# DISTRIBUTED CONWAY'S GAME OF LIFE SIMULATION

**CMPT431 Final Presentation** 

#### **Background Introduction**

- The Game of Life is a visualized cellular automaton devised by the British mathematician John Horton Conway in 1970.
- The universe (board) of the Game of Life is a twodimensional orthogonal grid of square cells.
- Each cell has three states:
  - Survival
  - Death
  - New Birth

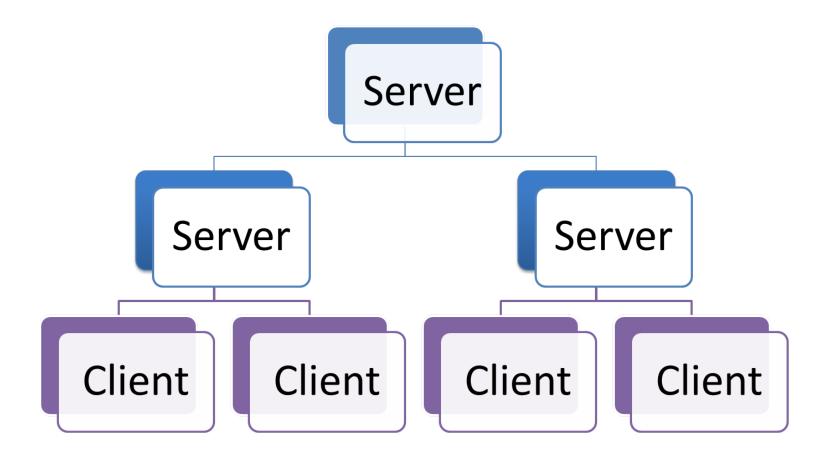
#### Objectives

- Implement a distributed Conway's Game of Life computational simulation
- Handle large-size packetized network messages
- Build a graphical interface to visualize the universe
- Support over 100,000,000 cells and 10,000,000 lives
- Support dynamically adding and deleting computing nodes
- Finish one cycle with 100,000,000 cells within 0.25 sec

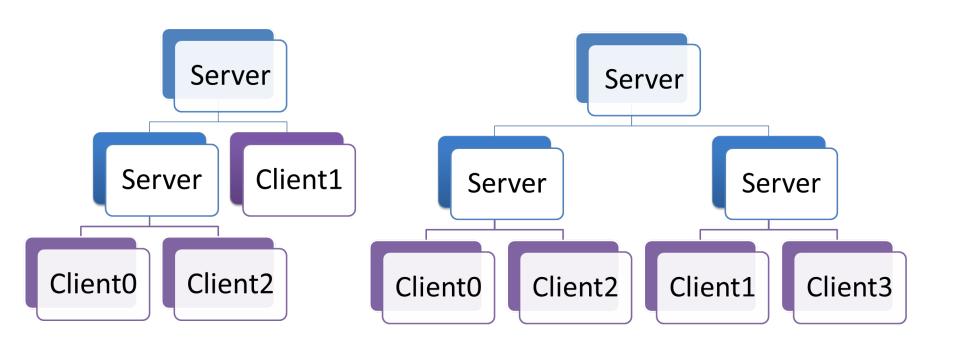
#### Challenges

- Computing nodes must exchange information to coordinate with each other
- Various messages must be designed to help implement dynamically adding and deleting nodes
- Server must take control to the computing nodes to guarantee transitions occur simultaneously
- Network messages should be minimized in order to avoid the system becomes network-bound with only a few computing nodes

#### Structures



### Adding



#### Adding (cont.)

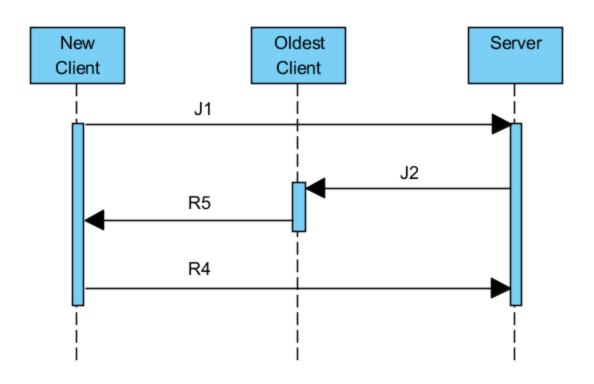
- A wants to join (first client)
  - A => Server: application
  - Server => A: accepted & data
  - A => Server: received

#### Adding (cont.)

- A wants to join (general)
  - A => Server: application
  - Server => B (\*): ask B to work with A
  - B => A: accepted & data
  - A => B: received
  - B => Server: everything is done

\*B is the first node who has no siblings; if everybody has, then choose first node.

