

Process algebras and network motifs 3

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- Introduction of reflective calculi
 - Syntax
 - Structural Equivalence
 - Semantics
 - A New Approach to Stochastics
 - Modeling in a reflective calculus

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Is working up to weak bisimulation sufficient?

If the algebraic notation is simply an alternative notation for the logic -- that's fine... but then we are still missing something, according to the proposition-as-types paradigm...

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```
class Reaction {
    Reagent[] _reagents;
    Reagent[] _resultants;
    float _basal_rate;
...

public Solution reduce(Solution s) {
    Solution ans = s.copy();
    for r in _reagents {
        ans.remove(r);
    }
    for r in _resultants {
        ans.add(r.copy());
    }
    return(ans);
}
```

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```
public Solution reduce(Solution s) {
    int i = random(10000);
    int p = nthDigitOfPi(i);
    Solution ans = s.copy();
    int i = random(10000);
    int p = nthDigitOfPi(i);
    for r in _reagents {
        int i = random(10000);
        int p = nthDigitOfPi(i);
        ans.remove(r);
        int i = random(10000);
        int p = nthDigitOfPi(i);
        int i = random(10000);
        int p = nthDigitOfPi(i);
    }
    ...
    for r in _resultants {
        ...
        ans.add(r.copy());
        ...
}
```

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- The two programs are weakly bisimilar
- Is the programmer who writes the first code including the second in her mental models of the first?
- Likewise is the biologist looking at a specific network including radical variants?
 - Under what conditions is occam's razor at work?
 - When might she be thinking of such classes of networks?

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Is working up to weak bisimulation sufficient?

No! A language of **individuals** -- or witnesses -- is still needed!

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Course check

That Michaelis-menten is a scheme or a type is obvious:

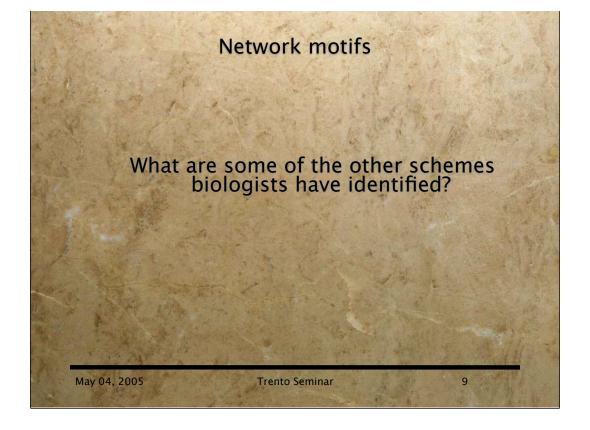
There are no elements or compounds mentioned in the equation

$$E + S \Leftrightarrow {}^{kl}_{k-1} ES \rightarrow {}^{k2} E + P$$

It is understood to be instantiated

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Why these motifs?

- Because they show up in physical systems?
- Because they are statistically overrepresented in physical networks?
- Because they are easy to analyze with existing tools?
- Because they cohere?
 - As a set of gadgets that may be combined they define some expressive class
 - Turing complete
 - CFG
 - PDA's
 - Regular languages

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Why these motifs?

- Analog between this question and this one:
 - Why these operators in our process calculus and not another set?
- Analog to an even more fundamental question:
 - Why this model of computation and not another?
- Must the answer come from features the domain?
 - Mobile process algebras are the only scale-invariant model of computation
 - Are there other invariants at work selecting network combinators?
 - How are these related to biologically relevant and realizable observations?

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Kinetic proofreading -- some biological processes

- Why kinetic proofreading?
 - dna replication
 - T-cell receptor signal transduction
 - dna damage detection
- Generalizes something we have already studied

$$E + S \Leftrightarrow {}^{kl}{}_{k-1} ES \to {}^{k2} E + P$$

$$E + S \Leftrightarrow {}^{up00}{}_{dn0} ES \to {}^{up01} ES^* \to {}^{up02} E + P$$

Generalizes to a family of schemes

$$E + S \Leftrightarrow up00 \atop dn0} ES \rightarrow up01 ES^* \rightarrow up02 E + P$$

$$E + S \iff up00 \longrightarrow ES \implies up00 \longrightarrow ES^* \implies up02 \longrightarrow ES^{**} \implies up03 \longrightarrow E + P$$

$$E \uparrow S \Longleftrightarrow^{up00} \underbrace{offl}_{offl} ES \longrightarrow^{up01} ES^{(1)} \longrightarrow^{up02} \cdots ES^{(n)} \longrightarrow^{up0n+1} E + P$$

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Kinetic proofreading -- some biological processes

- Why kinetic proofreading?
 - Because it provides another way to address 'why mobile process algebras...?'
 - The family of schemes may be described recursively via composition
 - Can we do this with ode's?

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Kinetic proofreading -- mass-action analysis

 $d[S]/dt = dn0([ES] + [ES^*]) - up00[E][S]$

 $d[ES]/dt = up00[E][S] - dn0[ES] - up01[ES^*]$

 $d[ES^*]/dt = up01[ES] - (dn0+up02)[ES^*]$

 $d[P]/dt = up02[ES^*]$

But this is not what justifies the term "proofreading"

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Kinetic proofreading -- the proofreading part

- Define Err as the ratio of formation of incorrect product to correct product
- In the MM case we calculate

 Err_{MM}

(up10/E](up12/up12+dn0))/up00/E](up01/up01 dn0)

(up02+dn0)/(up12+dn1)

In the KPR case we calculate

 $Err_{KPR} = (Err_{MM})^2$

• Why? What (in)equilibrium constraints must be enforced to have this work? What assumptions must be made regarding the relative magnitudes of dn0 and dn1?

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Kinetic proofreading -- course check

- What is the corresponding spatial logic formula?
- Can one formula capture the entire family of schemes?
- Can one set of equations?
- Let D(n) be the set of differential equations for n steps of delay, how does one capture D (n+1)?
- Let S(n) be the spatial logic formula for n steps of delay, how does one capture S(n+1)?

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