ROB311 Quiz 1

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Contents

1 Prologue

Summary:
- Variables:
  - State: $\mathbf{x}(t)$
  - Action(s): $\mathbf{u}(t)$
  - Measurement: $\mathbf{y}_k^{(i)}$
  - Context: $\mathbf{z}_k^{(i)}$
  - Old Context: $\mathbf{z}_{k-1}^{(i)}$
  - Plan: $\mathbf{p}_k^{(i)}$
  - (i): Ith agent
- Conversion to DT is necessary because robots are digitalized system and then converted back to CT for execution.

1.1 Setup of Planning Problems

Definition: In a planning problem, it is assumed that:
- the environment is representable using a discrete set of states, $\mathcal{S}$
- for each state, $s \in \mathcal{S}$, each agent, i, has a discrete set of actions, $\mathcal{A}_i(s)$, with $\mathcal{A}(s) := \times_i \mathcal{A}_i(s)$ (joint action set)
- **Move:** Any tuple, (s, a), where $s \in \mathcal{S}$ and $a \in \mathcal{A}(s)$
- **Transition:** Any 3-tuple, (s, a, s'), where $s, s' \in \mathcal{S}$ and $a \in \mathcal{A}(s)$
  - the transition resulting from a move may be deterministic/stochastic
- **Reward:** $\text{rwd}_i(s, a, s')$ is agent i's reward for the transition, (s, a, s')
- **Path:** Any sequence of transitions of the form.

$$p = \left\langle (s^{(0)}, a^{(1)}, s^{(1)}), (s^{(1)}, a^{(2)}, s^{(2)}), \dots \right\rangle$$

- **Objective:** Each agent wants to realize a path that maximizes its own reward

Warning: $\mathcal{A}(s)$ is the joint action set of all agents at state s.

1.2 Components of a Robotic System

Summary:

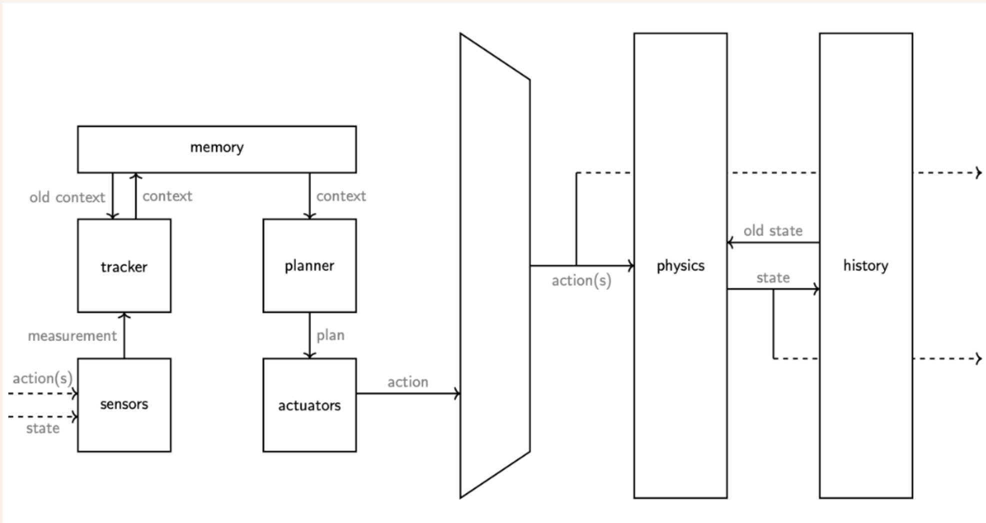


Figure 1: Components of a Robotic System (Words)

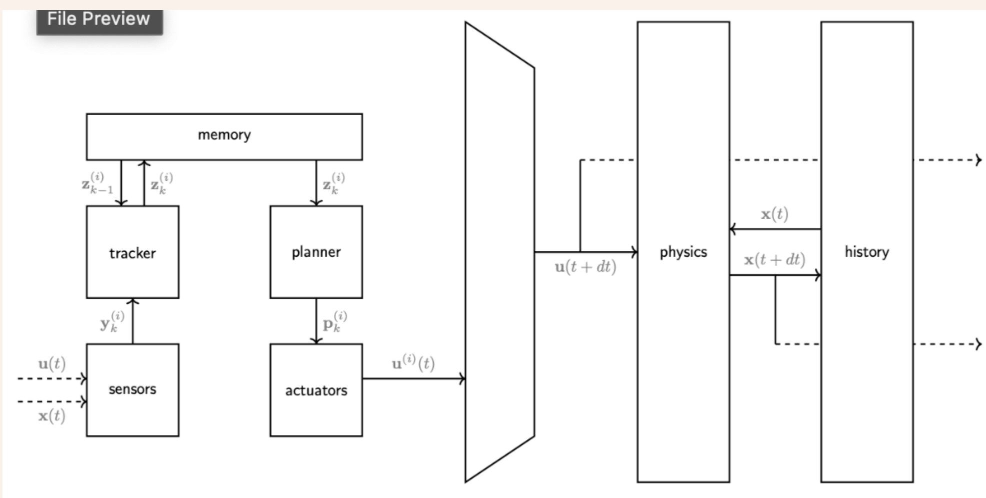


Figure 2: Components of a Robotic System (Math)

1.2.1 Overview (Robots, the Environment)

Definition:

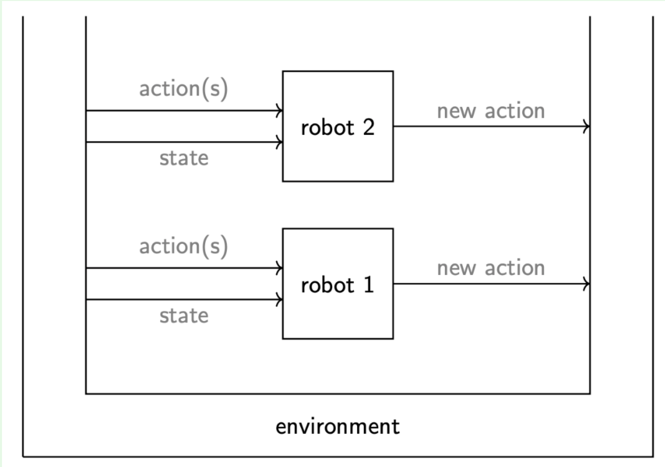


Figure 3: Overview (Robots, the Environment)

Notes:
- Environment → previous actions + current state → robot → new action → environment

1.2.2 Robot (Sensors, Actuators, the Brain)

Definition:

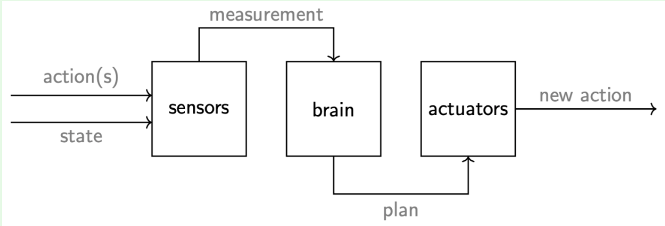


Figure 4: Robot (Sensors, Actuators, the Brain)

Notes:
- Measurements can be noisy and inaccurate if not a perfect sensor.
- Measurements go into the brain which can create a plan.

1.2.3 Brain (Tracker, Planner, Memory)

Definition:

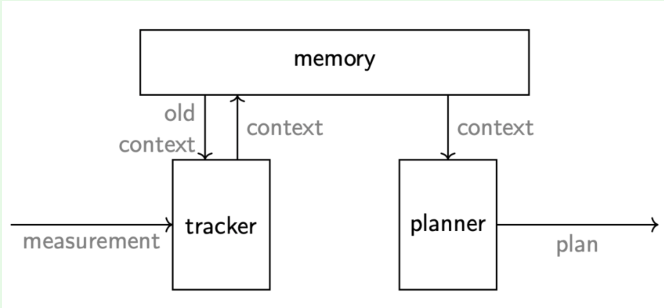


Figure 5: Brain (Tracker, Planner, Memory)

Notes:
- The tracker takes in the measurements and old context and updates the context.
- The planner takes in the context and creates a plan.
- The memory stores the context.

1.2.4 Environment (Physics, State)

Definition:

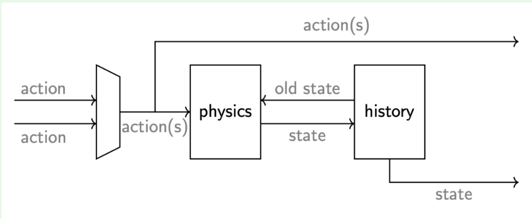


Figure 6: Environment (Physics, State)

1.3 Equations of a Robotic System

1.3.1 Sensing

Definition: Take a measurement:

$$\mathbf{y}^{(i)}(t) = \text{sns}^{(i)}\big(\mathbf{x}(t), \mathbf{u}(t), t\big)$$

Convert the measurement into a discrete-time signal using a sampling period of $T^{(i)}$:

$$\mathbf{y}_k^{(i)} = \text{dt}\big(\mathbf{y}^{(i)}(t), t, T^{(i)}\big)$$

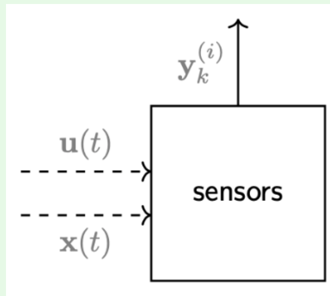


Figure 7: Sensing

1.3.2 Tracking

Definition: Track (update) the context:

$$\mathbf{z}_k^{(i)} = \text{trk}^{(i)}\big(\mathbf{z}_{k-1}^{(i)}, \mathbf{y}_k^{(i)}, k\big)$$

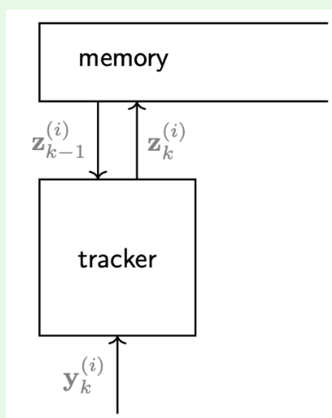


Figure 8: Tracking

1.3.3 Planning

Definition: Make a plan:

$$\mathbf{p}_k^{(i)} = \mathrm{pln}^{(i)}\big(\mathbf{z}_k^{(i)}, k\big)$$

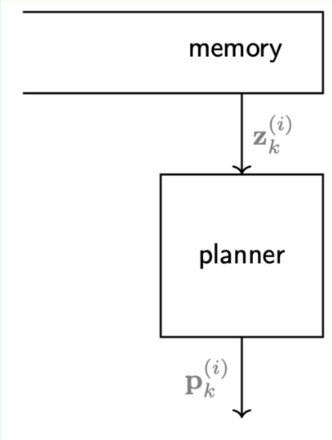


Figure 9: Planning

1.3.4 Acting

Definition: Convert the plan into a continuous-time signal using a sampling period of $T^{(i)}$:

$$\mathbf{p}(t) = \mathrm{ct}\big(\mathbf{p}_k^{(i)}, t, T^{(i)}\big)$$

Execute the plan:

$$\mathbf{u}^{(i)}(t) = \mathrm{act}^{(i)}\big(\mathbf{p}^{(i)}(t), t\big)$$

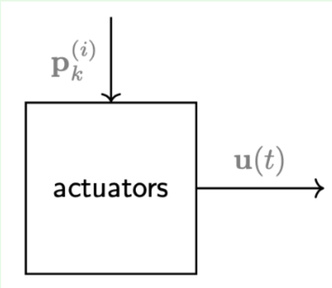


Figure 10: Acting

1.3.5 Simulating

Definition: Simulate the environment's response:

$$\dot{\mathbf{x}}(t) = \text{phy}\big(\mathbf{x}(t), \mathbf{u}(t), t\big)$$

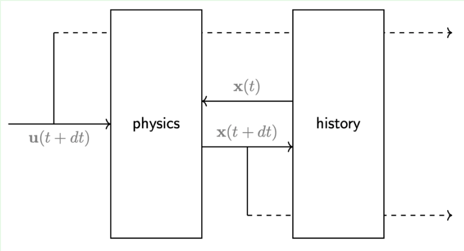


Figure 11: Simulating

2 Search Algorithms

Summary:

Alg.	Halting	Sound	Complete	Optimal	Time	Space
BFS	$d < \infty$ | non-NULL soln.	always	always	constant cst	b^{l^*}	b^{l^*+1}
DFS	$d < \infty$	always	$d < \infty$	never	b^d	bd
IDDFS	always	always	always	constant cst	b^{l^*}	bl^*
CFS	$d < \infty$ | non-NULL soln.	yes	$\epsilon > 0$	$\epsilon > 0$	$b^{c^*/\epsilon}$	$b^{c^*/\epsilon+1}$
HFS	$d < \infty$	never	never	never	-	-
A*	hur admissible, $\epsilon > 0$	always	hur admissible, $\epsilon > 0$	hur admissible, $\epsilon > 0$	$O\left(b^{c^*/\epsilon}\right)$	$O\left(b^{c^*/\epsilon+1}\right)$
IIA*	always	always	always	always	b^{l^*}	bl^*
WA*	-	-	-	-	-	-

Choose `REMOVE(·)` so algo. exhibits the characteristics:

- **Halting:** Terminates after finitely many nodes explored | **Sound:** Returned (possibly NULL) soln. is correct
- **Complete:** Halting & sound when a non-NULL soln. exists | **Opt.:** Returns an opt. soln. when mult. exist
- **Time:** Minimizes nodes **explored**/expanded/exported | **Space:** Minimizes nodes simultaneously open

Choose `REMOVE(·)` so algo. exhibits the characteristics for as many path trees as possible.

- b ($b < \infty$): Branching factor (the maximum number of children a node can have)
- d: Depth (the length of the longest path), l^*: Length of the shortest solution
- c^*: Cost of the cheapest solution, ϵ: Cost of the cheapest edge

Uninformed Search Algorithms

BFS

- Explores the least-recently expanded open node first.

DFS

- Explores the most-recently expanded open node first.

IDDFS

- Same as DFS but with iterative deepening.

CFS

- Explores the cheapest open node first.

Informed Search Algorithms

HFS

- Explores the node with the smallest hur-value first, $\text{ecst}(p) = \text{hur}(p)$

A*

- Explores the node with the smallest ecst-value first, $\text{ecst}(p) = \text{cst}(p) + \text{hur}(p)$

IIA*

- Same as A* but with iterative inflating on ecst.

