

System Programming

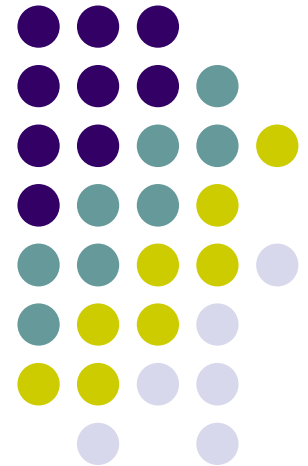
03. C. Memory II (ch 2.1)

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Compiler Complaint



- <http://xkcd.com/371/>



Roadmap



C:

```
car *c = malloc(sizeof(car));  
c->miles = 100;  
c->gals = 17;  
float mpg = get_mpg(c);  
free(c);
```

Java:

```
Car c = new Car();  
c.setMiles(100);  
c.setGals(17);  
mpg =  
    c.getMPG();
```

Memory & data

Integers & floats
x86 assembly
Procedures & stacks
Executables
Arrays & structs
Memory & caches
Processes
Virtual memory
Memory allocation
Java vs. C

Assembly language:

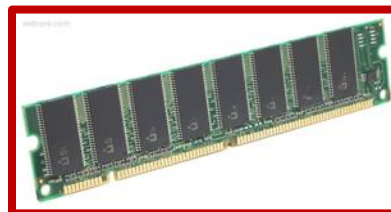
```
get_mpg:  
    pushq    %rbp  
    movq     %rsp, %rbp  
    ...  
    popq     %rbp  
    ret
```

Machine code:

```
0111010000011000  
100011010000010000000010  
1000100111000010  
110000011111101000011111
```

Computer system:

OS:



Memory, Data, and Addressing



- Representing information as bits and bytes
- Organizing and addressing data in memory
- Manipulating data in memory using C
- Boolean algebra and bit-level manipulations

Addresses and Pointers in C



- `&` = “address of” operator
- `*` = “value at address” or “dereference” operator

```
int* ptr;
```

Declares a variable, `ptr`, that is a pointer to (i.e. holds the address of) an `int` in memory

```
int x = 5;
```

```
int y = 2;
```

Declares two variables, `x` and `y`, that hold `ints`, and sets them to 5 and 2, respectively

```
ptr = &x;
```

Sets `ptr` to the address of `x` (“`ptr` points to `x`”)

```
y = 1 + *ptr;
```

“Dereference `ptr`”

Sets `y` to “1 plus the value stored at the address held by `ptr`. Because `ptr` points to `x`, this is equivalent to `y=1+x`;

What is `* (&y)` ?

Assignment in C (1)



- A variable is represented by a memory location
- Declaration \neq initialization (initially holds “garbage”)
- `int x, y;`
 - `x` is at address `0x04`, `y` is at `0x18`

0x00	0x01	0x02	0x03	
A7	00	32	00	0x00
00	01	29	F3	0x04 ✗
EE	EE	EE	EE	0x08
FA	CE	CA	FE	0x0C
26	00	00	00	0x10
00	00	10	00	0x14
01	00	00	00	0x18 y
FF	00	F4	96	0x1C
DE	AD	BE	EF	0x20
00	00	00	00	0x24

Assignment in C (2)

32-bit example
(pointers are
32-bits wide)



- A variable is represented by a memory location
- Declaration \neq initialization (initially holds “garbage”)
- `int x, y;`
 - `x` is at address `0x04`, `y` is at `0x18`

0x00	0x01	0x02	0x03	
				0x00
00	01	29	F3	0x04 ×
				0x08
				0x0C
				0x10
				0x14
01	00	00	00	0x18 y
				0x1C
				0x20
				0x24

Assignment in C (3)

32-bit example
(pointers are
32-bits wide)



- left-hand side = right-hand side;
 - LHS must evaluate to a memory *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location

● `int x, y;`

● `x = 0;`

0x00	0x01	0x02	0x03	
				0x00
00	00	00	00	0x04 x
				0x08
				0x0C
				0x10
				0x14
01	00	00	00	0x18 y
				0x1C
				0x20
				0x24

Assignment in C (3)

32-bit example
(pointers are
32-bits wide)



- left-hand side = right-hand side;
 - LHS must evaluate to a memory *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location

● `int x, y;`

● `x = 0;`

● `y = 0x3CD02700;`

little endian!