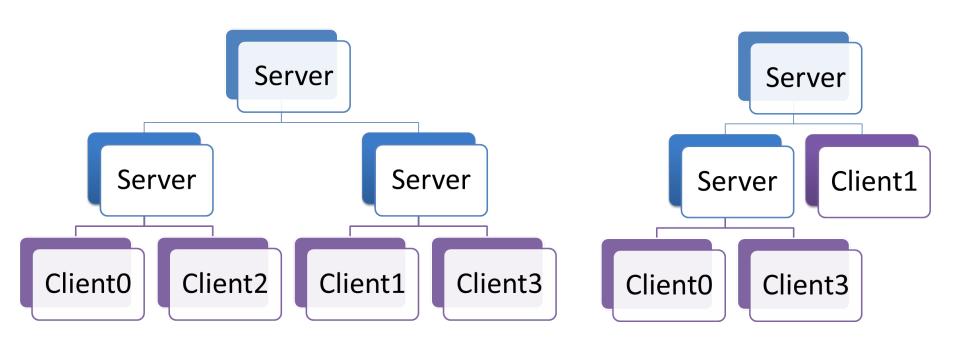
# Adding (cont.)



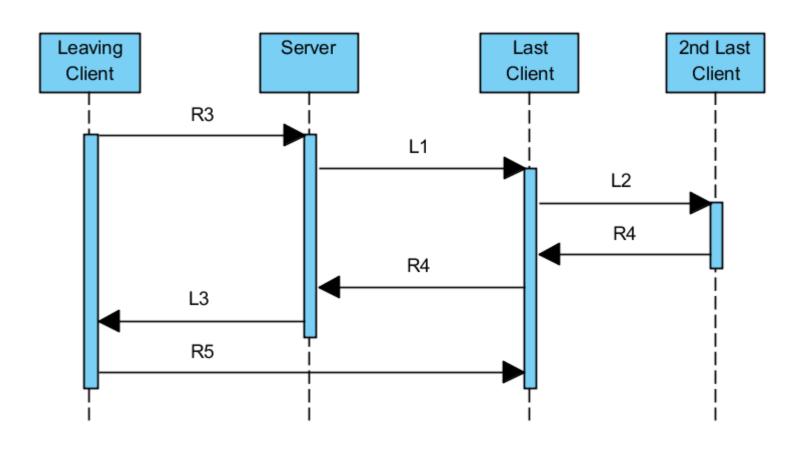
#### Message Passing

Message		Direction	Payload	Description	
R1	Next Clock	S→C	new clock number	Inform all the clients to start new computation round	
R2	Border	C→C	border status vector	Send border information to clients that is in charge of adjacent area	
R3	Report	C→S	whether is leaving, top and left coordinate, status bit map	Report finished computation, and return the computation result to the server in interactive mode	
R4	Confirm	$A \rightarrow A$	none	Confirm received last message	
R5	New Outfit	C→C	c o n n e c t i o n information, outfit information	Send all the computing information to a new added client	
R6	N e i g h b o r Update	C→C	c o n n e c t i o n information, relative position	Inform all the neighbors about the position change	
J1	Request Join	C→S	connection information	Send request to server for joining the computation system	
J2	Split	S→C	c o n n e c t i o n information, new client id, split direction	Inform the right client to split its area, construct the new outfit and send to the new client	
L1	Last Merge	S→C	connection information	Inform the last two client to merge into one, the idle one waits for new command	
L2	Outfit Merge	C→C	connection information, outfit	Send outfit to pair client for merging	
L3	L e a v e Permit	S→C	connection information	Send outfit to the specified client	

