#### PALEA

Know the unknown

#### NETWORK / DATAPROTECTION

- Deny hackers entry
- Keep sensitive data inside
- Not through front door (known corporate gateway)
- Not through back door (the unknown)

#### WHY SCAN?

- · A hacker break in can lead to several forms of damage.
  - Stealing customer records, financial forecasts
  - Undisclosed business plans and other property
- Your systems might be used in an attack
- · Legal obligations to report a hacker break-in
- Compliancy regulations or consequences

#### BACKDOORS TO THE INTERNET

- Unauthorized routes to the internet might enable compromised machines to talk to their hacker (botnets)
- Stealing computer resources
- Unauthorized routes to the internet might be used as covert channel for stealing data
- Unauthorized routes might be used as secret channel to enter the company

#### HOW TO FIND THESE CONNECTIONS

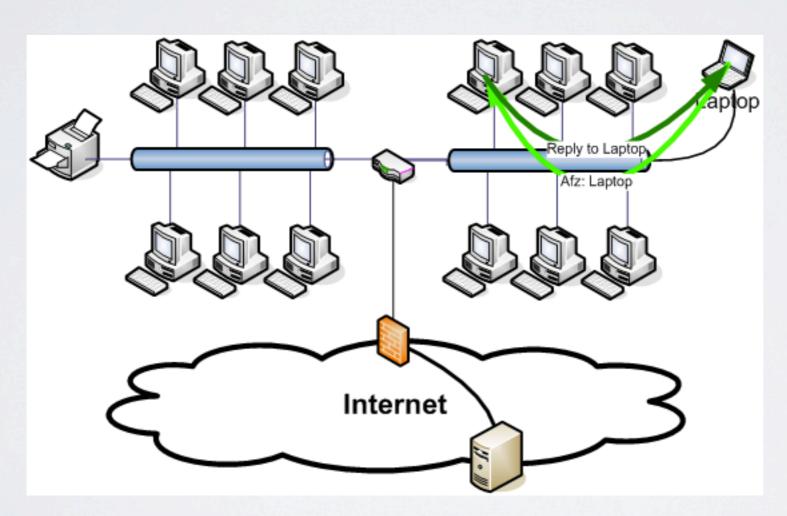
- We might do our very best to find external connections from the outsite. (a needle in a haystack)
- We better ask all the machines if they are a router or know a router to outside.
- · Let the sum of the machines do the work

#### HOW TO FIND THESE CONNECTIONS

• We simply need a system to try and deliver a package to the internet, something we can catch, identify and trace back to the specific system.

