Contents			<pre>typedef vector<ll> vi; typedef pair<ll, ll=""> pll;</ll,></ll></pre>
L	Setting 1.1 Default code	-	<pre>typedef vector<11> v1; #include <ext assoc_container.hpp="" pb_ds=""> #include <ext pb_ds="" tree_policy.hpp=""></ext></ext></pre>
2	Math 2.1 Extended Euclidean Algorithm 2.2 Primality Test 2.3 Integer Factorization (Pollard's rho) 2.4 Chinese Remainder Theorem 2.5 Query of nCr mod M in $O(Q + M)$ 2.6 pelindrome number 2.7 Matrix Pow 2.8 Catalan, Derangement, Partition, 2nd Stirling 2.9 Matrix Operations 2.10 Coursion Elimination	1 2 2 2 2 2 2 3	<pre>#include <ext detail="" pb_ds="" standard_policies.hpp=""> using namespacegnu_pbds; template <typename t=""> using ordered_set = tree<t, less<="" null_type,="">, rb_tree_tag,</t,></typename></ext></pre>
	2.10 Gaussian Elimination	4	<pre>#ifdef ONLINE_JUDGE constexpr bool ndebug = true; #else constexpr bool ndebug = false; #endif</pre>
3	Data Structure 3.1 Lazy Segment Tree	4	11 gcd(11 a, 11 b){return b?gcd(b,a%b):a;} 11 lcm(11 a, 11 b){if(a&&b)return a*(b/gcd(a,b)); return a+b;} 11 POW(11 a, 11 b, 11 rem){11 p=1;a%=rem;for(;b;b>>=1,a=(a*a)% rem)if(b&1)p=(p*a)%rem;return p;}
1	Graph 4.1 Dijkstra 4.2 LCA 4.3 Centroid Decomposition 4.4 Minimum Spanning Tree 4.5 Offline Dynamic Connectivity	5 5 5 5 6	<pre>ios_base::sync_with_stdio(0); cin.tie(0);</pre>
5	String 5.1 KMP 5.2 Z Algorithm	6 6 7	,
; L	Hash 6.1 Basic Hash Setting	7	}
tii tpi tpi tpi tpi tpi ts: us: us: yi	Include clude <format> ragma warning(disable:4996) ragma comment(linker, "/STACK:336777216") ragma GCC optimize("03,unroll-loops") ragma GCC target("avx,avx2,fma") ting namespace std; ting ll = long long; ting pll = pair<ll,ll; </ll,ll; ting tll = tuple<ll,ll,ll; </ll,ll,ll; ting ld = pair<ld,ld>; tedef unsigned long long ull; deedefint128 LL; pedef pair<ll, ll=""> pii;</ll,></ld,ld></format>		<pre>int main() { setup(); preprocess(); ll t = 1; // cin >> t; cin.ignore(); for (ll testcase = 1; testcase <= t; testcase++) { solve(testcase); } return 0; } 2 Math 2.1 Extended Euclidean Algorithm // ax+by=g, return (g,x,y)</pre>

```
KU - unipass
```

```
tuple<11, 11, 11> extended_gcd(11 a, 11 b){
 if (a == 0) {b, 0, 1};
  auto [g, x, y] = extended gcd(b % a, a);
  return {g, y - (b / a) * x, x};
// find x in [0,m) s.t. ax === qcd(a, m) (mod m)
11 modinverse(ll a, ll m) {
    return (get<1>(extended gcd(a, m))%m+m)%m;
2.2 Primality Test
// O(logn*logn)
bool is prime(ll n) {
 if (n < 2 | n % 2 == 0 | n % 3 == 0) return n == 2 | n == 3;
  ll k = builtin ctzll(n - 1), d = n - 1 >> k;
  for (ll a: { 2, 325, 9375, 28178, 450775, 9780504, 1795265022 }) {
   11 p = modpow(a % n, d, n), i = k;
    while (p != 1 && p != n - 1 && a % n && i--) p = modmul(p, p, n);
    if (p != n - 1 && i != k) return 0;
  return 1;
2.3 Integer Factorization (Pollard's rho)
11 pollard(ll n) {
 auto f = [n](11 x) \{ return modadd(modmul(x, x, n), 3, n); \};
 11 \times 0, y = 0, t = 30, p = 2, i = 1, q;
  while (t++ \% 40 \mid | gcd(p, n) == 1) {
    if (x == y) x = ++i, y = f(x);
    if (q = modmul(p, abs(x - y), n)) p = q;
    x = f(x), y = f(f(y));
  return gcd(p, n);
// integer factorization
// O(n^0.25 * Logn)
vector<ll> factor(ll n) {
 if (n == 1) return {};
  if (is_prime(n)) return { n };
 11 x = pollard(n);
  auto 1 = factor(x), r = factor(n / x);
 1.insert(l.end(), r.begin(), r.end());
  sort(1.begin(), 1.end());
 return 1;
2.4 Chinese Remainder Theorem
// x = r_i \mod m_i
// (y, m) 'x = y mod m' 'm = lcm(m i)', if not exists return (0, 0)
auto crt = [](auto r, auto m) {
  const int n = r.size(); i64 r0 = 0, m0 = 1;
  for (int i = 0; i < n; i++) {
    i64 r1 = r[i], m1 = m[i];
    if (m0 < m1) swap(r0, r1), swap(m0, m1);</pre>
    if (m0 % m1 == 0 && r0 % m1 != r1) return pair(OLL, OLL);
    if (m0 % m1 == 0) continue;
    i64 g = gcd(m0, m1);
    if ((r1 - r0) % g) return pair(0LL, 0LL);
    i64 u0 = m0 / g, u1 = m1 / g;
    i64 \times = (r1 - r0) / g \% u1 * modinv(u0, u1) % u1;
    r0 += x * m0, m0 *= u1; if (r0 < 0) r0 += m0;
  return pair(r0, m0);
```