# Leaving



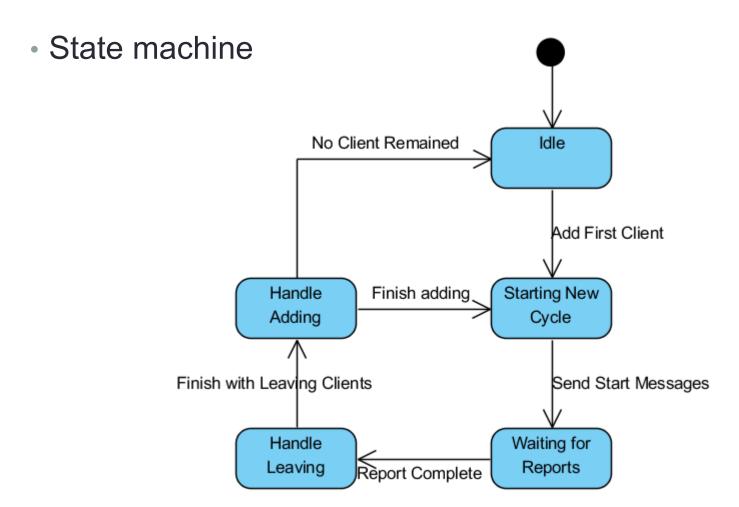
# Leaving (cont.)



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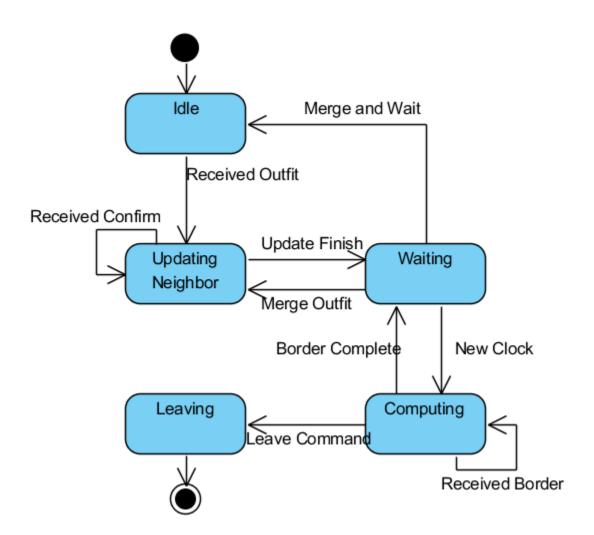
#### Server



### Work Cycle

- Server
  - Cycle start
  - Synchronization
  - Wait for results & Handle missing
  - Merge & Refresh
  - Respond to requests(join, leave)
  - Cycle end

#### Client



### Work Cycle (cont.)

#### Client

- Send joining application to server
- wait for accepting
- Cycle start
- Respond to server's request of sibling change, replacement or new sibling
- Synchronization
- Compute
- Compress if needed & Send
- Back up if necessary
- Send leaving application if necessary
- Cycle end

#### Neighbor Updating

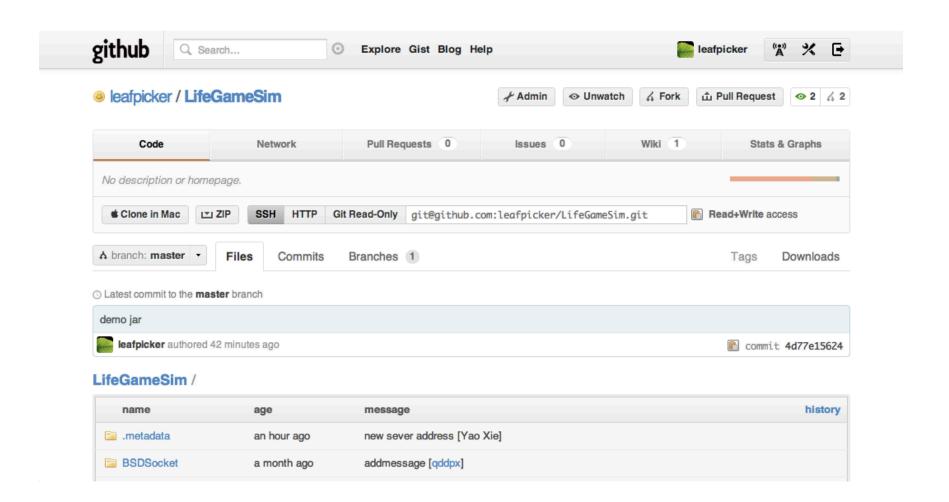
- Every node has up to 12 neighbors
- Send message to notice all neighbors

