

Exercises: Constrained Optimization

Exercise 1 (Penalization). We consider the following problem

$$(P) \quad \min_{x \in \mathbb{R}^n} \quad f(x) \\ \text{s.t.} \quad Ax = b, \quad x \leq 0$$

with value v and the following penalized versions

$$(P_t^{in}) \quad \min_{x \in \mathbb{R}^n} f(x) - t \sum_{i=1}^n \ln(-x_i) \\ \text{s.t.} \quad Ax = b, \quad x < 0$$

and

$$(P_t^{out}) \quad \min_{x \in \mathbb{R}^n} f(x) + t \sum_{i=1}^n (x_i)^+ \\ \text{s.t.} \quad Ax = b$$

with associated value v_t^{in} and v_t^{out} , and an optimal solution x_t^{in} and x_t^{out} .

1. Intuitively, assuming that f is "well behaved", for t going to which value does (P_t^{in}) tend to the original problem (P) ? In which sense?
2. What can you say about x_t^{in} ?
3. Can you compare v_t^{in} and v ?
4. Same questions for (P_t^{out}) .

Answers: For t going to 0 we have that (P_t^{in}) tends toward (P) : in the sense that $v_t^{(in)} \rightarrow v$ and x_t goes toward an optimal solution . For t small enough we have $v_t^{in} \geq v$. In any case x_t^{in} is admissible.

For t going to $+\infty$, we have that (P_t^{out}) tends toward (P) in the sense that $v_t^{(out)} \rightarrow v$ and x_t^{out} goes toward an optimal solution. For t large enough, x_t^{out} is optimal for (P) . We always have $v_t^{(out)} \leq v$.

Exercise 2 (Decomposition by prices). We consider the following energy problem:

- you are an energy producer with N production units
- you have to satisfy a given demand planning for the next 24h (i.e. the total output at time t should be equal to d_t)
- the time step is the hour, and each unit have a production cost for each planning given as a convex quadratic function of the planning
- For each unit i , the production planning $u^i = (u_t^i)_{t \in [24]}$ has to satisfy polyhedral constraints $u^i \in U^i$.

1. Model this problem as an optimization problem. In which class does it belongs ? How many variables ?
2. Apply Uzawa's algorithm to this problem. Why could this be an interesting idea ?
3. Give an economic interpretation to this method.
4. What would happen if each unit had production constraints ?

Exercise 3 (Kelley's convergence). We are going to prove that, if $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is convex, and X a non-empty polytope (bounded polyhedron) then Kelley's cutting plane algorithm is converging.

Consider $x_1 \in X$. We consider a sequence of points $(x^{(k)})_{k \in \mathbb{N}}$ such that $x^{(k+1)}$ is an optimal solution to

$$(\mathcal{P}^{(k)}) \quad \underline{v}^{(k+1)} = \min_{x \in X} z \\ \text{s.t.} \quad f(x^{(\kappa)}) + \langle g^{(\kappa)}, x - x^{(\kappa)} \rangle \leq z \quad \forall \kappa \in [k]$$

where $g^{(k)} \in \partial f(x^{(k)})$.

Denote $v = \min_{x \in X} f(x)$.

1. Show that v exists and is finite, and that there exists a sequence $x^{(k)}$.
2. Show that there exists L such that, for all k_1 and k_2 , we have $\|f(x^{(k_1)}) - f(x^{(k_2)})\| \leq L\|x^{(k_1)} - x^{(k_2)}\|$, and $\|g^{(k)}\| \leq L$.
3. Let $K_\varepsilon = \{k \in \mathbb{N} \mid f(x^{(k)}) > v + \varepsilon\}$ be the set of index such that $x^{(k)}$ is not an ε -optimal solution. Show that $f(x_k) \rightarrow v$ if and only if K_ε is finite for all $\varepsilon > 0$
4. Consider $k_1, k_2 \in K_\varepsilon$, such that $k_2 > k_1$. Show that
$$f(x^{(k_1)}) + \langle g^{(k_1)}, x^{(k_2)} - x^{(k_1)} \rangle \leq \underline{v}^{(k_2)} \leq v$$
5. Show that $\varepsilon + f(x^{(k_1)}) + \langle g^{(k_1)}, x^{(k_2)} - x^{(k_1)} \rangle < f(x^{(k_2)})$
6. Show that $\varepsilon < 2L\|x^{(k_2)} - x^{(k_1)}\|$.
7. Prove that $f(x^{(k)}) \rightarrow v$.
8. (Optional - hard) Find a complexity bound for the method (that is a number of iteration N_ε after which you are sure to have obtained a ε -optimal solution).

Hence the cutting-plane model underestimates f on X , and therefore $\underline{v}^{(k)} \leq v$ for all k . Moreover, since $(x^{(k_2)}, \underline{v}^{(k_2)})$ is feasible for $\mathcal{P}^{(k_2-1)}$, it satisfies in particular the cut $\kappa = k_1$, giving

$$f(x^{(k_1)}) + \langle g^{(k_1)}, x^{(k_2)} - x^{(k_1)} \rangle \leq \underline{v}^{(k_2)} \leq v.$$

4. As $k_2 \in K_\varepsilon$, we have $f(x^{(k_2)}) = v^{(k_2)} > v + \varepsilon \geq f(x^{(k_1)}) + \langle g^{(k_1)}, x^{(k_2)} - x^{(k_1)} \rangle + \varepsilon$ by the previous question.

5. We have

$$\begin{aligned} \varepsilon &< |f(x^{(k_2)}) - f(x^{(k_1)})| + |\langle g^{(k_1)}, x^{(k_2)} - x^{(k_1)} \rangle| \\ &\leq 2L\|x^{(k_2)} - x^{(k_1)}\| \end{aligned}$$

by Cauchy-Schwartz and question 2.

6. If $f(x^{(k)}) \not\rightarrow v$, then there exists $\varepsilon > 0$ such that $(x^{(k)})_{k \in K_\varepsilon}$ is not finite. As X is compact we can extract a converging subsequence, that is $x^{(\sigma(k))}$ such that $x^{(\sigma(k))} \rightarrow x^*$ and $\sigma(k) \in K_\varepsilon$, which is in contradiction with the result of 6.

Answers:

1. f is finite convex and thus continuous on X which is compact, yielding the existence and finiteness of v .
 f is subdifferentiable, thus we have the existence of $g^{(k)}$, and an optimal solution to $\mathcal{P}^{(k)}$ exists as the solution of a bounded linear program.
2. We have seen that on any compact K included in the domain of a convex function f , f is L -Lipschitz. Here $\text{dom}(f) = \mathbb{R}^n$, so on the compact $K = X + B(0, \varepsilon)$ f is L -Lipschitz, and on X any subgradient g is of norm lower than L .
3. $f(x_k) \rightarrow v$ iff $\forall \varepsilon > 0, \exists N_\varepsilon \in \mathbb{N}, k \geq K \implies f(x_k) \leq v + \varepsilon$. Hence $K_\varepsilon \subset [N_\varepsilon]$. By the subgradient inequality, for any $\kappa \leq k$ and any $x \in X$,

$$f(x) \geq f(x^{(\kappa)}) + \langle g^{(\kappa)}, x - x^{(\kappa)} \rangle.$$