

# Newton and Quasi-Newton algorithms

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May 2nd, 2025

# Why should I bother to learn this stuff?

- Newton algorithm is, in theory, the best black-box algorithm for smooth strongly convex function. It is used in practice as well as a stepping step for more advanced algorithm.
- Quasi-Newton algorithms (in particular L-BFGS) are the de facto by default algorithm for most smooth black-box optimization library. Used in large scale application (e.g. weather forecast) for decades.
- $\implies$  useful for
  - ▶ understanding the optimization software you might use as an engineer
  - ▶ understanding more advanced methods (e.g. interior points methods)
  - ▶ getting an idea of why the convergence might behave strangely in practice

## Oriented sum-up of previous courses

- There are two large classes of unconstrained, exact, black-box, optimization algorithms:
  - ▶ descent direction algorithm:  $x^{(k+1)} = x^{(k)} + t^{(k)} d^{(k)}$ ;
  - ▶ model based approach:  $x^{(k+1)} = \arg \min_x f^{(k)}(x)$ .
- We saw that defining a descent direction algorithm requires:
  - ▶ a direction  $d^{(k)}$ ;
  - ▶ a step  $t^{(k)}$ ;
  - ▶ a stopping test (e.g.  $\|\nabla f(x^{(k)})\|_2 \ll 1$ )
- We discussed gradient and conjugate gradient algorithms defined by  $d^{(k)} = -\nabla f(x^{(k)}) + \beta^{(k)} d^{(k-1)}$ :
  - ▶ convergence speed is sensitive to conditioning of the problem (i.e. if level sets are almost spherical);
  - ▶ you can precondition the problem through a change of coordinates;
  - ▶ can be interpreted as steepest descent method:  $d^{(k)} = \arg \min_{\|d\|_P \leq 1} \nabla f(x^{(k)})^\top d$

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## 1 Newton algorithm [BV 9.5]

- Algorithm presentation, intuition and property
- (Damped) Newton algorithm convergence

## 2 Quasi Newton [JCG - 11.2]

- Quasi-Newton methods
- BFGS algorithm

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Let  $f$  be  $\mathcal{C}^2$  such that  $\nabla^2 f(x) \succ 0^1$  for all  $x$  (so in particular strictly convex).

The Newton algorithm is a descent direction algorithm with:

- $d^{(k)} = -[\nabla^2 f(x^{(k)})]^{-1} \nabla f(x^{(k)})$
- $t^{(k)} = 1$

Note that

$$\nabla f(x^{(k)})^\top d^{(k)} = -\nabla f(x^{(k)})^\top [\nabla^2 f(x^{(k)})]^{-1} \nabla f(x^{(k)}) < 0$$

(unless  $\nabla f(x^{(k)}) = 0$ )

$\leadsto d^{(k)}$  is a descent direction.

We are now going to give multiple justifications for this direction choice.

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We have

$$f(x^{(k)} + d) = f(x^{(k)}) + \nabla f(x^{(k)})^\top d + \frac{1}{2} d^\top \nabla^2 f(x^{(k)}) d + o(\|d\|^2)$$

The Newton method chooses the direction  $d$  (with step 1) that minimizes this second-order approximation, which is given by

$$\nabla f(x^{(k)}) + \nabla^2 f(x^{(k)}) d^{(k)} = 0$$

↪ The Newton method can be seen as a **model-based** method, where the model at iteration  $k$  is simply the second-order approximation.

↪ A trust region method with confidence radius  $+\infty$  is simply the Newton method.



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- The Newton direction  $d^{(k)}$  is the steepest descent direction for the quadratic norm associated to  $\nabla^2 f(x^{(k)})$ :

$$d^{(k)} = \arg \min_d \left\{ \nabla f(x^{(k)})^\top d \mid \|d\|_{\nabla^2 f(x^{(k)})} \leq 1 \right\}$$

- Recall that the steepest gradient descent for a quadratic norm  $\|\cdot\|_P$  converges rapidly if the condition number of the Hessian, after a change of coordinate, is small.
- In particular a good choice near  $x^\#$  is  $P = \nabla^2 f(x^\#)$ .