1	2	3	4
12			5
11			6
10	9	8	7

#### Message Queue

- Java NIO
- Multi-thread
  - Monitoring new connection
  - Monitoring new message and push to a queue with thread protection

### Github Management



#### Performance Testing

Ensure the system will perform well and be reliable:

- a) Under high user load
- b) With a large dataset
- c) Over an extended period time
- d) As the load/dataset increases

#### Test Objectives

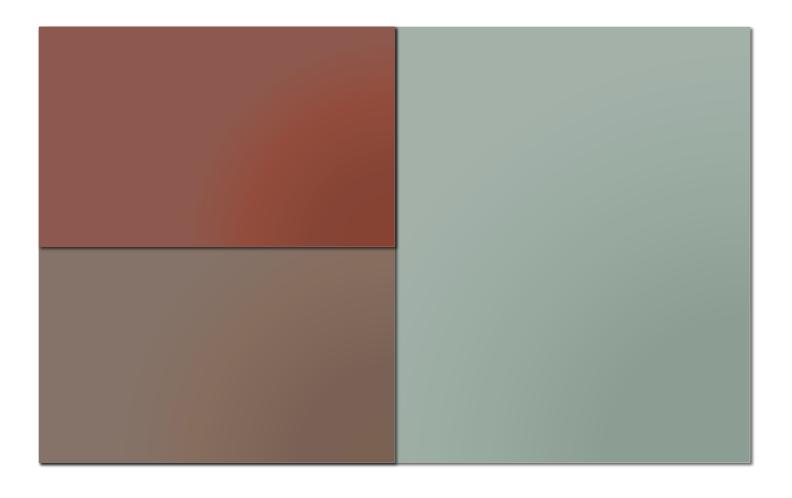
- With the increase of clients, we want to observe the updates(cycles) per second
- The longest computing time in each round
- The average computing time of each client in each round
- The network flow (including sending message and receiving message)

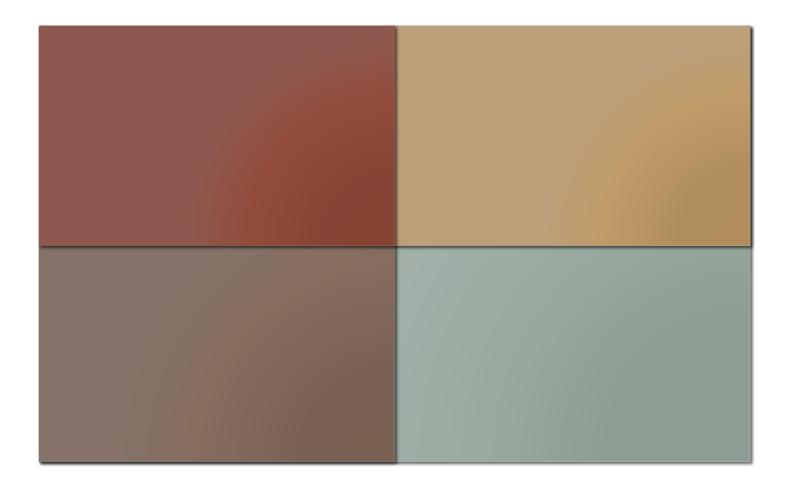
#### **Test Strategy**

- Append incoming clients to a pending list
- Add a new client per 20 cycles as a new round
- cycles / second = 20 \* 1000 / (endTime startTime)
- Each round:
  - Weakest-link principle (Cannikin Law)
  - compared with the longest computing time.
  - Monitor the sender and receiver and record the size of each package it send or receive
  - Total network flow =  $\sum$  sender +  $\sum$  receiver