0x00	0x01	0x02	0x03	
				0x00
00	00	00	00	0x04 ×
				0x08
				0x0C
				0x10
				0x14
00	27	D0	3C	0x18 y
				0x1C
				0x20
				0x24

Assignment in C (4)

32-bit example
(pointers are
32-bits wide)



- left-hand side = right-hand side;
 - LHS must evaluate to a memory *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location

● `int x, y;`

● `x = 0;`

● `y = 0x3CD02700;`

● `x = y + 3;`

- Get value at `y`, add 3, store in

0x00	0x01	0x02	0x03	
				0x00
03	27	D0	3C	0x04 x
				0x08
				0x0C
				0x10
				0x14
00	27	D0	3C	0x18 y
				0x1C
				0x20
				0x24

Assignment in C (5)

32-bit example
(pointers are
32-bits wide)



& = “address of”

* = “dereference”

- left-hand side = right-hand side;
 - LHS must evaluate to a memory *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location

- `int x, y;`

- `x = 0;`

- `y = 0x3CD02700;`

- `x = y + 3;`

- Get value at `y`, add 3, store in `x`

- `int* z;`

- `z` is at address `0x20`

0x00	0x01	0x02	0x03	
				0x00
03	27	D0	3C	0x04 x
				0x08
				0x0C
				0x10
				0x14
00	27	D0	3C	0x18 y
				0x1C
DE	AD	BE	EF	0x20 z
				0x24

Declaration: initially contains garbage

Assignment in C (6)

32-bit example
(pointers are
32-bits wide)



& = “address of”

* = “dereference”

- left-hand side = right-hand side;
 - LHS must evaluate to a memory *location*
 - RHS must evaluate to a *value* (could be an address)
 - Store RHS value at LHS location

- `int x, y;`

- `x = 0;`

- `y = 0x3CD02700;`

- `x = y + 3;`

- Get value at `y`, add 3, store in `x`

- `int* z = &y + 3;`

- Get address of `y`, “add 3”, store in `z`

0x00	0x01	0x02	0x03	
				0x00
03	27	D0	3C	0x04 x
				0x08
				0x0C
				0x10
				0x14
00	27	D0	3C	0x18 y
				0x1C
24	00	00	00	0x20 z
				0x24

Pointer arithmetic

Pointer Arithmetic



- Pointer arithmetic is scaled by the size of target type
 - In this example, `sizeof(int) = 4`
- `int* z = &y + 3;`
 - Get address of `y`, add `3* sizeof(int)`, store in `z`
 - `&y = 0x18`
 - `24 + 3 * (4) = 36`
- **Pointer arithmetic can be dangerous!**
 - Can easily lead to bad memory accesses
 - Be careful with data types and *casting*

Assignment in C (7)

32-bit example
(pointers are
32-bits wide)



& = “address of”

* = “dereference”

- `int x, y;`
- `x = 0;`
- `y = 0x3CD02700;`
- `x = y + 3;`
 - Get value at `y`, add 3, store in `x`
- `int* z = &y + 3;`
 - Get address of `y`, add **12**, store in `z`
- `*z = y;`
 - What does this do?

0x00	0x01	0x02	0x03	
				0x00
03	27	D0	3C	0x04 x
				0x08
				0x0C
				0x10
				0x14
00	27	D0	3C	0x18 y
				0x1C
24	00	00	00	0x20 z
				0x24

Assignment in C (8)

32-bit example
(pointers are
32-bits wide)



& = “address of”
* = “dereference”

- `int x, y;`
- `x = 0;`
- `y = 0x3CD02700;`
- `x = y + 3;`
 - Get value at `y`, add 3, store in `x`
- `int* z = &y + 3;`
 - Get address of `y`, add **12**, store in `z`

