A formalization of Dedekind domains and class groups of global fields

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Introduction

Our project is the first formalization of several essential notions of algebraic number theory, in the Lean 3 prover as part of mathlib.

Goal: lay a useful foundation for theory-building.

Developing with mathlib means updating your code regularly in exchange for frequent new results and improvements.

Background: global fields

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Theorem: rings of integers are Dedekind domains.

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Principal fractional ideals $\left\langle \frac{a}{b} \right\rangle = \frac{a}{b}\mathcal{O}_K$ for $\frac{a}{b} \in K$ form a subgroup of the fractional ideals; the quotient is the class group $\mathcal{C}l_{\mathcal{O}_K}$.

Theorem: if $\mathcal{O}_K$ is a ring of integers, $\mathcal{C}l_{\mathcal{O}_K}$ is a finite abelian group.
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Theorem: A Dedekind domain is a UFD $\iff$ it is a PID $\iff$ $\mathcal{C}l_{\mathcal{O}_K}$ is trivial $\iff$ class number of $K = 1$.

Number fields; field extensions

mathlib typically uses typeclasses for algebraic structures, e.g.

```
class is_number_field (K : Type*) [field K] : Prop :=
[cz : char_zero K] [fd : finite_dimensional ℚ K]
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Typeclass inference automates the implications `char_zero K →` `algebra ℚ K → module ℚ K` required for `finite_dimensional ℚ K`.

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A tower $L/K/F$ is given by inclusions `[algebra F K]` `[algebra K L]` `[algebra F L]` and an instance `[is_scalar_tower F K L]` stating the maps commute.
Coherence proof obligations are automated through typeclass search.

Representing $\mathbb{Q}(\alpha)$

Number fiels have the form $\mathbb{Q}(\alpha)$, where α is algebraic:
the minimal polynomial $f_\alpha \in \mathbb{Q}[x]$ is irreducible and $f_\alpha(\alpha) = 0$.

Several constructions of $\mathbb{Q}(\alpha)$: subtype of $\mathbb{C}$, quotient type $\mathbb{Q}[x]/P$, ...
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We used the power basis: $\mathbb{Q}(\alpha)$ has a $\mathbb{Q}$-basis $1, \alpha, \ldots, \alpha^{n-1}$.
Theorem: each construction of $\mathbb{Q}(\alpha)$ is a field with power basis generated by α.

Dedekind domains

We defined Dedekind domains as integral domains R with an `is_dedekind_domain R` instance:

```
class is_dedekind_domain (R : Type*) [integral_domain R] :
    Prop :=
(to_is_noetherian_ring : is_noetherian_ring R)
(dimension_le_one : ∀ (P : ideal R), P ≠ ⊥ →
  is_prime P → is_maximal P)
(is_integrally_closed :
  integral_closure R (fraction_ring R) = ⊥)
```

Fractional ideals

We formalized fractional ideals of R as
R-submodules I of $\mathsf{Frac}(R)$ such that $\exists a : R, aI \subseteq R$.

Fractional ideals have a semiring structure (like submodules):

- $0 = \{0\}$
- $1 = \{x \mid x \in R\}$
- $I + J = \{x + y \mid x \in I, y \in J\}$
- $I * J$ is generated by $x * y$, $x \in I$, $y \in J$
- $x \in I/J \iff \forall y \in J, x * y \in I$

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- $x \in I/J \iff \forall y \in J, x * y \in I$

Theorem: $I * (1/I) = 1$ for all $I \neq 0$ iff R is a Dedekind domain.

Re-defining division

The `group_with_zero` typeclass used to define $x/y := x * y^{-1}$. For fractional ideals we want to define $I^{-1} := 1/I$. How to deal with this circularity?

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Solution: turn defeq into propositional equality by adding a new operation $(/)$ to `group` (`_with_zero`) and an axiom $x/y = x * y^{-1}$.

This required about 500 changes in mathlib.

Dedekind domain theorems

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Difficulties:

- Showing $x \in I * J$ implies $x = \sum_k y_k z_k$ for $y_k \in I$, $z_k \in J$.
- Coercions: I can be an integral ideal or set $\subseteq R$ or a fractional ideal or submodule or set $\subseteq \mathsf{Frac}(R)$.

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Theorem: Principal ideal domains are Dedekind domains.
Corollary: $\mathbb{Z}$ and $\mathbb{F}_q[t]$ are Dedekind domains.

Rings of integers are Dedekind domains

Theorem: The integral closure of a Dedekind domain R in a finite separable extension $K/\operatorname{Frac}(R)$ is a Dedekind domain.

Corollary: Rings of integers, closures of PIDs in finite separable extensions, are Dedekind domains.

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"Accidental" collaboration with the Berkeley Galois theory group:

- We defined `intermediate_field`
- They used it to formalize the primitive element theorem
- We used that to show finite separable field extensions have a power basis
- They used that to show conjugate roots of α correspond to images $\sigma(\alpha)$ for $\sigma : F(\alpha) \to K$ fixing F
- We used that to show the trace form is nondegenerate

Finiteness of the class group

Theorem: If K is a global field, the class group of $\mathcal{O}_K$ is finite.

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We introduced a new notion of admissible absolute value,
and proved if abs : $R \to \mathbb{Z}$ is admissible,
this intermediate step in the classical proof holds:

```
theorem exists_mem_finset_approx'
  (a b : integral_closure R L) :=
  ∃ (q : integral_closure R L) (r ∈ finset_approx L f abs),
  abs_norm f abs (r • a - q * b) < abs_norm f abs b
```

Finiteness of the class group

After formalizing the remainder of the classical proof, it remained to find admissible absolute values.

For $\mathbb{Z}$, the usual absolute value is admissible.
For $\mathbb{F}_q[t]$, $|f|_{\deg} := q^{\deg f}$ is admissible.

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```
def class_group (f : fraction_map R K) :=
quotient_group.quotient (to_principal_ideal f).range

noncomputable def number_field.class_number (K : Type*)
  [field K] [is_number_field K] : ℕ :=
card (class_group (ring_of_integers.fraction_map K))
```

Conclusions

Total contribution: $\pm$ 5000 lines of project-specific code, $\pm$ 2500 lines background work.
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Rules of thumb for our work:

- Parametrize (`is_scalar_tower`, `power_basis`, ...) instead of choosing a canonical construction.
- Refactoring allows deep integration between different viewpoints.
- Contribute quickly and often, so others do your work for you.