

# System Programming

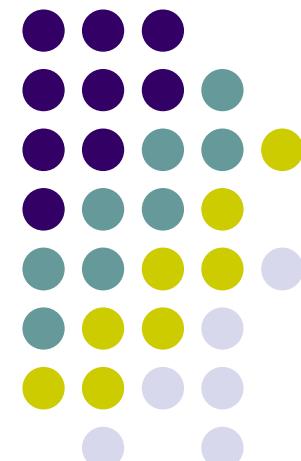
## 03. C. Memory II (ch 2.1)

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Data Science Lab @ PNU



**data**lab

data science laboratory

# Compiler Complaint



- <http://xkcd.com/371/>





# Roadmap

C:

```
car *c = malloc(sizeof(car));  
c->miles = 100;  
c->gals = 17;  
float mpg = get_mpg(c);  
free(c);
```

Java:

```
Car c = new Car();  
c.setMiles(100);  
c.setGals(17);  
mpg =  
    c.getMPG();
```

Assembly language:

```
get_mpg:  
    pushq  %rbp  
    movq   %rsp, %rbp  
    ...  
    popq  %rbp  
    ret
```

Machine code:

```
0111010000011000  
10001101000010000000010  
1000100111000010  
11000001111101000011111
```

Computer system:



OS:



Memory & data

Integers & floats

x86 assembly

Procedures & stacks

Executables

Arrays & structs

Memory & caches

Processes

Virtual memory

Memory allocation

Java vs. C

# Memory, Data, and Addressing



- Representing information as bits and bytes
- Organizing and addressing data in memory
- Manipulating data in memory using C
- Boolean algebra and bit-level manipulations

# Addresses and Pointers in C



- `&` = “address of” operator
- `*` = “value at address” or “dereference” operator

```
int* ptr;
```

Declares a variable, `ptr`, that is a pointer to (i.e. holds the address of) an `int` in memory

```
int x = 5;
```

```
int y = 2;
```

Declares two variables, `x` and `y`, that hold `ints`, and sets them to 5 and 2, respectively

```
ptr = &x;
```

Sets `ptr` to the address of `x` (“`ptr` points to `x`”)

```
y = 1 + *ptr;
```

“Dereference `ptr`”

Sets `y` to “1 plus the value stored at the address held by `ptr`. Because `ptr` points to `x`, this is equivalent to `y=1+x`;

What is `*(&y)` ?



# Assignment in C (1)

- A variable is represented by a memory location
- Declaration ≠ initialization (initially holds “garbage”)
- `int x, y;`
  - `x` is at address `0x04`, `y` is at `0x18`

0x00	0x01	0x02	0x03	0x00	0x04	0x08	0x0C	0x10	0x14	0x18	0x1C	0x20	0x24
A7	00	32	00										
00	01	29	F3										
EE	EE	EE	EE										
FA	CE	CA	FE										
26	00	00	00										
00	00	10	00										
01	00	00	00										
FF	00	F4	96										
DE	AD	BE	EF										
00	00	00	00										

X

y

# Assignment in C (2)

32-bit example  
(pointers are  
32-bits wide)



- A variable is represented by a memory location
- Declaration ≠ initialization (initially holds “garbage”)
- `int x, y;`
  - `x` is at address `0x04`, `y` is at `0x18`

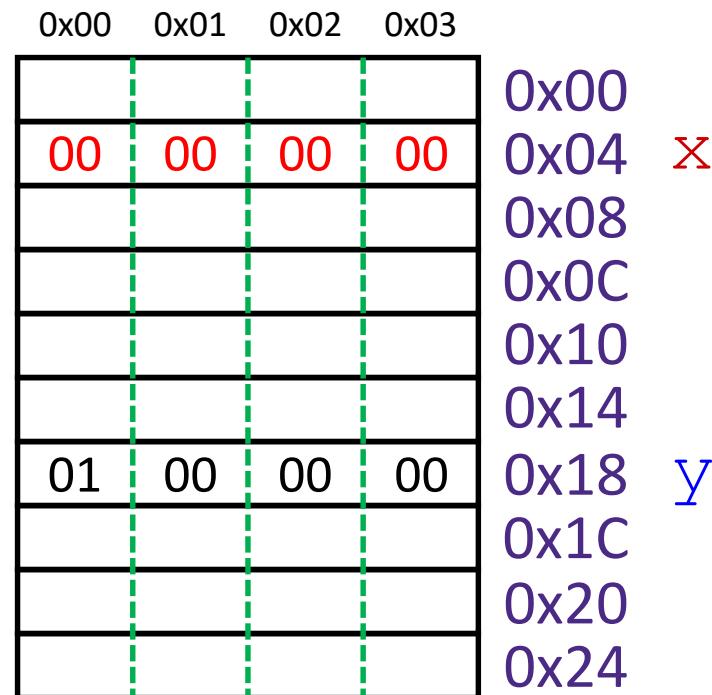
0x00	0x01	0x02	0x03	0x00	0x04	0x08	0x0C	0x10	0x14	0x18	0x1C	0x20	0x24
00	01	29	F3										
01	00	00	00										