2.5 Query of nCr mod M in O(Q+M)auto sol p e = [](int q, const auto& qs, const int p, const int e, const int mod) { $// qs[i] = \{n, r\}, nCr mod p^e in O(p^e)$ vector dp(mod, 1); for (int i = 0; i < mod; i++) {</pre> if (i) dp[i] = dp[i - 1]; if (i % p == 0) continue; dp[i] = mul(dp[i], i); $auto f = [&](i64 n) {$ i64 res = 0;while (n /= p) res += n; return res; }; auto $g = [\&](i64 n) {$ auto rec = [&](const auto& self, i64 n) -> int { if (n == 0) return 1; int q = n / mod, r = n % mod; int ret = mul(self(self, n / p), dp[r]); if (q & 1) ret = mul(ret, dp[mod - 1]); return ret: }; return rec(rec, n); }; auto bino = $[\&](i64 \text{ n}, i64 \text{ r}) \{$ if (n < r) return 0;</pre> if (r == 0 || r == n) return 1; i64 a = f(n) - f(r) - f(n - r);if (a >= e) return 0; int b = mul(g(n), modinv(mul(g(r), g(n - r)), mod));return mul(pow(p, a), b); }; vector res(q, 0); for (int i = 0; i < q; i++) { auto [n, r] = qs[i]; res[i] = bino(n, r);return res; }; auto sol = [](int q, const auto& qs, const int mod) { vector fac = factor(mod); vector r(q, vector(fac.size(), 0)); vector m(fac.size(), 1); for (int i = 0; i < fac.size(); i++) {</pre> auto [p, e] = fac[i]; for (int j = 0; j < e; j++) m[i] *= p; auto res = sol_p_e(q, qs, p, e, m[i]); for (int j = 0; j < q; j++) r[j][i] = res[j]; vector res(q, 0); for (int i = 0; i < q; i++) res[i] = crt(r[i], m).first; return res; }; 2.6 pelindrome number ll peli(string n) { 11 len = n.size(), cnt = 0; for (int i = 1; i < len; i++) cnt += 9 * pow(10, (i - 1) / 2); string half = n.substr(0, (len + 1) / 2); ll halfNum = stoll(half), base = pow(10, (len - 1) / 2); cnt += halfNum - base; string rev = half.substr(0, len / 2); reverse(rev.begin(), rev.end()); string full = half + rev;

```
if (full <= n) cnt++;</pre>
    return cnt;
2.7 Matrix Pow
void mulmat(vector<vector<11>> &a, vector<vector<11>> b) {
    11 n = a.size();
    11 m = a[0].size();
    11 k = b[0].size();
    vector ret(n, vector<11>(k, 0));
     for (ll i = 0; i < n; i++) {
         for (11 j = 0; j < k; j++) {
             for (11 \ 1 = 0; \ 1 < m; \ 1++) {
                 ret[i][j] += a[i][l] * b[l][j];
                 ret[i][j] %= mod;
    a = ret;
void powmat(vector<vector<11>> &ret, vector<vector<11>> &a, 11 n) {
     while (n) {
         if (n & 1) mulmat(ret, a);
         mulmat(a, a);
        n >>= 1;
}
2.8 Catalan, Derangement, Partition, 2nd Stirling
C_n = \frac{1}{n+1} {2n \choose n}, C_0 = 1, C_{n+1} = \sum_{i=0}^n C_i C_{n-i}, C_{n+1} = \frac{2(2n+1)}{n+2} C_n
D_n = (n-1)(D_{n-1} + D_{n-2}) = n! \sum_{i=1}^n \frac{(-1)^{i+1}}{i!}
P(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k+1} P(n - k(3k - 1)/2)
= P(n-1) + P(n-2) - P(n-5) - P(n-7) + P(n-12) + P(n-15) - P(n-22) - \cdots
P(n,k) = P(n-1,k-1) + P(n-k,k), S(n,k) = S(n-1,k-1) + k \cdot S(n-1,k)
2.9 Matrix Operations
inline bool is_zero(ld a) { return abs(a) < eps; }</pre>
// returns {det(A), A^-1, rank(A), tr(A)}
// A becomes invalid after call this O(n^3)