$\leadsto$  fast around  $x^\#$

# Solution of linearized optimality condition



The optimality condition is given by

$$\nabla f(x^\#) = 0$$

We can linearize it as

$$\nabla f(x^{(k)} + d) \approx \nabla f(x^{(k)}) + \nabla^2 f(x^{(k)})d = 0$$

And the Newton step  $d^{(k)}$  is the solution of this linearization.



- Recall that gradient and conjugate gradient methods can be accelerated through smart affine changes of variables (pre-conditioning).
- It is not the same for the Newton method:
  - ▶ Let  $A$  be an invertible matrix, and denote  $y = Ax + b$ , and  $\tilde{f} : x \mapsto f(Ax + b)$ .
  - ▶  $\nabla \tilde{f}(y) = A \nabla f(x)$  and  $\nabla^2 \tilde{f}(y) = A^\top \nabla^2 f(x) A$
  - ▶ The Newton step for  $\tilde{f}$  is thus

$$d_y = -(A^\top \nabla^2 f(x) A)^{-1} A \nabla f(x) = -A^{-1} (\nabla^2 f(x))^{-1} \nabla f(x) = A^{-1} d_x$$

- ▶ Consequently

$$x^{(k+1)} - x^{(k)} = A(y^{(k+1)} - y^{(k)})$$

➡ The Newton method does not really benefit from preconditioning!



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**Data:** Initial point  $x^{(0)}$ , second-order oracle, error  $\varepsilon > 0$ .

**while**  $\|\nabla f(x^{(k)})\| \geq \varepsilon$  **do**

    Solve for  $d^{(k)}$

$$\nabla^2 f(x^{(k)}) d^{(k)} = -\nabla f(x^{(k)})$$

    Compute  $t^{(k)}$  by backtracking line-search, starting from  $t = 1$ ;

$x^{(k+1)} = x^{(k)} + t^{(k)} d^{(k)}$

## Algorithm 1: Damped Newton algorithm

- The Newton algorithm with fixed step size  $t = 1$  may fail far from the optimum, and you should always use a backtracking line-search.
- If the function is not strictly convex the Newton direction is not necessarily a descent direction, and you should check for it (and default to a gradient step).

## Convergence idea



Assume that  $f$  is strongly convex, such that  $mI \preceq \nabla^2 f(x) \preceq LI$ , and that the Hessian  $\nabla^2 f$  is  $L$ -Lipschitz.

We have the following two phases of convergence:

- 1 **Damped phase:** far from the optimum, the step  $t^{(k)}$  might be less than 1. Each iteration yields an absolute improvement of  $-\gamma < 0$ .
- 2 **Quadratic phase:** close to the optimum, the step  $t^{(k)} = 1$ . Each iteration squares the error.

More precisely, we can show that there exists  $0 < \eta \leq m^2/L$  and  $\gamma > 0$  such that

- If  $\|\nabla f(x^{(k)})\|_2 \geq \eta$ , then

$$f(x^{(k+1)}) - f(x^{(k)}) \leq -\gamma$$

- If  $\|\nabla f(x^{(k)})\|_2 < \eta$ , then  $t^{(k)} = 1$  and

$$\frac{L}{2m^2} \|\nabla f(x^{(k+1)})\|_2 \leq \left( \frac{L}{2m^2} \|\nabla f(x^{(k)})\|_2 \right)^2$$

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## Newton is fast around the solution



We have, if  $\|\nabla f(x^{(k)})\|_2 < \eta$ , then  $t^{(k)} = 1$  and

$$\frac{L}{2m^2} \|\nabla f(x^{(k+1)})\|_2 \leq \left( \frac{L}{2m^2} \|\nabla f(x^{(k)})\|_2 \right)^2$$