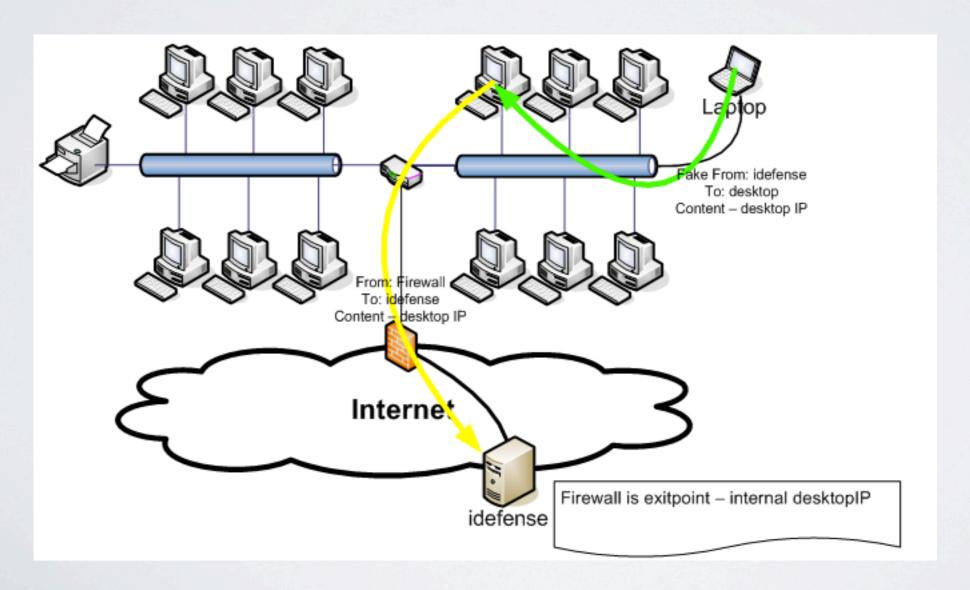
#### INTERNET INTERFACE AS WE KNOW IT

· Laptop sends a 'HELLO' to a system, the system replies



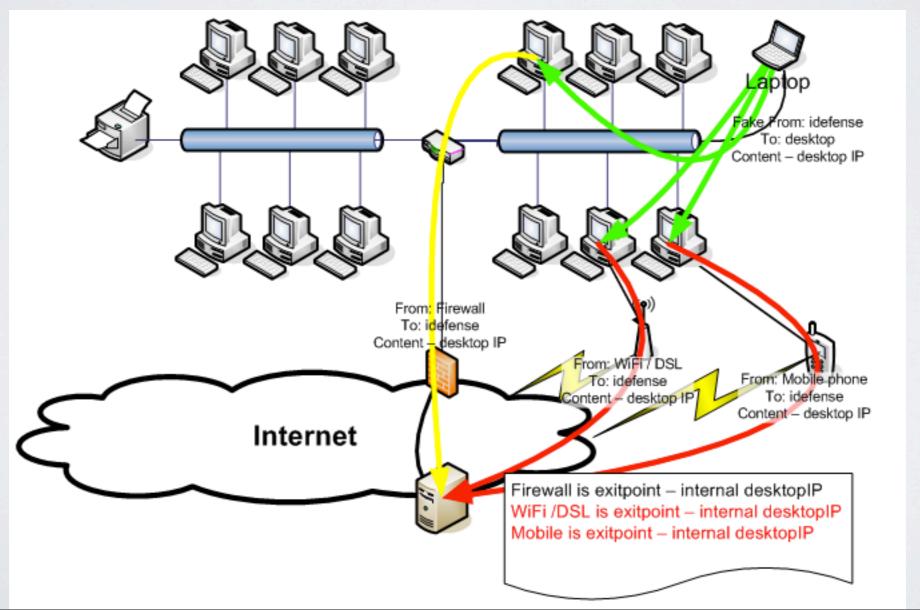
#### INTERNET INTERFACE AS WE KNOW IT

• Laptop sends a 'HELLO' to with a fake FROM address to a system, the reply gets sent to a server on the internet. Firewall MIGHT stop this.



# INTERNET INTERFACES AS THEY CAN BE

 Laptop sends packets to several systems with fake FROM address, some desktops might know alternatives ways to the outside.

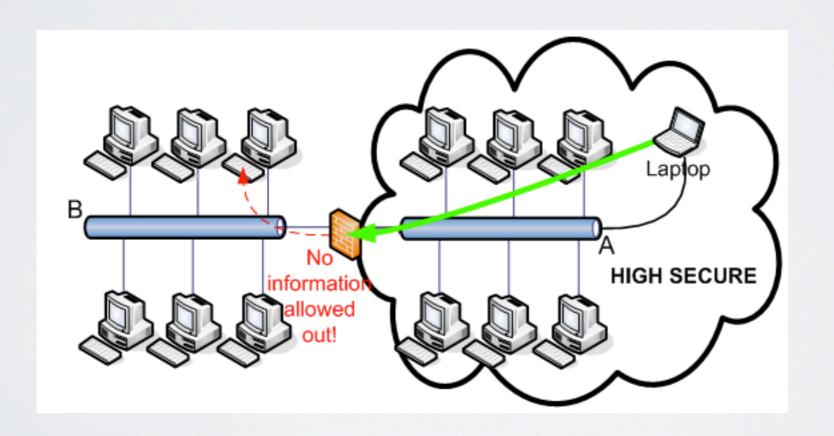


#### HIGHLY SECURE ENVIRONMENT

• Information in highly secure environments is not supposed to leak out to the rest of the network