tuple<ld, vector<vector<ld>>>,ll,ll> inv_det_rnk(auto A) {
  ld n=A.size(); ld det = 1; vector out(n, vector<ld>(n)); ld tr=0;
  for (int i = 0; i < n; i++) {
    out[i][i] = 1; tr+=A[i][i];
  for (int i = 0; i < n; i++) {
    if (is_zero(A[i][i])) {
      1d \max v = 0;
      int maxid = -1;
       for (int j = i + 1; j < n; j++) {
         auto cur = abs(A[j][i]);
        if (maxv < cur) {</pre>
          maxv = cur;
           maxid = j;
       if (maxid == -1 || is_zero(A[maxid][i])) return {0, out, i, tr};
      for (int k = 0; k < n; k++) {
        A[i][k] += A[maxid][k]; out[i][k] += out[maxid][k];
```

```
det *= A[i][i];
   ld coeff = 1.0 / A[i][i];
    for (int j = 0; j < n; j++) A[i][j] *= coeff,out[i][j] *= coeff;</pre>
    for (int j = 0; j < n; j++) if (j != i) {
     1d mp = A[j][i];
      for (int k = 0; k < n; k++) A[j][k] -= A[i][k] * mp;
      for (int k = 0; k < n; k++) out[j][k] -= out[i][k] * mp;
 return {det, out, n, tr};
      Gaussian Elimination
const double EPS = 1e-10;
typedef vector<vector<double>> VVD;
// Gauss-Jordan elimination with full pivoting.
// solving systems of linear equations (AX=B)
           a[][] = an n*n matrix
// INPUT:
             b[][] = an n*m matrix
// OUTPUT:
                 = an n*m matrix (stored in b[][])
             A^{-1} = an n*n matrix (stored in a[][])
// O(n^3)
bool gauss jordan(VVD& a, VVD& b) {
    const int n = a.size();
    const int m = b[0].size();
    vector<int> irow(n), icol(n), ipiv(n);
    for (int i = 0; i < n; i++) {
        int pj = -1, pk = -1;
        for (int j = 0; j < n; j++) if (!ipiv[j])</pre>
            for (int k = 0; k < n; k++) if (!ipiv[k])
                if (p_j == -1 \mid fabs(a_{j|k}) > fabs(a_{pj|pk})) \{ p_j = j; p_k = k; \}
        if (fabs(a[pj][pk]) < EPS) return false; // matrix is singular</pre>
        ipiv[pk]++;
        swap(a[pj], a[pk]);
        swap(b[pj], b[pk]);
        irow[i] = pj;
        icol[i] = pk;
        double c = 1.0 / a[pk][pk];
        a[pk][pk] = 1.0;
        for (int p = 0; p < n; p++) a[pk][p] *= c;
        for (int p = 0; p < m; p++) b[pk][p] *= c;</pre>
        for (int p = 0; p < n; p++) if (p != pk) {
            c = a[p][pk];
            a[p][pk] = 0;
            for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;
            for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;
    for (int p = n - 1; p >= 0; p --) if (irow[p] != icol[p]) {
        for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);
   return true;
2.11 Permutation and Combination
//Permutation
int arr[5] = {1,2,3,4,5};
    for(int i=0;i<5;i++)</pre>
```

```
cout << arr[i] << ' ';
cout <<'\n';
}while(next_permutation(arr,arr+5));
//also prev_permutation exist

//Combination
int arr[5] = {0, 0, 0, 1, 1}; // total : total cnt, 0 cnt : choose cnt
do{
    for(int i=0;i<5;i++)
        if(arr[i] == 0)
            cout << i+1 << ' ';
    cout <<'\n';
}while(next_permutation(arr,arr+5));</pre>
```