Let  $k = k_0 + \ell$ ,  $\ell \geq 1$ , with  $k_0$  such that  $\|\nabla f(x^{(k_0)})\|_2 < \eta$ . Then  $\|\nabla f(x^{(k)})\|_2 < \eta$ , and,

$$\frac{L}{2m^2} \|\nabla f(x^{(k)})\|_2 \leq \left( \frac{L}{2m^2} \|\nabla f(x^{(k-1)})\|_2 \right)^2$$

Recursively,

$$\frac{L}{2m^2} \|\nabla f(x^{(k)})\|_2 \leq \left( \frac{L}{2m^2} \|\nabla f(x^{(k_0)})\|_2 \right)^{2^\ell} \leq \frac{1}{2^{2^\ell}}$$

And thus

$$f(x^{(k)}) - v^\# \leq \frac{1}{2m} \|\nabla f(x^{(k)})\|_2^2 \leq \frac{2m^3}{L^2} \frac{1}{2^{2^\ell-1}}$$

$\leadsto$  in the quadratic convergence phase, Newton's algorithm gets the result in a few iterations (5 or 6).

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## Convergence speed - Wrap-up

The Newton algorithm, for strongly convex function, have two phases :

- The damped phase, where  $t^{(k)}$  can be less than 1. Each iteration yields an absolute improvement of  $-\gamma < 0$ .
- The quadratic phase, where each step  $t^{(k)} = 1$ .

Thus, the total number of iterations to get an  $\varepsilon$  solution is bounded above by

$$\frac{f(x^{(0)}) - v^\#}{\gamma} + \underbrace{\log_2(\log_2(\varepsilon_0/\varepsilon))}_{\lesssim 6}$$

where  $\varepsilon_0 = 2m^3/L^2$ .

Note that, in 6 iterations in the quadratic convergent phase we get an error  $\varepsilon \approx 5 \cdot 10^{-20} \varepsilon_0$ .

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- Full Newton step :  $x^{(k+1)} = x^{(k)} - [\nabla^2 f(x^{(k)})]^{-1} \nabla f(x^{(k)})$
- Can be seen through various lenses:
  - ①  $[\nabla^2 f(x^{(k)})]^{-1} \nabla f(x^{(k)})$  is a descent direction ( $f$  is strongly convex);
  - ② model-based algorithm where the model is the second-order approximation;
  - ③ preconditioned gradient algorithm, with adaptive preconditioning.
- Is incredibly fast around the optimal solution.
- Far from the optimum a full Newton step is a bad idea:
  - ▶ If  $f$  is not strongly convex the Newton direction might not be a descent direction<sup>2</sup>!
  - ▶  $\leadsto$  check if it is a descent direction, otherwise make a gradient step.
  - ▶ Even with convexity the step might be too aggressive,  $\leadsto$  receding step choice.
- Convergence of the (damped) Newton's algorithm is in two phases:
  - ▶ slow constant update far from the optimum,
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Newton's step is very efficient (near optimality) but has three drawbacks:

- 1 having a second-order oracle to compute the Hessian
- 2 storing the Hessian ( $n^2$  values)
- 3 solving a (dense) linear system :  $\nabla^2 f(x^{(k)})d = -\nabla f(x^{(k)})$

The main idea of Quasi Newton method is to define  $M^{(k)} \approx \nabla^2 f(x^{(k)})$  (or  $W^{(k)} \approx [\nabla^2 f(x^{(k)})]^{-1}$ ):

- 1 from first order information  $\leadsto$  no need to compute Hessian;
- 2 sparse  $\leadsto$  smaller storage requirements;
- 3  $d^{(k)} = -W^{(k)}\nabla f(x^{(k)}) \leadsto$  no linear system solving.



