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1. Why does the Offset field in the IP header measure the offset in 8-byte units? (Hint: Recall that the Offset field is 13 bits long.)

The Offset field is 13 bits long and maximum IP packet size is 2^16 bytes length. So, offset field can be measured in $2^16 - 13 = 8$ byte units.

2. Suppose a TCP message that contains 1024 bytes of data and 20 bytes of TCP header is passed to IP for delivery across two networks interconnected by a router (i.e., it travels from the source host to a router to the destination host). The first network has an MTU of 1020 bytes; the second has an MTU of 572 bytes. Each network's MTU gives the size of the largest IP datagram that can be carried in a link-layer frame. Give the sizes and offsets of the sequence of fragments delivered to the network layer at the destination host. Assume all IP headers are 20 bytes.

In the first router, the IP payload is 1024 + 20 = 1044 bytes. Maximum size of data allowed in first router is 1020 - 20 = 1000 bytes. The TCP payload can be fragmented as 1000, 48 bytes.

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1st fragment: 20 bytes header -1000 bytes of payload - offset 0 2nd fragment: 20 bytes header -44 bytes of payload - offset 1000/8 = 125
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In the second router, the maximum size of data allowed is 572 - 20 = 552 bytes. The second fragment from first router will not fragmented and will be delivered as such. The first fragment will be fragmented as 552, 448.

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3rd fragment: 20 bytes header -552 bytes of payload - offset 0 4th fragment: 20 bytes header -448 bytes of payload - offset 552/8 = 6
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So at the destination host, the 2nd, 3rd and 4th fragment will reach.

3. (a) Consider a subnet with prefix 128.226.40.128/26. Give an example of one IP address (of form a.b.c.d) that can be assigned in this network.

(b) Suppose an ISP owns the block of addresses of the form 128.226.40.64/26. Suppose it wants to create four subnets from this block, with each block having the same number of IP addresses. What are the prefixes (of form a.b.c.d/x) for the four subnets?

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Address: 128.226.40.64 10000000.11100010.00101000.01 000000
Netmask: 255.255.255.192 = 26 11111111.11111111.11111111.11 000000
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To create 4 subnets, we can add extra two bits to the netmask as

The prefix of four subnets are

- 1) 128.226.40.64/28
- 2) 128.226.40.80/28
- 3) 128.226.40.96/28
- 4) 128.226.40.112/28
- (c) How many hosts are allowed in each subnet created in (b)? (You do not need to consider hosts that are behind NATs.)

Number of hosts allowed in each subnet is $2^4 - 2 = 14$

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(a) Why is an ARP query sent within a broadcast frame?

The ARP query is sent within a broadcast frame because the querying host does not know the destination adapter corresponding to the destination IP address.

(b) Why is an ARP response sent within a frame with a specific destination MAC address?

The responding node knows the source adapter address to which the response to be sent, so there is no need to send a broadcast frame.

5) For the network below, show how Dijkstra's algorithm builds the routing table for router 3

From 3	0	1	2	4	5	6	7
3	1.3,3	0.9,3	∞	9.1,3	∞	1.6,3	∞
3,1	1.3,3		∞	9.1,3	∞	1.6,3	10.6,1
3,1,0			1.9,0	9.1,3	∞	1.6,3	2.3,0
3,1,0,6			1.9,0	4.3,6	∞		2.3,0
3,1,0,6,2				4.3,6	4.4,2		2.3,0
3,1,0,6,2,7				4.3,6	4.4,2		
3,1,0,6,2,7,4					4.4,2		

Routing Table for 3

Destination	Link	Next Hop
0	(3,0)	0
1	(3,1)	1
2	(3,0,2)	0
4	(3,6,4)	6

5	(3,0,2,5)	0
6	(3,6)	6
7	(3,0,7)	0

- 6) The distance vector of a router X is given below ((A,3,7), (B,2,4), (C,1,1), (D,2,2), (E,0,1), (W,3,8), (X,-1,0), (Y,3,1), (Z,2,5)) where an entry (Y,p,c) indicates X reaches Y through port p in c hops. X receives via port 2 a vector ((A,3,4),(B,2,3),(C,4,4),(D,1,10),(E,0,1),(F,3,3), (Y,3,1),(Z,-1,0)) Answer the following questions.
- (a) Give the vector of X after processing the incoming vector in the alphabetic order of routers.

((A,2,5)

(B,2,4)

(C,1,1)

(D,2,2)

(E,0,1)

(F,2,4)

(W,3,8)

(X,-1,0)

(Y,3,1)

(Z,2,1)

(b) Apparently, the input vector is advertised by Z. The communication link between X and Z is: $X2 \to Z1$. That is, the link connecting X and Z is called port 2 at X and port 1 at Z. Point out a routing loop between X and Z

Routing loop can happen between X and Z while the link between Z and A is broken and Z does not inform the link failure to X. If X sends data to A through node Z, Z sends back to X, knowing the failure between Z and A. The loop continues.