2.12 Lifting The Exponent

For any integers x, y a positive integer n, and a prime number p such that $p \nmid x$ and $p \nmid y$, the following statements hold:

- When p is odd:
 - If $p \mid x y$, then $\nu_p(x^n y^n) = \nu_p(x y) + \nu_p(n)$.
 - If n is odd and $p \mid x+y$, then $\nu_p(x^n+y^n) = \nu_p(x+y) + \nu_p(n)$.
- When p=2:
 - If $2 \mid x y$ and n is even, then $\nu_2(x^n y^n) = \nu_2(x y) + \nu_2(x + y) + \nu_2(n) 1$.
 - If 2 | x y and n is odd, then $\nu_2(x^n y^n) = \nu_2(x y)$.
 - Corollary
 - * If $4 \mid x y$, then $\nu_2(x + y) = 1$ and thus $\nu_2(x^n y^n) = \nu_2(x y) + \nu_2(n)$.
- For all *p*:
 - If gcd(n, p) = 1 and $p \mid x y$, then $\nu_p(x^n y^n) = \nu_p(x y)$.
 - If gcd(n, p) = 1, $p \mid x + y$ and n odd, then $\nu_p(x^n + y^n) = \nu_p(x + y)$.

3 Data Structure

3.1 Lazy Segment Tree

```
struct LazySeg {
    11 n;
    vector<ll> data, tree, lazy;
    LazySeg(11 n): n(n), data(n), tree(n<<2), lazy(n<<2) {}
    void seg_init(ll idx, ll s, ll e) {
        if (s == e) {
            tree[idx] = data[s];
            return;
        11 \text{ mid} = (s + e) >> 1;
        seg_init(idx<<1, s, mid);</pre>
        seg init(idx<<1|1, mid+1, e);
        tree[idx] = tree[idx<<1] + tree[idx<<1|1];</pre>
    void update_lazy(ll idx, ll s, ll e) {
        if (lazy[idx] != 0) {
            tree[idx] += (e-s+1) * lazy[idx];
            if (s != e) {
                 lazy[idx<<1] += lazy[idx];</pre>
                lazy[idx<<1|1] += lazy[idx];</pre>
            lazy[idx] = 0;
        }
    void seg_update(ll idx, ll s, ll e, ll l, ll r, ll d) {
        update_lazy(idx, s, e);
        if (1 > e \mid \mid r < s) return;
        if (1 <= s && e <= r) {
            tree[idx] += (e-s+1) * d;
```

```
if (s != e) {
                  lazv[idx<<1] += d;
                 lazy[idx<<1|1] += d;
             return;
         11 \text{ mid} = (s + e) >> 1;
         seg update(idx<<1, s, mid, l, r, d);</pre>
         seg update(idx<<1|1, mid+1, e, 1, r, d);
         tree[idx] = tree[idx<<1] + tree[idx<<1|1];</pre>
    ll seg_query(ll idx, ll s, ll e, ll l, ll r) {
         update_lazy(idx, s, e);
         if (1 \rightarrow e \mid \mid r < s) return 0;
        if (1 <= s && e <= r) return tree[idx];</pre>
        11 \text{ mid} = (s + e) >> 1;
        ll lsum = seg query(idx<<1, s, mid, l, r);
        ll rsum = seg_query(idx<<1|1, mid+1, e, 1, r);</pre>
         return lsum + rsum;
    // seg.init(v);
    void init(const vector<11>&v) {
         seg init(1, 0, n-1);
    // seg.update(l-1, r-1, d);
    void update(ll l, ll r, ll d) {
         seg update(1, 0, n-1, 1, r, d);
    // seq.query(l-1, r-1);
    11 query(11 1, 11 r) {
         if (1 > r) return 0;
         return seg_query(1, 0, n-1, 1, r);
};
```

