

# Faculty of Mathematics and Computer Science

The 
$$(n^2 - 1)$$
-Puzzle

Artificial Inteligence in Computer Games Final Project

Double Degree in Mathematics and Computer Science (University of Granada, Spain)

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## 1 Goals

- Apply the mathematical Group Theory to an AI problem.
- $\bullet\,$  Get familiar with working with data trees and complex data structures.
- ullet Offer a game-solving algorithm for the famous 15-Puzzle and generalize it to any board size.

## 2 Introduction

The construction of an algorithm which is able to solve the famous 15-Game is proposed.

The 15-Puzzle is a board game which consists of 15 numbered tiles and an empty one. The goal is to put every tile in order by sliding the tiles into the empty space.

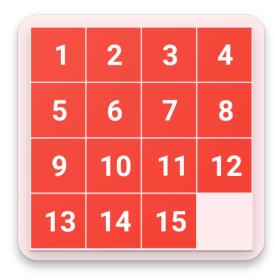


Figure 1: Example of a solved 15-Puzzle.

## 3 Mathematical Analysis

Let  $n \in \mathbb{N}$ ,  $n \geq 2$ . Our goal is to provide an algorithm which solves the  $(n^2 - 1)$ -Puzzle.

The first thing we need to consider is solvability. Can every situation on the board lead to a solution?

**Theorem 1** (Johnson and Story (1879)). On boards of size  $m \times n$ , where  $m, n \geq 2$ , all even permutations are solvable.

It can be proven by induction on m and n, starting with m = n = 2. This theorem proves that every situation on the board can be generated by 3-cycles.

**Corollary 1.** A board of size  $n^2$  can be represented by the Alternating Group  $A_{n^2-1}$ . This means that the total possible states which lead to a solution are reduced to the number  $|A_{n^2-1}| = \frac{(n^2-1)!}{2}$ 

On the other hand, we can also consider which is the solution with the fewest number of slides. This is, in fact, an NP-Hard problem, although there is a polynomial-time constant-factor approximation.

It is known that the 15-Game can be solved in no more than 80 moves, while the 8-Game in no more than 31.

15 11 10 14 15 11	10 13 11 9 13 14 9 12 15 10 14 2 5 3 7 6 2	15 9 13 11 12 10 14 3 7 6 2	12 9 13 15 11 10 14 3 7 6 2	12 14 13 15 11 9 10 3 7 6 2
4 8 6 1 4 8		4 8 5 1	4 8 5 1	4 8 5 1
	11 13 12 10 13	12 9 13 15 11 14 10	12 9 13	12 14 13
3 7 6 2 3 7	10 9 15 11 9 14 6 2 7 3 6 2	15 11 14 10 3 8 6 2	15 11 10 14 8 3 6 2	15 11 9 10 8 3 6 2
485148	5 1 4 8 5 1	4 7 5 1	4 7 5 1	4 7 5 1
	10 13 12 9 13	12 9 13	12 9 13	
intelligent and intelligent to the second	14 9 15 8 10 14	15 11 10 14	15 11 10 14	
	6 2 11 7 6 2 5 1 4 3 5 1	3 7 5 6 4 8 2 1	7 8 5 6 4 3 2 1	

Figure 2: The 17 positions of the 15-Puzzle which need 80 moves.

## 4 Computational Analysis

### 4.1 Board Generation

We first need to think how to generate a solvable situation for our game player. Knowing that every solvable distribution of the board is generated by 3-cycles, it is intelligent to start from the puzzle's idle position and apply any number of 3-cycles to shuffle it, obtaining a solvable distribution.

### 4.2 Algorithm Selection

On the other hand, it is necessary to think about what algorithm solves the problem in a more efficient way. However, we have to know that dealing with permutations and exploring a tree of possible solutions frequently involve the use of non-polynomial algorithms.

In this case, an easy option for our problem would be the A\* Algorithm.

