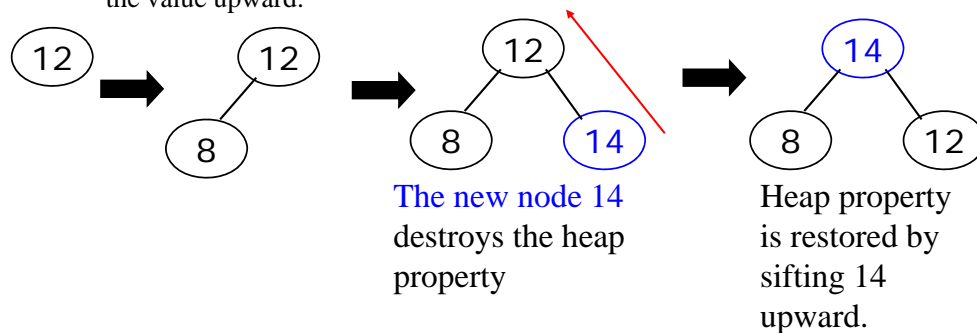


- All leaf nodes automatically have the heap property
- A binary tree is a heap if *all* nodes in it have the heap property

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Sifting Up

- Given a new node that destroys the heap property of a binary tree, you can restore the heap property by exchanging its value with its parent and cascade the value upward.

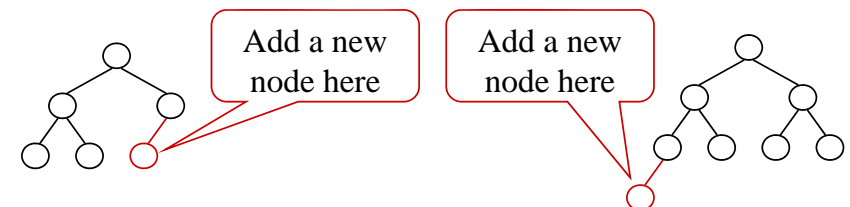


- This is called sifting up because the new value may have to go all the way up to the root.

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Constructing a heap I

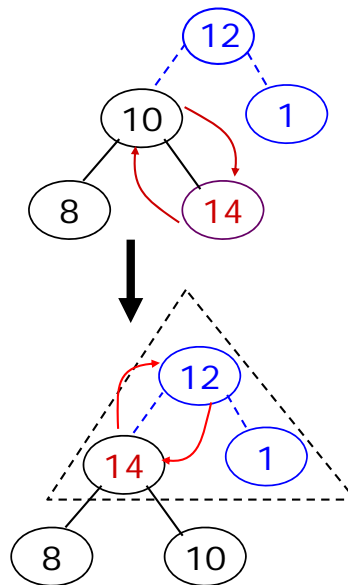
- A tree consisting of a single node is automatically a heap
- We construct a heap by adding nodes one at a time:
 - Add the node just to the right of the rightmost node in the deepest level
 - If the deepest level is full, start a new level
- Examples:



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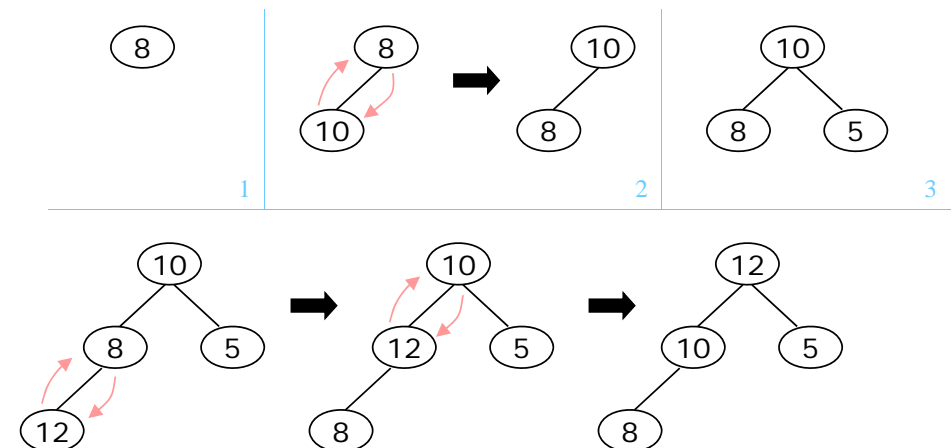
Constructing a heap II

