

Computing expectations by conditioning

Conditional PDF (continuous case)

- ❖ The conditional PDF of X given that $Y = y$

$$f_{X|Y}(x|y) = \frac{f(x, y)}{f_Y(y)}$$

$$E[X|Y = y] = \int_{-\infty}^{\infty} x f_{X|Y}(x|y) dx$$

- ❖ (Example 3.5) $f(x, y) = \begin{cases} 6xy(2 - x - y), & 0 < x < 1, 0 < y < 1 \\ 0, & \text{otherwise} \end{cases}$

- (Question) $E[X|Y = y] = ?$
- $f_Y(y) = \int_0^1 6xy(2 - x - y)dx = -6y \int_0^1 (x^2 + (y - 2)x)dx = -6y \left. \frac{x^3}{3} + \frac{(y-2)x^2}{2} \right|_0^1 = y(4 - 3y)$
- $f_{X|Y}(x|y) = \frac{f(x, y)}{f_Y(y)} = \frac{6xy(2 - x - y)}{y(4 - 3y)}$
- $E[X|Y = y] = \int_{-\infty}^{\infty} x f_{X|Y}(x|y) dx = \int_0^1 x \frac{6xy(2 - x - y)}{y(4 - 3y)} dx = \frac{6}{4 - 3y} \int_0^1 x^2 (2 - x - y) dx = \frac{5 - 4y}{8 - 6y}$

Conditional PDF – example

❖ (Example 3.7) $f(x, y) = \begin{cases} \frac{1}{2}ye^{-xy}, & 0 < x < \infty, 0 < y < 2 \\ 0, & \text{otherwise} \end{cases}$

▪ (Question) $E[e^{X/2}|Y = 1] = ?$

▪ (Solution) $f_Y(1) = \int_0^\infty \frac{1}{2}e^{-x}dx = \left. \frac{-e^{-x}}{2} \right|_0^\infty = \frac{1}{2}$

▪ $f_{X|Y}(x|1) = \frac{f(x,1)}{f_Y(1)} = \frac{\frac{1}{2}e^{-x}}{\frac{1}{2}} = e^{-x}$

▪ $E[e^{X/2}|Y = 1] = \int_0^\infty e^{\frac{x}{2}}f_{X|Y}(x|1)dx = \int_0^\infty e^{\frac{x}{2}}e^{-x}dx = -2 \int_0^\infty e^{-\frac{x}{2}}dx = 2$

Computing expectations by conditioning

❖ $E[X] = \sum_y E[X|Y = y]P[Y = y]$ (discrete case)

❖ $E[X] = \int_{-\infty}^{\infty} E[X|Y = y]f_Y(y) dy$

❖ (Example 3.9)

The number of misprints in Sam's probability chapter \sim Poisson (2), the number in his history chapter \sim Poisson (5), either book is equally likely chosen

❖ (Question) What is the expected number of misprints?

▪ (Solution) Let X denote the number of misprints and

$$Y = \begin{cases} 1, & \text{if his history book is chosen} \\ 2, & \text{if his probability book is chosen} \end{cases}$$

$$\begin{aligned} E[X] &= \sum_y E[X|Y = y]P[Y = y] = E[X|Y = 1]P[Y = 1] + E[X|Y = 2]P[Y = 2] \\ &= 5 * \frac{1}{2} + 2 * \frac{1}{2} = \frac{7}{2} \end{aligned}$$

Computing expectations by conditioning

- ❖ $E[X] = \sum_y E[X|Y = y]P[Y = y]$ (discrete case)
- ❖ $E[X] = \int_{-\infty}^{\infty} E[X|Y = y]f_Y(y) dy$
- ❖ (Example 3.11) a coin, having probability p of coming up heads. Let N denote the number of flips until the first head appears.
- ❖ (Question) $E[N] = ?$
 - (Solution)