### 4.2.1 Properties

- Completeness: If a solution exists, A\* will find it.
- Optimality: A\* always finds the optimal solution regarding cost. In our case, it finds the solution with the least number of movements, although not in a polynomial time.
- It uses a function to calculate costs: Let f(n) = g(n) + h(n), where g is a real-cost function, and h an heuristic one, which estimates the cost to a possible solution.
- In this case, both time and space complexities are factorial in the worst case, as we are dealing with permutations.

#### 4.2.2 Structures

- Explored Set: Nodes which have already been visited. It is frequent to use either a set, an unordered set or an unordered map. For this example, a simple set gives better results.
- Frontier Set: Nodes which are to be explored. The structure used is a priority queue which sorts based on cost.

### 4.2.3 Heuristic Selection

The heuristic function determines the number of nodes which are expanded. As a result, the selection of this function is an important step in the process. Heuristics also have some characteristics. Let h be an heuristic function.

**Definition 1.** We say that h is admissible  $\iff$ 

$$h(n) \le h^*(n) \ \forall n \ node$$

where  $h^*(n)$  is the true cost to reach node n. In other words, the heuristic does not overestimate costs.

**Definition 2.** We say that h is monotonic (or consistent)  $\iff$ 

$$h(n) - h(n') \le c(n, n') \ \forall n, n' \ nodes$$

where c(n, n') is the cost of going from node n to node n'. This guarantees the optimal solution is found without the need to update costs and search in the frontier set.

**Definition 3.** Let  $h_1$  and  $h_2$  be heuristics. We say that  $h_2$  dominates  $h_1 \iff$ 

$$h_2(n) \ge h_1(n) \ \forall n \ node$$

In practice, this means  $h_1$  will expand, at least, the same number of nodes as  $h_2$ .

**Proposition 1.** Monotonicity  $\implies$  Admissibility.

**Theorem 2.** Let h be an admissible heuristic used by the  $A^*$  Algorithm. Then, the  $A^*$  Algorithm is complete and optimal regarding cost.

Knowing this, we can now decide which heuristic to use.

**Manhattan Distance** One really appropriate heuristic for this problem is the Manhattan Distance, which is defined as:

$$d_{\text{Manhattan}} = \sum_{i=1}^{n} \sum_{j=1}^{n} |x_i^j - y_i^j|$$

where  $x_i^j$  is the tile i, j on the board and  $y_i^j$  is the position where  $x_i^j$  should be.

It will measure how far each tile is from its correct position, and then add all these values.

**Inverse Distance** This heuristic builds upon linear conflicts and uses the idea of inversions.

We define an inversion to be when a tile appears before another tile with a smaller number. The blank has no number and cannot contribute to inversions.

#### Considerations:

- When moving a tile horizontally, the total number of inversions never changes. This is due to the blank not affecting inversions.
- When moving a tile vertically, the total number of inversions can change by only -3, -1, +1, and +3.
  - Case 1: the three skipped tiles are all smaller (or larger) than the moved tile. Moving the tile will either add or fix three inversions, one for each skipped tile. So, the total number of inversions changes by +3 or -3.
  - Case 2: two of the tiles are larger and other is smaller (or vice versa). In this case, there's going to be a net change of +1 or -1 inversions.
- One vertical move can fix at most three inversions. If we assume the minimum number of vertical moves needed to fix the inversions, that results in  $\lfloor n\_inversions/3 \rfloor$ . If there is a remainder, the remaining inversions can be solved with at least one vertical move per remaining inversion.

As a result, going through the board left-to-right, top-to-bottom, we can define

$$vertical = \lfloor \frac{n\_inversions}{3} \rfloor + n\_inversions\%3$$

Analogously, we can define the horizontal count of inversion, but noticing we will go through the board top-to-bottom, left-to-right.

Finally, the heuristic is defined as:

$$d_{\text{Inverse}} = \text{vertical} + \text{horizontal}$$

Walking Distance We now define the walking distance as the sum of the Manhattan Distance and the Inverse Distance.

$$d_{\text{Walking}} = d_{\text{Manhattan}} + \alpha \cdot d_{\text{Inverse}}$$

where  $\alpha$  is a weighting factor.