# Assignment in C (3)

32-bit example  
(pointers are  
32-bits wide)



- left-hand side = right-hand side;
  - LHS must evaluate to a memory *location*
  - RHS must evaluate to a *value* (could be an address)
  - Store RHS value at LHS location
- `int x, y;`
- `x = 0;`



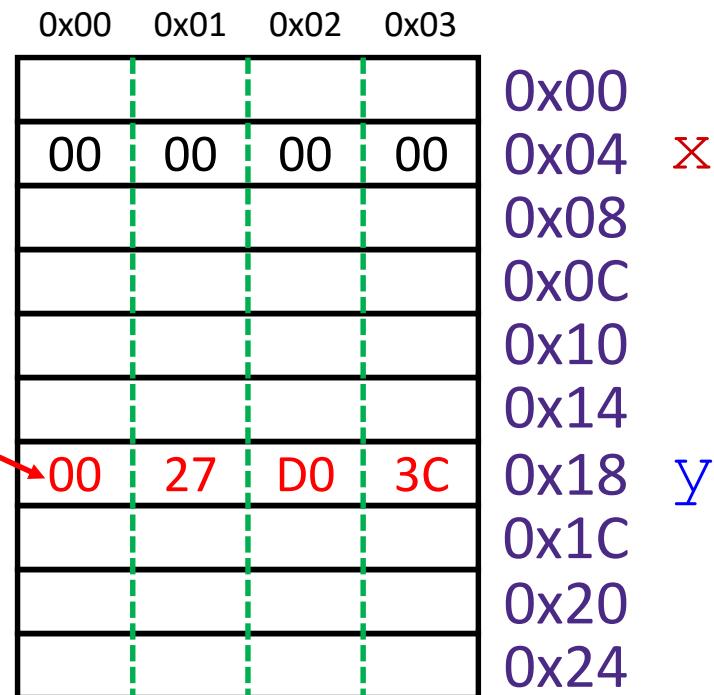
# Assignment in C (3)

32-bit example  
(pointers are  
32-bits wide)



- left-hand side = right-hand side;
  - LHS must evaluate to a memory *location*
  - RHS must evaluate to a *value* (could be an address)
  - Store RHS value at LHS location
- int x, y;
- x = 0;
- y = 0x3CD02700;

little endian!



# Assignment in C (4)

## 32-bit example (pointers are 32-bits wide)



- left-hand side = right-hand side;
    - LHS must evaluate to a memory *location*
    - RHS must evaluate to a *value* (could be an address)
    - Store RHS value at LHS location
  - int x, y;
  - x = 0;
  - y = 0x3CD02700;
  - x = y + 3;
    - Get value at y, add 3, store in

0x00	0x01	0x02	0x
03	27	D0	3
00	27	D0	3

0x00	0x01	0x02	0x03		
03	27	D0	3C	0x00	X
				0x04	
				0x08	
				0x0C	
				0x10	
				0x14	
00	27	D0	3C	0x18	Y
				0x1C	
				0x20	
				0x24	

# Assignment in C (5)

32-bit example  
(pointers are  
32-bits wide)



- left-hand side = right-hand side;
  - LHS must evaluate to a memory *location*
  - RHS must evaluate to a *value* (could be an address)
  - Store RHS value at LHS location
- `int x, y;`
- `x = 0;`
- `y = 0x3CD02700;`
- `x = y + 3;`
  - Get value at `y`, add 3, store in `x`
- `int* z;`
  - `z` is at address `0x20`

	0x00	0x01	0x02	0x03			
0x00	03	27	D0	3C			
0x04							X
0x08							
0x0C							
0x10							
0x14							
0x18	00	27	D0	3C			
0x1C							Y
0x20							
0x24	DE	AD	BE	EF			Z

Declaration: initially contains garbage

# Assignment in C (6)

32-bit example  
(pointers are  
32-bits wide)



- left-hand side = right-hand side;
  - LHS must evaluate to a memory *location*
  - RHS must evaluate to a *value* (could be an address)
  - Store RHS value at LHS location
- `int x, y;`
- `x = 0;`
- `y = 0x3CD02700;`
- `x = y + 3;`
  - Get value at `y`, add 3, store in `x`
- `int* z = &y + 3;`
  - Get address of `y`, “add 3”, store in `z`

	0x00	0x01	0x02	0x03		
0x00	03	27	D0	3C	X	
0x04						
0x08						
0x0C						
0x10						
0x14						
0x18	00	27	D0	3C	Y	
0x1C						
0x20						
0x24	24	00	00	00	Z	

# Pointer Arithmetic



- Pointer arithmetic is scaled by the size of target type
  - In this example, `sizeof(int) = 4`
- `int* z = &y + 3;`
  - Get address of `y`, add `3 * sizeof(int)`, store in `z`
  - `&y = 0x18`
  - `24 + 3 * (4) = 36`
- Pointer arithmetic can be dangerous!
  - Can easily lead to bad memory accesses
  - Be careful with data types and *casting*

# Assignment in C (7)

32-bit example  
(pointers are  
32-bits wide)



- int x, y;
- x = 0;
- y = 0x3CD02700;
- x = y + 3;
  - Get value at y, add 3, store in x
- int\* z = &y + 3;
  - Get address of y, add 12, store in z
- \*z = y;
  - What does this do?