- Each time we add a node, we may destroy the heap property of its parent node because the value of the new node may be smaller than the value of the parent node.
- To fix this, we have to sift up (now it is coasting upward)
- But each time we sift up, the value of the top node in the sift may increase, and this may destroy the heap property of *its* parent node (i.e., the new value of the parent may be greater than the value of the grandparent)
- We repeat the sifting up process, moving up in the tree, until either
 - We reach nodes whose values don't need to be swapped (because the parent is *still* larger than both children), or
 - We reach the root of the binary tree.



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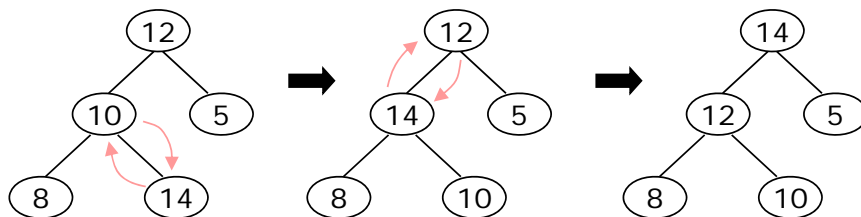
Constructing a heap III



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Other children are not affected

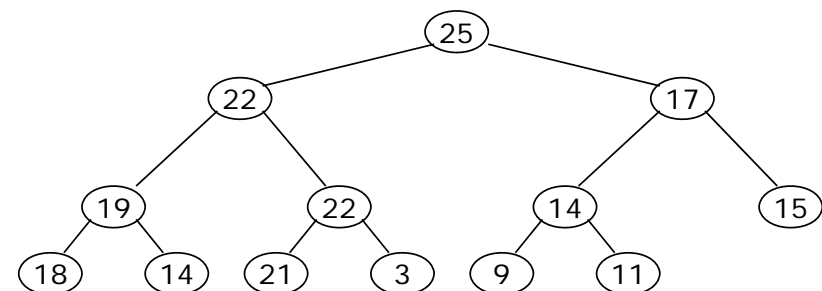


- The node containing 8 is not affected because its parent gets larger, not smaller
- The node containing 5 is not affected because its parent gets larger, not smaller
- The node containing 8 is still not affected because, although its parent got smaller, its parent is still greater than it was originally

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A sample heap

- Here's a sample binary tree after it has been heapified

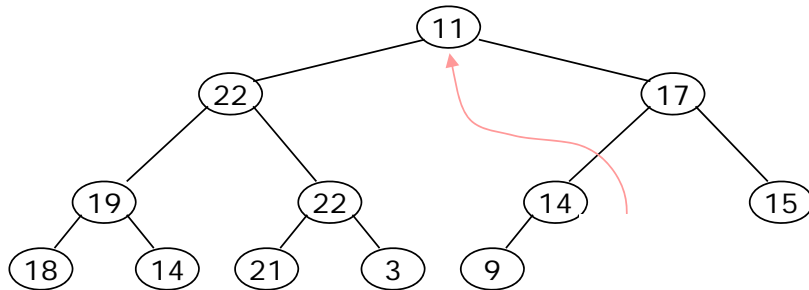


- Notice that heapified does *not* mean sorted
- Heapifying does *not* change the shape of the binary tree; this binary tree is balanced and left-justified because it started out that way

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Removing the root

- Notice that the largest number is now in the root
- Suppose we *discard* the root:

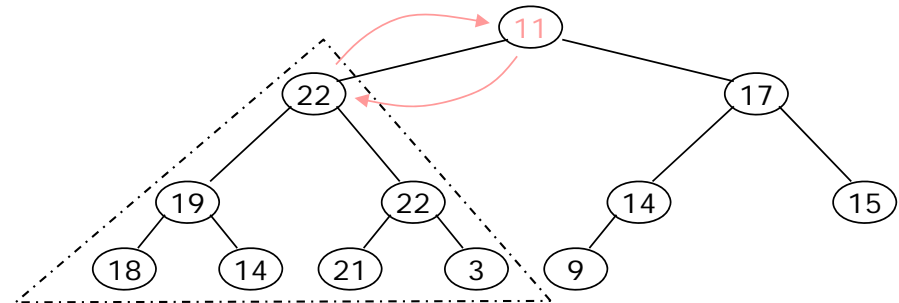


- How can we fix the binary tree so it is once again *balanced and left-justified*?
- Solution: remove the rightmost leaf at the deepest level and use it for the new root

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The Re-Heap method I

- Our tree is balanced and left-justified, **but is no longer a heap**
- However, *only the root* lacks the heap property

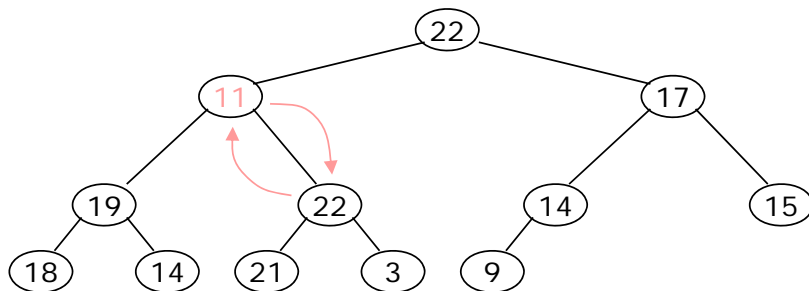


- We can **sift down** the root so that the heap property is restored.
- But after doing this, one of its children may have lost the heap property. If it is so, we continue to restore the heap property of the sub-tree downwardly.

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The reHeap method II

- Now the left child of the root (still the number 11) lacks the heap property

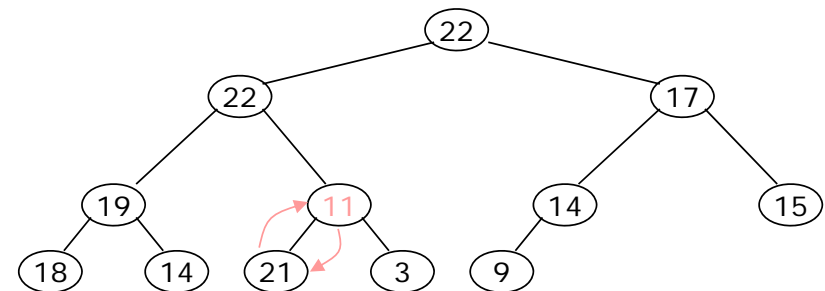


- We can **sift down** this node
- After doing this, one and only one of its children may have lost the heap property

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The reHeap method III

- Now the right child of the left child of the root (still the number 11) lacks the heap property:

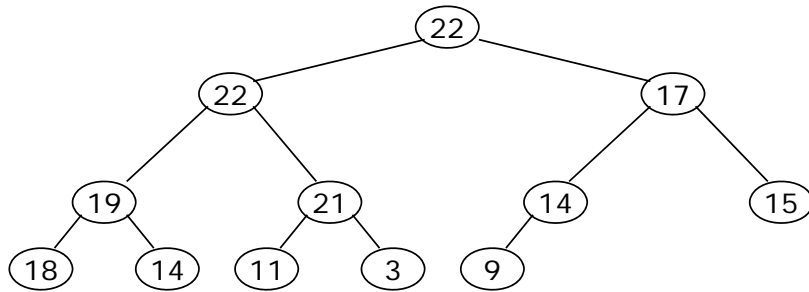


- We can **sift down** this node again
- After doing this, one and only one of its children may have lost the heap property — but it doesn't, because it's a leaf

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The reHeap method IV

- Our tree is once again a heap, because every node in it has the heap property



- Once again, the largest of the remaining values is in the root
- We can repeat this process until the tree becomes empty
- This produces a sequence of values in the order from the largest to the smallest