The target of a pointer is
also a memory location

- `*z = y;`
 - Get value of `y`, put in address
stored in `z`

0x00	0x01	0x02	0x03	
				0x00
03	27	D0	3C	0x04 x
				0x08
				0x0C
				0x10
				0x14
00	27	D0	3C	0x18 y
				0x1C
24	00	00	00	0x20 z
00	27	D0	3C	0x24

Arrays in C (1)

Arrays are adjacent locations in memory storing the same type of data object

`a` is a name for the array's address

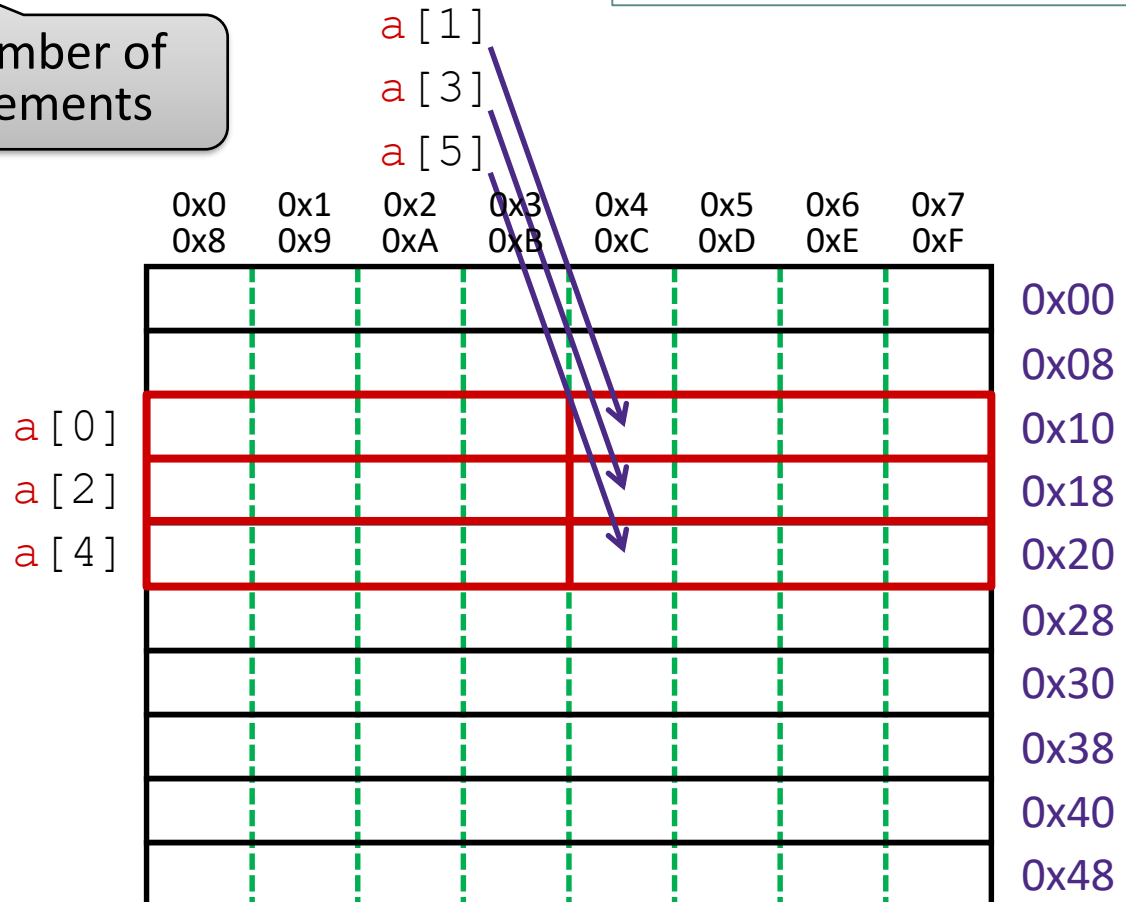
Declaration: `int a[6];`

element type

name

number of elements

64-bit example
(pointers are 64-bits wide)



Arrays in C (2)

Arrays are adjacent locations in memory storing the same type of data object

`a` is a name for the array's address

The address of `a[i]` is the address of `a[0]` plus `i` times the element size in bytes

Declaration: `int a[6];`

Indexing: `a[0]` = `0x015f`;
`a[5]` = `a[0]`;

