

COSC 3319-01 Lab 3 Report

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1 Option C — Relative File

1.1 Implementation Overview

1.2 Hash Functions

1.2.1 MyHash (Bad Hash)

The MyHash function provided is intentionally poor. It only uses a few characters from the key, which means that keys that differ outside the positions checked by the hash could produce the same hash, creating a collision. The float conversion, only to convert back to int at the end, seems unnecessary, and the constant divisors 517, 217, and 256 do not help spread the values out; they just scale the sum.

```
1 // === MY HASH (given; intentionally poor) ===
2 func MyHash(key string, mod int) int {
3     a := asciiCat(key, 1, 1) + asciiCat(key, 5, 5)
4     b := asciiCat(key, 3, 4)
5     c := asciiCat(key, 5, 6)
6     h := float64(a)/517.0 + float64(b)/217.0 + float64(c)/256.0
7     if h < 0 {
8         h = -h
9     }
10    return int(h) % mod
11 }
```

Listing 1: MyHash implementation

1.3 Insertion and Collision Handling

The insertion logic starts at the hashed index. If the slot is free, it inserts the key into that slot. If the slot is full, it uses the prober to find the next available slot. In linear probing, it will check the start slot plus one. This process repeats until an empty slot is found, and the probe counter is incremented to show how many attempts it took.

```
1 func InsertKeyC(tableFile *os.File, key []byte, start int, prober
2     Prober) (int, int, error) {
3     prober.Init(start, int(TableSize))
4     idx := start
5     probeCount := 0
6
7     for probeCount < TableSize {
8         var slot Slot
9         if err := ReadSlot(tableFile, idx, &slot); err != nil {
10             return -1, probeCount, err
11         }
12
13         if slot.Filled == 0 {
14             copy(slot.Key[:], key)
15             slot.Orig = int32(start)
16             slot.Final = int32(idx)
17             slot.Probes = int32(probeCount + 1)
18             slot.Filled = 1
19             if err := WriteSlot(tableFile, idx, &slot); err != nil {
20                 return -1, probeCount, err
21             }
22         }
23     }
24 }
```

```

21         return idx, probeCount + 1, nil
22     }
23
24     idx = prober.Next()
25     probeCount++
26 }
27
28 return -1, probeCount, fmt.Errorf("table full")
29 }
```

Listing 2: Linear probing insertion for 75 keys

1.4 Query Testing and Probe Statistics

To evaluate the hash function I gathered the min,max, and average of first 25 keys and the last 25 keys. I also compared the average of all 75 keys to the theoretical average.

Table 1: Probe Statistics for Linear Probing

Hash	Probe	F25 (min/max/avg)	L25 (min/max/avg)	All75 (avg)	Theory (avg)
MyHash	Linear	1/22/5.65	1/27/3.08	3.92	2.50

1.5 Theoretical vs. Measured Analysis

Given:

$$n = 75, \quad m = 100, \quad \alpha = \frac{n}{m} = 0.75$$

Successful Search (Linear Probing):

$$E_s = \frac{1}{2} \left(1 + \frac{1}{1 - \alpha} \right)$$

Evaluation with $\alpha = 0.75$:

$$E_s = \frac{1}{2} \left(1 + \frac{1}{1 - 0.75} \right) = \frac{1}{2} \left(1 + \frac{1}{0.25} \right) = \frac{1}{2}(1 + 4) = \frac{5}{2} = 2.5$$

Comparison with Theory:

Table 2: Measured vs. Theoretical Probes (Linear Probing, Successful Search)

Search Set	Measured Avg Probes	Theoretical Avg Probes
First 25 keys	5.65	2.5
Last 25 keys	3.08	2.5
All 75 keys	3.92	2.5

1.5.1 Explanation

The measured probes were higher than the theoretical average. This shows that the hash function used did a poor job at spreading the 75 keys over 100 slots. The first 25 took almost double the amount of probes on average than the theoretical average, while the last 25 were closer to the theoretical average.