& = “address of”  
\* = “dereference”

0x00	0x01	0x02	0x03	0x00	0x04	0x08	0x0C	0x10	0x14	0x18	0x1C	0x20	0x24
03	27	D0	3C										X
00	27	D0	3C										Y
													Z
24	00	00	00										

# Assignment in C (8)

32-bit example  
(pointers are  
32-bits wide)



- int x, y;
  - x = 0;
  - y = 0x3CD02700;
  - x = y + 3;
    - Get value at y, add 3, store in x
  - int\* z = &y + 3;
    - Get address of y, add 12, store in z
- The target of a pointer is  
also a memory location
- \*z = y;
    - Get value of y, put in address  
stored in z

& = “address of”  
\* = “dereference”

0x00	0x01	0x02	0x03	0x00	0x04	0x08	0x0C	0x10	0x14	0x18	0x1C	0x20	0x24
03	27	D0	3C										X
00	27	D0	3C										Y
24	00	00	00										Z
	00	27	D0	3C									

# Arrays in C (1)

Arrays are adjacent locations in memory storing the same type of data object

**a** is a name for the array's address

Declaration: `int a[6];`

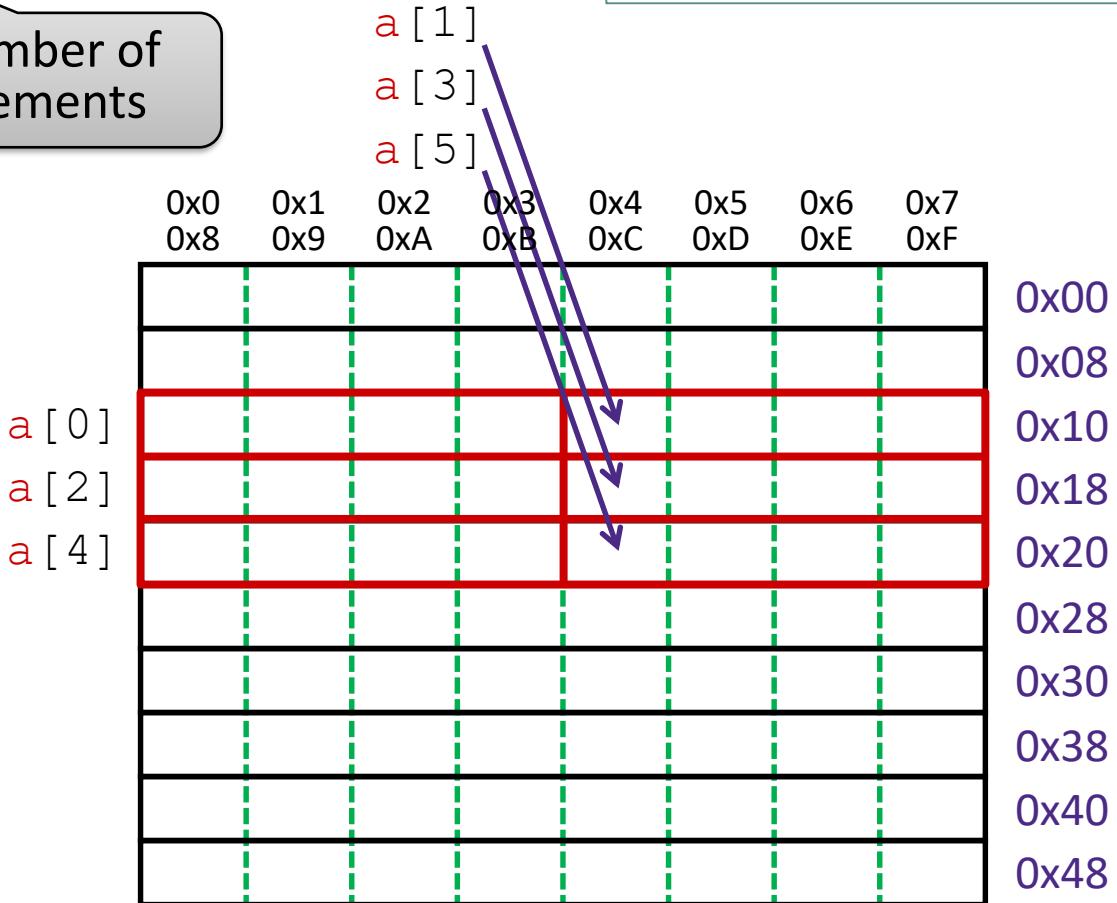
element type

name

number of elements

64-bit example

(pointers are 64-bits wide)