WA*

- Same as A* but $\text{ecst}(s) = w\text{cst}(s) + (1-w)\text{hur}(s)$ w/ $w \in [0,1]$
- $w = 0$: HFS, $w = 0.5$: A*, $w = 1$: CFS, iteratively increasing w from 0 to 1: anytime version of WA^*

2.1 Modifications to Search Algorithms:

Summary:

Modifications
Depth-Limiting • Enforce a depth limit, $d_{\max}$, to any search algorithm.
Iterative-Deepening • Iteratively increase the depth-limit to any search algorithm w/ depth-limiting.
Cost-Limiting • Enforce a cost limit of $c_{\max}$ to any search algorithm.
Iterative Inflating • Iteratively increase the cost limit, $c_{\max}$, to any search algorithm w/ cost-limiting.
Intra-Path Cycle Checking • Do not expand a path if it is cyclic.
Inter-Path Cycle Checking • Do not expand a path if its destination is that of an explored path.

2.2 Setup

Definition: In a search problem, it is assumed that:
- There is only one agent (us).
- For each state, $s \in S$, we have a discrete set of actions, $\mathcal{A}(s)$.
- The transition resulting from a move, (s, a), is deterministic; the resulting state is $tr(s, a)$.
- $cst(s, a, tr(s, a))$ is our cost for the transition, $(s, a, tr(s, a))$.
- We want to realize a path that minimizes our cost.

A search problem may have no solutions, in which case, we define the solution as `NULL`.

Warning: A NULL solution is not the same as $p = \langle\rangle$ (an empty solution w/ $s^{(0)} \in \mathcal{G}$).

2.3 Search Graphs

Definition: In a search graph (a graph representing a search problem):
- S is defined by the vertices.
- $\mathcal{G}$ is a subset of the vertices.
- $s^{(0)}$ is some vertex.
- $tr(\cdot, \cdot)$ and $\mathcal{T}$ are defined by the edges.
- $cst(\cdot, \cdot, \cdot)$ is defined by the edge weights.

2.4 Path Trees

Definition: A search algorithm explores a tree of possible paths.
- In such a tree, each node represents the path from the root to itself.
  - The node may also include other info (such as the path's origin, cost, etc).

2.5 Search Algorithms

Algorithm: All search algorithms follow the template below:

```
𝒪 ← {(⟨⟩, 0)}                                   ▷ initialize a set of open nodes
SEARCH(𝒪)
```

- $\langle\rangle$: Empty path, 0: Cost of empty path.

```
procedure SEARCH(𝒪)
    if 𝒪 = ∅ then
        return NULL                   ▷ the search algorithm failed to find a path to a goal
    n ← REMOVE(𝒪)                                              ▷ "explore" a node n
    if DST(n) ∈ 𝒢 then
        return n                             ▷ the search algorithm found a path to a goal
    for n′ ∈ CHL(n) do                              ▷ "expand" n and "export" its children
        𝒪 ← 𝒪 ∪ {n′}
    SEARCH(𝒪)
```

- Explore: Remove a node from the open set.
- Expand: Generate the children of the node.
- Export: Add the children to the open set.

Warning: The key difference is in the order that REMOVE$(\cdot)$ removes nodes.

2.6 Modifications to Search Algorithms

2.6.1 Depth-Limiting

Algorithm:

```
procedure SEARCHDL(𝒪, d_max):
    if 𝒪 = ∅ then
        return NULL                  ▷ the search algorithm failed to find a path to a goal
    n ← REMOVE(𝒪)                                                  ▷ "explore" a node, n
    if dst(n) ∈ 𝒢 then
        return n                            ▷ the search algorithm found a path to a goal
    for n′ ∈ chl(n) do                               ▷ "expand" n and "export" its children
        if len(n′) ≤ d_max then                              ▷ unless the child is too long
            𝒪 ← 𝒪 ∪ {n′}
    SEARCHDL(𝒪, d_max)
```

2.6.2 Iterative Deepening

Algorithm:

```
procedure SEARCHID():
    n ← NULL
    d_max = 0
    ▷ while a solution has not been found, reset the open set, run the search algorithm, then increase the
    depth-limit
    while n = NULL do
        𝒪 ← {(⟨⟩, 0)}
        n ← SEARCHDL(𝒪, d_max)
        d_max ← d_max + 1
    return n
```

Warning: Increasing d_{max} can be done in different ways.

2.6.3 Cost-Limiting

Algorithm:

```
procedure SEARCHCL(𝒪, c_max):
    if 𝒪 = ∅ then
        return NULL                  ▷ the search algorithm failed to find a path to a goal
    n ← REMOVE(𝒪)                                                  ▷ "explore" a node, n
    if dst(n) ∈ 𝒢 then
        return n                            ▷ the search algorithm found a path to a goal
    for n′ ∈ chl(n) do                               ▷ "expand" n and "export" its children
        if cst(n′) ≤ c_max then                         ▷ unless the child is too expensive
            𝒪 ← 𝒪 ∪ {n′}
    SEARCHCL(𝒪, c_max)
```

2.6.4 Iterative-Inflating

Algorithm:

```
procedure SEARCHII():
    n ← NULL
    c_max = 0
    ▷ while a solution has not been found, reset the open set, run the search algorithm, then increase the cost-limit
    while n = NULL do
        𝒪 ← {(⟨⟩, 0)}
        n ← SEARCHCL(𝒪, c_max)
        c_max ← c_max + ϵ
    return n
```

Warning: Increasing $c_{\max}$ can be done in different ways.

2.6.5 Intra-Path Cycle Checking

Algorithm:

```
procedure SEARCH(𝒪):
    if 𝒪 = ∅ then
        return NULL
    n ← REMOVE(𝒪)
    if dst(n) ∈ 𝒢 then
        return n
    for n′ ∈ chl(n) do                          ▷ "expand" n and "export" its children
        if not CYCLIC(n′) then                  ▷ unless the child is cyclic
            𝒪 ← 𝒪 ∪ {n′}
    SEARCH(𝒪)
```

- Optimately of an algorithm is preserved provided $\epsilon > 0$.

2.6.6 Inter-Path Cycle Checking

Algorithm:

```
procedure SEARCH(𝒪, 𝒞):
    if 𝒪 = ∅ then
        return NULL
    n ← REMOVE(𝒪)
    𝒞 ← 𝒞 ∪ {n}                                 ▷ add n to the closed set
    if dst(n) ∈ 𝒢 then
        return n
    for n′ ∈ chl(n) do                          ▷ "expand" n and "export" its children
        if n′ ∉ 𝒞 then                          ▷ unless the child's destination is closed
            𝒪 ← 𝒪 ∪ {n′}
    SEARCH(𝒪, 𝒞)
```

and then call the algorithm as follows:

```
𝒪 ← {(⟨⟩, 0)}
𝒞 ← {}                                          ▷ initialize a set of closed vertices
SEARCH(𝒪, 𝒞)
```

2.7 Informed Search Algorithms

2.7.1 Estimated Cost

Definition: $\text{ecst}(\cdot)$: estimate the total cost to a goal given a path, p, based on:

- $\text{cst}(p)$: Cost of path p
- $\text{hur}: S \to \mathbb{R}_+$: Estimate of the extra cost needed to get to a goal from $\text{dst}(p)$
  - $\text{hur}(s)$ estimates the cost to get to $\mathcal{G}$ from s and $\text{hur}(p)$ means $\text{hur}(\text{dst}(p))$.
  - $\text{hur}^*(s)$: The true cost to get to $\mathcal{G}$ from s.

2.7.2 Admissible

Motivation: We want to find a heuristic that under estimates (i.e. make paths look better than they are) the costs, rather than over estimate (i.e. make paths look worse than they are).

- Least useful heuristic: $\text{hur}(s) = 0$ for all $s \in \mathcal{S}$ or any other constant.
- Most useful heuristic: $\text{hur}(s) = \text{hur}^*(s)$ for all $s \in \mathcal{S}$.

Definition: A heuristic, $\text{hur}(\cdot)$, is said to be **admissible** if

$$\text{hur}(s) \leq \text{hur}^*(s)$$

for all $s \in \mathcal{S}$ and

$$\text{hur}(s) = 0$$

for all $s \in \mathcal{G}$.

Warning: Never over-estimates the overall cost, but may still estimate the cost of individual transition.

2.7.3 Consistent

Definition: A heuristic, $\text{hur}(\cdot)$, is said to be **consistent** if

$$\underbrace{\text{hur}(s) - \text{hur}(\text{tr}(s,a))}_{\text{estimated cost of the transition } (s,a,\text{tr}(s,a))} \leq \underbrace{\text{cst}(s,a,\text{tr}(s,a))}_{\text{true cost of the transition, } (s,a,\text{tr}(s,a))}$$

for all $s \in \mathcal{S}$, and $a \in \mathcal{A}(s)$, and

$$\text{hur}(s) = 0$$

for all $s \in \mathcal{G}$.

Warning: Never over-estimates the cost of individual transitions (and hence the overall cost).

Theorem: If a heuristic, $\text{hur}(\cdot)$, is consistent, then it is also admissible.

2.7.4 Domination

Definition: If hur_1 and hur_2 are admissible, then:

- hur_1 **strongly dominates** hur_2 if for all $s \in \mathcal{S} \setminus \mathcal{G}$:

$$\mathrm{hur}_1(s) > \mathrm{hur}_2(s)$$

- hur_1 **weakly dominates** hur_2 if for all $s \in \mathcal{S}$:

$$\mathrm{hur}_1(s) \geq \mathrm{hur}_2(s)$$

  and for some $s \in \mathcal{S}$:

$$\mathrm{hur}_1(s) > \mathrm{hur}_2(s)$$

Notes: Want the heuristic that dominates but is also admissible.

2.7.5 Designing Heuristics via Problem Relaxation

Definition: Let $\mathrm{hur}^*_{\mathrm{ori}}$ be the perfect heuristic for a search problem, and $\mathrm{cst}^*_{\mathrm{rel}}$ be the optimal cost for a relaxed version of the problem. Then

$$\mathrm{cst}^*_{\mathrm{rel}}(s) \leq \mathrm{hur}^*_{\mathrm{ori}}(s) \text{ for all } s \in \mathcal{S}.$$

2.7.6 Combining Heuristics

Definition: If $\{\mathrm{hur}_k(\cdot)\}_k$ are admissible (or consistent), then $\max_k\{\mathrm{hur}_k\}(\cdot)$ is also admissible (or consistent).

Definition: If $\mathrm{hur}_{\max} \equiv \max\{\mathrm{hur}_1, \mathrm{hur}_2\}$, then if hur_k is consistent:

$$\mathrm{hur}_k(s) - \mathrm{hur}_k(\mathrm{tr}(s,a)) \leq \mathrm{cst}(s, a, \mathrm{tr}(s,a))$$

$$\mathrm{hur}_{\max}(s) = \mathrm{hur}_{\max}(\mathrm{tr}(s,a)) - \mathrm{cst}^*(s, a, \mathrm{tr}(s,a))$$

2.7.7 Anytime Search Algorithms

Definition: An **anytime algorithm** finds a solution quickly (even if it is sub-optimal), and then iteratively improves it (if time permits).

2.8 Canonical Examples

2.8.1 How to setup a search problem?

Process:

1. Given a search graph, we need to define the following:
   - $\mathcal{S}$: set of vertices
   - $\mathcal{G}$: goal states (subset of $\mathcal{S}$)
   - $s^{(0)}$: initial state
   - $\mathcal{T}$: set of edges (defined by $\text{tr}(\cdot, \cdot)$)
     - $\text{tr}(\cdot, \cdot)$: transition function
   - $\text{cst}(\cdot, \cdot, \cdot)$: cost function (defined by edge weights)

Example:

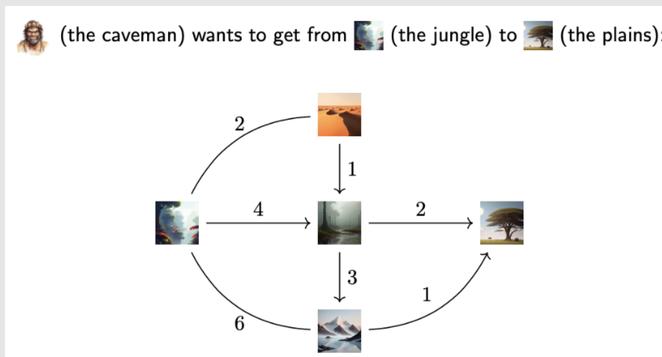


Figure 12

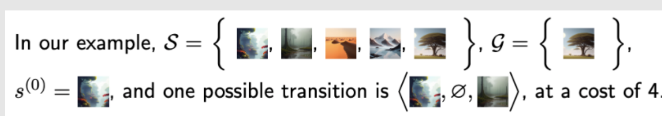


Figure 13

Example:

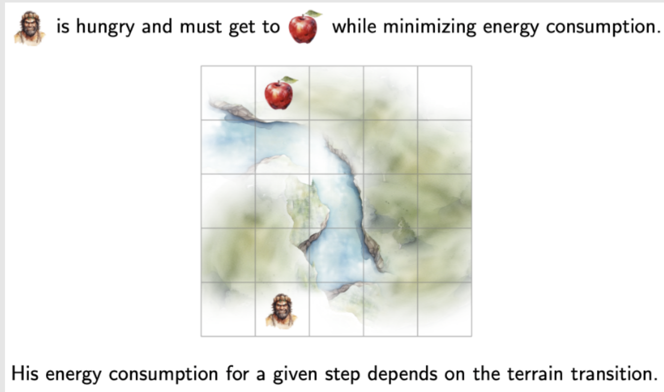


Figure 14

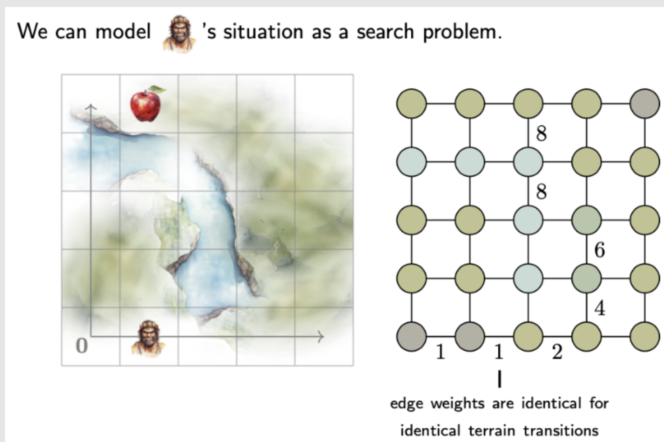


Figure 15

- $\mathcal{S} = \{0, \ldots, 4\}^2$
- $\mathcal{G} = \left\{ \begin{bmatrix} 1 \\ 4 \end{bmatrix} \right\}$
- $s^{(0)} = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$

2.8.2 How to setup a path tree?

Process:

1. Start at $s^{(0)}$
2. Choose a path until you reach a goal state.
3. Repeat until you have found all paths (probably infinite).

