

Simplex Algorithm Overview

- Simplex algorithm is an algorithm to solve linear programming problems.
- The linear programming problems we are to solve have a special constraint on the decision variables requiring them to be non-negative.

Example

$$\text{maximize } 3x_1 + 2x_2 + 5x_3$$

$$\begin{array}{lcl} & 2x_1 + 3x_2 + 6x_3 & \geq 5 \\ s.t. & 7x_1 + 5x_2 + 2x_3 & \leq 1 \\ & 3x_1 + 4x_2 & = 2 \end{array}$$

$$\forall i \in \{1, 2, 3\}, x_i \geq 0$$

Transformation

$$\text{maximize } 3x_1 + 2x_2 + 5x_3$$

$$\begin{array}{rcl} 2x_1 + 3x_2 + 6x_3 - x_4 + a_1 & = & 5 \\ s.t. \quad 7x_1 + 5x_2 + 2x_3 + x_5 & = & 1 \\ 3x_1 + 4x_2 + a_2 & = & 2 \end{array}$$

The added x variables are called slack variables.

The added a variables are called auxiliary/artificial variables.

Notations

- \vec{c}_b : Coefficient vector of Basic Variables (Variables that have non zero value) in the Objective Function
- \vec{c}_n : Coefficient vector of Non-Basic Variables (Variables are zero) in the Objective Function

- B : Coefficient matrix of Basic Variables in the constraints
- N : Coefficient matrix of Non-Basic Variables in the constraints
- \vec{b} : Vector of the RHS values in the constraints
- \vec{x}_b : Vector of Basic Variables

- \vec{x}_n : Vector of Non-Basic Variables

Matrix Representation

$$\text{maximize } \vec{c}_b^T \cdot \vec{x}_b + \vec{c}_n^T \cdot \vec{x}_n$$

$$\text{s.t. } B\vec{x}_b + N\vec{x}_n = \vec{b}$$

The Objective function can also be written as:

$$c_b^T B^{-1} \vec{b} + (c_n^T - c_b^T B^{-1} N) \vec{x}_n$$

The goal of the simplex algorithm is to set all dimensions of the vector $c_n^T - c_b^T B^{-1} N$ to be negative by swapping Basic and Non-Basic Variables (pivot).

The objective value is $c_b^T B^{-1} \vec{b}$

Each pivot will “improve” the objective value

Two Phases

- If \exists Artificial Variable(s), “Phase One” is required
- The objective of “Phase One” is to maximize $-\sum_i a_i$
- If the optimal value of “Phase One” $\neq 0$, the linear program is infeasible; otherwise, it is feasible

Implementation

- Single phase solver – Finished
- “Phase One” builder – Work in Process

```
struct Model {  
    int is_max;  
    int stat;  
    int num_non_basic;  
    int num_basic;  
    int num_art;  
    double** b_matrix;  
    double** n_matrix;  
    double* cb_vector;  
    double* cn_vector;
```

```
int* xb_index_vector;  
int* xn_index_vector;  
double* b_vector;  
double** art_matrix;  
int* art_ind_vector;  
};
```

```
int pivot(Model* model, int in, int out);  
int check(Model* model, int* in, int* out);  
int ratio_check(Model* model, int in_ind,  
                int* ind);
```

Sample I/O

Input

3 3

1 10 3 5

2 4 -4 0 4

1 -20 -5 0 7

4 2 3 0 10

$$\textit{maximize } 10x_1 + 3x_2 + 5x_3$$

$$\begin{array}{lcl} & 2x_1 + 4x_2 - 4x_3 & \leq 5 \\ s.t. & 1x_1 - 20x_2 - 5x_3 & \leq 1 \\ & 4x_1 + 2x_2 + 3x_3 & \leq 2 \end{array}$$

Output

$x_1 = 2.363636$

$x_3 = 0.181818$

$x_2 = 0.000000$

$\text{opt_val} = 24.545455$

Output from Xpress IVE

2

Optimum found

$x_1 = 2.36364$

$x_2 = 0$

$x_3 = 0.181818$

24.5455

Input

3 3

1 10 3 5

2 4 -4 0 4

1 -20 -5 0 7

4 2 -3 0 10

$$\textit{maximize } 10x_1 + 3x_2 + 5x_3$$

$$\begin{array}{lcl} & 2x_1 + 4x_2 - 4x_3 & \leq 5 \\ s.t. & 1x_1 - 20x_2 - 5x_3 & \leq 1 \\ & 4x_1 + 2x_2 - 3x_3 & \leq 2 \end{array}$$

Output

Unbounded

Output from Xpress IVE

8

Unbounded

$$x_1 = 0$$

$$x_2 = 0$$

$$x_3 = 0$$

0