# Arrays in C (2)

Arrays are adjacent locations in memory storing the same type of data object

a is a name for the array's address

The address of a [ i ] is the address of a [ 0 ] plus i times the element size in bytes

Declaration: int a [ 6 ] ;

Indexing: a [ 0 ] = 0x015f ;  
a [ 5 ] = a [ 0 ] ;

	0x0	0x1	0x2	0x3	0x4	0x5	0x6	0x7	0xF	
	0x8	0x9	0xA	0xB	0xC	0xD	0xE	0xF		
<u>a [ 0 ]</u>	5F	01	00	00					0x00	
<u>a [ 2 ]</u>									0x08	
<u>a [ 4 ]</u>					5F	01	00	00	0x10	
									0x18	
									0x20	
									0x28	
									0x30	
									0x38	
									0x40	
									0x48	

# Arrays in C (3)

Arrays are adjacent locations in memory storing the same type of data object

**a** is a name for the array's address

The address of **a [i]** is the address of **a [0]** plus **i** times the element size in bytes

Declaration: `int a[6];`

Indexing: `a[0] = 0x015f;`  
`a[5] = a[0];`

No bounds checking: `a[6] = 0xBAD;`  
`a[-1] = 0xBAD;`

	0x0	0x1	0x2	0x3	0x4	0x5	0x6	0x7	
	0x8	0x9	0xA	0xB	0xC	0xD	0xE	0xF	
						AD	0B	00	00
a[0]	5F	01	00	00					0x00
a[2]									0x08
a[4]					5F	01	00	00	0x10
	AD	0B	00	00					0x18
									0x20
									0x28
									0x30
									0x38
									0x40
									0x48

# Arrays in C (4)

Arrays are adjacent locations in memory storing the same type of data object

`a` is a name for the array's address

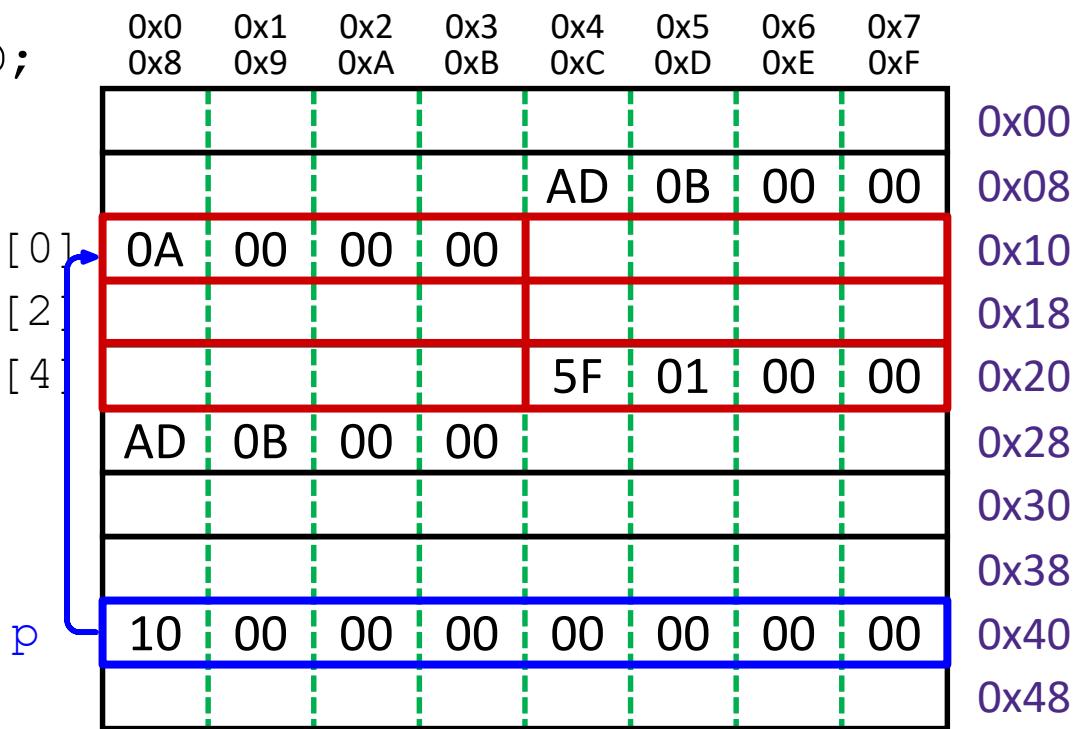
The address of `a[i]` is the address of `a[0]` plus `i` times the element size in bytes