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## Conditions on the approximate Hessian



We want to construct  $M^{(k)}$  an approximation of  $\nabla^2 f(x^{(k)})$ , leading to a quadratic model of  $f$  at iteration  $k$

$$f^{(k)}(x) := f(x^{(k)}) + \langle \nabla f(x^{(k)}), x - x^{(k)} \rangle + \frac{1}{2}(x - x^{(k)})^\top M^{(k)}(x - x^{(k)})$$

We ask that the gradient of the model  $f^{(k)}$  and the true function to match at the current and last iterates:

$$\begin{cases} \nabla f^{(k)}(x^{(k)}) = \nabla f(x^{(k)}) \\ \nabla f^{(k)}(x^{(k-1)}) = \nabla f(x^{(k-1)}) \end{cases}$$

This simply write as the **Quasi-Newton equation**

$$M^{(k)} \underbrace{(x^{(k)} - x^{(k-1)})}_{\delta_x^{(k-1)}} = \underbrace{\nabla f(x^{(k)}) - \nabla f(x^{(k-1)})}_{\delta_g^{(k-1)}}$$



Exercise: prove it

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$$\begin{cases} \nabla f^{(k)}(x^{(k)}) = \nabla f(x^{(k)}) \\ \nabla f^{(k)}(x^{(k-1)}) = \nabla f(x^{(k-1)}) \end{cases}$$

This simply write as the **Quasi-Newton equation**

$$M^{(k)} \underbrace{(x^{(k)} - x^{(k-1)})}_{\delta_x^{(k-1)}} = \underbrace{\nabla f(x^{(k)}) - \nabla f(x^{(k-1)})}_{\delta_g^{(k-1)}}$$



Exercise: prove it

## Conditions on the approximate Hessian



We want to construct  $M^{(k)}$  an approximation of  $\nabla^2 f(x^{(k)})$ , leading to a quadratic model of  $f$  at iteration  $k$

$$f^{(k)}(x) := f(x^{(k)}) + \langle \nabla f(x^{(k)}), x - x^{(k)} \rangle + \frac{1}{2}(x - x^{(k)})^\top M^{(k)}(x - x^{(k)})$$

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♣ Exercise: prove it

We are looking for a matrix  $M$  such that

- $M \succ 0$
- $M\delta_x = \delta_g$  (only possible if  $\delta_g^\top \delta_x > 0$  ♣ Exercise: prove it)
- $M^\top = M$
- $M$  is constructed from first order information only
- If possible,  $M$  is sparse

↪ an infinite number of solutions as we have  $n(n+1)/2$  variables and  $n$  constraints.

↪ Numerous quasi-Newton algorithms developed and tested between 1960-1980.

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## Choosing the approximate Hessian $M^{(k)}$



At the end of iteration  $k$  we have determined

- $x^{(k+1)}$  and  $\delta_x^{(k)} = x^{(k+1)} - x^{(k)}$
- $g^{(k+1)} = \nabla f(x^{(k+1)})$  and  $\delta_g^{(k)} = g^{(k+1)} - g^{(k)}$

and we are looking for  $M^{(k+1)} \approx \nabla^2 f(x^{(k+1)})$  satisfying the previous requirement.

The idea is to choose  $M^{(k+1)}$  close to  $M^{(k)}$ , that is to solve (analytically)

$$\begin{aligned} \text{Min}_{M \in S_{++}^n} \quad & d(M, M^{(k)}) \\ \text{s.t.} \quad & M \delta_x^{(k)} = \delta_g^{(k)} \end{aligned}$$

for some distance  $d$ .

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# Contents

- 1 Newton algorithm [BV 9.5]
  - Algorithm presentation, intuition and property
  - (Damped) Newton algorithm convergence
- 2 Quasi Newton [JCG - 11.2]
  - Quasi-Newton methods
  - BFGS algorithm



Broyden-Fletcher-Goldfarb-Shanno chose

$$d(A, B) := \operatorname{tr}(AB) - \ln \det(AB)$$

A few remarks

- $\Psi : M \mapsto \operatorname{tr} M - \ln \det(M)$  is convex on  $S_{++}^n$
- For  $M \in S_{++}^n$ ,  $\operatorname{tr} M - \ln \det(M) = \sum_{i=1}^n \lambda_i - \ln(\lambda_i)$
- $\Psi$  is minimized in the identity matrix
- $d(A, B) - n$  is the Kullback-Liebr divergence between  $\mathcal{N}(0, A)$  and  $\mathcal{N}(0, B)$



One of the pragmatic reasons for this choice of distance is that the optimal solution can be found analytically.