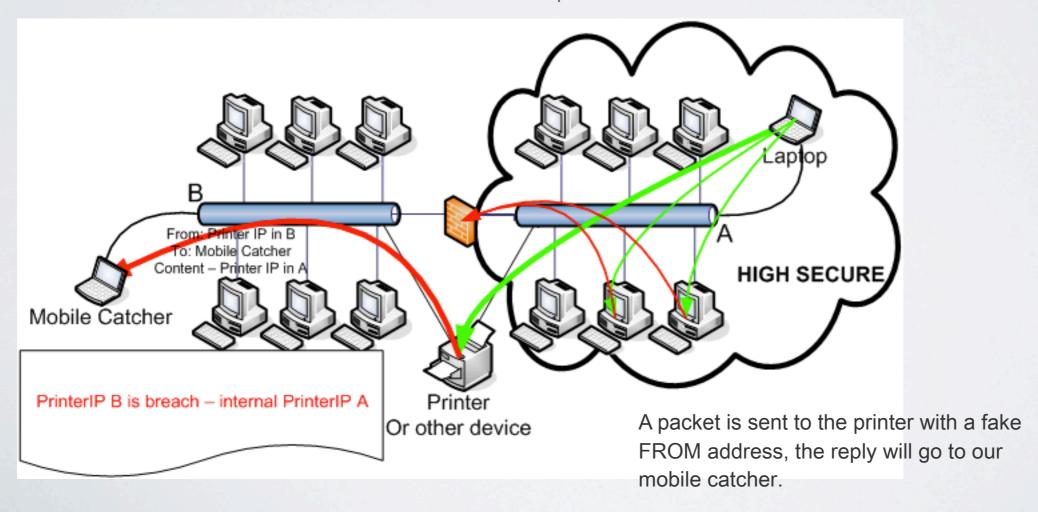
# INTERNAL INTERFACING AS WE KNOW IT

• Highly secure network is not allowed to leak data, a firewall is in place. This is the situation as documented



# INTERNAL INTERFACING ASTHEY MIGHT BE

- A printer is used to print to from the high secure environment but also from a user network.
- · This is can be abused as interface. The IT department never knew or realized.



# BEING IN CONTROL OF THE ENVIRONMENT

- These scans should be repeated 24/7 several times a day for full coverage
  - GSM / Modem connections aren't permanent, you need it to be active to find it
  - New interfaces might be created
  - IT is changing all the time
- · Security scans should be repeated as new vulnerabilities might develop

#### TIMETO GO TECH!

 Spoof IP to trick every system in an attempt to speak to the internet

· Tag all packets to recognize them once/if they arrive

#### THE PACKETS

- We will send out specially crafted ICMP and UDP packets. The structure of these packets are explained in the following slides.
- Some information is stored redundant to ensure it will be available after receiving the reply. This redundancy can also be used as sanity check for filtering out rogue packets floating on the net.

#### (THROW 1/3) ICMP ECHO REQUEST

0 1 2	3
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1	
Version  IHL  Type of Service  Total Length	
Identification   Flags  Fragment Offset   +-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+	
Time to Live   Protocol   Header Checksum	
Source Address (Spoofed to CATCHER )	
Destination Address (of potential gateway)	
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-	
Type	
Identifier (Session #)	
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-	
+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-+-	

- As data we place the Destination Address (the system we send the packet to) we also place it in the data in case another machine is the gateway or replying machine. We want to use a maximum of 64 bits of data since, if we get a DU of PU ICMP reply we will receive Internet Header + 64 bits of Original Data Datagram.
- [\*1] Unique is a unique number (counter) within a session, we also store locally what we send out, If we can ID the session and retrieve this unique number we can find the packet that was placed on the wire initially.

#### (THROW 2/3) UDP PACKET

```
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
|Version| IHL |Type of Service|
                Total Length
| Identification (NR in session)|Flags|
                 Fragment Offset
Time to Live I
        Protocol I
                Header Checksum
Source Address (Spoofed to CATCHER )
Destination Address (of potential gateway)
Source Port (session #)
               Destination port (53)
+-+-+-+-+-+-+-+-+-+-+ data segment -+-+-+-+-+-+-+-+-+
     Destination Address (of potential gateway)
I Identifier (Session #) | Sequence Number (unique *1) |
```

- As data we place the Destination Address (the system we send the packet to) we also place it in the data in case another machine is the gateway or replying machine. We want to use a maximum of 64 bits of data since, if we get a DU of PU ICMP reply we will receive Internet Header + 64 bits of Original Data Datagram.
- [1] (page 3, RFC 792) The address of the gateway or host that composes the ICMP message. Unless otherwise noted, this can be any of a gateway's addresses

#### (THROW 3/3) DATA IN PACKETS

- In ICMP and UDP data fields we can stuff at most 8 bytes that we can expect back as stated in RFC792 on page 3. Stuffing the IP address of the system we shoot the initial packet to will take up 4 bytes. This leaves us with 4 bytes, just enough to store session # and unique #
- It is necessary to store the system we shoot the packet to in case the packet leaves that machine from another interface or the reply is generated by another machine than the targeted machine (i.e. in case of Destination Unreachable reply). The from address will then contain the IP address of the replying interface. The data should therefore contain the system we targeted. Off course using session # and unique # we can recall the full initial packet from the logs in the 'thrower'. We include it anyway in order to generate live results without correlation phase.