Declaration: `int a[6];`

Indexing: `a[0] = 0x015f;`  
`a[5] = a[0];`

No bounds checking: `a[6] = 0xBAD;`  
`a[-1] = 0xBAD;`

Pointers: `int* p;`  
 equivalent {  
`p = a;`  
`p = &a[0];`  
`*p = 0xA;`



# Arrays in C (5)

Declaration: `int a[6];`

Indexing: `a[0] = 0x015f;`  
`a[5] = a[0];`

No bounds checking: `a[6] = 0xBAD;`  
`a[-1] = 0xBAD;`

Pointers: `int* p;`  
 equivalent { `p = a;`  
`p = &a[0];`  
`*p = 0xA;`

array indexing = address arithmetic  
 (both scaled by the size of the type)

equivalent { `p[1] = 0xB;`  
`* (p+1) = 0xB;`  
`p = p + 2;`

Arrays are adjacent locations in memory storing the same type of data object

`a` is a name for the array's address

The address of `a[i]` is the address of `a[0]` plus `i` times the element size in bytes

0x0	0x1	0x2	0x3	0x4	0x5	0x6	0x7	0x8	0x9	0xA	0xB	0xC	0xD	0xE	0xF	0x00	0x08	0x10	0x18	0x20	0x28	0x30	0x38	0x40	0x48	
																AD	0B	00	00							
<code>a[0]</code>	0A	00	00	00	0B	00	00	<code>p</code>	10	00	00	00	00	00	00	0x00	0x08	0x10	0x18	0x20	0x28	0x30	0x38	0x40	0x48	
<code>a[2]</code>																										
<code>a[4]</code>																5F	01	00	00							
	AD	0B	00	00																						

# Arrays in C (6)

Declaration: int a [6];

Indexing:      a [ 0 ] = 0x015f;  
                  a [ 5 ] = a [ 0 ];

No bounds checking: a[6] = 0xBAADF00D;  
a[-1] = 0xBAADF00D;

Pointers:      `int* p;`  
equivalent { `p = a;`  
              `p = &a[0];`  
              `*p = 0xA;`

array indexing = address arithmetic  
(both scaled by the size of the type)

equivalent { p[1] = 0xB;  
              \* (p+1) = 0xB;  
              p = p + 2;

`*p = a[1] + 1;`

Arrays are adjacent locations in memory storing the same type of data object

`a` is a name for the array's address

The address of  $a[i]$  is the address of  $a[0]$  plus  $i$  times the element size in bytes

# Fill in the blanks!

Declaration: `int a[6];`

Indexing: `a[0] = 0x015f;`  
`a[5] = a[0];`

No bounds checking: `a[6] = 0xBAD;`  
`a[-1] = 0xBAD;`

Pointers: `int* p;`  
 equivalent `{ p = a;`  
`p = &a[0];`  
`*p = 0xA;`

array indexing = address arithmetic  
 (both scaled by the size of the type)

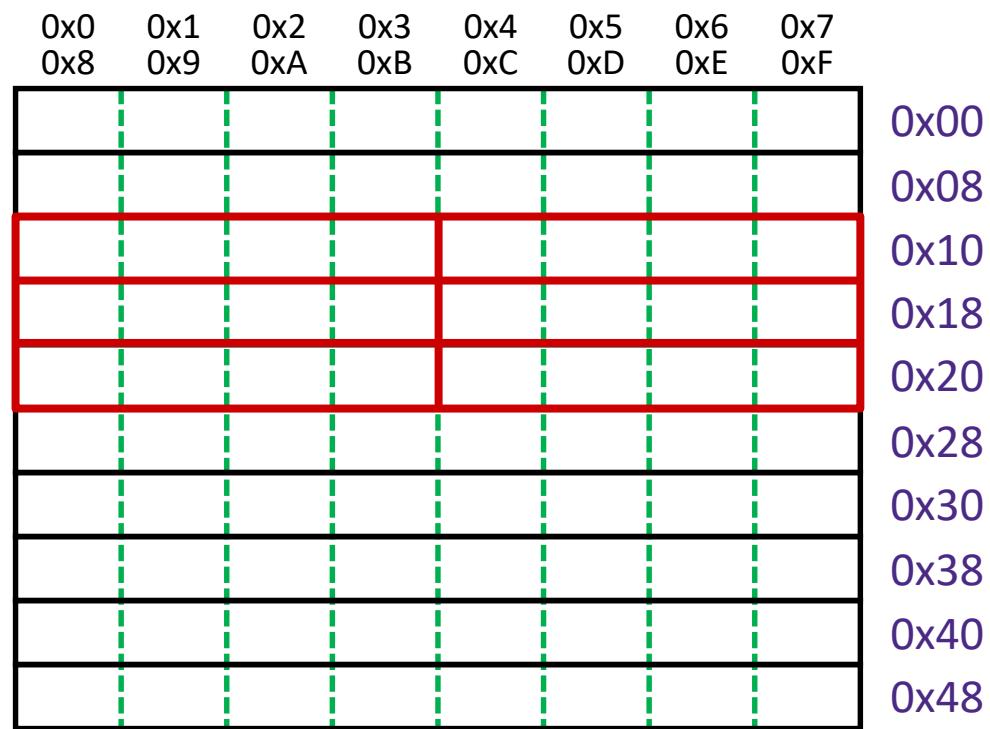
equivalent `{ p[1] = 0xB;`  
`* (p+1) = 0xB;`  
`p = p + 2;`