	0x0 0x8	0x1 0x9	0x2 0xA	0x3 0xB	0x4 0xC	0x5 0xD	0x6 0xE	0x7 0xF	
									0x00
									0x08
<code>a[0]</code>	5F	01	00	00					0x10
<code>a[2]</code>									0x18
<code>a[4]</code>					5F	01	00	00	0x20
									0x28
									0x30
									0x38
									0x40
									0x48

Arrays in C (3)

Arrays are adjacent locations in memory storing the same type of data object

`a` is a name for the array's address

The address of `a[i]` is the address of `a[0]` plus `i` times the element size in bytes

Declaration: `int a[6];`

Indexing: `a[0]` = `0x015f`;
`a[5]` = `a[0]`;

No bounds checking: `a[6]` = `0xBAD`;
`a[-1]` = `0xBAD`;

	0x0 0x8	0x1 0x9	0x2 0xA	0x3 0xB	0x4 0xC	0x5 0xD	0x6 0xE	0x7 0xF	
									0x00
					AD	0B	00	00	0x08
<code>a[0]</code>	5F	01	00	00					0x10
<code>a[2]</code>									0x18
<code>a[4]</code>					5F	01	00	00	0x20
	AD	0B	00	00					0x28
									0x30
									0x38
									0x40
									0x48

Arrays in C (4)

Arrays are adjacent locations in memory storing the same type of data object

`a` is a name for the array's address

The address of `a[i]` is the address of `a[0]` plus `i` times the element size in bytes

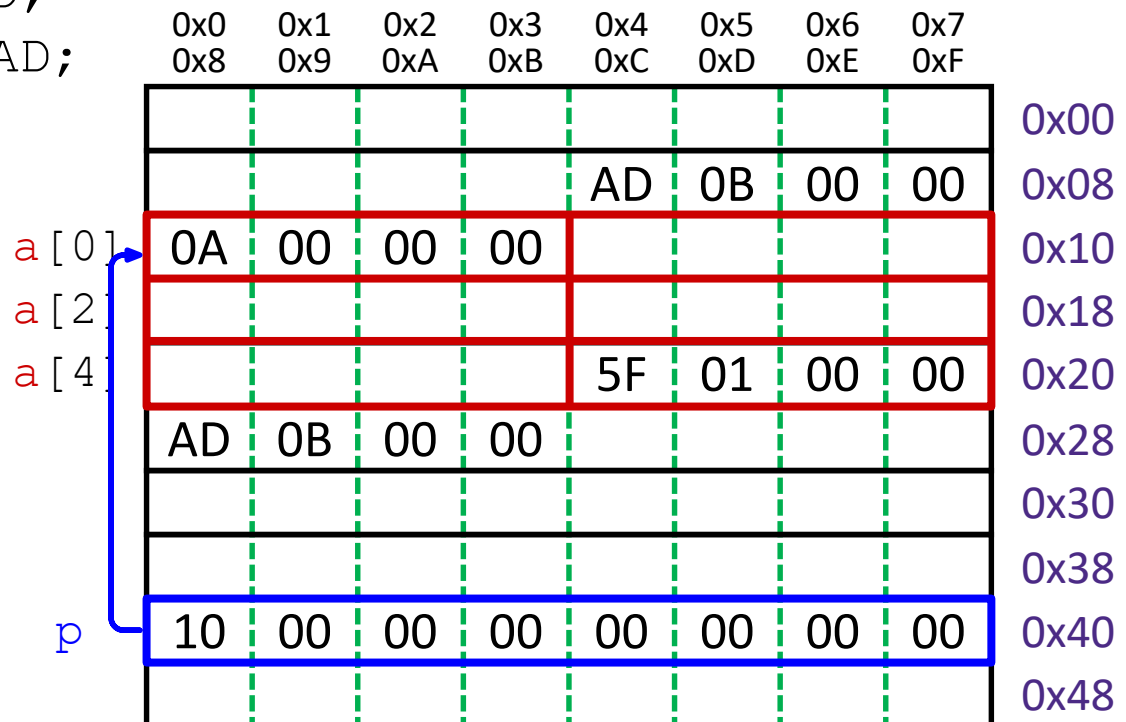
Declaration: `int a[6];`

Indexing: `a[0]` = 0x015f;
`a[5]` = `a[0]`;