3.2 Persistent Segment Tree

```
struct PST{
    11 n;
    vector<11> data;
    vector<vector<pll>>> tree;
    PST(11 n):n(n), data(n), tree(4*n) {}
    void seg init(ll idx, ll s, ll e){
            tree[idx].push back({0, data[s]});
            return;
        ll mid=(s+e)>>1;
        seg init(idx<<1, s, mid);</pre>
        seg init(idx<<1|1, mid+1, e);</pre>
        tree[idx].push_back({0, tree[idx<<1].back().second+tree[idx<<1|1].back().second});</pre>
    void seg update(ll idx, ll s, ll e, ll pos, ll val, ll ord){
        if(pos<s || pos>e) return;
        if(s==e){
            tree[idx].push_back({ord, val});
            return;
        11 mid=(s+e)>>1;
        seg_update(idx<<1, s, mid, pos, val, ord);</pre>
        seg update(idx<<1|1, mid+1, e, pos, val, ord);</pre>
        tree[idx].push_back({ord, tree[idx<<1].back().second+tree[idx<<1|1].back().second});</pre>
    11 seg query(11 idx, 11 s, 11 e, 11 1, 11 r, 11 ord){
        if(1>e || r<s)return 0;
```

}

```
if(1<=s && e<=r) {
            return prev(ranges::lower_bound(tree[idx], pll(ord, LLONG_MAX)))->second;
        ll mid=(s+e)>>1;
        return seg query(idx<<1, s, mid, l, r, ord)
                +seg query(idx<<1|1, mid+1, e, 1, r, ord);
    void init(const vector<11>&arr){
        data=arr:
        seg init(1, 0, n-1);
    void update(ll pos, ll val, ll ord){
        seg_update(1, 0, n-1, pos, val, ord);
    11 guery(11 1, 11 r, 11 ord){
        if(1>r)return 0;
        else return seg query(1, 0, n-1, 1, r, ord);
};
4 Graph
4.1 Diikstra
// O(ELogV)
vector<ll> dijk(ll n, ll s){
  vector<ll>dis(n,INF);
  priority_queue<pll, vector<pll>, greater<pll> > q; // pair(dist, v)
  dis[s] = 0;
  q.push({dis[s], s});
  while (!q.empty()){
    while (!q.empty() && visit[q.top().second]) q.pop();
    if (q.empty()) break;
    11 next = q.top().second; q.pop();
    visit[next] = 1:
    for (ll i = 0; i < adj[next].size(); i++)</pre>
      if (dis[adj[next][i].first] > dis[next] + adj[next][i].second){
        dis[adj[next][i].first] = dis[next] + adj[next][i].second;
        q.push({dis[adj[next][i].first], adj[next][i].first});}}
  for(ll i=0;i<n;i++)if(dis[i]==INF)dis[i]=-1;</pre>
  return dis;
4.2 LCA
const int MAXN = 100:
const int MAXLN = 9;
vector<int> tree[MAXN];
int depth[MAXN];
int par[MAXLN][MAXN];
void dfs(int nod, int parent) {
    for (int next : tree[nod]) {
        if (next == parent) continue;
        depth[next] = depth[nod] + 1;
        par[0][next] = nod;
        dfs(next, nod);
void prepare_lca() {
    const int root = 0;
    dfs(root, -1);
    par[0][root] = root;
    for (int i = 1; i < MAXLN; ++i)</pre>
        for (int j = 0; j < n; ++j)
```

```
par[i][j] = par[i - 1][par[i - 1][j]];
// find lowest common ancestor in tree between u & v
// assumption : must call 'prepare lca' once before call this
// O(LoaV)
int lca(int u, int v) {
    if (depth[u] < depth[v]) swap(u, v);</pre>
   if (depth[u] > depth[v]) {
        for (int i = MAXLN - 1; i >= 0; --i)
            if (depth[u] - (1 << i) >= depth[v])
               u = par[i][u];
    if (u == v) return u;
    for (int i = MAXLN - 1; i >= 0; --i) {
        if (par[i][u] != par[i][v]) {
           u = par[i][u];
           v = par[i][v];
   return par[0][u];
4.3 Centroid Decomposition
// O(n lg n) for centroid decomposition
auto cent_decom = [](const auto& adj) {
 const int n = adj.size() - 1;
 vector sz(n + 1, 1), dep(n + 1, 0), par(n + 1, 0);
 auto dfs = [&](const auto& self, int cur, int prv) -> void {
   for (auto [nxt, cost] : adj[cur]) {
     if (nxt == prv) continue;
      self(self, nxt, cur);
      sz[cur] += sz[nxt];
 };
  auto adjust = [&](int cur) {
   while (1) {
     int f = 0;
      for (auto [nxt, cost] : adj[cur]) {
       if (dep[nxt] || sz[cur] >= 2 * sz[nxt]) continue;
       sz[cur] -= sz[nxt], sz[nxt] += sz[cur];
       cur = nxt, f = 1;
        break;
      if (!f) return cur;
 };
  auto rec = [&](const auto& self, int cur, int prv) -> void {
   cur = adjust(cur);
   par[cur] = prv;
    dep[cur] = dep[prv] + 1;
    for (auto [nxt, cost] : adj[cur]) {
     if (dep[nxt]) continue;
     self(self, nxt, cur);
 };
 dfs(dfs, 1, 0);
 rec(rec, 1, 0);
  return pair(dep, par);
4.4 Minimum Spanning Tree
// O(ELogV)
```