Example:

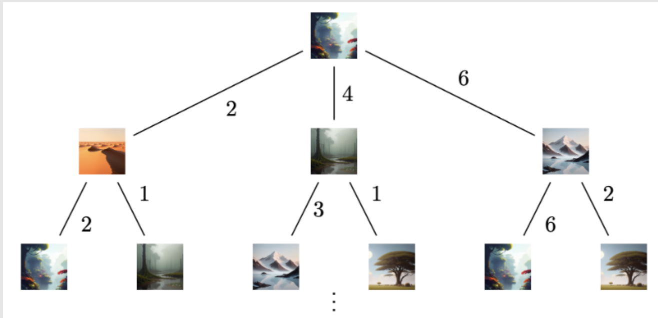


Figure 16

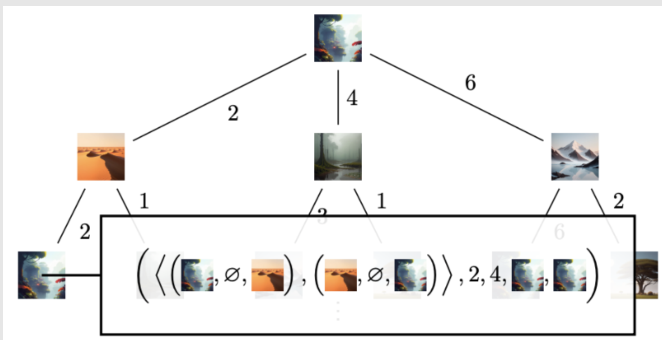


Figure 17

2.8.3 When to use each algorithm?

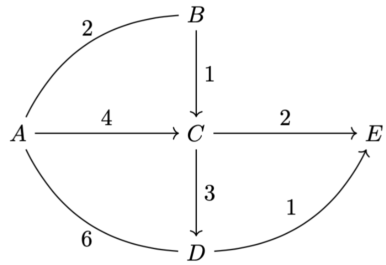
Process:

1. Do we have a heuristic?
   - **Yes:** Use informed search algorithms.
   - **No:** Use uninformed search algorithms.
2. Are path costs non-uniform?
   - **Yes:** Eliminate BFS.
   - **No:** Eliminate CFS, $A*$
3.
4. Is the search space finite or infinite?
   - **Finite:** Use any algorithm.
   - **Infinite:** Use BFS, IDDFS, CFS, or A*.
5. Do we need to guarantee finding a solution (completeness)?
   - **Yes:** Use BFS, IDDFS, IIA*, CFS (if $\epsilon > 0$).
   - **No:** Use DFS, HFS, WA*
6. Find properties needed for the problem and match them to the characteristics of the algorithm.
7. Choose the algorithm that best matches the properties.
   - **BFS:** Need shortest path in an unweighted graph.
   - **DFS:** Explore a deep path quickly, and completeness is not needed.
   - **IDDFS:** Want completeness of BFS but with the complexity of DFS.
   - **CFS:** Need the least-cost path in a weighted graph.
   - **HFS:**
   - **A*:**
   - **IIA*:**
   - **WA*:**

Example:

2.8.4 Tracing Search Algorithms

Example:



$$\text{hur}(s) = \begin{cases} 3 & s = A \\ 2 & s = B \\ 1 & s = C \\ 1 & s = D \\ 0 & s = E \end{cases}$$

Figure 18

Process: BFS

1. Start at s_0 as **current node**
2. Expand all neighboring nodes of the **current node** and add them to the open set (queue).
3. Remove the **current node** from the open set and add it to the path.
4. Choose the least-recently expanded node from the open set as the **current node**.
5. Repeat steps 2 and 4 until the goal state is reached or the open set is empty.

Example: BFS

Path	Open Set
	$\{A\}$
A	$\{AB, AC, AD\}$
AB	$\{AC, AD, ABA, ABC\}$
AC	$\{AD, ABA, ABC, ACD, ACE\}$
AD	$\{ABA, ABC, ACD, ACE, ADA, ADE\}$
ABA	$\{ABC, ACD, ACE, ADA, ADE, ABAB, ABAC, ABAD\}$
ABC	$\{ACD, ACE, ADA, ADE, ABAB, ABAC, ABAD, ABCD, ABCE\}$
ACD	$\{ACE, ADA, ADE, ABAB, ABAC, ABAD, ABCD, ABCE, ACDA, ACDE\}$
ACE	$\{ADA, ADE, ABAB, ABAC, ABAD, ABCD, ABCE, ACDA, ACDE\}$

Intra:

Path	Open Set
	$\{A\}$
A	$\{AB, AC, AD\}$
AB	$\{AC, AD, ABC, \cancel{ABA}\}$
AC	$\{AD, ABC, ACD, ACE\}$
AD	$\{ABC, ACD, ACE, ADE, \cancel{ADA}\}$
ABC	$\{ACD, ACE, ADE, ABCD, ABCE\}$
ACD	$\{ACE, ADE, ABCD, ABCE, ACDE, \cancel{ACBA}\}$
ACE	$\{ADE, ABCD, ABCE, ACDE\}$

Inter:

Path	Open Set	Closed Set
-	$\{A\}$	-
A	$\{AB, AC, AD\}$	$\{A\}$
AB	$\{AC, AD, ABC, \cancel{ABA}\}$	$\{A, B\}$
AC	$\{AD, ABC, ACD, ACE\}$	$\{A, B, C\}$
AD	$\{ABC, ACD, ACE, ADE, \cancel{ADA}\}$	$\{A, B, C, D\}$
ABC	$\{ACD, ACE, ADE, ABCE, \cancel{ABCD}\}$	$\{A, B, C, D\}$
ACD	$\{ACE, ADE, ABCE, ACDE, \cancel{ACDA}\}$	$\{A, B, C, D\}$
ACE	$\{ADE, ABCE, ACDE\}$	$\{A, B, C, D, E\}$

Process: DFS

1. Start at s_0 as **current node**
2. Expand all neighboring nodes of the **current node** and add them to the open set (stack).
3. Remove the **current node** from the open set and add it to the path.
4. Choose the most-recently expanded node from the open set as the **current node**.
5. Repeat steps 2 and 4 until the goal state is reached or the open set is empty.

Example: DFS

Path	**Open Set**
	$\{A\}$
A	$\{AB, AC, AD\}$
AD	$\{AB, AC, ADA, ADE\}$
ADE	$\{AB, AC, ADA\}$

Intra:

Path	**Open Set**
-	$\{A\}$
A	$\{AB, AC, AD\}$
AD	$\{AB, AC, ADE, \cancel{ADA}\}$
ADE	$\{AB, AC\}$

Inter:

Path	**Open Set**	**Closed Set**
-	$\{A\}$	-
A	$\{AB, AC, AD\}$	$\{A\}$
AD	$\{AB, AC, ADE, \cancel{ADA}\}$	$\{A, D\}$
ADE	$\{AB, AC\}$	$\{A, D, E\}$

Process: IDDFS

1. Start with a depth limit of 0.
2. Perform DFS up to the current depth limit.
3. If the goal state is not reached, increment the depth limit based on given fcn and repeat step 2.
4. Continue until the goal state is found or all nodes are explored.

Example: IDDFS

Depth	Path	Open Set
0		$\{A\}$
0	A	$\{\}$
1	A	$\{AB, AC, AD\}$
1	AD	$\{AB, AC\}$
1	AC	$\{AB\}$
1	AB	$\{\}$
2	AB	$\{ABA, ABC\}$
2	ABC	$\{ABA\}$
2	ABA	$\{\}$
3	ABA	$\{ABAB, ABAC, ABAD\}$
3	$ABAB$	$\{ABAC, ABAD\}$
3	$ABAC$	$\{ABAD\}$
3	$ABAD$	$\{\}$
4	$ABAD$	$\{ABADA, ABADE\}$
4	$ABADA$	$\{ABADE\}$
4	$ABADE$	$\{\}$

Process: CFS

1. Start at s_0 as **current node**
2. Expand all neighboring nodes of the **current node** and add them to the open set (priority queue).
3. Remove the **current node** from the open set and add it to the path.
4. Choose the cheapest expanded node from the open set as the **current node**.
5. Repeat steps 2 and 4 until the goal state is reached or the open set is empty.

Example: CFS

Path	Open Set
-	$\{A \mid 0\}$
A	$\{AB \mid 2,\ AC \mid 4,\ AD \mid 6\}$
AB	$\{AC \mid 4,\ AD \mid 6,\ ABC \mid 3,\ ABA \mid 4\}$
ABC	$\{AC \mid 4,\ AD \mid 6,\ ABA \mid 4,\ ABCE \mid 5,\ ABCD \mid 6\}$
AC	$\{AD \mid 6,\ ABA \mid 4,\ ABCE \mid 5,\ ABCD \mid 6,\ ACD \mid 7,\ ACE \mid 6\}$
ABA	$\{AD \mid 6,\ ABCE \mid 5,\ ABCD \mid 6,\ ACD \mid 7,\ ACE \mid 6,\ ABAB \mid 6,\ ABAC \mid 8,\ ABAD \mid 10\}$
$ABCE$	$\{AD \mid 6,\ ABCD \mid 6,\ ACD \mid 7,\ ACE \mid 6,\ ABAB \mid 6,\ ABAC \mid 8,\ ABAD \mid 10\}$

Intra:

Path	Open Set
-	$\{A \mid 0\}$
A	$\{AB \mid 2,\ AC \mid 4,\ AD \mid 6\}$
AB	$\{AC \mid 4,\ AD \mid 6,\ ABC \mid 3,\ \cancel{ABA}\}$
ABC	$\{AC \mid 4,\ AD \mid 6,\ ABCE \mid 5,\ ABCD \mid 6\}$
AC	$\{AD \mid 6,\ ABCE \mid 5,\ ABCD \mid 6,\ ACD \mid 7,\ ACE \mid 6\}$
$ABCE$	$\{AD \mid 6,\ ABCD \mid 6,\ ACD \mid 7,\ ACE \mid 6\}$

Inter:

Path	Open Set	Closed Set
-	$\{A \mid 0\}$	-
A	$\{AB \mid 2,\ AC \mid 4,\ AD \mid 6\}$	$\{A\}$
AB	$\{AC \mid 4,\ AD \mid 6,\ ABC \mid 3,\ \cancel{ABA}\}$	$\{A, B\}$
ABC	$\{AC \mid 4,\ AD \mid 6,\ ABCE \mid 5,\ ABCD \mid 6\}$	$\{A, B, C\}$
AC	$\{AD \mid 6,\ ABCE \mid 5,\ ABCD \mid 6,\ ACD \mid 7,\ ACE \mid 6\}$	$\{A, B, C\}$
$ABCE$	$\{AD \mid 6,\ ABCD \mid 6,\ ACD \mid 7,\ ACE \mid 6\}$	$\{A, B, C, E\}$

Process: HFS

1. Start at s_0 as **current node**
2. Expand all neighboring nodes of the **current node** and add them to the open set (priority queue).
3. Remove the **current node** from the open set and add it to the path.
4. Choose the lowest heuristic value expanded node from the open set as the **current node**.
5. Repeat steps 2 and 4 until the goal state is reached or the open set is empty.

Example: HFS

Path	**Open Set**
	$\{A \mid 3\}$
A	$\{AB \mid 2,\ AC \mid 1,\ AD \mid 1\}$
AC	$\{AB \mid 2,\ AD \mid 1,\ ACE \mid 0\}$
ACE	$\{AB \mid 2,\ AD \mid 1\}$

Process: A*

1. Start at s_0 as **current node**
2. Expand all neighboring nodes of the **current node** and add them to the open set (priority queue).
3. Remove the **current node** from the open set and add it to the path.
4. Choose the lowest $\text{esct}(p) = \text{cst}(p) + \text{hur}(p)$ expanded node from the open set as the **current node**.
5. Repeat steps 2 and 4 until the goal state is reached or the open set is empty.

Example: A*

Path	Open Set
-	$\{A \mid 3\}$
A	$\{AB \mid 2+2,\ AC \mid 4+1,\ AD \mid 6+1\}$
AB	$\{AC \mid 5,\ AD \mid 7,\ ABC \mid (2+1)+1,\ ABA \mid (2+2)+3\}$
ABC	$\{AC \mid 5,\ AD \mid 7,\ ABA \mid 7,\ ABCD \mid (2+1+3)+1,\ ABCE \mid (2+1+2)+0,\}$
AC	$\{AD \mid 7,\ ABA \mid 7,\ ABCD \mid 7,\ ABCE \mid 5,\ ACD \mid (4+3)+1,\ ACE \mid (4+2)+0\}$
$ABCE$	$\{AD \mid 7,\ ABA \mid 7,\ ABCD \mid 7,\ ACD \mid 8,\ ACE \mid 6\}$

Intra:

Path	Open Set
-	$\{A \mid 3\}$
A	$\{AB \mid 2+2,\ AC \mid 4+1,\ AD \mid 6+1\}$
AB	$\{AC \mid 5,\ AD \mid 7,\ ABC \mid (2+1)+1,\ \cancel{ABA}\}$
ABC	$\{AC \mid 5,\ AD \mid 7,\ ABCD \mid (2+1+3)+1,\ ABCE \mid (2+1+2)+0,\}$
AC	$\{AD \mid 7,\ ABCD \mid 7,\ ABCE \mid 5,\ ACD \mid (4+3)+1,\ ACE \mid (4+2)+0\}$
$ABCE$	$\{AD \mid 7,\ ABCD \mid 7,\ ACD \mid 8,\ ACE \mid 6\}$

Inter:

Path	Open Set	Closed Set
-	$\{A \mid 3\}$	-
A	$\{AB \mid 2+2,\ AC \mid 4+1,\ AD \mid 6+1\}$	$\{A\}$
AB	$\{AC \mid 5,\ AD \mid 7,\ ABC \mid (2+1)+1,\ \cancel{ABA}\}$	$\{A, B\}$
ABC	$\{AC \mid 5,\ AD \mid 7,\ ABCD \mid (2+1+3)+1,\ ABCE \mid (2+1+2)+0,\}$	$\{A, B, C\}$
AC	$\{AD \mid 7,\ ABCD \mid 7,\ ABCE \mid 5,\ ACD \mid (4+3)+1,\ ACE \mid (4+2)+0\}$	$\{A, B, C\}$
$ABCE$	$\{AD \mid 7,\ ABCD \mid 7,\ ACD \mid 8,\ ACE \mid 6\}$	$\{A, B, C, E\}$

Process: IIA*

1. Start with a cost limit of 0.
2. Perform A* up to the current cost limit.
3. If the goal state is not reached, increment the cost limit based on given fcn and repeat step 2.
4. Continue until the goal state is found or all nodes are explored.