No bounds checking: `a[6]` = 0xBAD;
`a[-1]` = 0xBAD;

Pointers: `int* p;`

equivalent $\left\{ \begin{array}{l} p = a; \\ p = \&a[0]; \\ *p = 0xA; \end{array} \right.$



Arrays in C (5)

Arrays are adjacent locations in memory storing the same type of data object

`a` is a name for the array's address

The address of `a[i]` is the address of `a[0]` plus `i` times the element size in bytes

Declaration: `int a[6];`

Indexing: `a[0] = 0x015f;`
`a[5] = a[0];`

No bounds checking: `a[6] = 0xBAD;`
`a[-1] = 0xBAD;`

Pointers: `int* p;`
equivalent $\begin{cases} p = a; \\ p = \&a[0]; \\ *p = 0xA; \end{cases}$

array indexing = address arithmetic
(both scaled by the size of the type)

equivalent $\begin{cases} p[1] = 0xB; \\ *(p+1) = 0xB; \\ p = p + 2; \end{cases}$

	0x0 0x8	0x1 0x9	0x2 0xA	0x3 0xB	0x4 0xC	0x5 0xD	0x6 0xE	0x7 0xF	
					AD	0B	00	00	0x08
<code>a[0]</code>	0A	00	00	00	0B	00	00	00	0x10
<code>a[2]</code>									0x18
<code>a[4]</code>					5F	01	00	00	0x20
	AD	0B	00	00					0x28
									0x30
									0x38
<code>p</code>	10	00	00	00	00	00	00	00	0x40
									0x48

Arrays in C (6)

Arrays are adjacent locations in memory storing the same type of data object

`a` is a name for the array's address

The address of `a[i]` is the address of `a[0]` plus `i` times the element size in bytes

Declaration: `int a[6];`

Indexing: `a[0] = 0x015f;`
`a[5] = a[0];`

No bounds checking: `a[6] = 0xBAD;`
`a[-1] = 0xBAD;`

Pointers: `int* p;`

equivalent $\begin{cases} p = a; \\ p = \&a[0]; \\ *p = 0xA; \end{cases}$

array indexing = address arithmetic
 (both scaled by the size of the type)

equivalent $\begin{cases} p[1] = 0xB; \\ *(p+1) = 0xB; \\ p = p + 2; \end{cases}$

`*p = a[1] + 1;`

	0x0 0x8	0x1 0x9	0x2 0xA	0x3 0xB	0x4 0xC	0x5 0xD	0x6 0xE	0x7 0xF	
					AD	0B	00	00	0x00
<code>a[0]</code>	0A	00	00	00	0B	00	00	00	0x08
<code>a[2]</code>	0C	00	00	00					0x10
<code>a[4]</code>					5F	01	00	00	0x18
	AD	0B	00	00					0x20
									0x28
									0x30
									0x38
<code>p</code>	18	00	00	00	00	00	00	00	0x40
									0x48

Fill in the blanks!

Arrays are adjacent locations in memory storing the same type of data object

`a` is a name for the array's address

The address of `a[i]` is the address of `a[0]` plus `i` times the element size in bytes

Declaration: `int a[6];`

Indexing: `a[0] = 0x015f;`
`a[5] = a[0];`

No bounds checking: `a[6] = 0xBAD;`
`a[-1] = 0xBAD;`

Pointers: `int* p;`

equivalent $\left\{ \begin{array}{l} p = a; \\ p = \&a[0]; \\ *p = 0xA; \end{array} \right.$ `a[0]`
`a[2]`
`a[4]`

array indexing = address arithmetic
(both scaled by the size of the type)

equivalent $\left\{ \begin{array}{l} p[1] = 0xB; \\ *(p+1) = 0xB; \\ p = p + 2; \end{array} \right.$

`*p = a[1] + 1;`

0x0 0x8	0x1 0x9	0x2 0xA	0x3 0xB	0x4 0xC	0x5 0xD	0x6 0xE	0x7 0xF	
								0x00
								0x08
								0x10
								0x18
								0x20
								0x28
								0x30
								0x38
								0x40
								0x48