**Proposition 2.** The Manhattan Distance for the  $(n^2 - 1)$ -Puzzle problem is a monotonic heuristic.

**Proposition 3.** The Inverse Distance for the  $(n^2 - 1)$ -Puzzle problem is an admissible heuristic.

**Proposition 4.** The Walking Distance for the  $(n^2 - 1)$ -Puzzle problem is an admissible heuristic for an appropriate value of  $\alpha$ , which dominates the Manhattan Distance and the Inverse Distance Heuristics. In particular, for  $\alpha = 2$ , this is true.

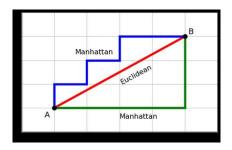


Figure 3: Comparison between the Euclidean and Manhattan distances.

## 5 Implementation

The implementation of the game has been done in C++11.

In this section, every function is defined in its respective header file. All these definitions have been omitted to avoid redundancy. Only the most relevant aspects of these headers are shown.

### 5.1 Generator

```
//generator.h
#ifndef GENERATOR_H
#define GENERATOR_H

#include <vector>
#include <iostream>

using namespace std;

const int N_PERMUTATIONS = 100; ///< Number of permutations to generate

const int LOWER_BOUND = 2; ///< Lower bound for board size

const int UPPER_BOUND = 4; ///< Upper bound for board size

#endif // GENERATOR_H</pre>
```

```
1 //generator.cpp
# include "../include/generator.h"
4 using namespace std;
_{7} * Obrief Function to get a random number from the vector
     without repetition.
_{\rm 8} * Cparam v The vector from which to get the random
    number
9 * Oreturn A random number from the vector
int getNum(vector<int>& v) {
      // Size of the vector
13
      int n = v.size();
14
15
      // Generate a random number
16
17
      // Make sure the number is within the index range
18
      int index = rand() % n;
19
20
      // Get random number from the vector
21
      int num = v[index];
22
23
      // Remove the number from the vector
      swap(v[index], v[n-1]);
25
      v.pop_back();
26
27
      // Return the removed number
28
      return num;
29
30 }
```

```
1 /**
_{2} * @brief Function to generate a random vector of n
    different numbers.
_{3} * Oparam n The number of different numbers to generate
* @return A vector of n different random numbers
5 */
6 vector < int > generateRandom(int n) {
      vector < int > v(n-1);
      // Fill the vector with the values
9
      // 1, 2, 3, ..., n-1
10
      for (int i = 1; i < n; i++) v[i-1] = i;
11
12
      vector<int> u;
      for(int i = 0; i < 3; i++) u.push_back(getNum(v));</pre>
15
     return u;
17
18 }
```

```
1 /**
  * Obrief Function to generate a permutated matrix.
   * Oparam n The size of the matrix (n x n)
   * @return A permutated matrix of size n x n
6 vector < vector < int >> generateBoard(int n) {
      srand(time(NULL));
8
      int n_squared = n*n;
10
      vector < int > permutation(n_squared);
11
12
      //Fill the board with the numbers from 0 to n^2-1
13
      for(int i = 0; i < n_squared; i++) permutation[i] =</pre>
14
     i;
      //Do the permutations
16
      for(int i = 0; i < N_PERMUTATIONS; i++){</pre>
17
18
           vector<int> swapper = generateRandom(n_squared);
19
           int swap = permutation[swapper[0]];
20
21
           permutation[swapper[0]] = permutation[swapper
22
     [1]];
           permutation[swapper[1]] = swap;
23
24
           swap = permutation[swapper[2]];
25
           permutation[swapper[2]] = permutation[swapper
26
     [0]];
           permutation[swapper[0]] = swap;
27
      }
28
      //Constructing the new permutated matrix
30
      int index = 0;
31
      vector < vector < int >> matrix(n, vector < int >(n));
32
      for(int i = 0; i < n; i++){</pre>
33
           for (int j = 0; j < n; j++) {
34
               matrix[i][j] = permutation[index++];
35
           }
36
      }
37
38
      return matrix;
39
40 }
```

```
/**
2 * @brief Function to show a matrix.
3 * @param matrix The matrix to show
4 */
5 void showMatrix(const vector<vector<int>>& matrix){
6    int n=matrix.size();
7    for(int i = 0; i < n; i++){
8        for(int j = 0; j < n; j++){
9            cout << matrix[i][j] << "\t";
10        }
11        cout << endl << endl;
12    }
13 }</pre>
```