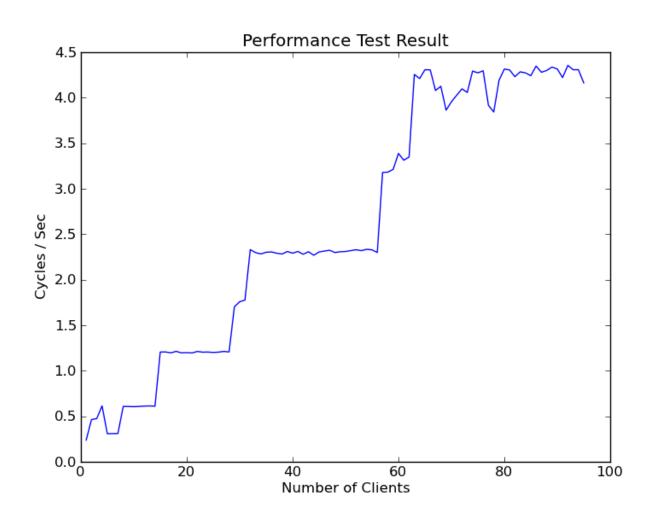




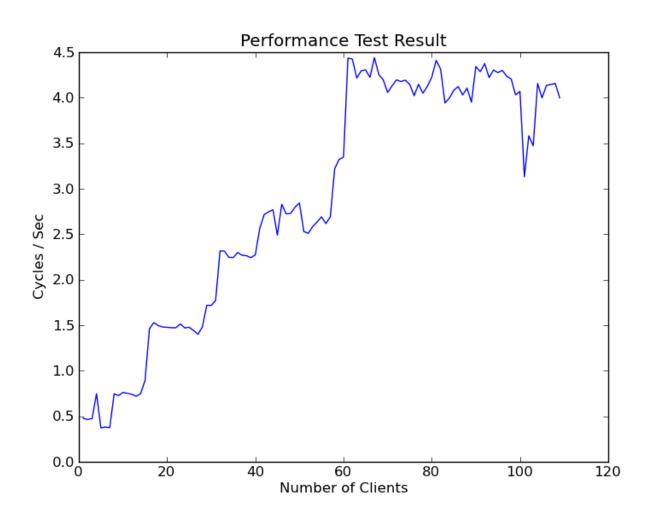




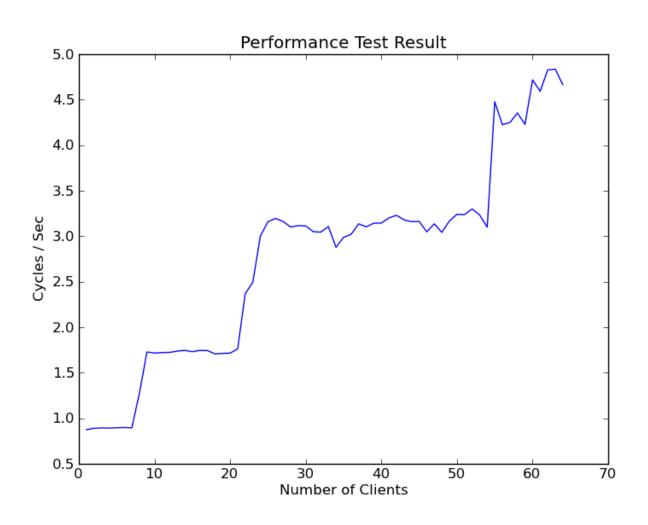
# Sample Run 1



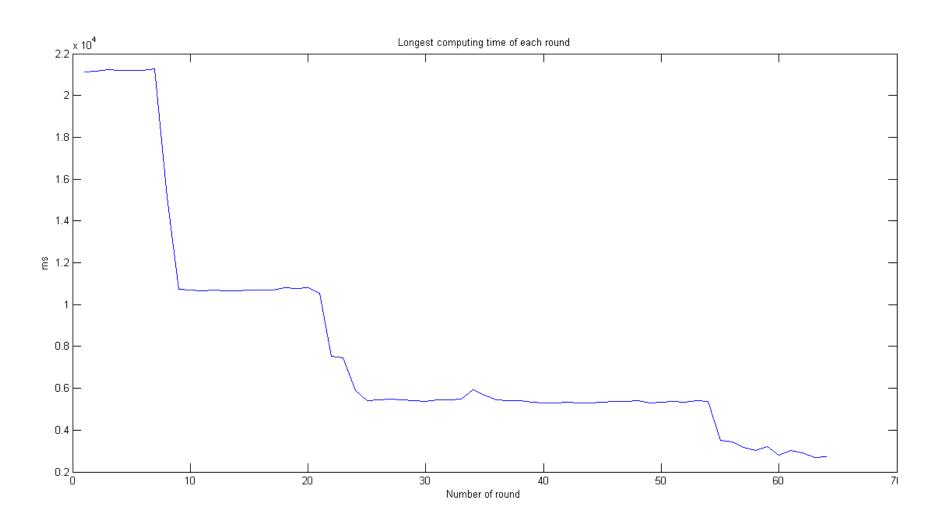
# Sample Run 2



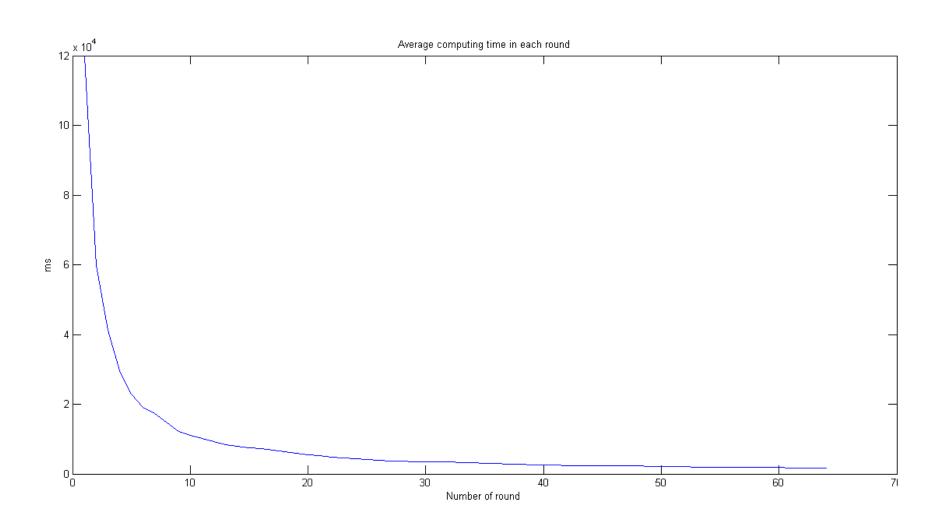
# Sample Run 3



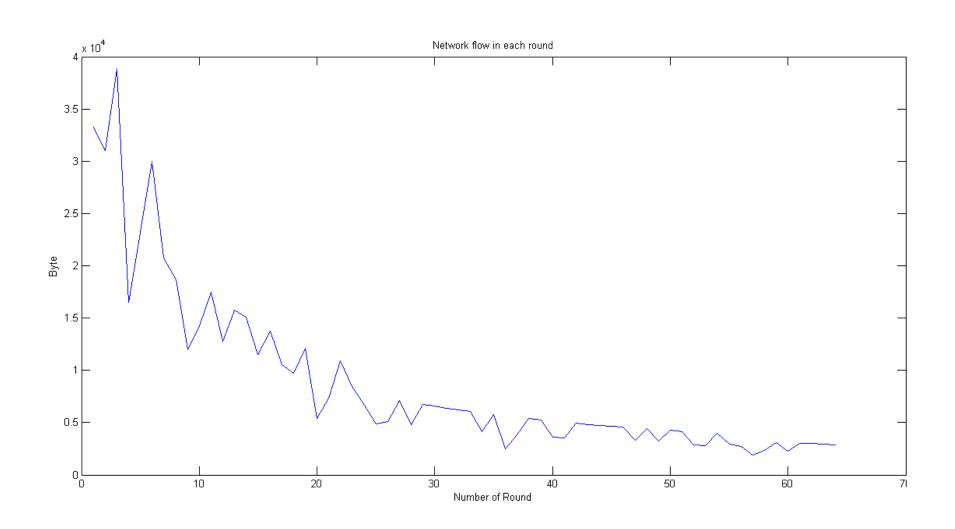
# **Longest Time**



# **Average Time**



#### **Network Bandwidth**



#### Conclusion

- The system can support dynamically adding and deleting computing nodes.
- It can support over 100,000,000 cells and 10,000,000 lives.
- With 100,000,000 cells, the system finishes about
  4.5 cycles / sec, better than our expectation.

Questions?