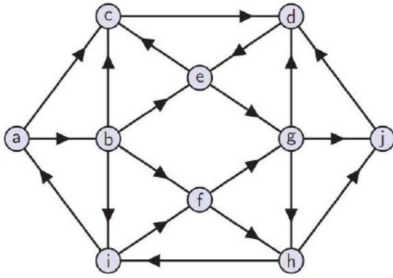


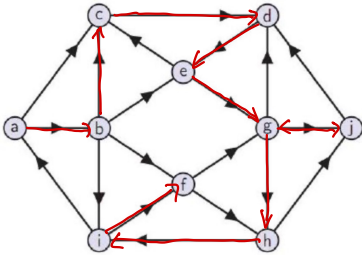
# GNT EXERCISES

## Task 03



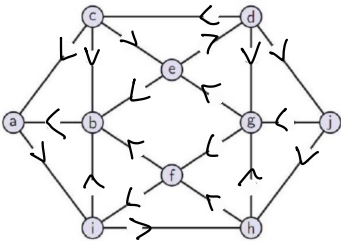
Is the graph in the figure strongly connected? Check using the algorithm from the presentation.

1) DFS (G) and store in order of visiting.



$L = \{A, B, C, D, E, G, J, H, I, F\}$

2) Compute  $G^T$

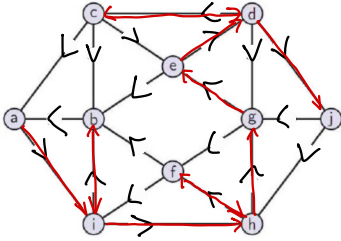


### 3) DFS( $G^T$ )

3.1) If we reach a cycle, we have a SCC

3.2) Pop from the stack  $L$  the elements from SCC:  $L = L \setminus \text{SCC}$

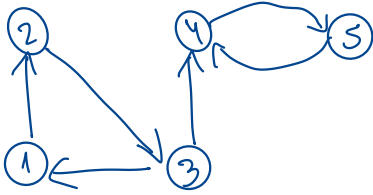
3.3) Repeat until  $L = \emptyset$ , over  $G^T = G^T \setminus \text{SCC}$



$$L = \{A, B, C, D, E, G, J, H, I, F\}$$

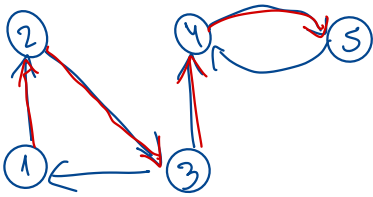
$$\text{SCC}_1 = \{A, I, B, H, F, G, E, D, C, J\}$$

We can see  $G$  is strongly connected, as we found 1 SCC.



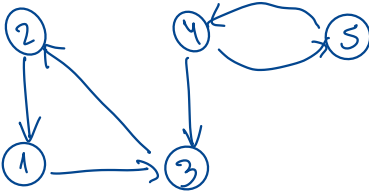
check if the graph is strongly connected.

1)

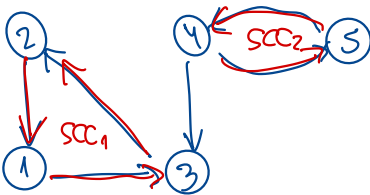


$$L = \{1, 2, 3, 4, 5\}$$

2)



3)



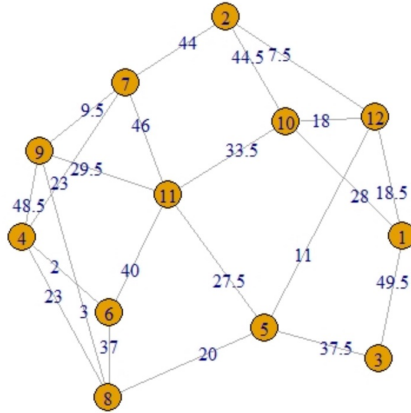
$$L = \{1, 2, 3, 4, 5\}$$

$$SCC_1 = \{1, 2, 3\}$$

$$SCC_2 = \{4, 5\}$$

## Task 04

Find the MST in the graph. Use two methods.



**Prim:** When done in paper, the algorithm can be formulated the following way.

$\text{Prim}(G, w, r) \{$

$S = V, A = \emptyset$

$A \cup \{e\}$ , where  $e$  is the edge of less weight from the root

$S = S \setminus \{r\}$

$\text{while}(S \neq \emptyset) \{$

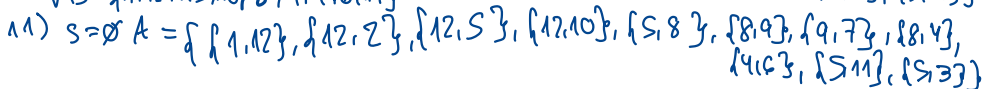
$\forall v \in V$  take edge of less weight  $e$  and mark it if it's not marked, and if it doesn't form a cycle.