### 5.2 Solver

```
//A_star.h
      #ifndef A_STAR_H
      #define A_STAR_H
3
      #include <queue>
5
      #include <vector>
6
      #include <list>
     #include <unordered_set>
      #include <set>
9
      #include <cmath>
10
      #include <string>
11
12
     using namespace std;
14 /**
_{\rm 15} * Obrief Enum to represent the possible moves in the
    puzzle.
17 enum class Action {
      UP,
            ///< Move the empty tile down
      DOWN, ///< Move the empty tile up
      LEFT,
            ///< Move the empty tile right
20
      RIGHT ///< Move the empty tile left
21
22 };
```

```
1 /**
2 * @brief Node structure to store the state of the board,
     the cost, and the moves made.
4 struct node {
5
      vector<vector<int>> board; ///< The current state of</pre>
6
      the board
      int cost; ///< The real cost to reach this node</pre>
      pair<int, int> pos_0; ///< The position of the empty</pre>
8
      tile on the board
      list < Action > moves; /// < The list of moves made to
9
     reach this node
      /**
11
      * @brief Default constructor for the node.
12
      */
13
      node() {
14
          this -> cost = 0;
15
      }
16
17
      /**
18
      * @brief Parameterized constructor for the node.
19
      \ast @param board The state of the board
20
      * Oparam cost The cost to reach this node
21
      * @param pos_0 The position of the empty tile
      * @param moves The list of moves made to reach this
23
     node
24
      node(const vector<vector<int>>& board, int cost,
25
     const pair<int, int>& pos_0, const list<Action>&
     moves) {
          this->board = board;
26
          this->cost = cost;
27
          this \rightarrow pos_0 = pos_0;
28
          this->moves = moves;
29
      }
30
31
32
33
34
35
36
```

```
37
      * Obrief Copy constructor for the node.
38
      * Oparam other The node to copy from
40
      node(const node& other) {*this=other;}
41
42
      /**
43
      * Obrief Overloading the assignment operator to copy
44
      the values of a node.
      * @param other The node to copy from
45
      * Oreturn A reference to the assigned node
46
47
      node& operator=(const node& other) {
48
          if (this != &other) {
49
              this->board = other.board;
               this->cost = other.cost;
               this->pos_0 = other.pos_0;
52
               this->moves = other.moves;
          }
54
          return *this;
      }
56
57
      /**
58
      * Obrief Overloading the equality operator to
59
     compare two nodes.
      * Oparam other The node to compare with
      * Oreturn True if the boards are equal, false
61
     otherwise
62
      bool operator == (const node& other) const {return
63
     this->board == other.board;}
      /**
65
      * @brief Overloading the less-than operator to
66
     compare two nodes.
      * Oparam other The node to compare with
67
      * Oreturn True if this node is less than the other
68
     node, false otherwise
      */
69
      bool operator < (const node& other) const {return this
     ->board < other.board;}
71 };
```

```
* @brief Functor for the priority queue to compare
   nodes based on their cost.
4 class Comparer {
5 public:
     /**
      * @brief Comparison operator to compare two nodes.
      * @param a The first node
      * @param b The second node
      * Oreturn True if the cost of node a is greater
10
    than the cost of node b, false otherwise
      */
11
    bool operator()(const node& a, const node& b) {
        return a.cost + WalkingDistanceHeuristic(a.board
    ) > b.cost + WalkingDistanceHeuristic(b.board);
14
15 };
```

