

Algorithms: Homework 2

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Problem 1

Using Figure 8.3 as a model, illustrate the operation of RADIX-SORT on the following list of English words: NOD, HOG, SHY, BAN, BAR, JET, EBB, PAR, ASH, PET, TED, ROT, FIG.

Solution.

$0(\text{initial})$	$\text{digit } 1$	$\text{digit } 2$	$\text{digit } 3$
NOD	EBB	BAN	ASH
HOG	NOD	BAR	BAN
SHY	TED	PAR	BAR
BAN	HOG	EBB	EBB
BAR	FIG	TED	FIG
JET	ASH	JET	HOG
EBB	BAN	PET	JET
PAR	BAR	SHY	NOD
ASH	PAR	FIG	PAR
PET	JET	NOD	PET
TED	PET	HOG	ROT
ROT	ROT	ROT	SHY
FIG	SHY	ASH	TED

■

Problem 2

Show how to sort n integers in the range 0 to $n^2 - 1$ in $\mathcal{O}(n)$ time.

Solution.

■

Problem 3

Using Figure 8.4 as a model, illustrate the operation of BUCKET-SORT on the array $A = \langle .49, .30, .56, .74, .68, .50, .69, .37, .41, .72 \rangle$

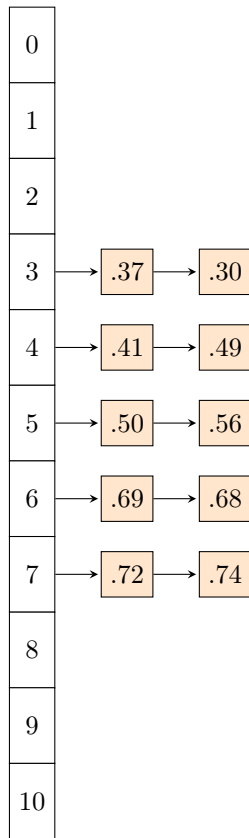


Figure 1: Insert data

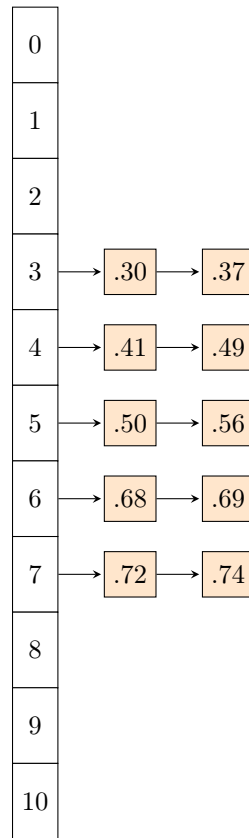


Figure 2: Sort each bucket

Solution. To get sorted data, chain all in Figure 2 from bucket 0 to bucket 10:
 $\langle .30, .37, .41, .49, .50, .56, .68, .69, .72, .74 \rangle$

Problem 4

Suppose that you have a “black-box” worst-case linear time median subroutine. Give a simple, linear-time algorithm that solves the selection problem for an arbitrary order statistic.

Solution. Psuedocode for finding the k -th smallest element:

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FIND-K-TH-SMALLEST( $X, k$ )
1:  if  $i > A.heap-size$ 
2:    error  $A[i]$  does not exist
3:   $key = A[i]$ 
4:   $A[i] = A[A.heap-size]$ 
5:   $A.heap-size = A.heap-size - 1$ 
6:  if  $A[i] < key$ 
7:    MAX-HEAPIFY( $A, i$ )
8:  else
9:     $key1 = A[i]$ 
10:    $A[i] = key$ 
11:   HEAP-INCREASE-KEY( $A, i, key1$ )

```

Problem 5

Let $X[1..n]$ and $Y[1..n]$ be two arrays, each containing n numbers already in sorted order. Give an $\mathcal{O}(\lg n)$ -time algorithm to find the median of all $2n$ elements in arrays X and Y .

Solution. Psuedocode for finding the median of arrays $X[1..n]$ and $Y[1..n]$, xi is X array's starting index and xj is X array's last index; yi is Y array's starting index and yj is Y array's last index:

FIND-MEDIAN(X, Y, xi, xj, yi, yj, n)