1.6 Table Dump Output

1.6.1 Complete Table Display (Slots 0–99)

	Hash address	Contents at address	Original hash	Final hash	Probes
1	0	G4iCoa3vqCmNxYpt	92	0	9
2	1	P6z6lJg68Qi4Qh5a	99	1	3
3	2	w6taYjy6apLID6VM	2	2	1
4	3	Xcrcg2aB1xVvowUx	93	3	11
5	4	VI5gBW2Du5KKCB3	94	4	11
6	5	GnNfs99CyXtt8ltD	5	5	1
7	6	O3oYAv9FHj3ntES3	6	6	1
8	7	rw25zuE6gQTcR9MQ	0	7	8
9	8	Rss0bMjTMdBCJRDi	92	8	17
10	9	URAq0d4i0tiCKJxf	88	9	22
11	10	STx0Ae50EGmfPiB1	10	10	1
12	11	M6rkIdzXL1v5wSXU	11	11	1
13	12	6PhrLOLEoSFUw12o	10	12	3
14	13	771YcuUELnP1gS1X	10	13	4
15	14	(empty)			
16	15	(empty)			
17	16	(empty)			
18	17	(empty)			
19	18	zTBmPxZPqwcj4bYJ	18	18	1
20	19	(empty)			
21	20	(empty)			
22	21	(empty)			
23	22	(empty)			
24	23	NodxCgQ10vtqkufT	23	23	1
25	24	MOGNdko4Wh8YNxFa	24	24	1
26	25	gGrsLvp1p9HefRNM	23	25	3
27	26	71ipe0g2kJPifs3R	24	26	3
28	27	(empty)			
29	28	(empty)			
30	29	(empty)			
31	30	(empty)			
32	31	(empty)			
33	32	vITHpRpZt6T2WFDk	32	32	1
34	33	(empty)			
35	34	(empty)			
36	35	(empty)			
37	36	(empty)			
38	37	Tmt2bzgUTrYh7L2R	37	37	1
39	38	v8MzHs6i7quhlbf	37	38	2
40	39	SP1rhCZhHiUw3eIy	39	39	1
41	40	(empty)			
42	41	(empty)			
43	42	RNczIjLhcT9nVZFb	42	42	1
44	43	rL0w62UqvaoJPCKP	43	43	1
45	44	AhZjGGhfRGWWeatr	43	44	2
46	45	uKqy42RdNMA5YnOm	42	45	4
47	46	(empty)			
48	47	(empty)			
49	48	(empty)			
50	49	(empty)			

54	50	(empty)			
55	51	B52h35qGLCwyUUFe	51	51	1
56	52	9o2d33dfNhkYEGd8	51	52	2
57	53	FLHP30tA4E8ZKqxd	53	53	1
58	54	R1PXjwk02vZpNLLU	52	54	3
59	55	X34HTtKZWPig4x09	53	55	3
60	56	5PMDQuWDBaXXMDs5	52	56	5
61	57	GZEicTpCdI5xndZJ	57	57	1
62	58	JaS41YZ1iYCfDAss	58	58	1
63	59	aoH1A69ayQtXkeks	59	59	1
64	60	6Mcx4fNznFLRNxVo	60	60	1
65	61	c1SC9LFErx0Ikdyh	61	61	1
66	62	Od9bTmGRYUBxZvWV	55	62	8
67	63	WT7SYUmqgyEV2JbU	61	63	3
68	64	PcHaDSFBYQA9eDwa	60	64	5
69	65	0V2DjJdgRcnW1xge	65	65	1
70	66	t8tpKaLT0yXZ6BtP	65	66	2
71	67	6FEWS3y9nd25BZg7	65	67	3
72	68	pf3tLUoKaJhoHkSE	65	68	4
73	69	p2M6T9vzfJuy5wP3	69	69	1
74	70	Ly8uLp270ZWQ3S1t	56	70	15
75	71	SjSVMCNUKJ9HaSeK	69	71	3
76	72	1z01LDvBMtP3RV7n	66	72	7
77	73	P13tCYqjSgW2gwNs	62	73	12
78	74	gJKD06MXSRW7STBJ	54	74	21
79	75	QMqSlhMrQ0fxTJ8M	75	75	1
80	76	oFupWNQhKXcMVZYV	74	76	3
81	77	7eEJiQNYk42EyCYP	73	77	5
82	78	f7NRjXKKkD7LqocM	78	78	1
83	79	pkocGBDOMU0elJMx	79	79	1
84	80	IWXDcG7iwFa06Kki	80	80	1
85	81	XOZGVmeodkS31VKb	78	81	4
86	82	K6N1rdL158Avqc81	82	82	1
87	83	TW4jpwQ1bQWx6KOA	78	83	6
88	84	OssRNRh0eMJVA2Jr	84	84	1
89	85	r58pD6asUZfqXMiZ	85	85	1
90	86	oX95VTkjrPk6lnXs	60	86	27
91	87	(empty)			
92	88	FqckjZ1Fmd0A7dbP	88	88	1
93	89	QRQ0Zsr7qv4mZuPw	89	89	1
94	90	1mwhjQD4093NIWs2	90	90	1
95	91	zER1ZfRtxewEgarL	90	91	2
96	92	BEcfZIGKOMGfyCav	92	92	1
97	93	S4iEWRIbYUOOXPEC	89	93	5
98	94	Pkc3Yo6dGHvsT6Tc	94	94	1
99	95	zVlsLvy4RCCTayyM	95	95	1
100	96	1Uv5Wp2c1eKyXzxz	95	96	2
101	97	xC7gnZz00wKqn3bI	97	97	1
102	98	sscXzYpa13LweUyP	94	98	5
103	99	JBxbnCbzRgHxtvaL	99	99	1