Example: IIA*

Cost	Path	Open Set
0	$\langle\rangle$	$\{\}$
1	$\langle\rangle$	$\{\}$
2	$\langle\rangle$	$\{\}$
3	$\langle\rangle$	$\{A \mid 3\}$
3	A	$\{\}$
4	A	$\{AB \mid 2+2\}$
4	AB	$\{ABC \mid 3+1\}$
4	ABC	$\{\}$
5	ABC	$\{ABCE \mid 5+0\}$
5	$ABCE$	$\{\}$

Process: WA*

1. Start at s_0 as **current node**
2. Expand all neighboring nodes of the **current node** and add them to the open set (priority queue).
3. Remove the **current node** from the open set and add it to the path.
4. Choose the lowest $\text{esct}(p) = w \cdot \text{cst}(p) + (1 - w) \cdot \text{hur}(p)$ expanded node from the open set as the **current node**.
5. Repeat steps 2 and 4 until the goal state is reached or the open set is empty.

Process: How to Prove Consistent/Admissible Given a Search Graph?
Admissible:

1. Given $\text{hur}(s)$ and search graph with $\text{cst}(s, a, \text{tr}(s, a))$. If consistent, then it is admissible.
2. Check $\forall s \in \mathcal{G}$, $\text{hur}(s) = 0$. If not, then it is not admissible.
3. For each $s \in \mathcal{S}$, calculate $\text{hur}^*(s)$ (i.e. actual cost of optimal soln.) using the search graph.
   (a) Start at s and choose path that gives the lowest cost to $s \in \mathcal{G}$.
4. Check if $\text{hur}(s) \leq \text{hur}^*(s)$ $\forall s \in \mathcal{S}$. If not, then it is not admissible.
5. Repeat $\forall s \in \mathcal{S}$.
6. If all are true, then it is admissible.

Consistent:

1. Given $\text{hur}(s)$ and search graph with $\text{cst}(s, a, \text{tr}(s, a))$.
2. Check $\forall s \in \mathcal{G}$, $\text{hur}(s) = 0$. If not, then it is not consistent.
3. For each $s \in \mathcal{S}$, calculate $\text{hur}(s) - \text{hur}(\text{tr}(s, a))$.
   (a) check if it is $\leq \text{cst}(s, a, \text{tr}(s, a))$. If not, then it is not consistent.
   (b) Repeat $\forall a \in \mathcal{A}(s)$
4. Repeat $\forall s \in \mathcal{S}$.
5. If all are true, then it is consistent.

Warning: Be careful of bidirectional edges bc for consistency you need compute the cost of the heuristic edge in both directions.

Example:

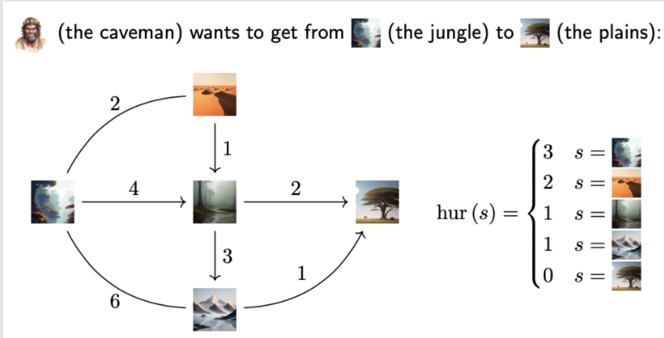


Figure 19: Jungle ($s^{(0)}$), Desert, Swamp, Mountain, Plains (Goal)

Admissible:

1. $s =$**Plains:** $\text{hur}(\text{Plains}) = 0$
2. $s =$**Jungle:** $\text{hur}(\text{Jungle}) = 3 \leq \text{hur}^*(\text{Jungle}) = 2 + 1 + 2 = 5$
3. $s =$**Desert:** $\text{hur}(\text{Desert}) = 2 \leq \text{hur}^*(\text{Desert}) = 1 + 2$
4. $s =$**Swamp:** $\text{hur}(\text{Swamp}) = 1 \leq \text{hur}^*(\text{Swamp}) = 2$
5. $s =$**Mountain:** $\text{hur}(\text{Mountain}) = 1 \leq \text{hur}^*(\text{Mountain}) = 1$
6. Therefore, it is admissible.

Consistent:

1. $s =$**Plains:** $\text{hur}(\text{Plains}) = 0$
2. $s =$**Jungle:**
   (a) $\text{hur}(\text{Jungle}) - \text{hur}(\text{Desert}) = 3 - 2 = 1 \leq \text{cst}(\text{Jungle}, \cdot, \text{Desert}) = 2$
   (b) $\text{hur}(\text{Jungle}) - \text{hur}(\text{Swamp}) = 3 - 1 = 2 \leq \text{cst}(\text{Jungle}, \cdot, \text{Swamp}) = 4$
   (c) $\text{hur}(\text{Jungle}) - \text{hur}(\text{Mountain}) = 3 - 1 = 2 \leq \text{cst}(\text{Jungle}, \cdot, \text{Mountain}) = 6$
3. $s =$**Desert:**
   (a) $\text{hur}(\text{Desert}) - \text{hur}(\text{Jungle}) = 2 - 3 = -1 \leq \text{cst}(\text{Desert}, \cdot, \text{Jungle}) = 2$
   (b) $\text{hur}(\text{Desert}) - \text{hur}(\text{Swamp}) = 2 - 1 = 1 \leq \text{cst}(\text{Desert}, \cdot, \text{Swamp}) = 1$
4. $s =$**Swamp:**
   (a) $\text{hur}(\text{Swamp}) - \text{hur}(\text{Mountain}) = 1 - 1 = 0 \leq \text{cst}(\text{Swamp}, \cdot, \text{Mountain}) = 3$
   (b) $\text{hur}(\text{Swamp}) - \text{hur}(\text{Plains}) = 1 - 0 = 1 \leq \text{cst}(\text{Swamp}, \cdot, \text{Plains}) = 2$
5. $s =$**Mountain:**
   (a) $\text{hur}(\text{Mountain}) - \text{hur}(\text{Jungle}) = 1 - 3 = -2 \leq \text{cst}(\text{Mountain}, \cdot, \text{Desert}) = 6$
   (b) $\text{hur}(\text{Mountain}) - \text{hur}(\text{Plains}) = 1 - 0 = 1 \leq \text{cst}(\text{Mountain}, \cdot, \text{Plains}) = 1$
6. Therefore, it is consistent.

Process: How to Design Heuristic via Problem Relaxation?

1. Make an assumption to simplify the problem as a relaxed problem.
2. Find the cost of the optimal solution of the relaxed problem, $\mathrm{cst}_{\mathrm{rel}}(s)$.
3. HOW TO FIND THE COST OF THE OPTIMAL SOLUTION?

Example:

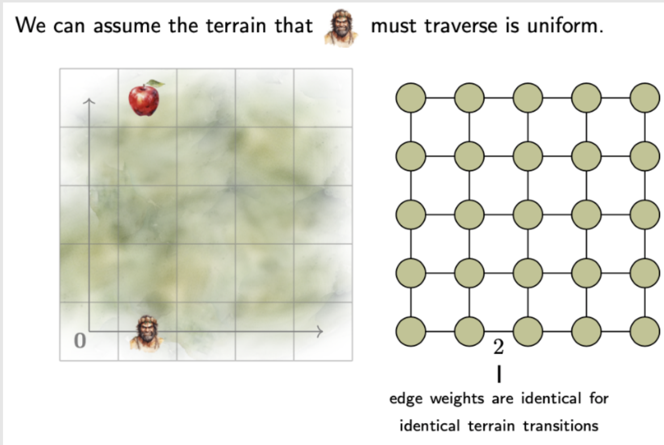


Figure 20

3 Constraint Satisfaction Problems

3.1 Setup of CSP

Definition: A **constraint satisfaction problem (CSP)** consists of:
- a set of **variables**, $\mathcal{V}$, where the domain of $V \in \mathcal{V}$ is $\mathrm{dom}(V)$
- a set of **constraints**, $\mathcal{C}$, where the scope of $C \in \mathcal{C}$ is $\mathrm{scp}(C) \subseteq \mathcal{V}$

3.2 Assignment

Definition: An **assignment** is a set of pairs, $\{(V, v)\}_{V \in \tilde{\mathcal{V}}}$, where $v \in \mathrm{dom}(V)$, and $\tilde{\mathcal{V}} \subseteq \mathcal{V}$. It is **complete** if $\tilde{\mathcal{V}} = \mathcal{V}$, and **partial** otherwise.

3.3 Formulating a CSP as a Search Problem

Motivation: We don't formulate a CSP as a search problem because the path tree of all possible ways to build a complete assignment is too large. The number of paths in the three is

$$\mathcal{O}\left(|\mathcal{V}|! \times b^d\right)$$

- $b = \max_{V \in \mathcal{V}} |\mathrm{dom}(V)|$
- $d = |\mathcal{V}|$

3.4 Consistent

3.4.1 Complete Assignment

Definition: A complete assignment, A, is **consistent** if it satisfies every constraint $\mathcal{C}$ with $\mathrm{scp}(\mathcal{C}) \subseteq \tilde{\mathcal{V}}$.

Warning: A solution to a CSP is any complete and consistent assignment.

3.4.2 Partial Assignment

Definition: A (possibly partial) assignment, $\{(V, v)\}_{V \in \tilde{\mathcal{V}}}$, is **consistent** if it satisfies every constraint, $C \in \mathcal{C}$ such that $\mathrm{scp}(C) \subseteq \tilde{\mathcal{V}}$.

3.4.3 k-Consistent

Definition: A CSP is ***k*-consistent** if for any consistent assignment of $k-1$ variables, $\{(V, v)\}_{V \in \tilde{\mathcal{V}}}$, and any k^{th} variable, V', there is a value, $v' \in \mathrm{dom}(V')$, so the assignment, $\{(V, v)\}_{V \in \tilde{\mathcal{V}}} \cup \{(V', v')\}$ is consistent.

Notes:
- Edge/Arc Consistent: $k = 2$

3.5 Constraint Satisfaction Algorithm

Algorithm:

```
A ← {}                                          ▷ initialize an empty assignment
for V ∈ 𝒱 do 𝒟(V) ← COPY(dom(V))
SATISFY(𝒱, 𝒞, 𝒟, A)
```

3.5.1 Satisfy

Algorithm:

```
procedure SATISFY(𝒱, 𝒞, 𝒟, A):
    if COMPLETE(A, 𝒱) then
        return A                                ▷ a solution was found
    V ← REMOVE(𝒱, A)
    for v ∈ 𝒟(V) do                             ▷ try each value in V's current domain
        𝒟′ ← COPY(𝒟)                            ▷ cache the current domains for backtracking
        A ← A ∪ {(V, v)}
        𝒟(V) ← {v}
        𝒟, success ← ENFORCE(𝒱, 𝒞, 𝒟, V, k)
        if success then                         ▷ enforce k consistency
            A ← SATISFY(𝒱, 𝒞, 𝒟, A)             ▷ recursively continue if possible
            if A ≠ NULL then
                return A
        𝒟 ← 𝒟′                                  ▷ backtrack if not possible
        A ← A \ {(V, v)}
    return NULL                                 ▷ No solution found in this branch
```

3.5.2 Enforce: Enforcing k-Consistency

Algorithm: Pre-pruning: Enforce without assigning any variables.

```
procedure ENFORCE(𝒱, 𝒞, 𝒟, V, k):
    Q ← {C ∈ 𝒞 s.t. V ∈ scp(C)}                 ▷ initialize affected constraints
    while Q ≠ ∅ do
        C ← REMOVE(Q)
        for V′ ∈ scp(C) do                      ▷ enforce consistency on each variable w.r.t. each affected constraint
            success ← ENFORCEVAR(k, V′, 𝒱, 𝒞, 𝒟)
            if not success then
                return False                    ▷ consistency could not be enforced
            Q ← Q ∪ {C′ ∈ 𝒞 | V′ ∈ scp(C′)}
    return True                                 ▷ consistency was enforced
```

3.5.3 EnforceVar: Enforcing k-Consistency

Algorithm:

```
procedure ENFORCEVAR(𝒱, 𝒞, 𝒟, V, k):
    for v ∈ 𝒟(V) do
        for C ∈ 𝒞 do
            if V ∈ scp(C) and |scp(C)| ≤ k then
                flag ← False
                for A ∈ 𝒳 × 𝒟(V′) do
                    if A ∪ {(V, v)} ∈ 𝒞 then
                        flag ← True
                        break
                if not flag then
                    𝒟(V) ← 𝒟(V) \ {v}
                if 𝒟(V) = ∅ then
                    return False                ▷ no valid domain values remain for V
    return True
```

3.6 Canonical Problems

Process: Setup of CSP:

1. Determine variables to track, domain of each variable, and constraints.

Example:

now wants to find food to meet his nutritional requirements:

	Supply	Nutrients (25 g) Carbs	Fat	Protein	Vitamins
Minimum	- - - -	8	3	2	1
Nuts	4	2	1	1	0
Fruits	5	2	0	0	0
Legumes	4	2	0	1	1
Grains	6	3	0	0	0
Meat	3	0	1	2	0
Maximum	- - - -	10	4	5	1

Figure 21: Information

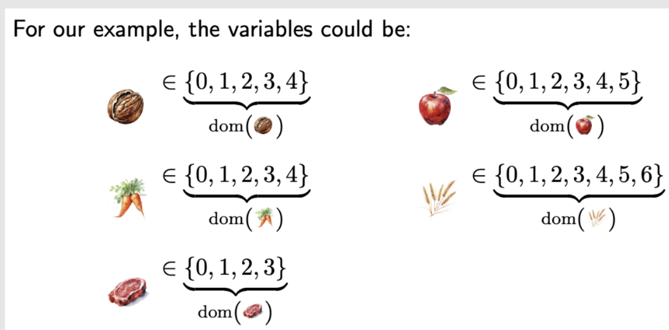


Figure 22: Variables

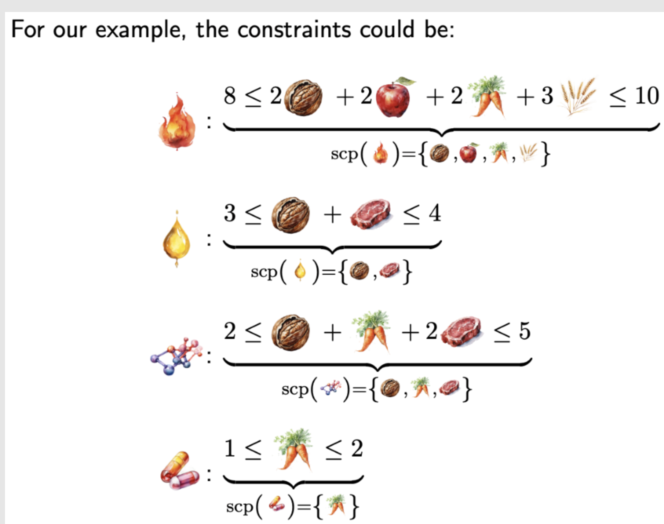


Figure 23: Constraints

Process: How to build a hyper-graph?