ll prim() {

```
priority queue<pll, vector<pll>, greater<pll> > q;
  11 count = 0; 11 ret = 0;
  q.push(make pair(0, 0)); // (cost, vertex)
  while (!q.empty()){
    11 x = q.top().second; // also able to get edges
    visit[x] = 1; ret += q.top().first; q.pop(); count++;
    for (ll i = 0; i < adj[x].size(); i++)</pre>
      q.push({adj[x][i].second, adj[x][i].first});
    while (!q.empty() && visit[q.top().second]) q.pop();
  if (count != n) return -1;
  else return ret;
11 Kruskal(){
 11 ret = 0;vector<11>par;
  iota(par.beging(),par.end(),0);
  vector<pair<11, pll>> e;
  for(ll i= 0; i < n; i++)
    for(ll j=0; j < adj[i].size(); j++)</pre>
      e.push_back({adj[i][j].second, {i, adj[i][j].first}});
  sort(e.begin(), e.end());
  for(ll i=0; i < e.size(); i++){</pre>
    11 x = e[i].second.first,y = e[i].second.second;
    if(find(x) != find(y)){
     union(x, y);
      ret += e[i].first;
  11 p=find(0);
  for(ll i=1;i<n;i++){</pre>
    if(find(i)!=p)return -1;
  else return ret;
     Offline Dynamic Connectivity
struct OFDC {
    vector<tlll> query;
    vector<ll> grp, sz;
    vector<vector<pll>>> tree;
    map<pll, 11> conn;
    11 n, q;
    OFDC(11 n, 11 q): n(n), q(q), query(q+1), grp(n+1), sz(n+1, 1), tree(4*(q+1)) {
        iota(grp.begin(), grp.end(), 0);
    void update(ll node, ll s, ll e, ll l, ll r, pll edge) {
        if (r < s || e < 1) return;
        if (1 <= s && e <= r) {</pre>
            tree[node].push back(edge);
            return;
        11 \text{ mid} = (s + e) >> 1;
        update(node << 1, s, mid, l, r, edge);</pre>
        update(node << 1 | 1, mid + 1, e, l, r, edge);
    11 _find(ll x) {
        if (grp[x] == x) return x;
        return find(grp[x]);
    pll _union(ll x, ll y) {
        x = _find(x), y = _find(y);
        if (x == y) return {-1, -1};
        if (sz[x] < sz[y]) swap(x, y);</pre>
        grp[y] = x;
```

```
sz[x] += sz[y];
        return {x, y};
    void _delete(ll u, ll v) {
        sz[u] -= sz[v];
        grp[v] = v;
    void dfs(ll node, ll s, ll e) {
        vector<pll> rconn;
        for (auto& [u, v]: tree[node]) {
            auto [x, y] = union(u, v);
            if (x != -1) rconn.push_back({x, y});
        if (s == e) {
            if (get<0>(query[s]) == 3) {
                cout << (_find(get<1>(query[s])) ==
                    _find(get<2>(query[s]))) << '\n';
        } else {