Listing 3: Full table dump for Option C

2 Option B — Memory Table

2.1 Implementation Overview

2.2 Array-Based Data Structure Design

2.2.1 Memory Layout and Slot Organization

2.2.2 Constraints (No Maps on Hashing Path)

2.3 Hash Functions

2.3.1 YourHash (Redesigned Hash)

```
1 func YourHash(key string, mod int) int {
2     h := fnv.New32a()
3     h.Write([]byte(key))
4     return int(h.Sum32() % uint32(mod))
5 }
```

Listing 4: YourHash implementation in memory

The YourHash That is provided improved upon MyHash by using FNV-1a over 16 byte keys. YourHash creates a sum with all 16 bytes that can then spread the 75 keys uniformly over the 100 slots with less collisions. With fewer collisions, the insertion algorithm performs fewer probe loops, which makes the overall insertion more efficient.

2.4 Insertion and Collision Handling

```
1 func InsertKeyB(memTable []MemSlot, key []byte, start int, prober
2     Prober) (int, int, error) {
3     prober.Init(start, M)
4     idx := start
5     probeCount := 0
6
7     for probeCount < M {
8         slot := &memTable[idx] // pointer to memTable
9
10        if !slot.Filled {
11            copy(slot.Key[:], key)
12            slot.Orig = start
13            slot.Final = idx
14            slot.Probes = probeCount + 1
15            slot.Filled = true
16            return idx, probeCount + 1, nil
17        }
18
19        idx = prober.Next()
20        probeCount++
21    }
22    return -1, probeCount, fmt.Errorf("table full")
```

Listing 5: Linear probing insertion for 75 keys in memory

2.5 Query Testing and Probe Statistics

Hash	Probe	F25 (min/max/avg)	L25 (min/max/avg)	All75 (avg)	Theory (avg)
YourHash	Linear	1/5/1.61	1/4/1.57	1.88	2.50

2.6 Table Dump Output

2.6.1 Complete Array Display (Indices 0–99)

1	==== Table Dump ====	2	Hash address	Contents at address	Original hash	Final hash	Probes
3							
4	0 GZEicTpCdI5xndZJ	0	0	0	0	1	
5	1 JBxbnCbzRgHxtvaL	0	0	1	1	2	
6	2 pf3tLUoKaJhoHkSE	2	2	2	2	1	
7	3 w6taYjy6apLID6VM	2	3	3	3	2	
8	4 (empty)						
9	5 (empty)						
10	6 (empty)						
11	7 G4iCoa3vqcMNXypt	7	7	7	7	1	
12	8 (empty)						
13	9 (empty)						
14	10 6Mcx4fNznFLRNxVo	10	10	10	10	1	
15	11 Rss0bMjTMdBCJRDi	11	11	11	11	1	
16	12 (empty)						
17	13 (empty)						
18	14 VIs5gBW2Du5KKCB3	14	14	14	14	1	
19	15 Pkc3Yo6dGHvsT6Tc	15	15	15	15	1	
20	16 B52h35qGLCwyUUfy	16	16	16	16	1	
21	17 1Uv5Wp2c1eKyXZxz	16	17	17	17	2	
22	18 03oYAv9FHj3ntES3	18	18	18	18	1	
23	19 771YcuUELnP1gS1X	17	19	19	19	3	
24	20 t8tpKa1T0yXZ6BtP	20	20	20	20	1	
25	21 NodxCgQ10vtqkufT	21	21	21	21	1	
26	22 zTBmPxZPqwcj4bYJ	20	22	22	22	3	
27	23 pkocGBDOMU0e1JMx	23	23	23	23	1	
28	24 SjSVMCNUKJ9HaSeK	20	24	24	24	5	
29	25 K6N1rdL158Avqc81	25	25	25	25	1	
30	26 Od9bTmGRYUBxZvWV	26	26	26	26	1	
31	27 sscXzYpa13LweUyP	25	27	27	27	3	
32	28 r58pD6asUZfqXMiZ	28	28	28	28	1	
33	29 RNczIjLhcT9nVZFb	27	29	29	29	3	
34	30 X34HTtKZpig4x09	30	30	30	30	1	
35	31 X0ZGVmeodkS31VKb	30	31	31	31	2	
36	32 WT7SYUmqgyEV2JbU	30	32	32	32	3	
37	33 IWXdCg7iwFa06Kki	33	33	33	33	1	
38	34 QRQ0Zsr7qv4mZuPw	27	34	34	34	8	
39	35 1z01LDvBMtP3RV7n	34	35	35	35	2	
40	36 5PMDQuWDBaXXMDs5	36	36	36	36	1	
41	37 MOGNdko4Wh8YNxFa	37	37	37	37	1	
42	38 AhZjGGhfRGWWeatr	37	38	38	38	2	
43	39 (empty)						
44	40 Tmt2bzgUTrYh7L2R	40	40	40	40	1	
45	41 TW4jpwQ1bQWx6K0A	40	41	41	41	2	
46	42 f7NRjXKkkD7LqocM	42	42	42	42	1	
47	43 OssRNRh0eMJVA2Jr	43	43	43	43	1	