#### (CATCH 1/4) WHATTO EXPECT

- We've send out ICMP Echo Request and UPD packets. We might get back:
  - ICMP reply
  - DU as reply to ICMP
  - DU as reply to UDP
  - PU as reply to UDP

#### (CATCH 2/4) ICMP REPLY

```
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
|Version| IHL |Type of Service|
                 Total Length
Identification
            |Flags|
                 Fragment Offset
Time to Live I
        Protocol I
                Header Checksum
Source Address (Sending Interface )
Destination Address (of potential gateway)
Type
        Code
                 Checksum
I Identifier (Session #) | Sequence Number (unique *1) | I
+-+-+-+-+-+-+-+-+-+-+ data segment -+-+-+-+-+-+-+-+-+
      Destination Address (of potential gateway)
I Identifier (Session #) | Sequence Number (unique *1) |
```

- As data we place the Destination Address (the system we send the packet to) we also place it in the data in case another machine is the gateway or replying machine. We want to use a maximum of 64 bits of data since, if we get a DU of PU ICMP reply we will receive Internet Header + 64 bits of Original Data Datagram.
- [\*1] Unique is a unique number (counter) within a session, we also store locally what we send out, If we can ID the session and retrieve this unique number we can find the packet that was placed on the wire initially.

#### (CATCH 3/4) ICMP DU

```
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
|Version| IHL |Type of Service|
            Total Length
Identification |Flags|
              Fragment Offset
Time to Live |
       Protocol |
              Header Checksum
Source Address (Sending Interface ) | < ----- A
Destination Address (CATCHER)
Type DU | Code DU |
              Checksum
DATA IN ICMP DU
|Version| IHL |Type of Service|
             Total Length
Identification
          |Flags| Fragment Offset
| Time to Live |
       Protocol |
              Header Checksum
Source Address (Spoofed to CATCHER )
< ----- B if not equals to A we've</pre>
    Destination Address (of potential gateway)
found a server replying over NAT
Type
       Code
              Checksum
| Identifier (Session #) | Sequence Number (unique *1) |
Destination Address (of potential gateway)
I Identifier (Session #) | Sequence Number (unique *1) |
```

#### (CATCH 4/4) ICMP DU

```
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
|Version| IHL |Type of Service| Total Length
Identification |Flags|
              Fragment Offset
| Time to Live |
       Protocol | Header Checksum
Source Address (Sending Interface ) | < ----- A
Destination Address (CATCHER )
Type DU | Code DU/PU |
              Checksum
DATA IN UDP DU/PU
|Version| IHL |Type of Service|
             Total Length
| Identification (NR in session)|Flags| Fragment Offset
| Time to Live |
      Protocol |
              Header Checksum
Source Address (Spoofed to CATCHER )
Destination Address (of potential gateway)
                    | < ----- B if not equals to A we've</pre>
Source Port (Session #) | Destination port (53)
Checksum
Destination Address (of potential gateway)
Identifier (Session #)
         | Sequence Number (unique *1) |
```

The setup, vmware fusion with

- 2 Linux VM's, sender and catcher
- Windows VM, our dual homed victim

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#### PROOF OF CONCEPT

The setup, vmware fusion with

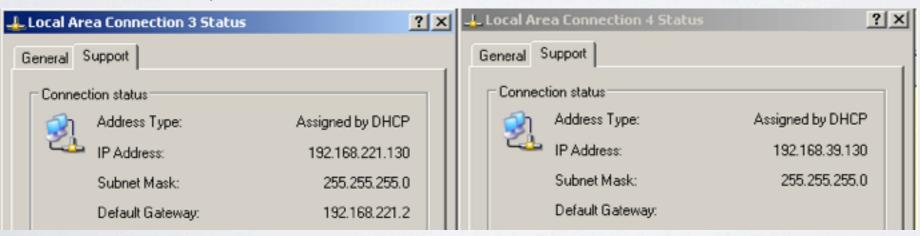
Linux, Is thrower this POC.