Question: The variable values after Line 3 executes are shown on the right. What are they after Line 4 & 5?



```

1  void main() {
2      int a[] = {5, 10};
3      int *p = a;
4      p = p + 1;
5      *p = *p + 1;
6  }
    
```

	Data (decimal)	Address (decimal)
a[0]	5	100
a[1]	10	
	⋮	
p	100	

	p	*p	a[0]	a[1]	then	p	*p	a[0]	a[1]
(A)	101	10	5	10		101	11	5	11
(B)	104	10	5	10		104	11	5	11
(C)	100	6	6	10		101	6	6	10
(D)	100	6	6	10		104	6	6	10

Representing strings



- C-style string stored as an array of bytes (`char *`)
 - Elements are one-byte **ASCII codes** for each character
 - No “String” keyword, unlike Java

32	space	48	0	64	@	80	P	96	`	112	p
33	!	49	1	65	A	81	Q	97	a	113	q
34	"	50	2	66	B	82	R	98	b	114	r
35	#	51	3	67	C	83	S	99	c	115	s
36	\$	52	4	68	D	84	T	100	d	116	t
37	%	53	5	69	E	85	U	101	e	117	u
38	&	54	6	70	F	86	V	102	f	118	v
39	'	55	7	71	G	87	W	103	g	119	w
40	(56	8	72	H	88	X	104	h	120	x
41)	57	9	73	I	89	Y	105	i	121	y
42	*	58	:	74	J	90	Z	106	j	122	z
43	+	59	;	75	K	91	[107	k	123	{
44	,	60	<	76	L	92	\	108	l	124	
45	-	61	=	77	M	93]	109	m	125	}
46	.	62	>	78	N	94	^	110	n	126	~
47	/	63	?	79	O	95	_	111	o	127	del

ASCII: American Standard Code for Information Interchange

Null-terminated Strings



- **Example:** “Life is good” stored as a 13-byte array

<i>Decimal:</i>	76	105	102	101	32	105	115	32	103	111	111	100	0
<i>Hex:</i>	0x4c	0x69	0x66	0x65	0x20	0x69	0x73	0x20	0x67	0x6f	0x6f	0x64	0x00
<i>Text:</i>	L	i	f	e		i	s		g	o	o	d	\0

- Last character followed by a 0 byte (`'\0'`)
(a.k.a. “**null terminator**”)
 - Must take into account when allocating space in memory
 - Note that `'0' ≠ '\0'` (i.e. character 0 has non-zero value)
- How do we compute the length of a string?
 - Traverse array until null terminator encountered

Endianness and Strings

C (char = 1 byte)

```
char s[6] = "12345";
```

String literal

0x31 = 49 decimal = ASCII '1'

IA32, x86-64
(little endian)

SPARC
(big endian)

0x00	31	↔	31	0x00	'1'
0x01	32	↔	32	0x01	'2'
0x02	33	↔	33	0x02	'3'
0x03	34	↔	34	0x03	'4'
0x04	35	↔	35	0x04	'5'
0x05	00	↔	00	0x05	'\0'

- Byte ordering (endianness) is not an issue for 1-byte values
 - The whole array does not constitute a single value
 - Individual elements are values; chars are single bytes
- Unicode characters – up to 4 bytes/character
 - ASCII codes still work (just add leading zeros)
 - Unicode can support the many characters in all languages in the world
 - Java and C have libraries for Unicode (Java commonly uses 2 bytes/char)

Examining Data Representations



- Code to print byte representation of data
 - Any data type can be treated as a *byte array* by **casting** it to `char`
 - C has **unchecked** casts **!! DANGER !!**

```
void show_bytes(char* start, int len) {  
    int i;  
    for (i = 0; i < len; i++)  
        printf("%p\t0x%.2x\n", start+i, *(start+i));  
    printf("\n");  
}
```

printf directives:

<code>%p</code>	Print pointer
<code>\t</code>	Tab
<code>%x</code>	Print value as hex
<code>\n</code>	New line