The following code is the implementation of a hash function in case we want to use an unordered set to manage the explored set of nodes. However, in this case, the explored set is managed with a normal set structure.

```
* @brief Hash function for the unordered_set to hash
     nodes.
  class Hash {
      public:
6
           * Obrief Hash function to generate a hash value
      for a node.
           * Oparam node The node to hash
           * Oreturn The hash value of the node
           */
10
          size_t operator()(const node& node) const {
               size_t hashValue = 0;
12
               hash<int> hashInt; // Hash Function for
13
     integers
14
               for (int i = 0; i < node.board.size(); i++)</pre>
     {
                   for (int j = 0; j < node.board[i].size()</pre>
16
     ; j++) {
                       hashValue ^= hashInt(node.board[i][j
17
     ]) + 0x9e3779b9 + (hashValue << 6) + (hashValue >> 2)
                   }
18
19
               return hashValue;
20
          }
#endif // A_STAR_H
```

```
1 //A_star.cpp
#include "../include/A_star.h"
4 using namespace std;
6 /**
_{7} * @brief Function to calculate the Manhattan distance
    heuristic.
  * Oparam board The state of the board
  * @return The Manhattan distance heuristic value
10 */
11 int ManhattanDistanceHeuristic(const vector<vector<int
    >>& board){
12
     int cost = 0;
13
   int n = board.size();
14
      for(int i=0; i<n; i++){</pre>
          for(int j=0; j<n; j++){</pre>
16
              if(board[i][j] != 0)
17
                   cost += abs(i - board[i][j]/n) + abs(j -
18
      board[i][j]%n);
          }
19
      }
20
      return cost;
21
22 }
```

```
1 /**
* @brief Function to calculate the Inversion distance
    heuristic.
* @param board The state of the board
* @return The Inversion distance heuristic value
6 int InversionDistanceHeuristic(const vector<vector<int
    >>& board){
     int n= board.size();
      int inversion_count = 0;
      for(int i=0; i<n; i++){</pre>
10
          for(int j=0; j<n-1; j++){</pre>
11
              if(board[i][j] != 0){
                  if(board[i][j] > board[i][j+1])
13
     inversion_count++;
                  if(board[j][i] > board[j+1][i])
14
    inversion_count++;
              }
15
          }
      }
17
      return inversion_count/3 + inversion_count%3;
19
20 }
* @brief Function to calculate the Walking distance
    heuristic.
* @param board The state of the board
* @return The Walking distance heuristic value
5 */
6 double WalkingDistanceHeuristic(const vector<vector<int
    >>& board, double weight){
   return ManhattanDistanceHeuristic(board) + weight*
    InversionDistanceHeuristic(board);
9 }
```

```
1 /**
2 * @brief Function to find the empty tile on the board.
  * Oparam board The state of the board
* @return The position of the empty tile
6 pair < int , int > findEmptyTile(const vector < vector < int >>&
     board){
    int n = board.size();
   for(int i=0; i<n; i++){</pre>
9
      for(int j=0; j<n; j++){</pre>
10
        if(board[i][j] == 0)
11
          return pair < int , int > (i, j);
12
      }
13
    }
14
15 }
1 /**
2 * Obrief Function to determine if a board is a solution
  * Oparam board The state of the board
   * Oreturn True if the board is a solution, false
    otherwise
  */
6 bool isSolution(const vector<vector<int>>& board){
      int n = board.size();
8
      bool isSolution = true;
10
      for(int i=0; i<n && isSolution; i++){</pre>
11
           for(int j=0; j<n && isSolution; j++){</pre>
12
               if(board[i][j] != i*board.size() + j)
13
                   isSolution = false;
14
          }
      }
16
17
      return isSolution;
19 }
```

```
1 /**
   * @brief Function to permutate the tiles of the board.
   * Oparam board The state of the board
   * @param pos_0 The position of the empty tile
  * @param action The action to perform
   * @return The new position of the empty tile
8 pair < int , int > permutateBoard(vector < vector < int >> & board ,
      const pair < int , int > % pos_0, const Action & action) {
9
    int i = pos_0.first;
10
    int j = pos_0.second;
11
    switch(action){
13
      case Action::DOWN:
        if(i!=0){
           swap(board[i][j], board[i-1][j]);
16
           i=i-1:
17
        }
18
        break;
19