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Sorting

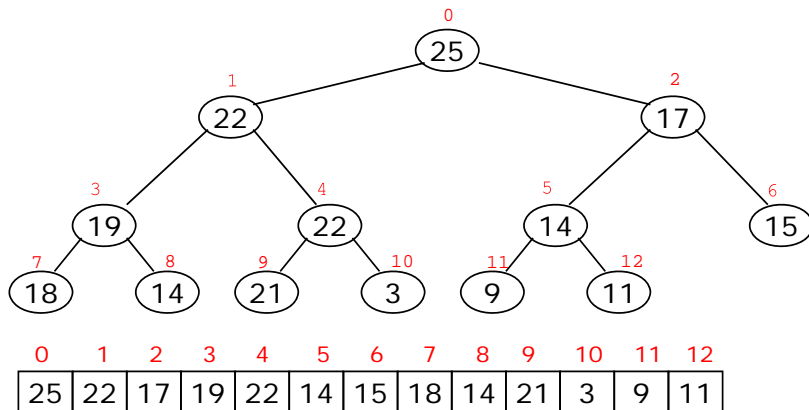
- What do heaps have to do with sorting an array?
- Here's the neat part:
 - Because the binary tree is *balanced* and *left justified*, it can be represented as an array
 - All our operations on binary trees can be represented as operations on *arrays*
 - To sort:

```

heapify the array;
while the array isn't empty
{
    remove and replace the root;
    reheap the new root node;
}
  
```

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Mapping into an array

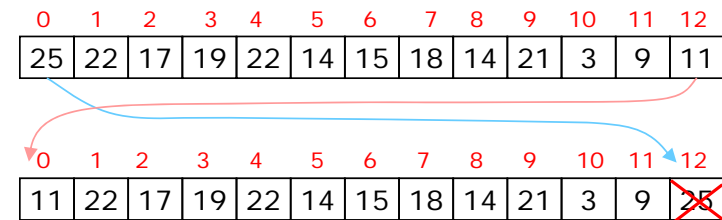


- Notice:
 - The left child of index i is at index $2*i+1$
 - The right child of index i is at index $2*i+2$
 - Example: the children of node 3 (19) are 7 (18) and 8 (14)

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Removing and replacing the root

- The “root” is the first element in the array
- The “rightmost node at the deepest level” is the last element
- Swap them...



- ...And pretend that the last element in the array no longer exists—that is, the “last index” is 11 (9)

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Reheap and repeat

- Reheap the root node (index 0, containing 11)...

- ...And again, remove and replace the root node
- Remember, though, that the “last” array index is changed
- Repeat until the last becomes first, and the array is sorted!

```
/* Heap Sort */
#include<stdio.h>
#include<stdlib.h>

#define N 100
#define BASE_START_INDEX 63
// pow(2,0) + pow (2,1) + ... + pow (2, (floor[lg(N)] -1) )
// 1 + 2 + 4 + 8 + 16 + 32 = 63

int x[N];

main()
{
    int i, done;
    int parent, this1, temp;
    long checksum;
    int max_index;
    int j, sorted;

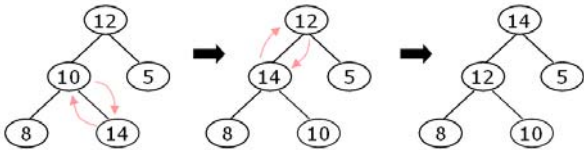
    printf("\n\nHeap Sort:\n\n");
```

```
srand(17); // fix the seed
for(i=0;i<N;i++) x[i]= rand() % 250;

printf("\nBefore:");
for(i=0;i<N;i++) printf(" %d ",x[i]);

checksum = 0;
for(i = 0; i <N; i++ ){ checksum +=x[i];}
printf("\nChecksum before sorting: %ld\n", checksum);
```

```
// Here I build the real heap.
for (i=0;i<N;i++) // inspect the array element one by one
{
    done = 0;
    this1=i; // the element under investigation
              // this1 refers to the child when sifting up
    while (!done)
    {
        if (this1==0) //reach the root thus it is done
            done=1;
        else
        {
            // start sifting up
            parent = (this1-1)/2;
            if (x[parent] < x[this1])
            {
                temp = x[parent]; // exchange child with parent
                x[parent] = x[this1];
                x[this1] = temp;
                this1=parent; // child is sifted up
            }
            else
                done =1; // no exchange so sifting completed
        }
    } //else
} // while
} // for i
```



```
printf ("\nAfter heapify : ");
for(j=0;j<N;j++) printf(" %d ",x[j]);
```