48	44	c1SC9LFErx0Ikdyh	44	44	1
49	45	(empty)			
50	46	(empty)			
51	47	BEcfZIGK0MGfyCav	47	47	1
52	48	zVlsLvy4RCCTayyM	47	48	2
53	49	p2M6T9vzfJuy5wP3	48	49	2
54	50	rw25zuE6gQTcR9MQ	47	50	4
55	51	URAq0d4i0tiCKJxf	51	51	1
56	52	oX95VTkjrPk6lnXs	47	52	6
57	53	JaS41YZliYCfDAss	53	53	1
58	54	71ipe0g2kJPifs3R	53	54	2
59	55	QMqSlhMrQ0fxTJ8M	55	55	1
60	56	6PhrLOLeoSfuw12o	55	56	2
61	57	uKqy42RdNMA5Yn0m	55	57	3
62	58	gJKD06MXSRW7STBJ	56	58	3
63	59	xC7gnZz00wKqn3bI	59	59	1
64	60	PcHaDSFBYQA9eDwa	59	60	2
65	61	(empty)			
66	62	1mwhjQD4093NIWs2	62	62	1
67	63	zER1ZfRtxewEgarL	62	63	2
68	64	Ly8uLp270ZWQ3Slt	62	64	3
69	65	FqcKjZ1Fmd0A7dbP	62	65	4
70	66	v8MzHs6i7quhlbxr	65	66	2
71	67	7eEJiQNYk42EyCYP	62	67	6
72	68	(empty)			
73	69	(empty)			
74	70	(empty)			
75	71	aoH1A69ayQtXkeks	71	71	1
76	72	vITHpRpZt6T2WFDk	72	72	1
77	73	M6rkIdzXL1v5wSXU	73	73	1
78	74	9o2d33dfNhkYEGd8	73	74	2
79	75	gGrsLvp1p9HefRNM	72	75	4
80	76	oFupWNQhKXcMVZYV	73	76	4
81	77	STx0Ae50EGmfPiB1	77	77	1
82	78	(empty)			
83	79	(empty)			
84	80	(empty)			
85	81	R1PXjwk02vZpNLLU	81	81	1
86	82	SP1rhCZhHiUw3eIy	82	82	1
87	83	(empty)			
88	84	6FEWS3y9nd25BZg7	84	84	1
89	85	(empty)			
90	86	(empty)			
91	87	(empty)			
92	88	(empty)			
93	89	P6z6lJg68Qi4Qh5a	89	89	1
94	90	FLHP30tA4E8ZKqxd	90	90	1
95	91	GnNfs99CyXtt81tD	89	91	3
96	92	(empty)			
97	93	(empty)			
98	94	rL0w62UqvoaJPCKP	94	94	1
99	95	P13tCYqjSgW2gwNs	95	95	1
100	96	S4iEWribyU00XPEC	96	96	1
101	97	0V2DjJdgRcnW1xge	97	97	1
102	98	(empty)			
103	99	Xcrcg2aBIxVvowUx	99	99	1

Listing 6: Full memory table dump for Option B

2.7 Theoretical vs. Measured Comparison

Search Set	Measured Avg Probes C	Measured Avg Probes B	Theoretical
First 25 keys	5.65	1.61	2.5
Last 25 keys	3.08	1.57	2.5
All 75 keys	3.92	1.88	2.5

2.8 Conclusion

Together, Option B and Option C show the clear difference between the two hash functions. MyHash was intentionally poor, while YourHash was a better implementation. MyHash produced higher average probes compared to the theoretical average, while YourHash stayed below the theoretical average. Both options used the same table size, the same keys, and the same linear probing logic, but Option C's hash required roughly twice as many probes on average as the hash used in Option B.

A well-designed hash keeps probe lengths short and makes access patterns more efficient, while a poorly designed hash leads to clustering and an inefficient number of probes.