Examining Data Representations



- Code to print byte representation of data
 - Any data type can be treated as a *byte array* by **casting** it to `char`
 - C has **unchecked** casts **!! DANGER !!**

```
void show_bytes(char* start, int len) {  
    int i;  
    for (i = 0; i < len; i++)  
        printf("%p\t0x%.2x\n", start+i, *(start+i));  
    printf("\n");  
}
```

```
void show_int(int x) {  
    show_bytes( (char *) &x, sizeof(int));  
}
```

show_bytes Execution Example



```
int a = 12345; // 0x00003039
printf("int a = 12345;\n");
show_int(a); // show_bytes((char *) &a, sizeof(int));
```

- Result (Linux x86-64):

- **Note:** The addresses will change on each run (try it!), but fall in same general range

int a = 12345;	
0x7fffb7f71dbc	0x39
0x7fffb7f71dbd	0x30
0x7fffb7f71dbe	0x00
0x7fffb7f71dbf	0x00

Summary



- Assignment in C results in value being put in memory location
- Pointer is a C representation of a data address
 - $\&$ = “address of” operator
 - $*$ = “value at address” or “dereference” operator
- Pointer arithmetic scales by size of target type
 - Convenient when accessing array-like structures in memory
 - Be careful when using – particularly when *casting* variables
- Arrays are adjacent locations in memory storing the same type of data object
 - Strings are null-terminated arrays of characters (ASCII)

Memory, Data, and Addressing



- Representing information as bits and bytes
- Organizing and addressing data in memory
- Manipulating data in memory using C
- **Boolean algebra and bit-level manipulations**



Boolean Algebra

- Developed by George Boole in 19th Century
 - Algebraic representation of logic (True \rightarrow 1, False \rightarrow 0)
 - AND: $A \& B = 1$ when both A is 1 and B is 1
 - OR: $A | B = 1$ when either A is 1 or B is 1
 - XOR: $A \wedge B = 1$ when either A is 1 or B is 1, but not both
 - NOT: $\sim A = 1$ when A is 0 and vice-versa
 - DeMorgan's Law: $\sim (A | B) = \sim A \& \sim B$
 $\sim (A \& B) = \sim A | \sim B$

AND		
$\&$	0	1
0	0	0
1	0	1

OR		
	0	1
0	0	1
1	1	1

XOR		
\wedge	0	1
0	0	1
1	1	0

NOT	
\sim	
0	1
1	0

General Boolean Algebras



- Operate on bit vectors
 - Operations applied bitwise
 - All of the properties of Boolean algebra apply

01101001	01101001	01101001	
& 01010101	01010101	^ 01010101	~ 01010101

- Examples of useful operations:

$$x \wedge x = 0$$

01010101
^ 01010101

00000000

$$x | 1 = 1, \quad x | 0 = x$$

01010101
11110000

11110101

Representing & Manipulating Sets



- Representation

- A w -bit vector represents subsets of $\{0, \dots, w-1\}$
- $a_j = 1$ iff $j \in A$

01101001 { 0, 3, 5, 6 }
76543210

01010101 { 0, 2, 4, 6 }
76543210

- Operations

- & Intersection 01000001 { 0, 6 }
- | Union 01111101 { 0, 2, 3, 4, 5, 6 }
- ^ Symmetric difference 00111100 { 2, 3, 4, 5 }
- ~ Complement 10101010 { 1, 3, 5, 7 }

Bit-Level Operations in C



- $\&$ (AND), $|$ (OR), \wedge (XOR), \sim (NOT)
 - View arguments as bit vectors, apply operations bitwise
 - Apply to any “integral” data type
 - long, int, short, char, unsigned

- Examples with char a, b, c;