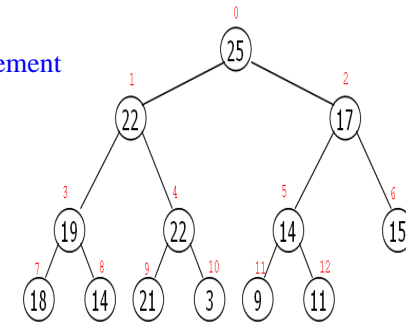
```
checksum = 0;
for(j = 0; j < N; j++){ checksum +=x[j];}
printf("\nChecksum after heapify: %ld\n", checksum);
```

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```
// insertion and heap restoration
for(i=N-1;i>=1;i--)
{
    temp = x[0];
    x[0] = x[i];
    x[i] = temp; // exchange the top with the current_last_element
                // x[i] to x[n-1] are already committed
```

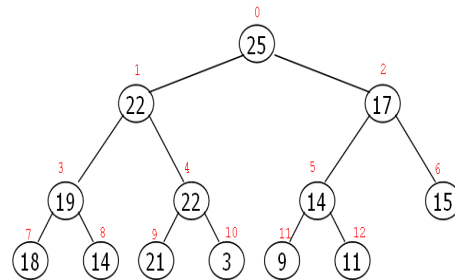
// have to sift down because the top has been disturbed

```
done = 0;
this1 = 0;    // store from root
              // this1 refers to the parent
              // when sifting down
```



```
while (!done)
{
    if (this1 >= BASE_START_INDEX) //reach the bottom
        done=1;                    // thus it is done as it has no child
    else if ( (2*this1+1)>= i )      //if the left child is already committed
        done =1;                  // done bcos right child is also committed
```

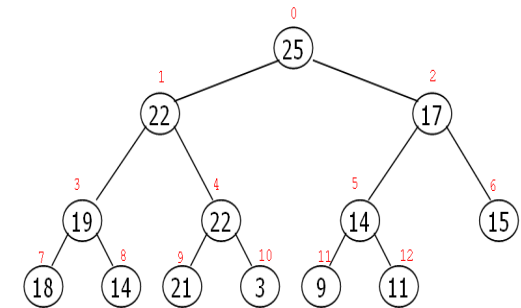
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```
else if ( (2*this1+2)>= i )
{ // left child not committed, only right child is committed
    if (x[this1] < x[2*this1+1])
    {
        temp = x[this1]; // exchange parent with left child
        x[this1] = x[2*this1+1];
        x[2*this1+1] = temp;
        this1 = 2*this1+1; // sifting down
    }
    else
        done=1; // no exchange with left child, restoration is completed
}
```