      case Action::UP:
20
        if (i!=board.size()-1){
21
           swap(board[i][j], board[i+1][j]);
22
           i=i+1;
        }
24
        break;
      case Action::RIGHT:
26
        if(j!=0){
27
           swap(board[i][j], board[i][j-1]);
28
           j = j - 1;
29
        }
30
        break;
      case Action::LEFT:
32
        if(j!=board.size()-1){
33
           swap(board[i][j], board[i][j+1]);
34
           j=j+1;
35
        }
36
        break;
37
    }
38
39
    return pair<int,int>(i,j);
40
41 }
```

```
1 /**
2 * @brief Function to generate a child node by applying
    an action to the current node.
3 * @param current_node The current node
* @param action The action to apply
* @return The generated child node
7 node Apply(const node& n, const Action& action){
      vector < vector < int >> board = n.board;
9
      int cost = n.cost;
10
      pair < int , int > pos_0 = n.pos_0;
11
     int i = pos_0.first;
12
      int j = pos_0.second;
      list < Action > moves = n.moves;
     if ((action == Action::UP && i != board.size() - 1)
     \Pi
          (action == Action::DOWN && i != 0) ||
17
          (action == Action::LEFT && j != board.size() -
18
     1) ||
          (action == Action::RIGHT && j != 0)) {
19
          pos_0 = permutateBoard(board, pos_0, action);
20
          cost++;
21
          moves.push_back(action);
22
      }
      return node(board, cost, pos_0, moves);
25
26 }
```

```
* Obrief A* algorithm to solve the puzzle.
   * @param board The initial state of the board
  * @return The list of actions to solve the puzzle
6 list<Action> A_star(const vector<vector<int>>& board){
      node current_node(board, 0, findEmptyTile(board),
8
     list < Action > ());
      priority_queue < node, vector < node > , Comparer > frontier;
9
      set < node > explored;
    list < Action > moves;
11
    bool solutionFound = isSolution(current_node.board);
13
    frontier.push(current_node);
14
15
    while(!solutionFound && !frontier.empty()){
16
      frontier.pop();
17
      explored.insert(current_node);
18
19
      if(isSolution(current_node.board)) solutionFound =
20
     true;
21
      if (!solutionFound) {
        // Generate child UP
23
        node child_up = Apply(current_node, Action::UP);
        if(explored.find(child_up) == explored.end())
25
     frontier.push(child_up);
26
27
      if (!solutionFound) {
28
        // Generate child DOWN
        node child_down = Apply(current_node, Action::DOWN
30
     );
        if(explored.find(child_down) == explored.end())
31
     frontier.push(child_down);
      }
32
      if(!solutionFound){
34
        // Generate child LEFT
35
        node child_left = Apply(current_node, Action::LEFT
36
     );
37
```

```
if(explored.find(child_left) == explored.end())
38
     frontier.push(child_left);
      }
39
40
      if(!solutionFound){
41
        // Generate child RIGHT
42
        node child_right = Apply(current_node, Action::
43
     RIGHT);
        if(explored.find(child_right) == explored.end())
44
     frontier.push(child_right);
      }
45
46
      if (!solutionFound and !frontier.empty()){
47
        current_node = frontier.top();
        while(!frontier.empty() && explored.find(current_
     node) != explored.end()){
          frontier.pop();
50
          if(!frontier.empty())
51
             current _node=frontier.top();
        }
53
      }
55
      if(solutionFound) moves=current_node.moves;
56
57
58
      return moves;
60 }
```

```
1 /**
* @brief Function to convert the string to actions.
* Oparam action The string to convert
_4 * @return The action representation of the string
6 Action stringToAction(char action) {
     switch (action) {
          case 'w': return Action::UP;
          case 's': return Action::DOWN;
          case 'a': return Action::LEFT;
          case 'd': return Action::RIGHT;
     }
12
13 }
* @brief Function to convert the actions to string.
* @param action The action to convert
  * Oreturn The string representation of the action
6 string actionToString(const Action& action) {
     switch (action) {
          case Action::UP:
                              return "UP";
          case Action::DOWN: return "DOWN";
         case Action::LEFT: return "LEFT";
10
         case Action::RIGHT: return "RIGHT";
11
     }
12
13 }
```