1. Circle the variables that appear in constraint C_i $\forall i$.

Example:

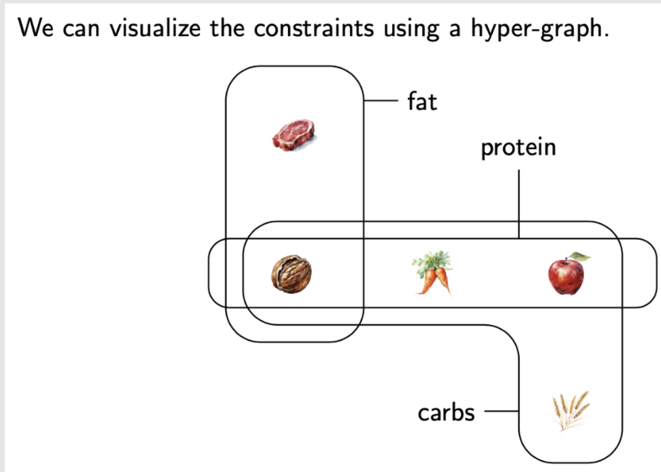


Figure 24

Process: How to Enforce k-Consistency?

1. Given $\mathcal{V}$ w/ $\text{dom}(V) = \{v_1, \ldots, v_{|\text{dom}(V)|}\}$ $\forall V \in \mathcal{V}$ and $\mathcal{C}$ w/ $\text{scp}(C) = \{V_1, \ldots, V_{|\text{scp}(C)|}\}$ $\forall C \in \mathcal{C}$.
2. Remove all constraints that have $k+1$ or more variables and add the rest to a queue.
3. **Pre-pruning:** For each remaining $C \in \mathcal{C}$, do the following:
   (a) For each $V \in \text{scp}(C)$, do the following:
      i. For each $v \in \text{dom}(V)$, do the following:
      - Fix V to v.
      - For the other $V \in \text{scp}(C)$, check if the constraint is satisfied by trying all combinations (need only one).
      - **Key:** If the constraint is not satisfied, then remove the value from $\text{dom}(V)$. Add any affected constraints back to the queue.
4. Repeat until the queue is empty.

Warning: Can think of checking as picking $k-1$ variables, then choosing any value for the k^{th} variable that satisfies all constraints. While enforcing is fixing a variable to a value, then checking if there is a combination for the other variables that satisfies all constraints.

Warning: Enforcing k-consistency is enforcing $k-1, \ldots, 1$-consistency.

Process: How to determine a solution to a CSP?

1. After pre-pruning the domains.
2. Assign variables in alphabetical order and values in numerical order.
3. Prune the pre-pruned domains.
4. If you can assign all variables, then you have a solution. If you have domain wipeout, backtrack.
5. Repeat the process until you find all solutions.

Process: Checking k-Consistency

1. Enforce k-consistency.
2. If you have to pre-prune, then not k-consistent.

Example: Pre-Pruning Domains

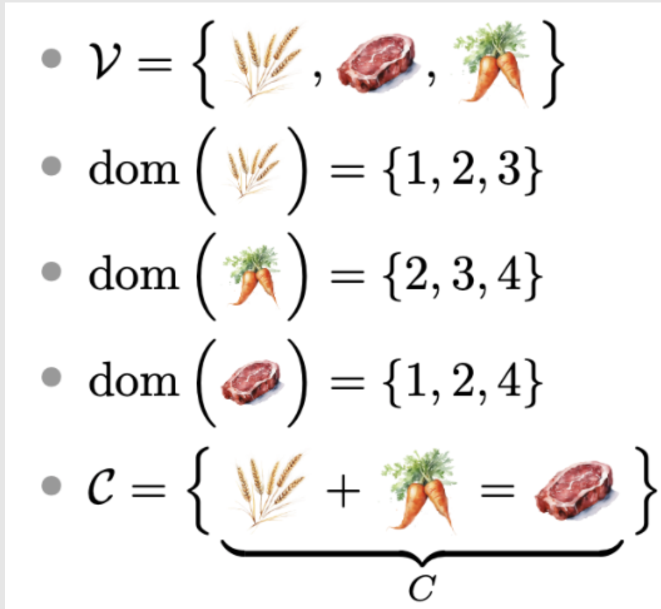


Figure 25

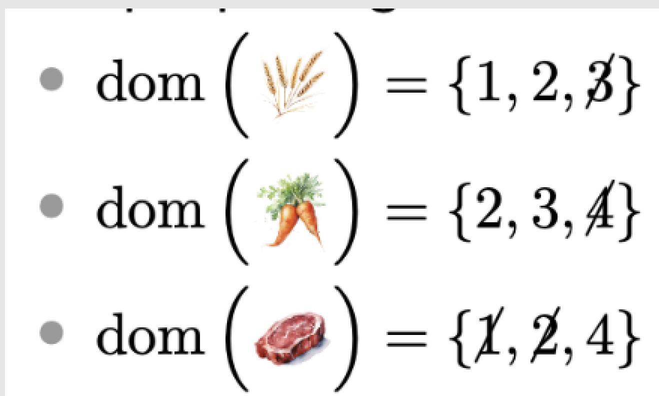


Figure 26: Pre-pruning. Since only one constraint, it is also pruning.

Example:

1. **Given:** Consider a CSP in which $\mathcal{V} = \{A, B, C, D, E\}$, where:

$$\begin{aligned}
\text{dom}(A) &= \{0, 1, 2, 3, 4\} \\
\text{dom}(B) &= \{0, 1, 2, 3, 4\} \\
\text{dom}(C) &= \{0, 1, 2, 3\} \\
\text{dom}(D) &= \{0, 1, 2, 3, 4, 5\} \\
\text{dom}(E) &= \{0, 1, 2, 3, 4, 5, 6\}
\end{aligned}$$

   and $\mathcal{C} = \{C_1, C_2, C_3, C_4\}$, where:

$$\begin{aligned}
&C_1 : 8 \leq 2a + 2b + 2d + 3e \leq 10 \\
&C_2 : 3 \leq a + c \leq 4 \\
&C_3 : 2 \leq a + b + 2c \leq 5 \\
&C_4 : 1 \leq b \leq 2
\end{aligned}$$

2. **Problem:** Solve the following CSP using $k = 4$ consistency. Pre-prune the domains using $k = 4$ consistency. Assign variables in alphabetical order and values in numerical order.

Example: 4-Consistency Pre-Pruning

Queue	**Fixed Value**	**Satisfactory Combination?**
		$C_4 : 1 \leq b \leq 2$
$\{C_2, C_3, C_1\}$	$b = 0$, $b = 1$, $b = 2$, $b = 3$, $b = 4$	No, Yes, Yes, No, No

- $\text{dom}(A) = \{0, 1, 2, 3, 4\}$
- $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, 2, 3\}$
- $\text{dom}(D) = \{0, 1, 2, 3, 4, 5\}$
- $\text{dom}(E) = \{0, 1, 2, 3, 4, 5, 6\}$
- $\{C_2, C_3, C_1\}$

		$C_2 : 3 \leq a + c \leq 4$
$\{C_3, C_1\}$	$a = 0$, $a = 1$, $a = 2$, $a = 3$, $a = 4$	Yes, Yes, Yes, Yes, Yes
-	$c = 0$, $c = 1$, $c = 2$, $c = 3$	Yes, Yes, Yes, Yes

- $\text{dom}(A) = \{0, 1, 2, 3, 4\}$
- $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, 2, 3\}$
- $\text{dom}(D) = \{0, 1, 2, 3, 4, 5\}$
- $\text{dom}(E) = \{0, 1, 2, 3, 4, 5, 6\}$
- $\{C_3, C_1\}$

		$C_3 : 2 \leq a + b + 2c \leq 5$
$\{C_1\}$	$a = 0$, $a = 1$, $a = 2$, $a = 3$, $a = 4$	Yes, Yes, Yes, Yes, Yes
-	$b = 1$, $b = 2$	Yes, Yes
-	$c = 0$, $c = 1$, $c = 2$, $c = 3$	Yes, Yes, Yes, No

- $\text{dom}(A) = \{0, 1, 2, 3, 4\}$
- $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, 2, \cancel{3}\}$
- $\text{dom}(D) = \{0, 1, 2, 3, 4, 5\}$
- $\text{dom}(E) = \{0, 1, 2, 3, 4, 5, 6\}$
- $\{C_1, C_2\}$

Example: 4-Consistency Continued:

Queue	Fixed Value	Satisfactory Combination?
	$C_1 : 8 \leq 2a + 2b + 2d + 3e \leq 10$	
$\{C_2\}$	$a = 0$, $a = 1$, $a = 2$, $a = 3$, $a = 4$	Yes, Yes, Yes, Yes, Yes
-	$b = 1$, $b = 2$	Yes, Yes
-	$d = 0$, $d = 1$, $d = 2$, $d = 3$, $d = 4$, $d = 5$	Yes, Yes, Yes, Yes, Yes, No
-	$e = 0$, $e = 1$, $e = 2$, $e = 3$, $e = 4$, $e = 5$, $e = 6$	Yes, Yes, Yes, No, No, No, No

- $\text{dom}(A) = \{0, 1, 2, 3, 4\}$
- $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, 2, \cancel{3}\}$
- $\text{dom}(D) = \{0, 1, 2, 3, 4, \cancel{5}\}$
- $\text{dom}(E) = \{0, 1, 2, \cancel{3}, \cancel{4}, \cancel{5}, \cancel{6}\}$
- $\{C_2\}$

Queue	Fixed Value	Satisfactory Combination?
	$C_2 : 3 \leq a + c \leq 4$	
$\{\}$	$a = 0$, $a = 1$, $a = 2$, $a = 3$, $a = 4$	No, Yes, Yes, Yes, Yes
-	$c = 0$, $c = 1$, $c = 2$	Yes, Yes, Yes

- $\text{dom}(A) = \{\cancel{0}, 1, 2, 3, 4\}$
- $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, 2, \cancel{3}\}$
- $\text{dom}(D) = \{0, 1, 2, 3, 4, \cancel{5}\}$
- $\text{dom}(E) = \{0, 1, 2, \cancel{3}, \cancel{4}, \cancel{5}, \cancel{6}\}$
- $\{C_3, C_1\}$

Queue	Fixed Value	Satisfactory Combination?
	$C_3 : 2 \leq a + b + 2c \leq 5$	
$\{C_1\}$	$a = 1$, $a = 2$, $a = 3$, $a = 4$	Yes, Yes, Yes, Yes
-	$b = 1$, $b = 2$	Yes, Yes
-	$c = 0$, $c = 1$, $c = 2$	Yes, Yes, No

- $\text{dom}(A) = \{\cancel{0}, 1, 2, 3, 4\}$
- $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, \cancel{2}, \cancel{3}\}$
- $\text{dom}(D) = \{0, 1, 2, 3, 4, \cancel{5}\}$
- $\text{dom}(E) = \{0, 1, 2, \cancel{3}, \cancel{4}, \cancel{5}, \cancel{6}\}$
- $\{C_1, C_2\}$

Example: 4-Consistency Continued:

Queue	**Fixed Value**	**Satisfactory Combination?**

$$C_1 : 8 \leq 2a + 2b + 2d + 3e \leq 10$$

$\{C_2\}$	$a = 1$, $a = 2$, $a = 3$, $a = 4$	Yes, Yes, Yes, Yes
-	$b = 1$, $b = 2$	Yes, Yes
-	$d = 0$, $d = 1$, $d = 2$, $d = 3$, $d = 4$	Yes, Yes, Yes, Yes, No
-	$e = 0$, $e = 1$, $e = 2$	Yes, Yes, Yes

- $\text{dom}(A) = \{\cancel{0}, 1, 2, 3, 4\}$
- $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, \cancel{2}, \cancel{3}\}$
- $\text{dom}(D) = \{0, 1, 2, 3, \cancel{4}, \cancel{5}\}$
- $\text{dom}(E) = \{0, 1, 2, \cancel{3}, \cancel{4}, \cancel{5}, \cancel{6}\}$
- $\{C_2\}$

$$C_2 : 3 \leq a + c \leq 4$$

$\{\}$	$a = 1$, $a = 2$, $a = 3$, $a = 4$	No, Yes, Yes, Yes
-	$c = 0$, $c = 1$	Yes, Yes

- $\text{dom}(A) = \{\cancel{0}, \cancel{1}, 2, 3, 4\}$
- $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, \cancel{2}, \cancel{3}\}$
- $\text{dom}(D) = \{0, 1, 2, 3, \cancel{4}, \cancel{5}\}$
- $\text{dom}(E) = \{0, 1, 2, \cancel{3}, \cancel{4}, \cancel{5}, \cancel{6}\}$
- $\{C_3, C_1\}$

$$C_3 : 2 \leq a + b + 2c \leq 5$$

$\{C_1\}$	$a = 2$, $a = 3$, $a = 4$	Yes, Yes, Yes
-	$b = 1$, $b = 2$	Yes, Yes
-	$c = 0$, $c = 1$	Yes, Yes

- $\text{dom}(A) = \{\cancel{0}, \cancel{1}, 2, 3, 4\}$, $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, \cancel{2}, \cancel{3}\}$, $\text{dom}(D) = \{0, 1, 2, 3, \cancel{4}, \cancel{5}\}$
- $\text{dom}(E) = \{0, 1, 2, \cancel{3}, \cancel{4}, \cancel{5}, \cancel{6}\}$
- $\{C_1\}$

$$C_1 : 8 \leq 2a + 2b + 2d + 3e \leq 10$$

$\{\}$	$a = 2$, $a = 3$, $a = 4$	Yes, Yes, Yes
-	$b = 1$, $b = 2$	Yes, Yes
-	$d = 0$, $d = 1$, $d = 2$, $d = 3$	Yes, Yes, Yes, No
-	$e = 0$, $e = 1$, $e = 2$	Yes, Yes, No