```

1:  if  $n == 1$ 
2:    // Both arrays have only one element, pick the smaller one as median
3:    return MIN( $X[1], Y[1]$ )
4:   $xmid = X[(xi + xj)/2]$  //  $X$ 's median
5:   $ymid = Y[(yi + yj)/2]$  //  $Y$ 's median
6:  if  $xmid < ymid$ 
7:    // median must be greater than  $xmid$  and less than  $ymid$ 
8:    return FIND-MEDIAN( $X, Y, xmid + 1, xj, yi, ymid, n/2$ )
9:  else
10:   // median must be less than  $xmid$  and greater than  $ymid$ 
11:   return FIND-MEDIAN( $X, Y, xi, xmid, ymid + 1, yj, n/2$ )

```

This algorithm reduce the array size by half every recursion call, similar with binary search. Thus, it's running time is $\mathcal{O}(\lg n)$. ■

Problem 6

Demonstrate the insertion of the keys 23, 38, 19, 40, 29, 43, 12, 27, 11 into a hash tale with collisions resolved by chaining. Let the table have 9 slots, and let the hash function be $h(k) = k \pmod{9}$.

Solution. ■

- (1) insert 23 $\rightarrow 23 \pmod{9} = 5$
- (2) insert 38 $\rightarrow 38 \pmod{9} = 2$
- (3) insert 19 $\rightarrow 19 \pmod{9} = 1$
- (4) insert 40 $\rightarrow 40 \pmod{9} = 4$
- (5) insert 29 $\rightarrow 29 \pmod{9} = 2$, insert head
- (6) insert 43 $\rightarrow 43 \pmod{9} = 7$
- (7) insert 12 $\rightarrow 12 \pmod{9} = 3$
- (8) insert 27 $\rightarrow 27 \pmod{9} = 0$
- (9) insert 11 $\rightarrow 11 \pmod{9} = 2$, insert head

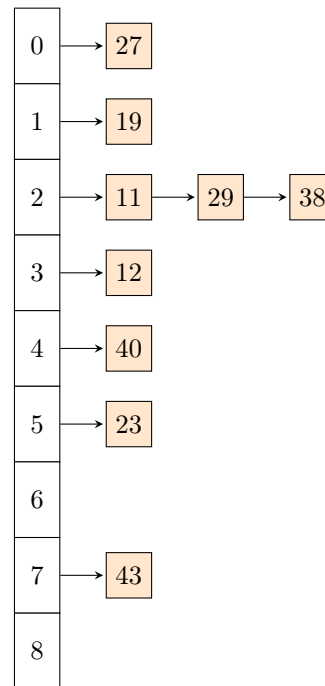


Figure 3: Hash table with chaining

Problem 7

Consider inserting the keys 10, 23, 31, 4, 12, 28, 17, 87, 58 into a hash table of length $m = 11$ using open addressing with the auxiliary hash function $h'(k) = k \pmod{m}$. Illustrate the result of inserting these keys using **linear probing**, using **quadratic probing** with $c_1 = 1$, and $c_2 = 3$, and using **double hashing** with $h_2(k) = 1 + (k \pmod{m-1})$.

- (1) insert 10 $\rightarrow 10 \pmod{11} = 10$
- (2) insert 23 $\rightarrow 23 \pmod{11} = 1$
- (3) insert 31 $\rightarrow 31 \pmod{11} = 9$
- (4) insert 4 $\rightarrow 4 \pmod{11} = 4$
- (5) insert 12 $\rightarrow 12 \pmod{11} = 1 \times$ collision
insert 12 to slot 2 because slot 1 is occupied
- (6) insert 28 $\rightarrow 28 \pmod{11} = 6$
- (7) insert 17 $\rightarrow 17 \pmod{11} = 6 \times$
insert 17 to index 7 because slot 6 is occupied
- (8) insert 87 $\rightarrow 87 \pmod{11} = 10 \times$
insert 87 to index 0 because slot 10 is occupied