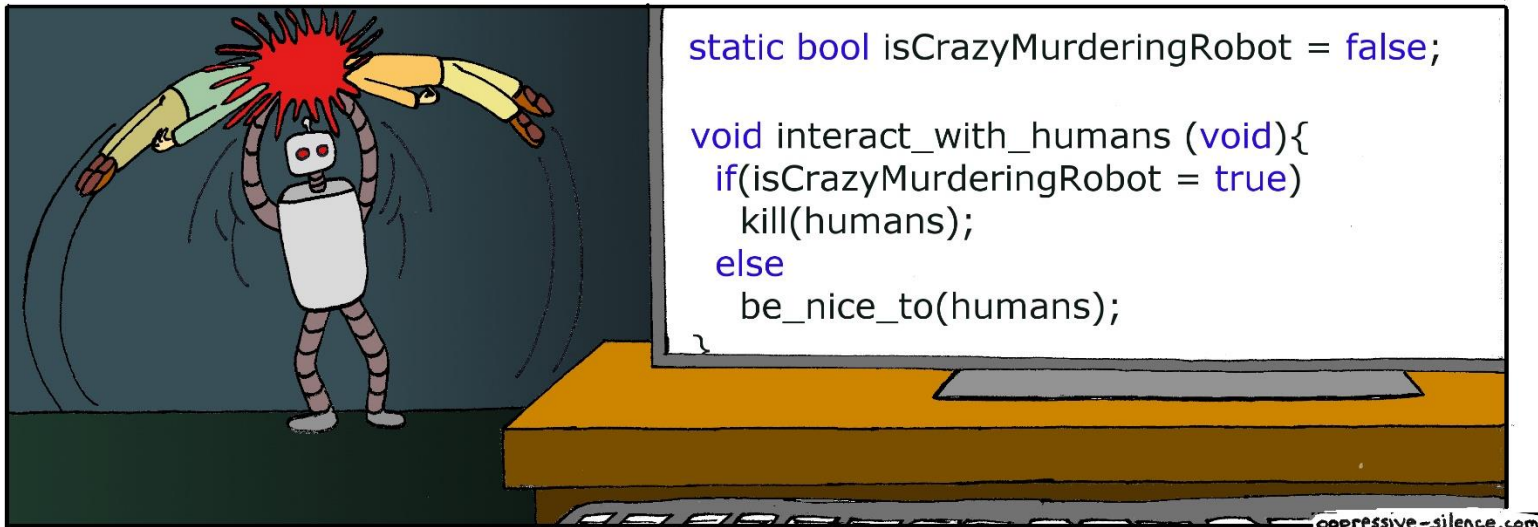
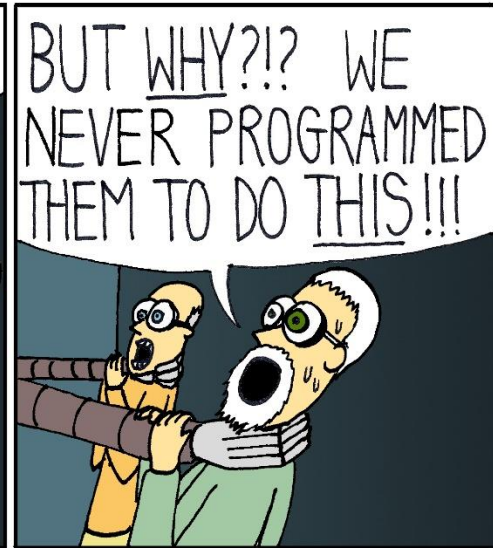
- ```
a = (char) 0x41; // 0x41->0b 0100 0001
b = ~a; // 0b ->0x
```
- ```
a = (char) 0x69;      // 0x69->0b 0110 1001
b = (char) 0x55;      // 0x55->0b 0101 0101
c = a & b;            //          0b          ->0x
```
- ```
a = (char) 0x41; // 0x41->0b 0100 0001
b = a; // 0b 0100 0001
c = a ^ b; // 0b ->0x
```

# Contrast: Logic Operations in C



- Logical operators in C: `&&` (AND), `||` (OR), `!` (NOT)
  - 0 is False, anything nonzero is True
  - Always return 0 or 1
  - **Early termination** (a.k.a. short-circuit evaluation) of `&&`, `||`
- Examples (`char` data type)
  - `!0x41 -> 0x00`
  - `!0x00 -> 0x01`
  - `!!0x41 -> 0x01`
  - `p && *p++`
    - Avoids **null pointer** (0x0) access via *early termination*
    - Short for: `if (p) { *p++; }`

# Logic Operations!!!



# Q&A