### 5.3 Main

```
#include "../include/generator.h"
#include "../include/A_star.h"
using namespace std;
6 int main() {
      int n; //Length of the board
9
      //Presentation
      cout << "Welcome to the (n^2-1)-puzzle solver!!!!!"</pre>
11
     << endl;
      do{
13
           cout << "Introduce a board size between " <<</pre>
14
     LOWER_BOUND << " and " << UPPER_BOUND << ": ";
          cin >> n;
15
      }while(n<LOWER_BOUND || n>UPPER_BOUND);
16
17
      //Generate the board
19
      cout << "You chose the board size: " << n << "x" <<
20
     n \ll endl;
      if(n==4) cout << "You will play the famous 15-puzzle</pre>
21
     !!!!" << endl;
      cout << "Your puzzle to solve is: " << endl << endl;</pre>
22
23
24
      vector < vector < int >> board;
25
      //Generate a board that is not a solution
      do{
           board = generateBoard(n);
29
      }while (isSolution(board));
30
31
      pair < int , int > pos_0 = findEmptyTile(board);
32
      showMatrix(board);
34
35
36
37
```

```
int option;
38
      do{
39
           cout << "Press 0 to solve the puzzle by yourself</pre>
      or 1 to solve it automatically: ";
           cin >> option;
41
      }while(option!=0 && option!=1);
42
       cout << endl;
43
44
      //Manual Solver
      if(option == 0){
46
           char move;
48
49
           do{
50
                cout << "You chose to solve the puzzle by</pre>
     yourself!!!!" << endl;</pre>
                cout << "Use the following keys to move the</pre>
     empty tile:" << endl;</pre>
                cout << "W: Move the empty tile down" <<</pre>
     endl;
                cout << "S: Move the empty tile up" << endl;</pre>
                cout << "A: Move the empty tile right" <<</pre>
     endl;
                cout << "D: Move the empty tile left" <<</pre>
56
     endl;
                cout << "Q: Quit the game" << endl;</pre>
                cout << endl;</pre>
                showMatrix(board);
59
60
                cout << "Select a key: ";</pre>
61
                cin >> move;
                cout << endl;</pre>
63
                if (move == 'w' || move == 'a' || move == 's' ||
65
     move == 'd') {
                     Action move_s=stringToAction(move);
66
                    pos_0 = permutateBoard(board, pos_0,
67
     move_s);
                     actionToString(move_s);
                }
69
                else if (move!='q') cout << "Select a valid</pre>
     key!!!!" << endl << endl;
           }while(!isSolution(board) && move != 'q');
```

```
72
           //Not playing anymore
73
           if(move=='q') cout << "You quit the game!!!!" <<</pre>
      endl;
           else cout << "The puzzle has been solved!!!!" <<</pre>
75
      endl;
      }
76
77
      //Automatic Solver
       else{
79
80
           cout << "You chose to solve the puzzle</pre>
81
     automatically!!!!" << endl;</pre>
           cout << endl << "Solving the puzzle..." << endl;</pre>
82
83
           list < Action > moves = A_star(board);
85
           cout << "The solution is: " << endl << endl;</pre>
86
87
           showMatrix(board);
88
           for(auto it = moves.begin(); it != moves.end();
90
     it++){
               cout << actionToString(*it) << endl;</pre>
91
               pos_0 = permutateBoard(board, pos_0, *it);
92
               showMatrix(board);
93
           }
94
95
           cout << endl << "The puzzle has been solved!!!!"</pre>
96
      << endl;
      }
97
      cout << "
99
     " << endl:
100
      return 0;
101
102
103 }
```

## 6 Conclusions

- How different heuristics can alter the execution time of an algorithm.
- How complex it is to provide a good solution to a problem based on heuristics.
- The need to get familiar with data structures and their particularities.

### References

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