- $\text{dom}(A) = \{\cancel{0}, \cancel{1}, 2, 3, 4\}$, $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, \cancel{2}, \cancel{3}\}$, $\text{dom}(D) = \{0, 1, 2, \cancel{3}, \cancel{4}, \cancel{5}\}$
- $\text{dom}(E) = \{0, 1, \cancel{2}, \cancel{3}, \cancel{4}, \cancel{5}, \cancel{6}\}$
- $\{\}$

Example: 4-Consistency Post-Pre-Pruning:

$$C_1 : 8 \leq 2a + 2b + 2d + 3e \leq 10$$
$$C_2 : 3 \leq a + c \leq 4$$
$$C_3 : 2 \leq a + b + 2c \leq 5$$
$$C_4 : 1 \leq b \leq 2$$

Solution	**Updated Necessary Domains After Assignment**
$A = 2$	$\text{dom}(B) = \{1,2\}$, $\text{dom}(C) = \{\cancel{0},1\}$, $\text{dom}(D) = \{0,1,2\}$, $\text{dom}(E) = \{0,1\}$
$A = 2, B = 1$	$\text{dom}(C) = \{\cancel{0},1\}$, $\text{dom}(D) = \{0,1,2\}$, $\text{dom}(E) = \{0,1\}$
$A = 2, B = 1, C = 1$	$\text{dom}(D) = \{0,1,2\}$, $\text{dom}(E) = \{0,1\}$
$A = 2, B = 1, C = 1, D = 0$	$\text{dom}(E) = \{\cancel{0},1\}$
$A = 2, B = 1, C = 1, D = 0, E = 1$	**Solution Found**
$A = 2$	$\text{dom}(B) = \{1,2\}$, $\text{dom}(C) = \{\cancel{0},1\}$, $\text{dom}(D) = \{0,1,2\}$, $\text{dom}(E) = \{0,1\}$
$A = 2, B = 1$	$\text{dom}(C) = \{\cancel{0},1\}$, $\text{dom}(D) = \{0,1,2\}$, $\text{dom}(E) = \{0,1\}$
$A = 2, B = 1, C = 1$	$\text{dom}(D) = \{0,1,2\}$, $\text{dom}(E) = \{0,1\}$
$A = 2, B = 1, C = 1, D = 1$	$\text{dom}(E) = \{0,\cancel{1}\}$
$A = 2, B = 1, C = 1, D = 1, E = 0$	**Solution Found**
$A = 2$	$\text{dom}(B) = \{1,2\}$, $\text{dom}(C) = \{\cancel{0},1\}$, $\text{dom}(D) = \{0,1,2\}$, $\text{dom}(E) = \{0,1\}$
$A = 2, B = 1$	$\text{dom}(C) = \{\cancel{0},1\}$, $\text{dom}(D) = \{0,1,2\}$, $\text{dom}(E) = \{0,1\}$
$A = 2, B = 1, C = 1$	$\text{dom}(D) = \{0,1,2\}$, $\text{dom}(E) = \{0,1\}$
$A = 2, B = 1, C = 1, D = 2$	$\text{dom}(E) = \{0,\cancel{1}\}$
$A = 2, B = 1, C = 1, D = 2, E = 0$	**Solution Found**
$A = 2$	$\text{dom}(B) = \{1,2\}$, $\text{dom}(C) = \{\cancel{0},1\}$, $\text{dom}(D) = \{0,1,2\}$, $\text{dom}(E) = \{0,1\}$
$A = 2, B = 2$	$\text{dom}(C) = \{\cancel{0},\cancel{1}\}$, $\text{dom}(D) = \{0,1,\cancel{2}\}$, $\text{dom}(E) = \{0,\cancel{1}\}$
-	**No Solution Found**
$A = 3$	$\text{dom}(B) = \{1,2\}$, $\text{dom}(C) = \{0,1\}$, $\text{dom}(D) = \{0,1,2\}$, $\text{dom}(E) = \{0,1\}$
$A = 3, B = 1$	$\text{dom}(C) = \{0,\cancel{1}\}$, $\text{dom}(D) = \{0,1,\cancel{2}\}$, $\text{dom}(E) = \{0,\cancel{1}\}$
$A = 3, B = 1, C = 0$	$\text{dom}(D) = \{0,1,\cancel{2}\}$, $\text{dom}(E) = \{0,\cancel{1}\}$
$A = 3, B = 1, C = 0, D = 0$	$\text{dom}(E) = \{0,\cancel{1}\}$
$A = 3, B = 1, C = 0, D = 0, E = 0$	**Solution Found**
$A = 3$	$\text{dom}(B) = \{1,2\}$, $\text{dom}(C) = \{0,1\}$, $\text{dom}(D) = \{0,1,2\}$, $\text{dom}(E) = \{0,1\}$
$A = 3, B = 1$	$\text{dom}(C) = \{0,\cancel{1}\}$, $\text{dom}(D) = \{0,1,\cancel{2}\}$, $\text{dom}(E) = \{0,\cancel{1}\}$
$A = 3, B = 1, C = 0$	$\text{dom}(D) = \{0,1,\cancel{2}\}$, $\text{dom}(E) = \{0,\cancel{1}\}$
$A = 3, B = 1, C = 0, D = 1$	$\text{dom}(E) = \{0,\cancel{1}\}$
$A = 3, B = 1, C = 0, D = 1, E = 0$	**Solution Found**
$A = 3$	$\text{dom}(B) = \{1,2\}$, $\text{dom}(C) = \{0,1\}$, $\text{dom}(D) = \{0,1,2\}$, $\text{dom}(E) = \{0,1\}$
$A = 3, B = 2$	$\text{dom}(C) = \{0,\cancel{1}\}$, $\text{dom}(D) = \{0,\cancel{1},\cancel{2}\}$, $\text{dom}(E) = \{0,\cancel{1}\}$
$A = 3, B = 2, C = 0$	$\text{dom}(D) = \{0,\cancel{1},\cancel{2}\}$, $\text{dom}(E) = \{0,\cancel{1}\}$
$A = 3, B = 2, C = 0, D = 0$	$\text{dom}(E) = \{0,\cancel{1}\}$
$A = 3, B = 2, C = 0, D = 0, E = 0$	**Solution Found**
$A = 4$	$\text{dom}(B) = \{1,\cancel{2}\}$, $\text{dom}(C) = \{0,\cancel{1}\}$, $\text{dom}(D) = \{0,1,\cancel{2}\}$, $\text{dom}(E) = \{0,\cancel{1}\}$
$A = 4, B = 1$	$\text{dom}(C) = \{0,\cancel{1}\}$, $\text{dom}(D) = \{0,\cancel{1},\cancel{2}\}$, $\text{dom}(E) = \{0,\cancel{1}\}$
$A = 4, B = 1, C = 0$	$\text{dom}(D) = \{0,\cancel{1},\cancel{2}\}$, $\text{dom}(E) = \{0,\cancel{1}\}$
$A = 4, B = 1, C = 0, D = 0$	$\text{dom}(E) = \{0,\cancel{1}\}$
$A = 4, B = 1, C = 0, D = 0, E = 0$	**Solution Found**

Example: 3-Consistency

Fixed Value	Satisfactory Combination?
$C_4 : 1 \leq b \leq 2$	
$b = 0$, $b = 1$, $b = 2$, $b = 3$, $b = 4$	No, Yes, Yes, No, No

- $\text{dom}(A) = \{0, 1, 2, 3, 4\}$
- $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, 2, 3\}$
- $\text{dom}(D) = \{0, 1, 2, 3, 4, 5\}$
- $\text{dom}(E) = \{0, 1, 2, 3, 4, 5, 6\}$

Fixed Value	Satisfactory Combination?
$C_2 : 3 \leq a + c \leq 4$	
$a = 0$, $a = 1$, $a = 2$, $a = 3$, $a = 4$	Yes, Yes, Yes, Yes, Yes
$c = 0$, $c = 1$, $c = 2$, $c = 3$	Yes, Yes, Yes, Yes

- $\text{dom}(A) = \{0, 1, 2, 3, 4\}$
- $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, 2, 3\}$
- $\text{dom}(D) = \{0, 1, 2, 3, 4, 5\}$
- $\text{dom}(E) = \{0, 1, 2, 3, 4, 5, 6\}$

Fixed Value	Satisfactory Combination?
$C_3 : 2 \leq a + b + 2c \leq 5$	
$a = 0$, $a = 1$, $a = 2$, $a = 3$, $a = 4$	Yes, Yes, Yes, Yes, Yes
$b = 1$, $b = 2$	Yes, Yes
$c = 0$, $c = 1$, $c = 2$, $c = 3$	Yes, Yes, Yes, No

- $\text{dom}(A) = \{0, 1, 2, 3, 4\}$
- $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, 2, \cancel{3}\}$
- $\text{dom}(D) = \{0, 1, 2, 3, 4, 5\}$
- $\text{dom}(E) = \{0, 1, 2, 3, 4, 5, 6\}$

Example:

Fixed Value	**Satisfactory Combination?**

$$C_2 : 3 \leq a + c \leq 4$$

$a = 0$, $a = 1$, $a = 2$, $a = 3$, $a = 4$	No, Yes, Yes, Yes, Yes
$c = 0$, $c = 1$, $c = 2$	Yes, Yes, Yes

- $\text{dom}(A) = \{\cancel{0}, 1, 2, 3, 4\}$
- $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, 2, \cancel{3}\}$
- $\text{dom}(D) = \{0, 1, 2, 3, 4, 5\}$
- $\text{dom}(E) = \{0, 1, 2, 3, 4, 5, 6\}$

$$C_3 : 2 \leq a + b + 2c \leq 5$$

$a = 1$, $a = 2$, $a = 3$, $a = 4$	Yes, Yes, Yes, Yes
$b = 1$, $b = 2$	Yes, Yes
$c = 0$, $c = 1$, $c = 2$	Yes, Yes, No

- $\text{dom}(A) = \{\cancel{0}, 1, 2, 3, 4\}$
- $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, \cancel{2}, \cancel{3}\}$
- $\text{dom}(D) = \{0, 1, 2, 3, 4, 5\}$
- $\text{dom}(E) = \{0, 1, 2, 3, 4, 5, 6\}$

$$C_2 : 3 \leq a + c \leq 4$$

$a = 1$, $a = 2$, $a = 3$, $a = 4$	No, Yes, Yes, Yes
$c = 0$, $c = 1$	Yes, Yes

- $\text{dom}(A) = \{\cancel{0}, \cancel{1}, 2, 3, 4\}$
- $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, \cancel{2}, \cancel{3}\}$
- $\text{dom}(D) = \{0, 1, 2, 3, 4, 5\}$
- $\text{dom}(E) = \{0, 1, 2, 3, 4, 5, 6\}$

$$C_3 : 2 \leq a + b + 2c \leq 5$$

$a = 2$, $a = 3$, $a = 4$	Yes, Yes, Yes
$b = 1$, $b = 2$	Yes, Yes
$c = 0$, $c = 1$	Yes, Yes

- $\text{dom}(A) = \{\cancel{0}, \cancel{1}, 2, 3, 4\}$
- $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, \cancel{2}, \cancel{3}\}$
- $\text{dom}(D) = \{0, 1, 2, 3, 4, 5\}$
- $\text{dom}(E) = \{0, 1, 2, 3, 4, 5, 6\}$

4. **Conclusion:** $\text{dom}(A) = \{2, 3, 4\}$, $\text{dom}(B) = \{1, 2\}$, $\text{dom}(C) = \{0, 1\}$, $\text{dom}(D) = \{0, 1, 2, 3, 4, 5\}$, $\text{dom}(E) = \{0, 1, 2, 3, 4, 5, 6\}$

Example: **2-Consistency**

Fixed Value	**Satisfactory Combination?**
$C_4 : 1 \leq b \leq 2$	
$b=0$, $b=1$, $b=2$, $b=3$, $b=4$	No, Yes, Yes, No, No

- $\text{dom}(A) = \{0, 1, 2, 3, 4\}$
- $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, 2, 3\}$
- $\text{dom}(D) = \{0, 1, 2, 3, 4, 5\}$
- $\text{dom}(E) = \{0, 1, 2, 3, 4, 5, 6\}$

$C_2 : 3 \leq a + c \leq 4$	
$a=0$, $a=1$, $a=2$, $a=3$, $a=4$	Yes, Yes, Yes, Yes, Yes
$c=0$, $c=1$, $c=2$, $c=3$	Yes, Yes, Yes, Yes

- $\text{dom}(A) = \{0, 1, 2, 3, 4\}$
- $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, 2, 3\}$
- $\text{dom}(D) = \{0, 1, 2, 3, 4, 5\}$
- $\text{dom}(E) = \{0, 1, 2, 3, 4, 5, 6\}$

4. **Conclusion:** $\text{dom}(A) = \{0, 1, 2, 3, 4\}$, $\text{dom}(B) = \{1, 2\}$, $\text{dom}(C) = \{0, 1, 2, 3\}$, $\text{dom}(D) = \{0, 1, 2, 3, 4, 5\}$, $\text{dom}(E) = \{0, 1, 2, 3, 4, 5, 6\}$

Example: **1-Consistency**

Fixed Value	**Satisfactory Combination?**
$C_4 : 1 \leq b \leq 2$	
$b=0$, $b=1$, $b=2$, $b=3$, $b=4$	No, Yes, Yes, No, No

- $\text{dom}(A) = \{0, 1, 2, 3, 4\}$
- $\text{dom}(B) = \{\cancel{0}, 1, 2, \cancel{3}, \cancel{4}\}$
- $\text{dom}(C) = \{0, 1, 2, 3\}$
- $\text{dom}(D) = \{0, 1, 2, 3, 4, 5\}$
- $\text{dom}(E) = \{0, 1, 2, 3, 4, 5, 6\}$

4. **Conclusion:** $\text{dom}(A) = \{0, 1, 2, 3, 4\}$, $\text{dom}(B) = \{1, 2\}$, $\text{dom}(C) = \{0, 1, 2, 3\}$, $\text{dom}(D) = \{0, 1, 2, 3, 4, 5\}$, $\text{dom}(E) = \{0, 1, 2, 3, 4, 5, 6\}$

Learning Problems

Definition: Assume that there is some (unknown) relationship,

$$f : \mathcal{X} \rightarrow \mathcal{Y} \text{ s.t. } x \mapsto_f y$$

- $\mathcal{X}$: Input Space
- $\mathcal{Y}$: Output Space (i.e. information we desire about input)

Find $h : \mathcal{X} \rightarrow \mathcal{Y}$ (hypothesis) s.t. $h \approx f$, given some data about f:

$$\mathcal{D} = \left\{ \left(x^{(i)}, y_i \right), x^{(i)} \in \mathcal{X}, y_i = f\left(x^{(i)} \right) \in \mathcal{Y}, i = 1 \dots N \right\}$$

- $\text{in}(\mathcal{D}) = \{x \text{ s.t. } (x, y) \in \mathcal{D}\}$
- $\text{out}(\mathcal{D}) = \{y \text{ s.t. } (x, y) \in \mathcal{D}\}$

3.7 Classification vs. Regression Problems

Definition:
- **Classification Problems:** $\mathcal{X} \subseteq \mathbb{R}^M$ and $\mathcal{Y} \subseteq \mathbb{N}$
- **Regression Problems:** $\mathcal{X} \subseteq \mathbb{R}^M$ and $\mathcal{Y} \subseteq \mathbb{R}$

3.8 Feature Spaces

Definition: Easier to learn relationships from high-level features (instead of the raw input). Need mapping b/w input space and feature space:

$$\phi : \mathcal{X} \rightarrow \mathcal{F}$$

4 PAC Learning

4.1 Probably Approximately Correct (PAC) Estimations

Motivation: More than one fcn may be consistent w/ the data, how to find the best one?