- (9) insert 58 $\rightarrow 58 \pmod{11} = 3$

0	87
1	23
2	12
3	58
4	4
5	0
6	28
7	17
8	0
9	31
10	10

Figure 4: linear probing

resolve collision: $h(k) + 1 * i + 3 * i^2 \pmod{11}$

- (1) insert 10 $\rightarrow 10 \pmod{11} = 10$
- (2) insert 23 $\rightarrow 23 \pmod{11} = 1$
- (3) insert 31 $\rightarrow 31 \pmod{11} = 9$
- (4) insert 4 $\rightarrow 4 \pmod{11} = 4$
- (5) insert 12 $\rightarrow 12 \pmod{11} = 1 \times$ collision
 $(1 + 12 + 3 * 144) \pmod{11} = (1 + 1 + 3) \pmod{11} = 5$
insert 12 to slot 5
- (6) insert 28 $\rightarrow 28 \pmod{11} = 6$
- (7) insert 17 $\rightarrow 17 \pmod{11} = 6 \times$
 $(6 + 17 + 3 * 289) \pmod{11} = (6 + 6 + 9) \pmod{11} = 10 \times$
 $(10 + 17 + 3 * 289) \pmod{11} = (10 + 6 + 9) \pmod{11} = 3$
insert 17 to slot 3
- (8) insert 87 $\rightarrow 87 \pmod{11} = 10 \times$
 $(10 + 87 + 3 * 87 * 87) \pmod{11} = (10 + 10 + 3) \pmod{11} = 1 \times$
 $(1 + 87 + 3 * 87 * 87) \pmod{11} = (1 + 10 + 3) \pmod{11} = 3 \times$
 $(3 + 87 + 3 * 87 * 87) \pmod{11} = (3 + 10 + 3) \pmod{11} = 5 \times$
 $(5 + 87 + 3 * 87 * 87) \pmod{11} = (5 + 10 + 3) \pmod{11} = 7$
insert 87 to slot 7
- (9) insert 58 $\rightarrow 58 \pmod{11} = 3 \times$
 $(7 + 58 + 3 * 58 * 58) \pmod{11} = (7 + 3 + 5) \pmod{11} = 4 \times$
 $(4 + 58 + 3 * 58 * 58) \pmod{11} = (4 + 3 + 5) \pmod{11} = 1 \times$
 $(1 + 58 + 3 * 58 * 58) \pmod{11} = (1 + 3 + 5) \pmod{11} = 9 \times$
 $(9 + 58 + 3 * 58 * 58) \pmod{11} = (9 + 3 + 5) \pmod{11} = 6 \times$
 $(6 + 58 + 3 * 58 * 58) \pmod{11} = (6 + 3 + 5) \pmod{11} = 3 \times$
 $(3 + 58 + 3 * 58 * 58) \pmod{11} = (3 + 3 + 5) \pmod{11} = 0$
insert 58 to slot 0

0	58
1	23
2	0
3	17
4	4
5	12
6	28
7	87
8	0
9	31
10	10

Figure 5: quadratic probing

Problem 8

Solve the following assembly-line problem:

$$e_1 = 1, e_2 = 1, x_1 = 7, x_2 = 3$$

$$a_{1,1} = 3, a_{1,2} = 2, a_{1,3} = 5, a_{1,4} = 4, a_{1,5} = 2$$

$$a_{2,1} = 4, a_{2,2} = 2, a_{2,3} = 4, a_{2,4} = 6, a_{2,5} = 4$$

$$t_{1,1} = 1, t_{1,2} = 2, t_{1,3} = 3, t_{1,4} = 1$$

$$t_{2,1} = 2, t_{2,2} = 3, t_{2,3} = 1, t_{2,4} = 2.$$

Solution. ■

Problem 9

Find an optimal parenthesization of a matrix-chain product whose sequence of dimension is $\langle 5, 7, 3, 2, 4, 5 \rangle$

Solution. ■

Problem 10

Determine an LCS of $\langle 1, 1, 0, 1, 0, 1 \rangle$ and $\langle 0, 0, 1, 1, 0, 1, 1 \rangle$.

Solution. ■