Let $Y = \begin{cases} 1, & \text{if the first flip results in a head} \\ 0, & \text{if the first flip results in a tail} \end{cases}$

$$E[N] = \sum_y E[N|Y = y]P[Y = y] = E[N|Y = 0]P[Y = 0] + E[N|Y = 1]P[Y = 1] = pE[N|Y = 1] + (1 - p)E[N|Y = 0]$$

$$E[N|Y = 1] = 1, E[N|Y = 0] = 1 + E[N]$$

$$E[N] = p + (1 - p)(1 + E[N]) \rightarrow E[N] = \frac{1}{p}$$

Computing expectations by conditioning – example

- ❖ $E[X] = \sum_y E[X|Y = y]P[Y = y]$ (discrete case)
- ❖ $E[X] = \int_{-\infty}^{\infty} E[X|Y = y]f_Y(y) dy$
- ❖ (Example 3.12) a miner, having three doors. The first door \rightarrow safety after two hours of travel. The second \rightarrow the mine after three hours of travel. The third \rightarrow the mine after five hours of travel. The miner chooses any one of the doors equally likely. Let X denote the time until the miner reaches safety and Y the door he initially chooses.
- ❖ (Question) $E[X] = ?$
 - (Solution)
$$E[X] = \sum_y E[X|Y = y]P[Y = y] = E[X|Y = 1]P[Y = 1] + E[X|Y = 2]P[Y = 2] + E[X|Y = 3]P[Y = 3] = \frac{1}{3} * 2 + \frac{1}{3} (3 + E[X]) + \frac{1}{3} (5 + E[X])$$
$$\Leftrightarrow \frac{1}{3}E[X] = \frac{10}{3} \Leftrightarrow E[X] = 10$$

Example 3.15

❖ Let N_k denote the number of necessary trials to obtain k consecutive successes, and let $M_k = E[N_k]$. Each trial \sim Bernoulli (p)

- $N_k = N_{k-1} + A_{k-1,k}$ where $A_{k-1,k}$ is the number of additional trials needed to go from $k - 1$ successes to k in a row

- $M_k = M_{k-1} + E[A_{k-1,k}]$

- $E[A_{k-1,k}] = p * 1 + (1 - p)(1 + M_k)$

- Thus, $M_k = M_{k-1} + p + (1 - p)(1 + M_k)$

$$\Leftrightarrow pM_k = M_{k-1} + 1$$

$$\Leftrightarrow M_k = \frac{1}{p}M_{k-1} + \frac{1}{p}$$

- $N_1 \sim \text{Geometric}(p)$, so $M_1 = \frac{1}{p}$

- $M_2 = \frac{1}{p^2} + \frac{1}{p}, \dots, M_k = \frac{1}{p^k} + \frac{1}{p^{k-1}} + \dots + \frac{1}{p}$

Analyzing Quick-Sort Algorithm (Example 3.16)

Quick-Sort Algorithm

- ❖ Quick-Sort (left, right)
 - Step 1 – Choose a pivot value
 - Step 2 – Partition the array using the pivot value
 - Step 3 – Quick-Sort the left partition recursively
 - Step 4 – Quick-Sort the right partition recursively

- ❖ Pseudocode

Quick-Sort(left, right)