4.1.1 Hoeffding's Inequality

Motivation: Bound $|\mu - \nu|$ w.r.t. N.

Definition: For any $\epsilon > 0$,

$$\mathbb{P}(|\nu - \mu| \geq \epsilon) \leq 2e^{-2\epsilon^2 N} \tag{1}$$

- μ: Probabillity of an event.
- ν: Relative frequency in a sample size N.
- ϵ: Tolerance (i.e. how close we want ν to be to μ).
  - $\epsilon \to 0$: $\nu = \mu$
- $\mu \overset{?}{\approx} \nu$: μ is probably approximately equal to ν. As $N \to \infty$: $\nu \to \mu$

Warning: Approx. the true dist. w/ high prob. by taking a large enough N (i.e. empirical dist. converges to true dist.).
- i.e. Probability of a sig. deviation shrinks exp. w/ N.

4.2 PAC Learning

4.2.1 Error

Definition:
- **Out-Sample Error:**
$$E_{\text{out}} = \mathbb{P}[f \neq h]$$
- **In-Sample Error:**
$$E_{\text{in}} = \frac{1}{N}\sum_{i=1}^{N} \mathbb{I}[f(x^{(i)}) \neq h(x^{(i)})]$$

4.2.2 Union Bound Theorem

Theorem: The prob. of at least one of the events $E_1, \ldots, E_M$ occurring is bounded by the sum of the prob. of each event occurring:

$$\mathbb{P}\left[E_1 \vee \cdots \vee E_M\right] \leq \sum_{i=1}^{M} \mathbb{P}[E_i]$$

Notes:
- If the events are mutually exclusive, then the union bound is tight (i.e. equality holds).
- If the events are highly correlated, then the union bound is loose (i.e. inequality holds)
  - Some events may be more likely to occur together.

4.2.3 Generalization of Hoeffding's Inequality

Definition: Assuming that h is chosen from a set of hypotheses $\mathcal{H}$, derive a (loose) upper-bound on $|E_{\text{out}} - E_{\text{in}}|$:

$$\begin{aligned}
\mathbb{P}\left[\bigvee_{h\in\mathcal{H}} (|E_{\text{out}} - E_{\text{in}}(h)| > \varepsilon)\right] &\leq \sum_{h\in\mathcal{H}} \mathbb{P}\left[|E_{\text{out}} - E_{\text{in}}(h)| > \varepsilon\right] \\
&\leq \sum_{h\in\mathcal{H}} 2e^{-2\varepsilon^2 N} \\
&= 2|\mathcal{H}|e^{-2\varepsilon^2 N}
\end{aligned}$$

- Endow $\mathcal{F}$ (i.e. fcn space) w/ prob. distribution, $P : \mathcal{X} \to [0, 1]$, then
  - E_{out} (i.e. true error of a hyp. over entire dist. of data) is analogous to μ
  - $E_{\text{in}}(h)$ (i.e. empirical error of hyp. on a finite sample) is analogous to ν.

Notes:
- $E_{\text{in}}(h) \overset{?}{\approx} E_{\text{out}}$ requires small $|\mathcal{H}|$ (generalization)
  - Look at inequality, small $|\mathcal{H}| \to$ small $E_{\text{out}} - E_{\text{in}}$ (i.e. prevents overfitting but leads to underfitting)
- $E_{\text{in}}(h) \approx 0$ requires large $|\mathcal{H}|$ (discrimination)
  - Need large $|\mathcal{H}|$ to capture the true dist. (i.e. prevents underfitting but leads to overfitting)

Example:

1. **Given:** An opaque box containing red and blue balls. Take N IID samples.
   - μ: Probability of drawing a blue balls (unknown).
   - ν: Relative frequency of blue balls in the sample (known).
2. **Problem 1:** What is ν in this case? 8 balls total, 5 are blue.
3. **Solution 1:** $\nu = \frac{5}{8}$
4. **Problem 2:** How to partition $\mathcal{F}$ into regions where $f = h$ and $f \neq h$?
5. **Solution 2:**

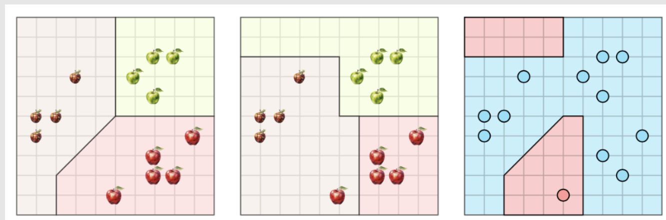


Figure 27: LS h, MS f

6. **Problem 3:** What is the out-sample error?
7. **Solution 3:** In words, the probability of the hypothesis being wrong.

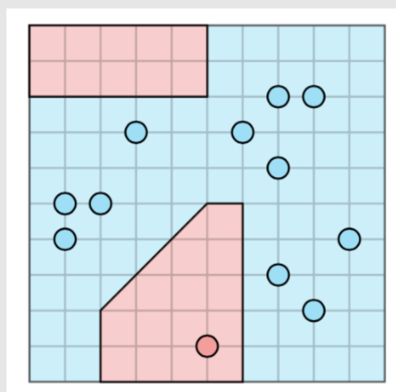


Figure 28

8. **Problem 4:** What is the in-sample error given this sample of 11 balls s.t. $f = h$, 1 ball s.t. $f \neq h$?
9. **Solution 4:** $E_{\text{in}} = \frac{1}{12}$

5 Decision Trees

5.1 Structure

Definition: Each vertex in a decision tree is either:

1. A **condition vertex**: a vertex that sorts points based on a question.
2. A **decision vertex**: a vertex that assigns all points a specific class.

Notes: We want to find the minimum # of condition vertices (or questions) needed to "sufficiently discriminate" (identify the class of every point in $\mathcal{D}$).

- More condition vertices improve discrimination.
- Less condition vertices improve generalization.

5.2 Building a Decision Tree

Definition: Consider determining the class of a randomly chosen target point.

- If we ask a K-ary question abt. the pts. in $\mathcal{D}$, we can form K subsets, $\mathcal{D}^{(1)}, \dots, \mathcal{D}^{(K)}$, using the answers s.t.
  - $|\mathcal{D}^{(k)}| \in \{0, \dots, |\mathcal{D}|\}$
  - $|\mathcal{D}| = \sum_{k=1}^{K} |\mathcal{D}^{(k)}|$

5.3 Special Case

Notes: Suppose each pt. belongs to a unique class (i.e. the # of classes is $|\mathcal{D}|$).

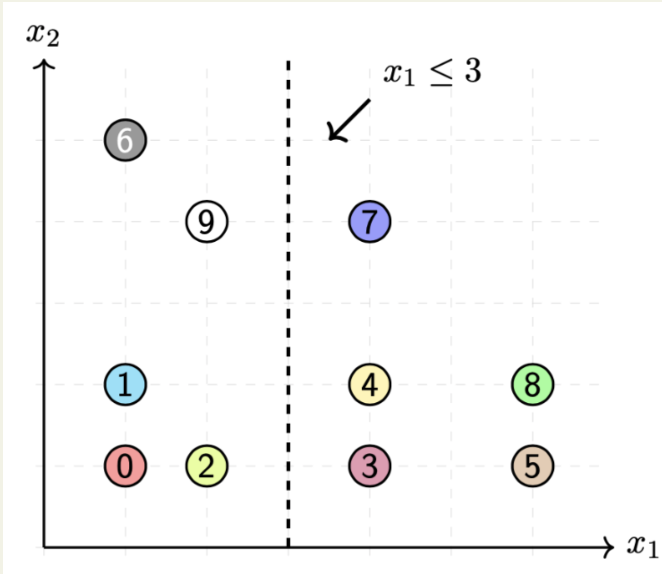


Figure 29

1. Before asking the question: $|\mathcal{D}|$ possible guesses for the target point's class.
2. After asking the question: Either
   - $|\mathcal{D}^{(1)}|, \ldots, |\mathcal{D}^{(K-1)}|$ or
   - $|\mathcal{D}^{(K)}|$

   guesses, depending on the answer for the target point.
   - i.e. $|\mathcal{D}^{(K)}|$ if the target point belongs to class K (Yes)
   - i.e. $|\mathcal{D}^{(1)}|, \ldots, |\mathcal{D}^{(K-1)}|$ if the target point belongs to class $1, \ldots, K-1$ (No)
3. Goal: Minimize the # of guesses needed in the worst-case, which would be
$$\max\{|\mathcal{D}^{(1)}|, \ldots, |\mathcal{D}^{(K)}|\}.$$
   - i.e. Target point falls into the largest subset after a question is asked.
4. Given the constraints on $|\mathcal{D}^{(1)}|, \ldots, |\mathcal{D}^{(K)}|$, we can show that $\max\{|\mathcal{D}^{(1)}|, \ldots, |\mathcal{D}^{(K)}|\}$ is minimized when
$$|\mathcal{D}^{(K)}| \in \left\{ \left\lfloor \frac{|\mathcal{D}|}{K} \right\rfloor, \left\lceil \frac{|\mathcal{D}|}{K} \right\rceil \right\}.$$
   Basically, the best question splits $\mathcal{D}$ into K sets of (roughly) the same size.

Warning: Roughly due to floor/ceil.

5.3.1 # of K-ary Questions Needed

Theorem: Given a classification data-set, $\mathcal{D}$, in which the class of each point is unique (i.e., $|\text{out}(\mathcal{D})| = |\mathcal{D}|$), the class of a randomly chosen target point can be determined within
$$\lceil \log_K(|\mathcal{D}|) \rceil$$
K-ary questions.

5.4 General Case

Motivation: Suppose points do not necessarily belong to a unique class.
- X is the class of a randomly chosen target point.
- Y is the answer to a K-ary question for X.

5.4.1 Expected # of Questions

Definition: Using the theorem above, since for each class, c, we can partition $\mathcal{D}$ into $\lceil 1/p_c \rceil$ subsets, with a subset containing all class c points
- p_c: Proportion of class c points.

If the target point's class is c, we can confirm it w/in $\lceil \log_K(\lceil 1/p_c \rceil) \rceil$ K-ary questions.

Thus, the expected # of Qs needed is

$$\sum_c p_c \lceil \log_2(\lceil 1/p_c \rceil) \rceil.$$

Notes: i.e. Reduces to special cases with each subset containing a unique class.

5.4.2 Entropy, Conditional Entropy, and Information Gain

Definition: The **entropy** of a random variable X (in K-its) is defined as

$$H(X) = -\sum_{\forall x \in X} p_X(x) \log_K(p_X(x)).$$

The **conditional entropy** of a random variable, X, given a random variable Y, is

$$H(X|Y) = -\sum_{\forall y \in Y} \sum_{\forall x \in X} p_{X|Y}(x|y) \log_K(p_{X|Y}(x|y)).$$

The **information gain** from Y is:

$$IG(X|Y) = H(X) - H(X|Y).$$

- Maximize $IG(X|Y)$ (i.e. choose the question to maximize the information gained).

Process:

1. Calculate $H(X)$ (i.e. entropy before the split).
2. Calculate $H(X|Y)$ (i.e. entropy after the split).
   (a) Calculate entropy for each subset of X based on the question, Y.
   (b) Calculate the weighted average of the entropies.
3. Calculate $IG(X|Y) = H(X) - H(X|Y)$.