`*p = a[1] + 1;`

Arrays are adjacent locations in memory storing the same type of data object

`a` is a name for the array's address

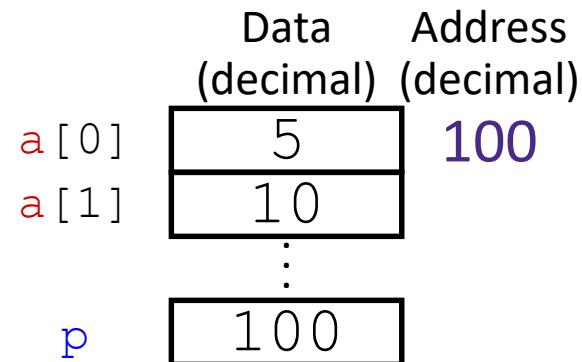
The address of `a[i]` is the address of `a[0]` plus `i` times the element size in bytes



**Question:** The variable values after Line 3 executes are shown on the right. What are they after Line 4 & 5?



```
1 void main () {  
2     int a[] = {5, 10};  
3     int *p = a;  
4     p = p + 1;  
5     *p = *p + 1;  
6 }
```



	p	*p	a[0]	a[1]		p	*p	a[0]	a[1]
--	---	----	------	------	--	---	----	------	------

- (A) 101 10 5 10 then 101 11 5 11
- (B) 104 10 5 10 then 104 11 5 11
- (C) 100 6 6 10 then 101 6 6 10
- (D) 100 6 6 10 then 104 6 6 10

# Representing strings



- C-style string stored as an array of bytes (`char *`)
  - Elements are one-byte **ASCII codes** for each character
  - No “String” keyword, unlike Java

32	<b>space</b>	48	<b>0</b>	64	<b>@</b>	80	<b>P</b>	96	<b>`</b>	112	<b>p</b>
33	<b>!</b>	49	<b>1</b>	65	<b>A</b>	81	<b>Q</b>	97	<b>a</b>	113	<b>q</b>
34	<b>"</b>	50	<b>2</b>	66	<b>B</b>	82	<b>R</b>	98	<b>b</b>	114	<b>r</b>
35	<b>#</b>	51	<b>3</b>	67	<b>C</b>	83	<b>S</b>	99	<b>c</b>	115	<b>s</b>
36	<b>\$</b>	52	<b>4</b>	68	<b>D</b>	84	<b>T</b>	100	<b>d</b>	116	<b>t</b>
37	<b>%</b>	53	<b>5</b>	69	<b>E</b>	85	<b>U</b>	101	<b>e</b>	117	<b>u</b>
38	<b>&amp;</b>	54	<b>6</b>	70	<b>F</b>	86	<b>V</b>	102	<b>f</b>	118	<b>v</b>
39	<b>,</b>	55	<b>7</b>	71	<b>G</b>	87	<b>W</b>	103	<b>g</b>	119	<b>w</b>
40	<b>(</b>	56	<b>8</b>	72	<b>H</b>	88	<b>X</b>	104	<b>h</b>	120	<b>x</b>
41	<b>)</b>	57	<b>9</b>	73	<b>I</b>	89	<b>Y</b>	105	<b>i</b>	121	<b>y</b>
42	<b>*</b>	58	<b>:</b>	74	<b>J</b>	90	<b>Z</b>	106	<b>j</b>	122	<b>z</b>
43	<b>+</b>	59	<b>;</b>	75	<b>K</b>	91	<b>[</b>	107	<b>k</b>	123	<b>{</b>
44	<b>,</b>	60	<b>&lt;</b>	76	<b>L</b>	92	<b>\</b>	108	<b>l</b>	124	<b> </b>
45	<b>-</b>	61	<b>=</b>	77	<b>M</b>	93	<b>]</b>	109	<b>m</b>	125	<b>}</b>
46	<b>.</b>	62	<b>&gt;</b>	78	<b>N</b>	94	<b>^</b>	110	<b>n</b>	126	<b>~</b>
47	<b>/</b>	63	<b>?</b>	79	<b>O</b>	95	<b>_</b>	111	<b>o</b>	127	<b>del</b>

ASCII: American Standard Code for Information Interchange

# Null-terminated Strings

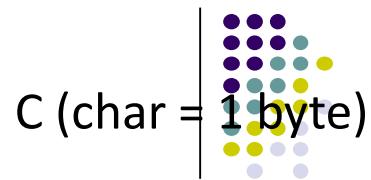


- Example: “Life is good” stored as a 13-byte array

Decimal:	76	105	102	101	32	105	115	32	103	111	111	100	0
Hex:	0x4c	0x69	0x66	0x65	0x20	0x69	0x73	0x20	0x67	0x6f	0x6f	0x64	0x00
Text:	L	i	f	e		i	s		g	o	o	d	\0

- Last character followed by a 0 byte (‘\0’)  
(a.k.a. “null terminator”)
  - Must take into account when allocating space in memory
  - Note that ‘0’ ≠ ‘\0’ (i.e. character 0 has non-zero value)
- How do we compute the length of a string?
  - Traverse array until null terminator encountered