if right \leq left, return

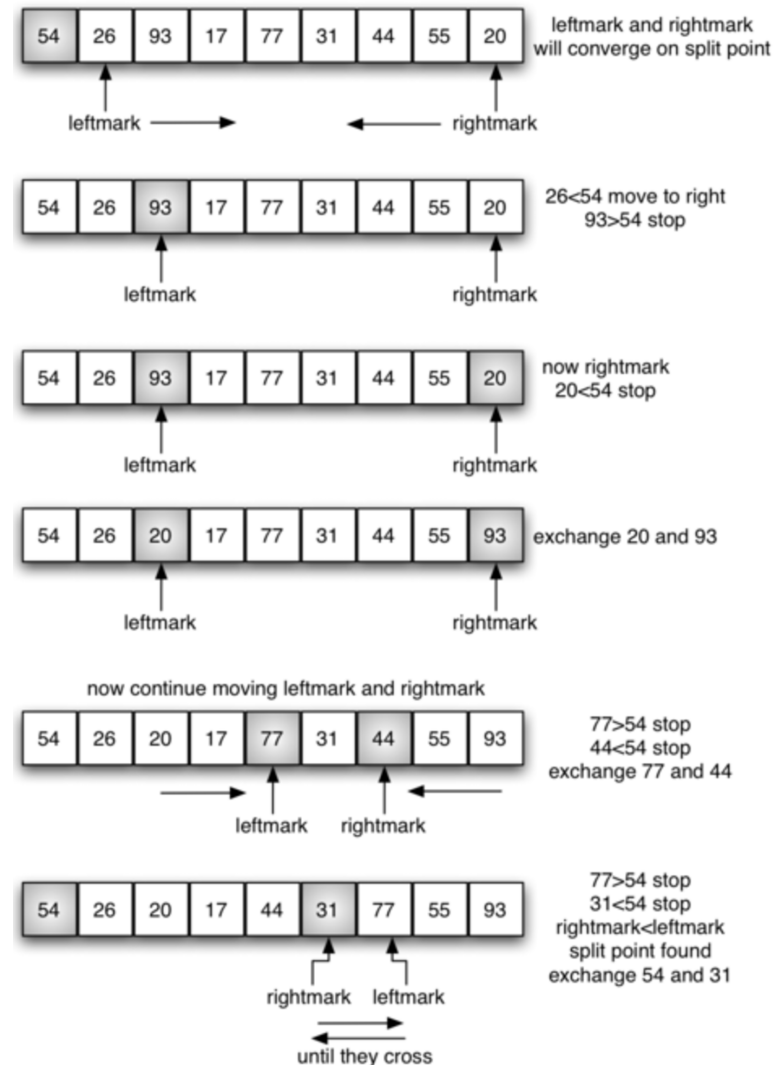
else

 pivot = A[left] // or choose at random

 partition = partitioning(left, right, pivot)

 Quick-Sort(left, partition-1)

 Quick-Sort(partition+1, right)




The expected number of comparisons

- ❖ $E[X] = \sum_y E[X|Y = y]P[Y = y]$ (discrete case)
- ❖ (Example 3.16) Analyzing the Quick-Sort Algorithm
 - Assume the pivot value chosen at random
 - The number of comparisons vary depending on the pivot value
 - Let M_n the expected number of comparisons needed by the Quick-Sort algorithm to sort a set of n distinct values
- $M_n = \sum_{j=1}^n E[\text{number of comparisons} | \text{value selected is } j\text{th smallest}] \frac{1}{n}$
- $\Leftrightarrow M_n = \sum_{j=1}^n (n-1 + M_{j-1} + M_{n-j}) \frac{1}{n}$
- $\Leftrightarrow M_n = n-1 + \frac{2}{n} \sum_{k=1}^{n-1} M_k$
- $\Leftrightarrow nM_n = n(n-1) + 2 \sum_{k=1}^{n-1} M_k$

The expected number of comparisons

- $M_n = \sum_{j=1}^n E[\text{number of comparisons} | \text{value selected is } j\text{th smallest}] \frac{1}{n}$
 - $\Leftrightarrow nM_n = n(n-1) + 2 \sum_{k=1}^{n-1} M_k \dots (1)$
 - $\Leftrightarrow (n+1)M_{n+1} = (n+1)n + 2 \sum_{k=1}^n M_k \dots (2)$
 - $(n+1)M_{n+1} - nM_n = 2n + 2M_n \dots \text{by } (2) - (1)$
 - $\Leftrightarrow (n+1)M_{n+1} = (n+2)M_n + 2n$
 - $\Leftrightarrow \frac{M_{n+1}}{n+2} = \frac{M_n}{n+1} + \frac{2n}{(n+1)(n+2)}$
 - $\Leftrightarrow \frac{M_{n+1}}{n+2} = \frac{M_{n-1}}{n} + \frac{2(n-1)}{n(n+1)} + \frac{2n}{(n+1)(n+2)} = \dots = 2 \sum_{k=0}^{n-1} \frac{n-k}{(n+1-k)(n+2-k)}$
 - $\Leftrightarrow M_{n+1} = 2(n+2) \sum_{k=0}^{n-1} \frac{n-k}{(n+1-k)(n+2-k)} = 2(n+2) \sum_{i=1}^n \frac{i}{(i+1)(i+2)}$
 - $\Leftrightarrow M_{n+1} = 2(n+2) \sum_{i=1}^n \left[\frac{2}{i+2} - \frac{1}{i+1} \right]$



Let $n-k = i$
- $$\sim 2(n+2) \left[\int_3^{n+2} \frac{2}{x} dx - \int_2^{n+1} \frac{1}{x} dx \right] = 2(n+2)[2 \log(n+2) - \log(n+1) + \log 2 - 2 \log 3]$$
- $$= 2(n+2)[\log(n+2) + \log \frac{n+2}{n+1} + \log 2 - 2 \log 3] \sim 2(n+2) \log(n+2)$$