th1 Link encap:Ethernet HWaddr 00:0c:29:08:56:8d inet addr:192.168.39.131 Bcast:192.168.39.255 Mask:255.255.255.0

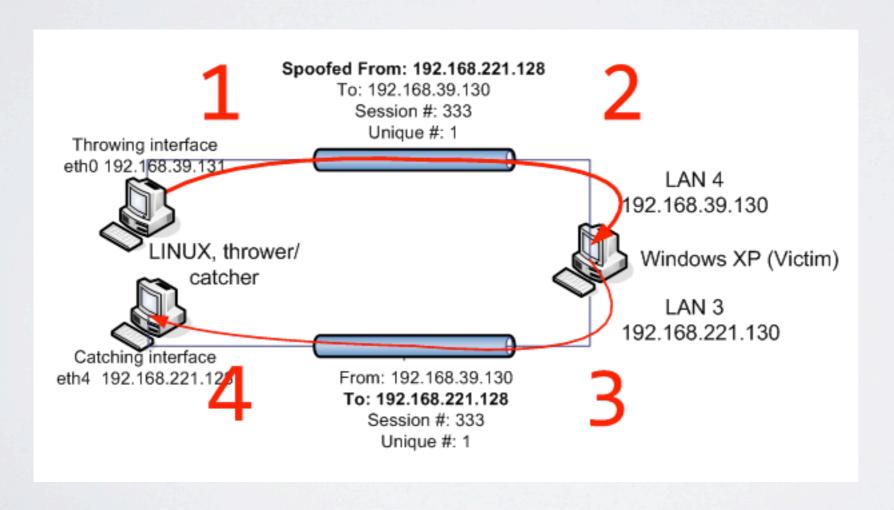
· Linux, Is cacher in this POC

eth4 Link encap:Ethernet HWaddr 00:0c:29:08:56:83 inet addr:192.168.221.128 Bcast:192.168.221.255 Mask:255.255.25

· Windows, dual homed



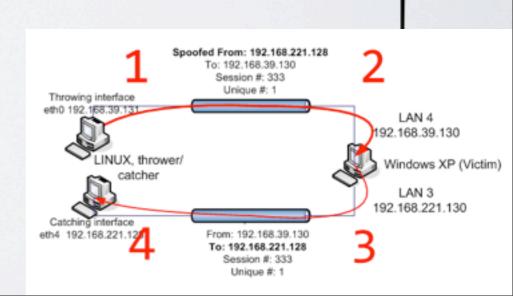
This setup gives the following network diagram. The numbers indicate the path and steps a specially crafted packet will take.



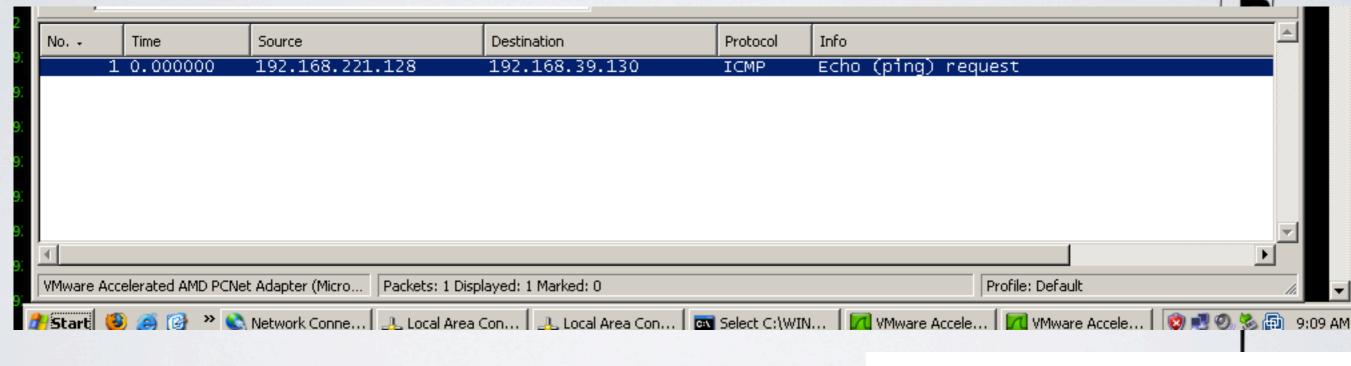
**Step I:** We send a packet from the Linux Thrower machine to the dual homed XP.We want to test if it has an interface on the 192.168.221.0/24