# Endianness and Strings



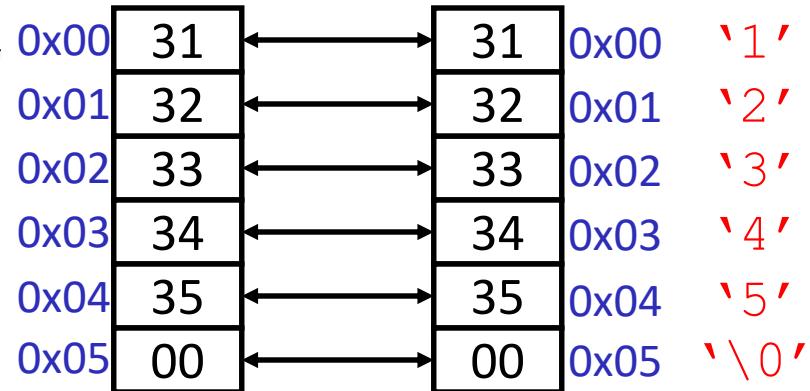
```
char s[6] = "12345";
```

String literal

0x31 = 49 decimal = ASCII '1'

IA32, x86-64  
(little endian)

SPARC  
(big endian)



- Byte ordering (endianness) is not an issue for 1-byte values
  - The whole array does not constitute a single value
  - Individual elements are values; chars are single bytes
- Unicode characters – up to 4 bytes/character
  - ASCII codes still work (just add leading zeros)
  - Unicode can support the many characters in all languages in the world
  - Java and C have libraries for Unicode (Java commonly uses 2 bytes/char)

# Examining Data Representations



- Code to print byte representation of data
  - Any data type can be treated as a *byte array* by **casting** it to char
  - C has **unchecked** casts **!! DANGER !!**

```
void show_bytes(char* start, int len) {  
    int i;  
    for (i = 0; i < len; i++)  
        printf("%p\t0x%.2x\n", start+i, *(start+i));  
    printf("\n");  
}
```

## printf directives:

%p	Print pointer
\t	Tab
%x	Print value as hex
\n	New line

# Examining Data Representations



- Code to print byte representation of data
  - Any data type can be treated as a *byte array* by **casting** it to char
  - C has **unchecked** casts **!! DANGER !!**

```
void show_bytes(char* start, int len) {  
    int i;  
    for (i = 0; i < len; i++)  
        printf("%p\t0x%.2x\n", start+i, *(start+i));  
    printf("\n");  
}
```

```
void show_int(int x) {  
    show_bytes( char * &x, sizeof(int));  
}
```

# show\_bytes Execution Example



```
int a = 12345; // 0x00003039
printf("int a = 12345;\n");
show_int(a); // show_bytes( (char *) &a, sizeof(int));
```

- Result (Linux x86-64):

- **Note:** The addresses will change on each run (try it!), but fall in same general range

```
int a = 12345;
0x7ffb7f71dbc      0x39
0x7ffb7f71dbd      0x30
0x7ffb7f71dbe      0x00
0x7ffb7f71dbf      0x00
```



# Summary

- Assignment in C results in value being put in memory location
- Pointer is a C representation of a data address
  - `&` = “**address of**” operator
  - `*` = “value at address” or “**dereference**” operator
- Pointer arithmetic scales by size of target type
  - Convenient when accessing array-like structures in memory
  - Be careful when using – particularly when *casting* variables
- Arrays are adjacent locations in memory storing the same type of data object
  - Strings are null-terminated arrays of characters (ASCII)

# Memory, Data, and Addressing



- Representing information as bits and bytes
- Organizing and addressing data in memory
- Manipulating data in memory using C
- **Boolean algebra and bit-level manipulations**



# Boolean Algebra

- Developed by George Boole in 19th Century
  - Algebraic representation of logic (True → 1, False → 0)
  - AND:  $A \& B = 1$  when both A is 1 and B is 1
  - OR:  $A | B = 1$  when either A is 1 or B is 1
  - XOR:  $A \hat{\wedge} B = 1$  when either A is 1 or B is 1, but not both
  - NOT:  $\sim A = 1$  when A is 0 and vice-versa
  - DeMorgan's Law:  
 $\sim (A | B) = \sim A \& \sim B$   
 $\sim (A \& B) = \sim A | \sim B$

AND		OR		XOR		NOT		
&	0	1		0	1	^	0	1
	0	0	0	0	1	0	0	1
	1	0	1	1	1	1	1	0

# General Boolean Algebras



- Operate on bit vectors
  - Operations applied bitwise
  - All of the properties of Boolean algebra apply

$$\begin{array}{r} 01101001 \\ \& 01010101 \\ \hline \end{array} \quad \begin{array}{r} 01101001 \\ \perp 01010101 \\ \hline \end{array} \quad \begin{array}{r} 01101001 \\ \wedge 01010101 \\ \hline \end{array} \quad \begin{array}{r} \sim 01010101 \\ \hline \end{array}$$