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```
else
{ // both children are not committed
    if (x[2*this1+1] > x[2*this1+2])
        max_index= 2*this1+1;
    else
        max_index = 2*this1+2;
    if (x[this1] < x[max_index])
    {
        temp = x[this1]; // exchange parent with the max child
        x[this1] = x[max_index];
        x[max_index] = temp;
        this1=max_index; // sifting down
    }
    else
        done=1; // no exchange thus heap restoration is complete
}
} // while
} // for i
```



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```
printf("\nEnd:");
for(i=0;i<N;i++) printf(" %d ",x[i]);

checksum = 0;
for(i = 0; i < N; i++ ){ checksum +=x[i];}
printf("\nChecksum after sorting: %ld\n", checksum);
```

```
sorted = 1;
for (i = 0; i < N-1; i++ )
{
    if (x[i] > x[i+1])
    {
        sorted = 0;
        break;
    }
}

if (sorted ==1)
    printf ("\n sorted!");
else
    printf ("Not sorted");

return 0;
}
```

```
Heap Sort:

Before: 137 198 22 162 119 28 152 139 89 145 3 123 130 133 80
167 204 209 123 104 8 44 134 159 212 244 106 189 223 92 11 196
17 105 165 189 17 84 182 84 71 146 158 9 80 64 135 94 182 60
86 48 177 0 40 105 7 100 234 86 91 161 91 93 55 234 249 68
135 23 221 87 83 113 117 128 2 118 103 173 77 245 114 162 40
76 197 231 112 220 96 213 11 189 138 168 232 218 91 27

Checksum before sorting: 12012
After heapify : 249 245 234 234 244 232 223 221 189 197 231 218
177 212 161 198 204 167 182 158 173 220 213 209 130 159 106 105
189 91 92 137 196 135 165 89 117 128 123 104 146 145 162 112
135 134 189 168 182 60 86 48 123 0 40 28 7 100 133 80 86 11
91 93 55 17 162 68 105 23 139 87 83 17 113 84 2 118 103 84 77
71 114 8 40 76 119 3 80 9 96 44 11 64 138 22 152 94 91 27
Checksum after heapify: 12012

End: 0 2 3 7 8 9 11 11 17 17 22 23 27 28 40 40 44 48 55 60
64 68 71 76 77 80 80 83 84 84 86 86 87 89 91 91 91 92 93 94
96 100 103 104 105 105 106 112 113 114 117 118 119 123 123 128
130 133 134 135 135 137 138 139 145 146 152 158 159 161 162 162
165 167 168 173 177 182 182 189 189 189 196 197 198 204 209 212
213 218 220 221 223 231 232 234 234 244 245 249
Checksum after sorting: 12012
sorted!
```

What is the best time complexity and worst time complexity of Heap Sort to sort n data?

We will now learn how to search an array. The first algorithm will be the sequential search.

```
// search1.c sequential search
# include <stdio.h>
# include <stdlib.h>
# define n 60
# define seed 29

int x[n];

main ()
{
    int search_key, found, i;
    srand(seed); // fix the seed
    for(i=0;i<n;i++) x[i]= rand() % 150;

    printf("\nEnter a search key: ");
    scanf ("%d", &search_key);

    found = 0;
    for(i = 0; i <n; i++ )
    {
        if (x[i] == search_key)
        {
            found=1;
            break;
        }
    }
}
```

```
if (found)
    printf ("\n %d is found in the array", search_key);
else
    printf ("\n %d is not found in the array", search_key);

printf("\n\nContents of array:\n");
for(i=0;i<n;i++) printf(" %d ",x[i]);

return 0;
}
```

--	--	--	--	--	--	--	--	--	--

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Binary Search: The array will have to be sorted first.

```
// search2.c binary search
```

```
# include <stdio.h>
# include <stdlib.h>

# define n 30
# define seed 29

int x[n];

main ()
{
    int i, height, exchange, temp;
    int search_key, found, left, right, mid;

    srand(seed); // fix the seed
    for(i=0;i<n;i++) x[i]= rand() % 300;

    exchange=1;
}
```

```
for (height=n-2 ; height>=0 && exchange==1 ; height--)
{
    exchange=0; // assume no exchange

    for (i=0; i<=height; i++) // rise the number up to height
    if (x[i] > x[i+1])
    {
        temp = x[i];
        x[i] = x[i+1];
        x[i+1] = temp;
        exchange=1;
    }
}
```

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```
printf ("\nEnter a search key: ");
scanf ("%d", &search_key);
```

```
left=0;
right= n-1;
found=0;
```

2	6	13	18	23	41	56	74	92	97
---	---	----	----	----	----	----	----	----	----

```
while ( (left <= right) && !found)
{
    mid = (left+right)/2;
    if (x[mid] == search_key)
        found=1;
    else if (x[mid] < search_key) left = mid+1;
    else right = mid -1;
}
```

```
if (found)
    printf ("\n %d is found in the array", search_key);
else
    printf ("\n %d is not found in the array", search_key);
```

```
printf("\n\nContents of array:");
for(i=0;i<n;i++) printf(" %d ",x[i]);
```

```
return 0;
}
```

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