

CONTEXT-FREE GRAMMAR and PARSING TABLE

Problem: Given a grammar, how do we construct a parsing table?

Context-free grammar:

$$G = (N, T, P, S)$$

where:

- N is finite set of non-terminals.
- T is finite set of terminals.
- P is finite set of productions in the form of $A \rightarrow \alpha$

where:

- $A \in N$
- $\alpha \in (N \cup T)^*$

is a sentential form

- $S \in N$ is the start symbol

Sentential form:
String $\in N \cup T$ that
can be derived from S

Context-free grammar vs Regular grammar

Context-free

- Exactly 1 non-terminal on LHS
- RHS can be a mix of terminals and non-terminals

Regular

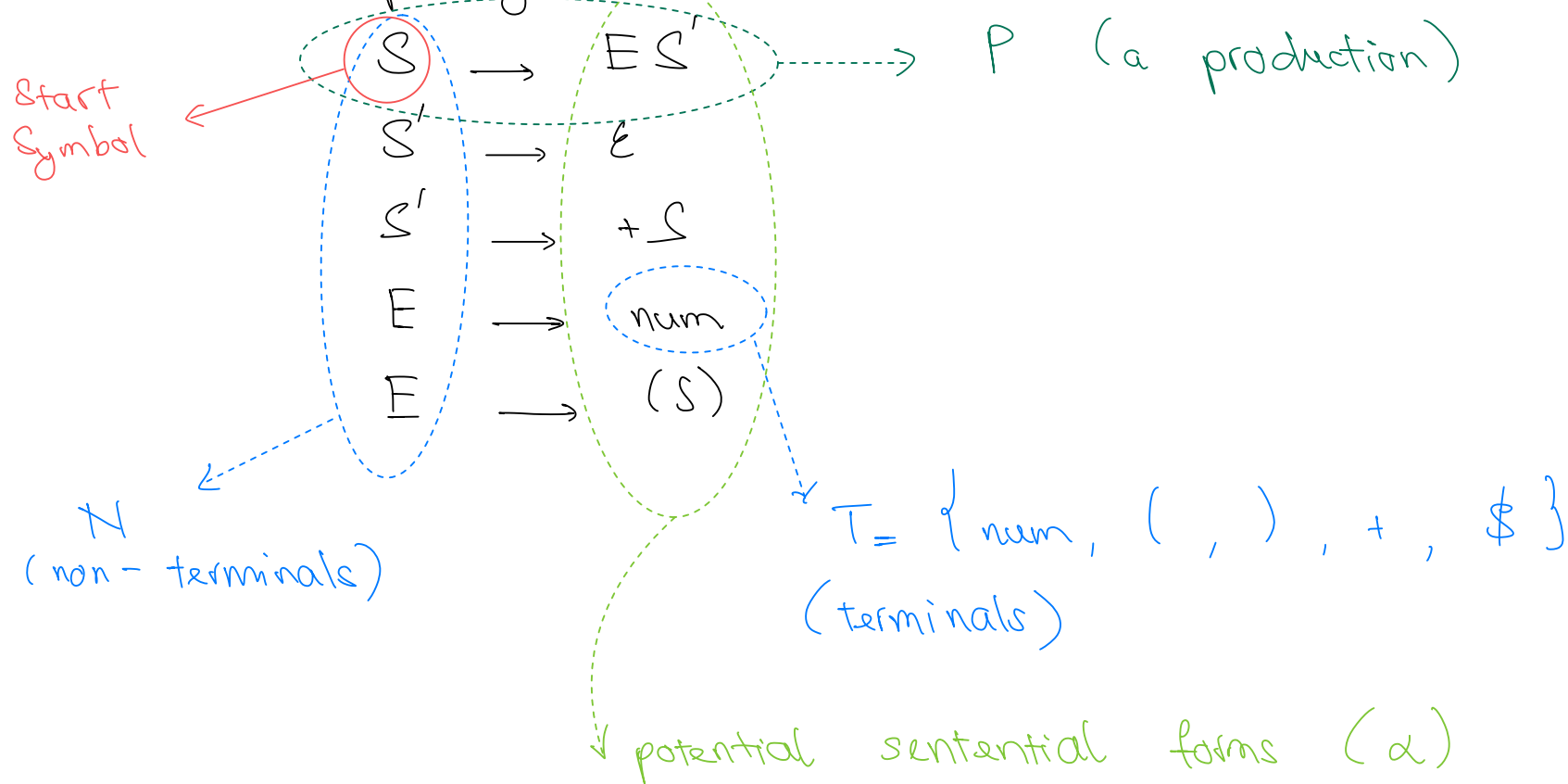
- Right-linear: $A \rightarrow \alpha B$ or $A \rightarrow \alpha$
- Left-linear: $A \rightarrow B\alpha$ or $A \rightarrow \alpha$
- Cannot mix left and right linear in 1 production

Some notation shortcuts

- $\Sigma = N \cup T$
- $F = T \cup \{\epsilon\}$
- $T' = T \cup \{\$\}$

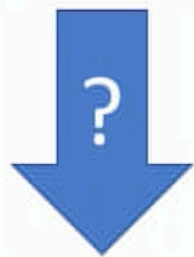
Example: Context-free grammar

Given this simple grammar:



Construct Parsing Table from Grammar

$S \rightarrow ES'$
 $S' \rightarrow \epsilon$
 $S' \rightarrow +S$
 $E \rightarrow \text{num}$
 $E \rightarrow (S)$



A {

	num	+	()	\$
S	$\rightarrow ES'$		$\rightarrow ES'$		
S'		$\rightarrow +S$		$\rightarrow \epsilon$	$\rightarrow \epsilon$
E	$\rightarrow \text{num}$		$\rightarrow (S)$		

t

Problem: Given a terminal t , what production $A \rightarrow \alpha$, should be filled in the parsing table cell $PT[A, t]$

→ Intuition: a production $A \rightarrow \alpha$ is a candidate if any of these true

① $t \in \text{FIRST}(\alpha)$?

Does this terminal t the first thing we see using this production $A \rightarrow \alpha$, where α is the RHS ?

Example: Given $t = "("$, what production goes in $PT[S, "("]$?

Consider production $S \rightarrow ES'$

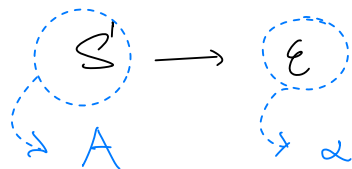
$\text{FIRST}(S)$
 $= \text{FIRST}(E)$
 $= "(" \rightarrow \text{Valid production}$

② $\text{NULLABLE}(\alpha)$ and $t \in \text{FOLLOW}(A)$?

If current RHS α can derive null (ϵ), then does the terminal t following the production A ?

Example: Given $t = ")"$, what production goes in $PT[S', ")"]$?

Consider production



$\text{FOLLOW}(S')$
 $= \text{FOLLOW}(S)$
 $= ")"$
 \checkmark

$\text{NULLABLE}(\epsilon)$
 \checkmark
 $\Rightarrow \text{Valid production}$

Construct Parsing Table Algorithm