We have<sup>3</sup> (to alleviate notation we drop the index  $k$  on  $\delta_x^{(k)}$  and  $\delta_g^{(k)}$ )

$$M^{(k+1)} = M^{(k)} + \frac{\delta_g \delta_g^\top}{\delta_g^\top \delta_g} - \frac{M^{(k)} \delta_x \delta_x^\top M^{(k)}}{\delta_x^\top M^{(k)} \delta_x}$$

Even better, denoting  $W = M^{-1}$ , we can show<sup>4</sup> that:

$$W^{(k+1)} = \left( I - \frac{\delta_x \delta_g^\top}{\delta_g^\top \delta_x} \right) W^{(k)} \left( I - \frac{\delta_g \delta_x^\top}{\delta_g^\top \delta_x} \right) + \frac{\delta_x \delta_x^\top}{\delta_g^\top \delta_x}$$

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**Data:** Initial point  $x^{(0)}$ , First order oracle, error  $\varepsilon > 0$ .

$W^{(0)} = I$ ;

**while**  $\|\nabla f(x^{(k)})\| \geq \varepsilon$  **do**

$g^{(k)} := \nabla f(x^{(k)})$ ;

$d^{(k)} := -W^{(k)}g^{(k)}$ ;

    Compute  $t^{(k)}$  by backtracking line-search, starting from  $t = 1$ ;

$x^{(k+1)} = x^{(k)} + t^{(k)}d^{(k)}$ ;

$\delta_g = g^{(k+1)} - g^{(k)}$ ,  $\delta_x = x^{(k+1)} - x^{(k)}$ ;

$W^{(k+1)} = \left(I - \frac{\delta_x \delta_g^\top}{\delta_g^\top \delta_x}\right) W^{(k)} \left(I - \frac{\delta_g \delta_x^\top}{\delta_g^\top \delta_x}\right) + \frac{\delta_x \delta_x^\top}{\delta_g^\top \delta_x}$ ;

$k = k + 1$ ;

## Algorithm 2: BFGS algorithm

- ✓ First order oracle only
- ✓ No need to solve a linear system
- ✗ Still large memory requirement
- ✓ Convergence comparable to Newton's algorithm

# Limited-memory BFGS (L-BFGS)



- For  $n \geq 10^3$  storing the matrices is a difficulty.
- Instead of storing and updating the matrix  $W^{(k)}$  we store  $(\delta_x, \delta_g)$  pairs.
- We can then compute  $d^{(k)} = -W^{(k)}g^{(k)}$  directly from the last 5 to 20 pairs, using recursively the update rule and never computing  $W^{(k)}$ .

→ An algorithm with:

- ✓ First order oracle only
- ✓ No need to solve a linear system
- ✓ Same storage requirement as gradient algorithm
- ✓ Convergence comparable to Newton's algorithm

→ this is the "go to" algorithm when you want high-level precision for strongly convex smooth problems. It is the default choice in a lot of optimization libraries.



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# What you have to know

- At least one idea behind Newton's algorithm.
- The Newton step.
- That quasi-Newton methods are almost as good as Newton, without requiring a second order oracle.

## What you really should know

- Newton's algorithm default step is 1, but you should use backtracking step anyway.
- Newton's algorithm converges in two phases : a slow damped phase, and a very fast quadratically convergent phase close to the optimum (at most 6 iterations).
- BFGS is the by default quasi-Newton method. It work by updating an approximation of the inverse of the Hessian close to the precedent approximation and satisfying some natural requirement.
- L-BFGS limit the memory requirement by never storing the matrix but only the step and gradient updates.

# What you have to be able to do

- Implement a damped Newton method.

## What you should be able to do

- Implement a BFGS method (with the update formula in front of your eyes)