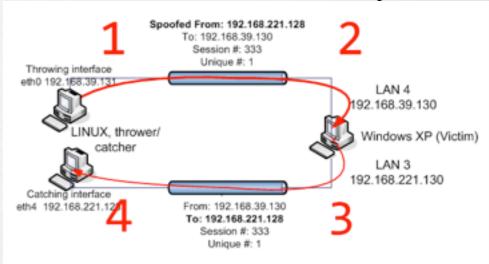
```
network. >>> from poc import SendUDP, SendICMP
          >>> SendICMP("192.168.221.128","192.168.39.130",333,1)
                86621.7597001
```

Note: in order to be able to use this code you need python and impacket installed (http://www.python.org/) (http://oss.coresecurity.com/projects/impacket.html)

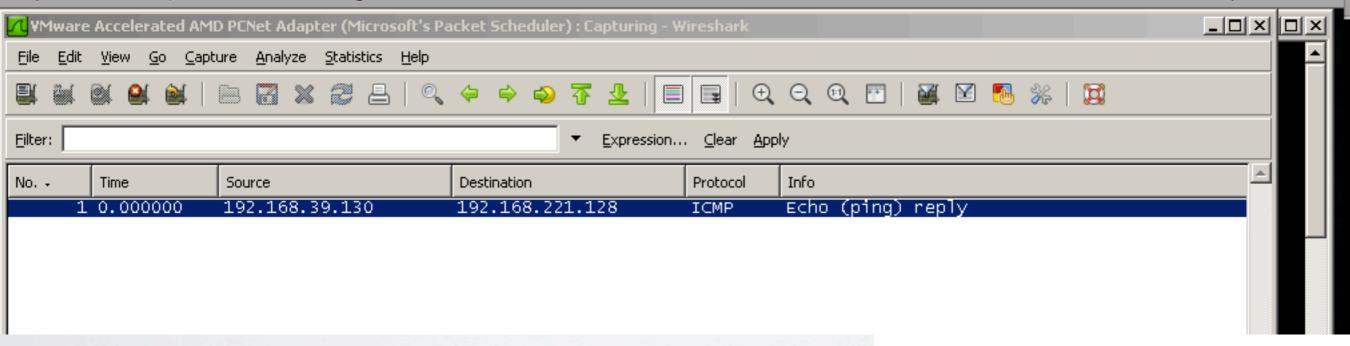


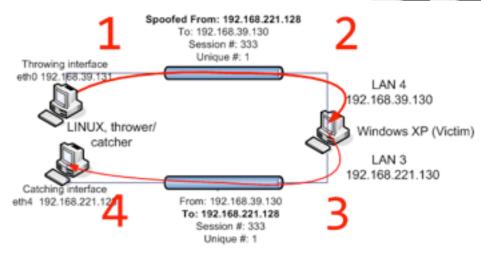
Step 2: The spoofed packet is received on interface LAN 4.





**Step 3:** The reply is put on interface LAN 3 determined by the route the packet needs to be delivered to 192.168.221.128

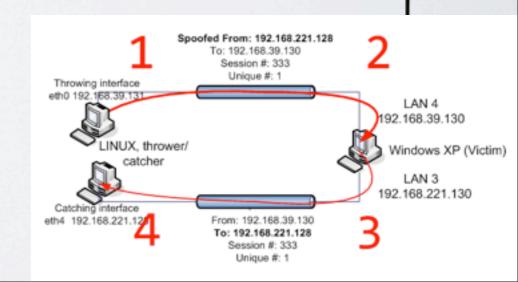




Step 4: The reply received by the Linux machine on eth4

```
xychix@pentest:~/cathro-dev/WIP/cathro-bin$ sudo tcpdump -i eth4
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on eth4, link-type EN10MB (Ethernet), capture size 96 bytes
08:22:30.472170 IP 192.168.39.130 > pentest.local: ICMP echo reply, id 333, seq 1, length 16
```

eth4 Link encap:Ethernet HWaddr 00:0c:29:08:56:83 inet addr:192.168.221.128 Bcast:192.168.221.255 Mask:255.255.255.0





# QUESTIONS?

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