$u = \text{vertex } x \text{ incident with } e \text{ in the path}$

$A \cup \{e\}$

$S = S \setminus \{u\}$

$\}$



# Kruskal:

$\text{Kruskal}(G, w)$

$A = \emptyset$

$\forall v \in V, \text{MakeSet}(v)$

Sort edges by weight in non-decreasing order in a queue  $Q$ .

while ( $Q \neq \emptyset$ )

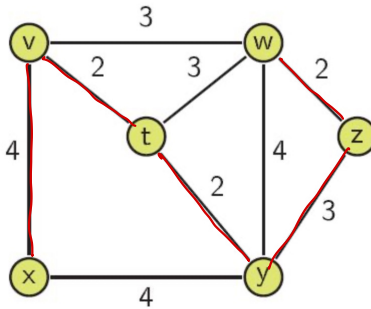
$\{u, v\} = Q.\text{pop}$

    if ( $u \notin \text{Set}(v)$ )

$A = A \cup \{u, v\}$

    } Union( $\text{Set}(u), \text{Set}(v)$ )

}



$Q = \{ \{v, t\}, \{t, y\}, \{w, z\}, \{y, z\}, \{t, w\}, \{v, w\}, \{x, v\}, \{x, y\}, \{y, w\} \}$

$\{v\} \{t\} \{y\} \{w\} \{z\} \{x\} \quad A = \emptyset$

1)  $\{v, t\} \{y\} \{w\} \{z\} \{x\}$

$A = \{ \{v, t\} \}$

2)  $\{v, t, y\} \{w\} \{z\} \{x\}$

$A = \{ \{v, t\}, \{t, y\} \}$

3)  $\{v, t, y\} \{w, z\} \{x\}$

$A = \{ \{v, t\}, \{t, y\}, \{w, z\} \}$

4)  $\{v, t, y, w, z\} \{x\}$

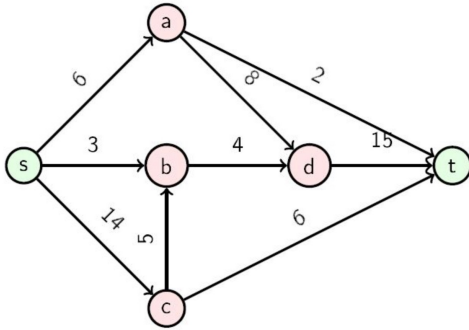
$A = \{ \{v, t\}, \{t, y\}, \{w, z\}, \{y, z\} \}$

5)  $\{v, t, y, w, z, x\}$

$A = \{ \{v, t\}, \{t, y\}, \{w, z\}, \{y, z\}, \{x, y\} \}$

### Task 09

For the network, find the maximum flow function and the value of the maximum flow, from vertex  $s$  to  $t$ .



Ford-Fulkerson ( $G, s, t$ ) {

$\forall (u, v) \in E, f(u, v) = 0$

Build residual network  $G_f$

while ( $\exists$  augmenting path  $p$  from  $s$  to  $t$  in  $G_f$ ) {

$C_f(p) = \min \{ C_f(u, v) \mid (u, v) \in p \}$

for  $(u, v) \in p$  {

if  $(u, v) \in E$   $f(u, v) = f(u, v) + C_f(p)$

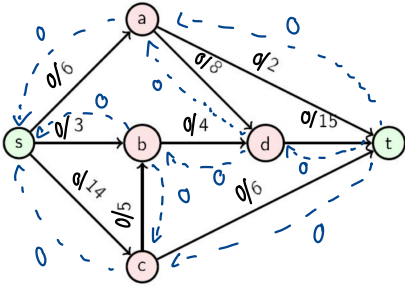
} else  $f(v, u) = f(v, u) - C_f(p)$

}

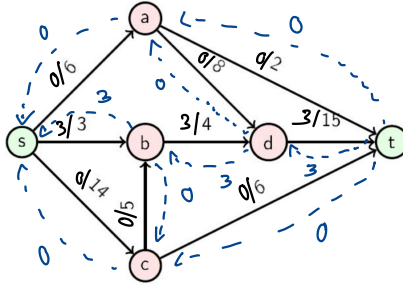
return  $\sum_p C_f(p) = \sum_{v \in V} f(s, v)$

}

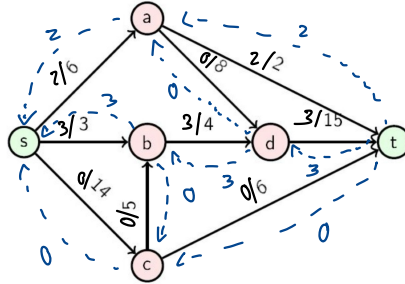
$G_8$



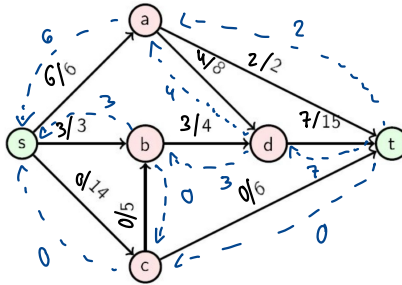
1)  $P = (s, b, d, t)$   
 $C_8(P) = 3$



2)  
 $P = (s, a, t)$   
 $C_8(P) = 2$



3)  
 $P = (s, a, d, t)$   
 $C_8(P) = 4$

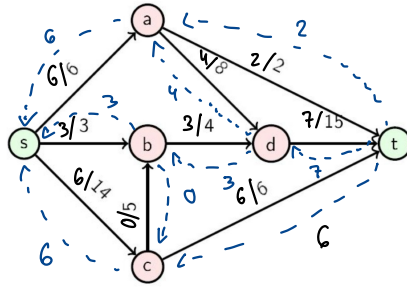




4)

$$P = (s, c, t)$$

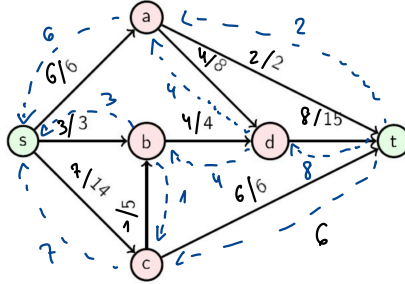
$$C_g(P) = 6$$



5)

$$P = (s, c, b, d, t)$$

$$C_g(P) = 1$$



We can see there are no more augmenting paths  $\Rightarrow$

$$|f| = \sum_P C_g(P) = 3 + 2 + 4 + 6 + 1 = 16$$