- Examples of useful operations:

$$x \wedge x = 0$$

$$\begin{array}{r} 01010101 \\ \wedge 01010101 \\ \hline 00000000 \end{array}$$

$$x \mid 1 = 1, \quad x \mid 0 = x$$

$$\begin{array}{r} 01010101 \\ \perp 11110000 \\ \hline 11110101 \end{array}$$

# Representing & Manipulating Sets



## • Representation

- A  $w$ -bit vector represents subsets of  $\{0, \dots, w-1\}$
- $a_j = 1$  iff  $j \in A$

01101001 { 0, 3, 5, 6 }  
76543210

01010101 { 0, 2, 4, 6 }  
76543210

## • Operations

• &	Intersection	01000001	{ 0, 6 }
•	Union	01111101	{ 0, 2, 3, 4, 5, 6 }
• ^	Symmetric difference	00111100	{ 2, 3, 4, 5 }
• ~	Complement	10101010	{ 1, 3, 5, 7 }

# Bit-Level Operations in C



- $\&$  (AND),  $|$  (OR),  $\wedge$  (XOR),  $\sim$  (NOT)
  - View arguments as bit vectors, apply operations bitwise
  - Apply to any “integral” data type
    - long, int, short, char, unsigned

## • Examples with `char a, b, c;`

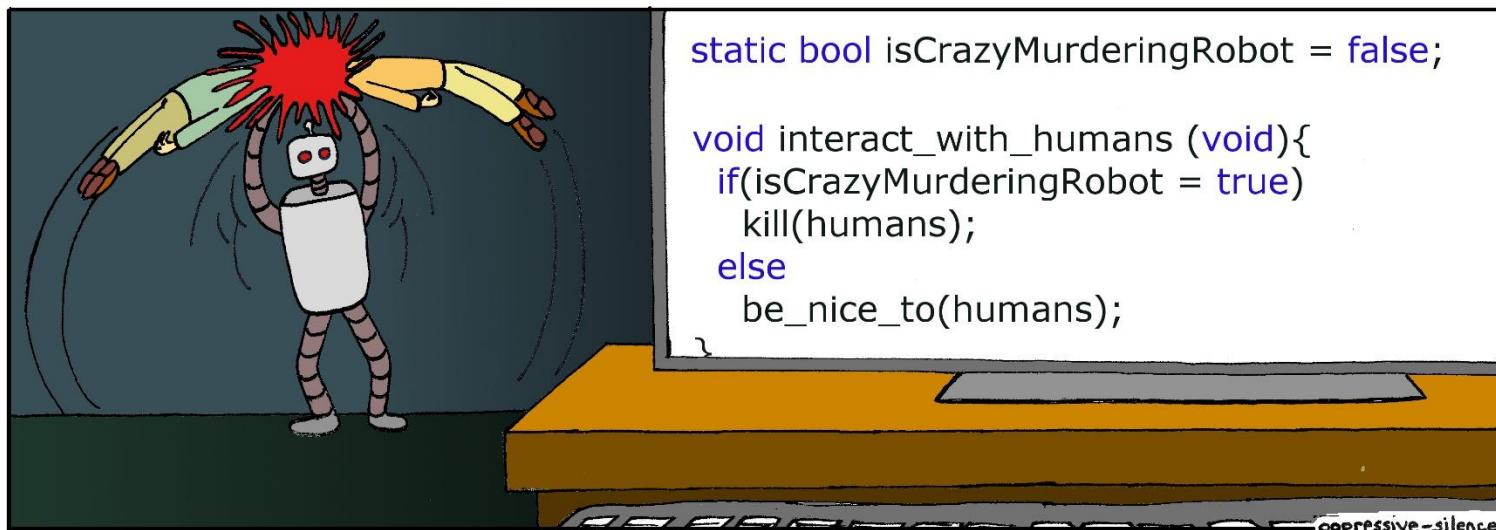
- `a = (char) 0x41;` // 0x41->0b 0100 0001
- `b = ~a;` // 0b ->0x
- `a = (char) 0x69;` // 0x69->0b 0110 1001
- `b = (char) 0x55;` // 0x55->0b 0101 0101
- `c = a & b;` // 0b ->0x
- `a = (char) 0x41;` // 0x41->0b 0100 0001
- `b = a;` // 0b 0100 0001
- `c = a ^ b;` // 0b ->0x



# Contrast: Logic Operations in C

- Logical operators in C: `&&` (AND), `||` (OR), `!` (NOT)
  - 0 is False, anything nonzero is True
  - Always return 0 or 1
  - **Early termination** (a.k.a. short-circuit evaluation) of `&&`, `||`
- Examples (`char` data type)
  - `! 0x41` → `0x00`
  - `! 0x00` → `0x01`
  - `! ! 0x41` → `0x01`
  - `p && *p++`
    - Avoids **null pointer** (`0x0`) access via *early termination*
    - Short for: `if (p) { *p++; }`

# Logic Operations!!!



# Q&A

