

Network Calculus

Exercise Series

Lecture 11

2005

Exercise 1 For a CBR connection, here are some values from an ATM operator:

peak cell rate (cells/s)	100	1000	10000	100000
CDVT (microseconds)	2900	1200	400	135

1. What are the (P, B) parameters in b/s and bits for each case ? How does T compare to τ ?
2. If a connection requires a peak cell rate of 1000 cells per second and a cell delay variation of 1400 microseconds, what can be done ?
3. Assume the operator allocates the peak rate to every connection at one buffer. What is the amount of buffer required to assure absence of loss ? Numerical Application for each of the following cases, where a number N of identical connections with peak cell rate P is multiplexed.

case	1	2	3	4
nb of connexions	3000	300	30	3
peak cell rate (c/s)	100	1000	10000	100000

Exercise 2 Is it true that offering a service curve β implies that, during any busy period of length t , the amount of service received rate is at least $\beta(t)$?

Exercise 3 We consider a buffer of size X cells, served at a constant rate of c cells per second. We put N identical connections into the buffer; each of the N connections is constrained both by $GCRA(T_1, \tau_1)$ and $GCRA(T_2, \tau_2)$. What is the maximum value of N which is possible if we want to guarantee that there is no cell loss at all ?

Give the numerical application for $T_1 = 0.5$ ms, $\tau_1 = 4.5$ ms, $T_2 = 5$ ms, $\tau_2 = 495$ ms, $c = 10^6$ cells/second, $X = 10^4$ cells

Exercise 4 We consider a flow defined by its function $R(t)$, with $R(t)$ = the number of bits observed since time $t = 0$.

1. The flow is fed into a buffer, served at a rate r . Call $q(t)$ the buffer content at time t . We do the same assumptions as in the lecture, namely, the buffer is large enough, and is initially empty. What is the expression of $q(t)$ assuming we know $R(t)$?
We assume now that, unlike what we saw in the lecture, the initial buffer content (at time $t = 0$) is not 0, but some value $q_0 \geq 0$. What is now the expression for $q(t)$?
2. The flow is put into a leaky bucket policer, with rate r and bucket size b . This is a policer, not a shaper, so nonconformant bits are discarded. We assume that the bucket is large enough, and is initially empty. What is the condition on R which ensures that no bit is discarded by the policer (in other words, that the flow is

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