

Problem: Given a singly-linked list of integers, determine whether or not it is a palindrome.

A **palindrome** is something that reads the same backward as forward.

**approach 1:** we know that we're going to be doing some sort of reversal here.

A straightforward solution would be to iterate through the linked list & push these values into a stack.

ex 1)  $r \rightarrow a \rightarrow c \rightarrow e \rightarrow c \rightarrow a \rightarrow r \rightarrow \emptyset$

stack:  $r, a, c, e, c, a, r$

ex 2)  $a \rightarrow b \rightarrow c \rightarrow \emptyset$

stack:  $a, b, c$

Then, iterate through the beginning of the stack again while popping the stack until it is empty.

The **fast runner** technique can also be used to cut the iteration in half, since one half of a palindrome should be equal to the other half reversed.

**approach 2:** To reduce space complexity, we can reverse the first half of the linked list, then iterate through the reversed half & the remaining un-reversed half to check if it is a valid palindrome.

1) The fast pointer travels 2 nodes for every 1 node of the slower pointer, its purpose is to reveal the midpoint of a linked list.

**Odd-numbered case:**

$a \rightarrow b \rightarrow a \rightarrow \emptyset$

$s \text{ --- } s$

$f \xrightarrow{\text{.next}} f$

**even-numbered case:**

$a \rightarrow b \rightarrow b \rightarrow a \rightarrow \emptyset$

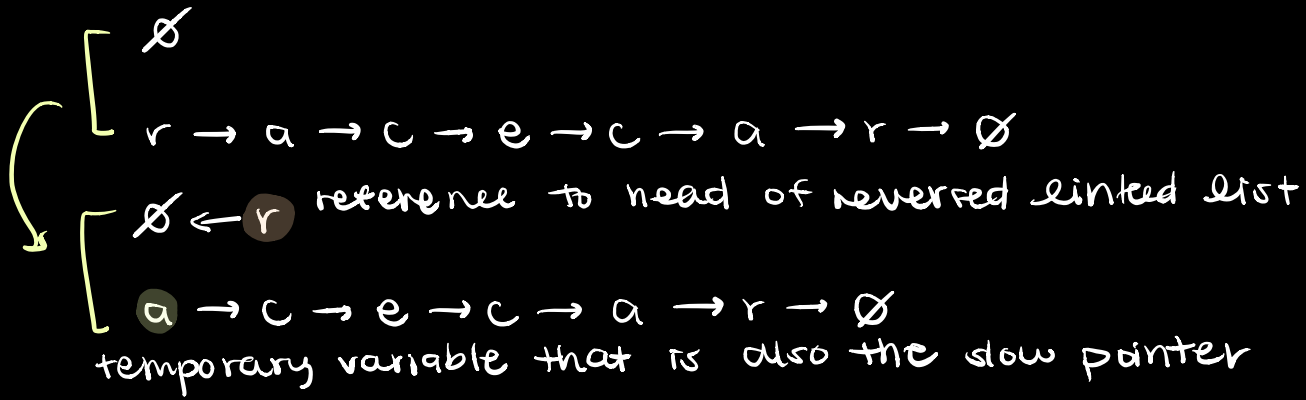
$s \text{ --- } s \text{ --- } s$

$f \xrightarrow{\text{.next}} f \xrightarrow{\text{.next}} f$

from these examples, we can conclude that we can continue iterating until  $f$  is null, or  $f.next$  is null.

2) The slow pointer moves 1 node at a time, and works as a temporary variable to overwrite the linked list & reverse it.

The first node of our linked list should be the last node of our reversed linked list. Therefore, this node's next node would be null.



## Implementation

```
function isPalindrome(head) {
```

```
  let reversedHead = null;
```

```
  let fastPointer = head;
```

```
  while (fastPointer && fastPointer.next) {
```

```
    fastPointer = fastPointer.next.next; // move the fast ptr first
```

```
    let temp = head.next; // store the next node
```

```
    head.next = reversedHead; // set the rev. list's next node
```

```
    reversedHead = head; // move rev. list's pointer
```

```
    head = temp; // move the slow pointer to next node to iterate
```

```
  }
```

```
  // the head should now be at the start of
```

```
  // the non-reversed half of the linked list
```

```
  // everything preceding it has been prepended
```

```
  // to the reversed linked list
```

```
  // if fast is not null, the linked list has an
```

```
  // odd-numbered length, so we should disregard
```

also ensures  
length  $\geq 2$

// the midpoint

```
if (fastPointer) {  
    head = head.next;  
}
```

// iterate through the original linked list until we hit null

```
while (head) {  
    if (head.value !== reversedHead.value) {  
        return false;  
    }  
    head = head.next;  
    reversedHead = reversedHead.next;  
}
```

```
return true;
```

```
}
```