Example: Consider a classification problem where $\mathcal{X} = \{0, \dots, 9\}^2$, $\mathcal{Y} = \{0, 1, 2\}$ and suppose we are given

$$\mathcal{D} = \left\{ \left(\begin{bmatrix} 1 \\ 1 \end{bmatrix}, 0 \right), \left(\begin{bmatrix} 1 \\ 2 \end{bmatrix}, 0 \right), \left(\begin{bmatrix} 2 \\ 1 \end{bmatrix}, 0 \right), \left(\begin{bmatrix} 4 \\ 1 \end{bmatrix}, 1 \right), \left(\begin{bmatrix} 4 \\ 2 \end{bmatrix}, 1 \right), \left(\begin{bmatrix} 6 \\ 1 \end{bmatrix}, 2 \right), \left(\begin{bmatrix} 1 \\ 5 \end{bmatrix}, 2 \right), \left(\begin{bmatrix} 4 \\ 4 \end{bmatrix}, 2 \right), \left(\begin{bmatrix} 6 \\ 2 \end{bmatrix}, 2 \right), \left(\begin{bmatrix} 2 \\ 4 \end{bmatrix}, 2 \right) \right\}.$$

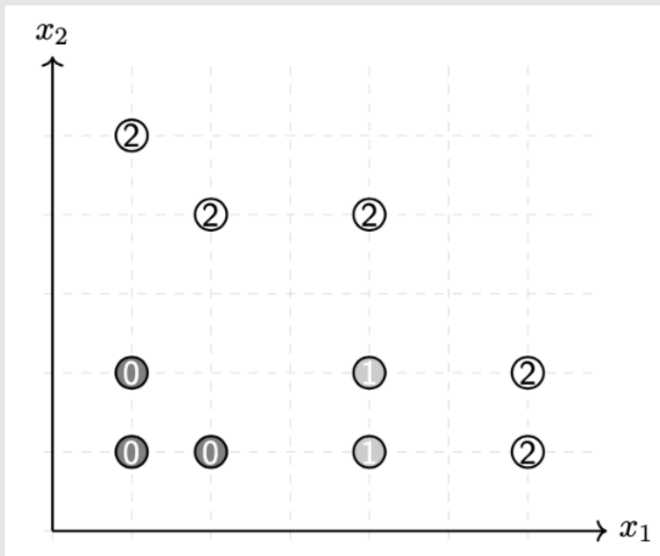


Figure 30

Example: **2-Ary Question**

1. **Given:** $X = \{0, 1, 2\}$, $Y = \begin{cases} 1, & \text{if } x_1 \leq 3 \quad \text{(Yes)} \\ 0, & \text{if } x_1 > 3 \quad \text{(No)} \end{cases}$,
2. **Problem:** $IG(X|Y) = ?$
3. **Solution:**
   (a) Entropy before the split: $H(X) = \frac{3}{10}\log_2\left(\frac{10}{3}\right) + \frac{2}{10}\log_2\left(\frac{10}{2}\right) + \frac{5}{10}\log_2\left(\frac{10}{5}\right)$
   (b) Entropy after the split:
      i. $H(X \mid x_1 \leq 3) = \frac{3}{5}\log_2\left(\frac{5}{3}\right) + \frac{2}{5}\log_2\left(\frac{5}{2}\right)$
      ii. $H(X \mid x_1 > 3) = \frac{2}{5}\log_2\left(\frac{5}{2}\right) + \frac{3}{5}\log_2\left(\frac{5}{3}\right)$.
      iii. Weighted Avg. Entropy: $H(X|Y) = \frac{5}{10}H(X \mid x_1 \leq 3) + \frac{5}{10}H(X \mid x_1 > 3)$
   (c) $IG(X|Y) = H(X) - H(X|Y)$

Example: **2-Ary Question**

1. **Given:** $X = \{0, 1, 2\}$, $Y = \begin{cases} 1, & \text{if } x_2 \leq 3 \quad \text{(Yes)} \\ 0, & \text{if } x_2 > 3 \quad \text{(No)} \end{cases}$,
2. **Problem:** $IG(X|Y) = ?$
3. **Solution:**
   (a) Entropy before the split: $H(X) = \frac{3}{10}\log_2\left(\frac{10}{3}\right) + \frac{2}{10}\log_2\left(\frac{10}{2}\right) + \frac{5}{10}\log_2\left(\frac{10}{5}\right)$
   (b) Entropy after the split:
      i. $H(X \mid x_2 > 3) = \frac{3}{3}\log_2\left(\frac{3}{3}\right)$
      ii. $H(X \mid x_2 \leq 3) = \frac{3}{5}\log_2\left(\frac{5}{3}\right) + \frac{2}{5}\log_2\left(\frac{5}{2}\right) + \frac{2}{5}\log_2\left(\frac{5}{2}\right)$.
      iii. Weighted Avg. Entropy: $H(X|Y) = \frac{3}{10}H(X \mid x_2 > 3) + \frac{7}{10}H(X \mid x_2 \leq 3)$
   (c) $IG(X|Y) = H(X) - H(X|Y)$

Example: **3-Ary Question**

1. **Given:** $X = \{0, 1, 2\}$, $Y = \begin{cases} 1, & \text{if } x_1 \leq 3 \text{ and } x_2 \leq 3 \\ 2, & \text{if } x_1 \leq 3 \text{ and } x_2 > 3 \\ 3, & \text{if } x_1 > 3 \end{cases}$
2. **Problem:** $IG(X|Y) = ?$
3. **Solution:**
   (a) Entropy before the split: $H(X) = \frac{3}{10}\log_2\left(\frac{10}{3}\right) + \frac{2}{10}\log_2\left(\frac{10}{2}\right) + \frac{5}{10}\log_2\left(\frac{10}{5}\right)$
   (b) Entropy after the split:
      i. $H(X \mid x_1 \leq 3 \text{ and } x_2 \leq 3) = \frac{3}{3}\log_2\left(\frac{3}{3}\right)$
      ii. $H(X \mid x_1 \leq 3 \text{ and } x_2 > 3) = \frac{2}{2}\log_2\left(\frac{2}{2}\right)$
      iii. $H(X \mid x_1 > 3) = \frac{2}{5}\log_2\left(\frac{5}{2}\right) + \frac{3}{5}\log_2\left(\frac{5}{3}\right)$
      iv. $H(X|Y) = \frac{3}{10}H(X \mid x_1 \leq 3 \text{ and } x_2 \leq 3) + \frac{2}{10}H(X \mid x_1 \leq 3 \text{ and } x_2 > 3) + \frac{5}{10}H(X \mid x_1 > 3)$
   (c) $IG(X|Y) = H(X) - H(X|Y)$

Example: Decision Tree

1. **Given:** $X = \{0, 1, 2\}$
2. **Problem:** Draw a decision tree using binary conditions of the form, $x_i \leq k$, where $i \in \{1, 2\}$ and $k \in \mathbb{Z}$, that maximizes the information gained at each level.
3. **Solution (Level 1):**
   (a) Entropy before the split: $H(X) = \frac{3}{10}\log_2\left(\frac{10}{3}\right) + \frac{2}{10}\log_2\left(\frac{10}{2}\right) + \frac{5}{10}\log_2\left(\frac{10}{5}\right) = 1.485[\text{bits}]$
   (b) Entropy after the split and information gain (everything in base 2 since 2-ary).

Split	**Entropy**
$x_1 \leq 1$	$H(X\|Y) = \frac{3}{10}\left[\frac{2}{3}\log\left(\frac{3}{2}\right) + \frac{1}{3}\log\left(\frac{3}{1}\right)\right] + \frac{7}{10}\left[\frac{1}{7}\log\left(\frac{7}{1}\right) + \frac{2}{7}\log\left(\frac{7}{2}\right) + \frac{4}{7}\log\left(\frac{7}{4}\right)\right] = 1.241[\text{bits}]$ • $IG(X\|Y) = 1.485 - 1.241 = 0.244[\text{bits}]$
$x_1 \leq 2, 3$	$H(X\|Y) = \frac{5}{10}\left[\frac{3}{5}\log\left(\frac{5}{3}\right) + \frac{2}{5}\log\left(\frac{5}{2}\right)\right] + \frac{5}{10}\left[\frac{2}{5}\log\left(\frac{5}{2}\right) + \frac{3}{5}\log\left(\frac{5}{3}\right)\right] = 0.971[\text{bits}]$ • $\boxed{IG(X\|Y) = 1.485 - 0.971 = 0.514[\text{bits}]}$
$x_1 \leq 4, 5$	$H(X\|Y) = \frac{8}{10}\left[\frac{3}{8}\log\left(\frac{8}{3}\right) + \frac{2}{8}\log\left(\frac{8}{2}\right) + \frac{3}{8}\log\left(\frac{8}{3}\right)\right] + \frac{2}{10}\left[\frac{2}{2}\log\left(\frac{2}{2}\right)\right] = 1.249[\text{bits}]$ • $IG(X\|Y) = 1.485 - 1.249 = 0.236[\text{bits}]$
$x_1 \leq 6$	$H(X\|Y) = \frac{10}{10}\left[\frac{3}{10}\log\left(\frac{10}{3}\right) + \frac{2}{10}\log\left(\frac{10}{2}\right) + \frac{5}{10}\log\left(\frac{10}{5}\right)\right] = 1.485[\text{bits}]$ • $IG(X\|Y) = 1.485 - 1.485 = 0[\text{bits}]$
$x_2 \leq 1$	$H(X\|Y) = \frac{4}{10}\left[\frac{2}{4}\log\left(\frac{4}{2}\right) + \frac{1}{4}\log\left(\frac{4}{1}\right) + \frac{1}{4}\log\left(\frac{4}{1}\right)\right] + \frac{6}{10}\left[2 \cdot \frac{1}{6}\log\left(\frac{6}{1}\right) + \frac{4}{6}\log\left(\frac{6}{4}\right)\right] = 1.351[\text{bits}]$ • $IG(X\|Y) = 1.485 - 1.351 = 0.134[\text{bits}]$
$x_2 \leq 2, 3$	$H(X\|Y) = \frac{7}{10}\left[\frac{3}{7}\log\left(\frac{7}{3}\right) + \frac{2}{7}\log\left(\frac{7}{2}\right) + \frac{2}{7}\log\left(\frac{7}{2}\right)\right] + \frac{3}{10}\left[\frac{3}{3}\log\left(\frac{3}{3}\right)\right] = 1.090[\text{bits}]$ • $IG(X\|Y) = 1.485 - 1.090 = 0.395[\text{bits}]$
$x_2 \leq 4$	$H(X\|Y) = \frac{9}{10}\left[\frac{3}{9}\log\left(\frac{9}{3}\right) + \frac{2}{9}\log\left(\frac{9}{2}\right) + \frac{4}{9}\log\left(\frac{9}{4}\right)\right] + \frac{1}{10}\left[\frac{1}{1}\log\left(\frac{1}{1}\right)\right] = 1.377[\text{bits}]$ • $IG(X\|Y) = 1.485 - 1.377 = 0.108[\text{bits}]$
$x_2 \leq 5$	$H(X\|Y) = \frac{10}{10}\left[\frac{3}{10}\log\left(\frac{10}{3}\right) + \frac{2}{10}\log\left(\frac{10}{2}\right) + \frac{5}{10}\log\left(\frac{10}{5}\right)\right] = 1.485[\text{bits}]$ • $IG(X\|Y) = 1.485 - 1.485 = 0[\text{bits}]$

Example: Decision Tree Continued:

4. **Solution (Level 2):** $x_1 \leq 2, 3$ has the highest information gain. For clarity, choose $x_1 \leq 3$ as the question.
   (a) Entropy before the split (treat as 2 indep. problems)
      i. $H(X_L) = \frac{3}{5}\log\left(\frac{5}{3}\right) + \frac{2}{5}\log\left(\frac{5}{2}\right) = 0.971$
      ii. $H(X_R) = \frac{2}{5}\log\left(\frac{5}{2}\right) + \frac{3}{5}\log\left(\frac{5}{3}\right) = 0.971$
   (b) Entropy after the split and information gain (everything in base 2 since 2-ary).

Split	**Entropy**
	Left Split
$x_1 \leq 1$	$H(X_L\|Y) = \frac{3}{5}\left[\frac{2}{3}\log\left(\frac{3}{2}\right) + \frac{1}{3}\log\left(\frac{3}{1}\right)\right] + \frac{2}{5}\left[\frac{1}{2}\log\left(\frac{1}{2}\right) + \frac{1}{2}\log\left(\frac{1}{2}\right)\right] = 0.151[\text{bits}]$ • $IG(X\|Y) = 0.971 - 0.151 = 0.820[\text{bits}]$
$x_2 \leq 1$	$H(X_L\|Y) = \frac{2}{5}\left[\frac{2}{2}\log\left(\frac{2}{2}\right)\right] + \frac{3}{5}\left[\frac{1}{3}\log\left(\frac{3}{1}\right) + \frac{2}{3}\log\left(\frac{3}{2}\right)\right] = 0.551[\text{bits}]$ • $IG(X\|Y) = 0.971 - 0.551 = 0.420[\text{bits}]$
$x_2 \leq 2, 3$	$H(X_L\|Y) = \frac{3}{5}\left[\frac{3}{3}\log\left(\frac{3}{3}\right)\right] + \frac{2}{5}\left[\frac{2}{2}\log\left(\frac{2}{2}\right)\right] = 0[\text{bits}]$ • $\boxed{IG(X_L\|Y) = 0.971 - 0 = 0.971[\text{bits}]}$
	Right Split
$x_1 \leq 4, 5$	$H(X_R\|Y) = \frac{3}{5}\left[\frac{2}{3}\log\left(\frac{3}{2}\right) + \frac{1}{3}\log\left(\frac{3}{1}\right)\right] + \frac{2}{5}\left[\frac{2}{2}\log\left(\frac{2}{2}\right)\right] = 0.551[\text{bits}]$ • $\boxed{IG(X_L\|Y) = 0.971 - 0.551 = 0.420[\text{bits}]}$
$x_2 \leq 1$	$H(X_R\|Y) = \frac{2}{5}\left[\frac{1}{2}\log\left(\frac{2}{1}\right) + \frac{1}{2}\log\left(\frac{2}{1}\right)\right] + \frac{3}{5}\left[\frac{2}{3}\log\left(\frac{3}{2}\right) + \frac{1}{3}\log\left(\frac{3}{1}\right)\right] = 0.951[\text{bits}]$ • $IG(X_L\|Y) = 0.971 - 0.951 = 0.020[\text{bits}]$
$x_2 \leq 2, 3$	$H(X_R\|Y) = \frac{4}{5}\left[\frac{2}{4}\log\left(\frac{4}{2}\right) + \frac{2}{4}\log\left(\frac{4}{2}\right)\right] + \frac{1}{5}\left[\frac{1}{1}\log\left(\frac{1}{1}\right)\right] = 0.8[\text{bits}]$ • $IG(X_L\|Y) = 0.971 - 0.8 = 0.171[\text{bits}]$

Example: Decision Tree Continued:

5. **Solution (Level 3):** $x_2 \leq 2, 3$ and $x_1 \leq 4, 5$ has the highest information gain. For clarity, choose $x_2 \leq 3$ as the question for the left split and choose $x_1 \leq 5$ as the question for the right split.
   (a) Since 3 are pure splits already, therefore, look at right-left side only.
   (b) Entropy before the split for the right-left side
      i. $H(X_{RL}) = \frac{2}{3}\log\left(\frac{3}{2}\right) + \frac{1}{3}\log\left(\frac{3}{1}\right) = 0.918[\text{bits}]$
   (c) Entropy after the split and information gain (everything in base 2 since 2-ary).

Split	**Entropy**
$x_2 \leq 1$	$H(X_{RL}\|Y) = \frac{1}{3}\left[\frac{1}{1}\log\left(\frac{1}{1}\right)\right] + \frac{2}{3}\left[\frac{1}{2}\log\left(\frac{2}{1}\right) + \frac{1}{2}\log\left(\frac{2}{1}\right)\right] = 0.667[\text{bits}]$
	• $IG(X\|Y) = 0.971 - 0.667 = 0.304[\text{bits}]$
$x_2 \leq 2, 3$	$H(X_{RL}\|Y) = \frac{1}{3}\left[\frac{1}{1}\log\left(\frac{1}{1}\right)\right] + \frac{2}{3}\left[\frac{2}{2}\log\left(\frac{2}{2}\right)\right] = 0[\text{bits}]$
	• $\boxed{IG(X\|Y) = 0.971 - 0 = 0.971[\text{bits}]}$

6. Now all regions in our graph contain a pure set (one class). Note this took more questions than needed, but IG is a heuristic so its not perfect.