

# Closure Conversion by explicit Lambda Abstraction

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We want to translate from the nonclosed `IL1` language

$$\begin{aligned} v &::= n \mid x \mid \lambda x k.c \mid \text{halt} \\ e &::= v \mid v_0 + v_1 \mid (v_i) \mid \pi_n v \\ c &::= \text{let } x = e \text{ in } c \mid v_0 \ v_1 \ v_2 \end{aligned}$$

into the closed restriction with an explicit abstraction  $\bar{\rho} \in \text{var list}$  over all of the variables.

$$\begin{aligned} v &::= n \mid x \mid \lambda x k.\bar{\rho}.c \mid \text{halt} \\ e &::= v \mid v_0 + v_1 \mid (v_i) \mid \pi_n v \\ c &::= \text{let } x = e \text{ in } c \mid v_0 \ v_1 \ v_2 \ \bar{\rho} \end{aligned}$$

Here, we can proceed as natural with the translation functions

$$\begin{aligned} \mathcal{V}[[v]] &: v \\ \mathcal{E}[[e]] &: e \\ \mathcal{C}[[c]] &: c \end{aligned}$$

## 1 Values

$$\begin{aligned} \mathcal{V}[[n]] &= n \\ \mathcal{V}[[x]] &= \rho_{\langle x \rangle} \\ \mathcal{V}[[\text{halt}]] &= \text{halt} \\ \mathcal{V}[[\lambda x k.c]] &= \lambda x k.\bar{\rho}. \text{let } \rho_{\langle x \rangle} = x, \rho_{\langle k \rangle} = k \text{ in } \mathcal{C}[[c]] \end{aligned}$$

## 2 Expressions

$$\begin{aligned} \mathcal{E}[[v]] &= \mathcal{V}[[v]] \\ \mathcal{E}[[v_0 + v_1]] &= \mathcal{V}[[v_0]] + \mathcal{V}[[v_1]] \\ \mathcal{E}[[v_i]] &= (\mathcal{V}[[v_i]]) \\ \mathcal{E}[[\pi_n v]] &= \pi_n \mathcal{V}[[v]] \end{aligned}$$

### 3 Commands

$$\begin{aligned}\mathcal{C}[\![v_0 \ v_1 \ v_2]\!] &= \mathcal{V}[\![v_0]\!] \ \mathcal{V}[\![v_1]\!] \ \mathcal{V}[\![v_2]\!] \ \bar{\rho} \\ \mathcal{C}[\![\text{let } x = e \text{ in } c]\!] &= \text{let } \rho_{\langle x \rangle} = \mathcal{E}[\![e]\!] \text{ in } \mathcal{C}[\![c]\!]\end{aligned}$$

For example, suppose I have the command

$$c \triangleq \text{let } f = \lambda x k. k \ (x + 1) \ k \text{ in } f \ 4 \ \text{halt}$$

there are three variables here,  $x, k, f$ , so we can construct the injection:

$$\langle \cdot \rangle = \{x \mapsto 1, k \mapsto 2, f \mapsto 3\}$$

so

$$\bar{\rho} = \rho_1, \rho_2, \rho_3$$

then the above would translate into

$$\mathcal{C}[\![c]\!] = \text{let } \rho_{\langle f \rangle} = \mathcal{E}[\![\lambda x k. k(x + 1)k]\!] \text{ in } \mathcal{C}[\![f \ 4 \ \text{halt}]\!]$$

expand the  $\mathcal{E}$  first

$$\begin{aligned}&= \text{let } \rho_{\langle f \rangle} = \mathcal{V}[\![\lambda x k. k(x + 1)k]\!] \text{ in } \mathcal{C}[\![f \ 4 \ \text{halt}]\!] \\&= \text{let } \rho_{\langle f \rangle} = \lambda x k \rho_1 \rho_2 \rho_3. \text{let } \rho_{\langle x \rangle} = x, \rho_{\langle k \rangle} = k \text{ in } \mathcal{C}[\![k \ (x + 1) \ k]\!] \text{ in } \mathcal{C}[\![f \ 4 \ \text{halt}]\!] \\&\vdots \\&= \text{let } \rho_{\langle f \rangle} = \lambda x k \rho_1 \rho_2 \rho_3. \text{let } \rho_{\langle x \rangle} = x, \rho_{\langle k \rangle} = k \text{ in } \rho_{\langle k \rangle}(\rho_{\langle x \rangle} + 1) \rho_{\langle k \rangle} \rho_1 \rho_2 \rho_3 \text{ in } \rho_{\langle f \rangle} \ 4 \ \text{halt} \ \rho_1 \rho_2 \rho_3 \\&= \text{let } \rho_3 = \lambda x k \rho_1 \rho_2 \rho_3. \text{let } \rho_1 = x, \rho_2 = k \text{ in } \rho_2(\rho_1 + 1) \rho_2 \rho_1 \rho_2 \rho_3 \text{ in } \rho_3 \ 4 \ \text{halt} \ \rho_1 \rho_2 \rho_3\end{aligned}$$