            11 \text{ mid} = (s + e) >> 1;
            dfs(node << 1, s, mid);</pre>
            dfs(node << 1 | 1, mid + 1, e);
        for (auto& [u, v]: rconn) {
            _delete(u, v);
    void run() {
        for (ll i = 0; i < q; i++) {
            auto& [type, u, v] = query[i];
            cin >> type >> u >> v;
            if (u > v) swap(u, v);
            if (type == 1) {
                conn[{u, v}] = i;
            } else if (type == 2) {
                update(1, 0, q, conn[{u, v}], i, {u, v});
                conn.erase({u, v});
        for (auto&[edge, time] : conn) {
            auto&[u, v] = edge;
            update(1, 0, q, time, q, {u, v});
        dfs(1, 0, q);
    }
};
5
    String
5.1 KMP
void calculate pi(vector<int>& pi, const string& str) {
  pi[0] = -1;
  for (int i = 1, j = -1; i < str.size(); i++) {
    while (j \ge 0 \&\& str[i] != str[j + 1]) j = pi[j];
    if (str[i] == str[j + 1]) pi[i] = ++j;
    else pi[i] = -1;
// returns all positions matched
// 0(|text|+|pattern|)
vector<int> kmp(const string& text, const string& pattern) {
  vector<int> pi(pattern.size()), ans;
  if (pattern.size() == 0) return ans;
  calculate_pi(pi, pattern);
  for (int i = 0, j = -1; i < text.size(); i++) {
    while (j >= 0 && text[i] != pattern[j + 1]) j = pi[j];
```

```
if (text[i] == pattern[j + 1]) {
    j++;
    if (j + 1 == pattern.size()) ans.push_back(i - j), j = pi[j];
}
return ans;
}

5.2 Z Algorithm

// Z[i]: maximum common prefix length of &s[0] and &s[i] with O(|s|)
auto get_z = [](const string& s) {
    const int n = s.size(); vector z(n, 0); z[0] = n;
    for (int i = 1, l = -1, r = -1; i < n; i++) {
    if (i <= r) z[i] = min(r - i + 1, z[i - 1]);
    while (i + z[i] < n && s[z[i]] == s[i + z[i]]) z[i]++;
    if (r < i + z[i] - 1) l = i, r = i + z[i] - 1;</pre>
```

6 Hash

};

return z;

6.1 Basic Hash

```
struct chash {
    size_t operator()(const pll& _x) const {
        auto [x, y] = _x;
        size_t hx = hash<ll>()(x);
        size_t hy = hash<ll>()(y);
        return ((hx<22) | (hx>>22)) ^ hy;
}
size_t operator()(const tuple<ll, string, ll>& _x) const {
        auto [x, y, z] = _x;
        size_t hx = hash<ll>()(x);
        size_t hy = hash<string>()(y);
        size_t hz = hash<ll>()(z);
        return ((hx<22) | (hx>>22)) ^ ((hy<17) | (hy>>17)) ^ hz;
}
int main() {
        unordered_map<pl1, ll, chash> a;
}
```