CHAPTER 4 - Memory

Questions 1

Consider the following page reference string: e, c, b, e, a, g, d, c, e, g, d, a

With 4 frames, how many page faults would occur with the following page replacement algorithms? Fill in the tables accordingly.

RS: reference strings; F0: frame 0, F1: frame 1, etc.

Hint: all frames are initially empty, so your first unique pages will all cost one fault each.

1. Optimal

| - r opania | | | | | | | | | | | | |
|-------------|---|---|---|---|---|---|---|---|---|----|----|----|
| Time | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 |
| RS | е | С | b | е | а | g | d | С | е | g | d | а |
| F0 | | | | | | | | | | | | |
| F1 | | | | | | | | | | | | |
| F2 | | | | | | | | | | | | |
| F3 | | | | | | | | | | | | |
| Page fault? | | | | | | | | | | | | |

Total page fault:

2 I RU

| Z. LIVO | | | | | | | | | | | | |
|-------------|---|---|---|---|---|---|---|---|---|----|----|----|
| Time | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 |
| RS | е | С | b | е | а | g | d | С | е | g | d | а |
| F0 | | | | | | | | | | | | |
| F1 | | | | | | | | | | | | |
| F2 | | | | | | | | | | | | |
| F3 | | | | | | | | | | | | |
| Page fault? | | | | | | | | | | | | |

Total page fault:

3.Second chance

| 0.0000110 | | | | | | | | | | | | |
|-------------|---|---|---|---|---|---|---|---|---|----|----|----|
| Time | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 |
| RS | е | С | b | е | а | g | d | С | е | g | d | а |
| F0 | | | | | | | | | | | | |
| F1 | | | | | | | | | | | | |
| F2 | | | | | | | | | | | | |
| F3 | | | | | | | | | | | | |
| Page fault? | | | | | | | | | | | | |

Total page fault:

Questions 2

- 1. True or False? A program does not need to be stored in memory in its entirety.
- 2. True or False? A physical address space is at least as large as a virtual address space.
- 3. When does a page fault occur?
- **4.** True or False? In a pure demand paged system a page is never brought into memory until it is needed.

1. A machine has 48 bit virtual addresses and 32 bit physical addresses. Pages are 8 KB. How many entries are needed for the page table?

2³⁵ pages

2. For each of the following decimal virtual addresses, compute the virtual page number and offset for a 4-KB page and for an 8 KB page: 20000, 32768, 60000.

20000 = page number * 4 * 1024 + offset

| Virtual address | Page number | Offset |
|-----------------|-------------|--------|
| 20000 | 4 | 3616 |
| 32768 | | |
| 60000 | | |

2.2 Page size = 8KB

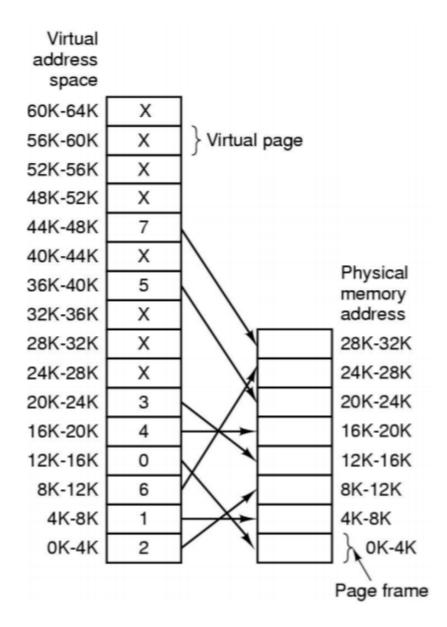
20000 = page number * 8 * 1024 + offset

| Virtual address | Page number | Offset |
|-----------------|-------------|--------|
| 20000 | 2 | 3616 |
| 32768 | | |
| 60000 | | |

3. The figure below shows a virtual address space from 0 to 64K and 32K of physical memory. There are 16 pages and 8 frames and transfers between memory and disk are in pages. Give the physical address corresponding to the following virtual addresses, explain how did you get the answer?:

a) 20 b) 4100 c) 8300

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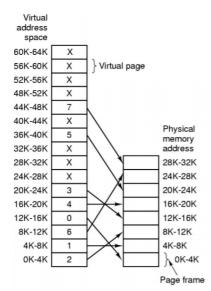


4. A memory free in 3 frames. How many page fault occur after running as the following page 7, 0, 1, 2, 0, 3, 0, 4, 2, 3, 0, 3, 2, 1, 2, 0, 1, 0, 7 using FIFO

| | 7 | 0 | 1 | 2 | 0 | 3 | 0 | 4 | 2 | 3 | 0 | 3 | 2 | 1 | 2 | 0 | 1 | 0 | 7 |
|----|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| F1 | | | | | | | | | | | | | | | | | | | |
| F2 | | | | | | | | | | | | | | | | | | | |
| F3 | | | | | | | | | | | | | | | | | | | |
| PF | | | | | | | | | | | | | | | | | | | |

2. With given page table and 16 bit virtual address, that is split 4 bit page numbers and 12 bit offset. If user references the address 15016, which physical address is in memory?

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Answer: ? 2728

3. A memory free in 4 frames. How many page faults do occur after running as the following page 2 3 2 0 1 5 2 4 5 3 2 5 2 using LRU

Answer:

| | 2 | 3 | 2 | 0 | 1 | 5 | 2 | 4 | 5 | 3 | 2 | 5 | 2 |
|----|---|---|---|---|---|---|---|---|---|---|---|---|---|
| F1 | | | | | | | | | | | | | |
| F2 | | | | | | | | | | | | | |
| F3 | | | | | | | | | | | | | |
| F4 | | | | | | | | | | | | | |
| PF | | | | | | | | | | | | | |

4. A memory free in 4 frames. Which state of the memory after the page 4 is accessed when the requested page as 2 3 2 0 1 5 2 4 5 3 2 5 2 using LRU

6. Assume that the Page Table below is in effect. The number of lines per page is 400. The actual memory location for line 1634 is $_{----}34_{--}$.

| Page Number | Page Frame Number |
|-------------|-------------------|
| 0 | 8 |
| 1 | 10 |
| 2 | 5 |
| 3 | 11 |
| 4 | 0 |

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7. A computer has four page frames. The time of loading, time of last access, and the R and M bits for each page are as shown below (the times are in clock ticks):

| Page | Loaded | Last ref | R | М |
|------|--------|----------|---|---|
| 0 | 226 | 280 | 0 | 0 |
| 1 | 160 | 265 | 0 | 1 |
| 2 | 110 | 270 | 1 | 0 |
| 3 | 120 | 285 | 1 | 1 |

Which page will LRU replace?

1

8. A computer has four page frames. The time of loading, time of last access, and the R and M bits for each page are as shown below (the times are in clock ticks):

| Page | Loaded | Last ref. | R | | М | |
|------|--------|-----------|-----|---|---|--|
| 0 | 226 | | 200 | | | |
| 0 | 226 | | 280 | 0 | Ü | |
| 1 | 160 | | 265 | 0 | 1 | |
| 2 | 110 | | 270 | 0 | 0 | |
| 3 | 120 | | 285 | 1 | 1 | |

Which page will Second Chance replace? (NRU)

9. A memory free in 3 frames. How many page hits do? Assume that the running as the following page 7, 0 , 1, 2 , 0, 3, 0 , 4, 2 , 3 , 0 , 3 , 2 , 1, 2, 0, 1, 0, 7 using LRU

| | 7 | 0 | 1 | 2 | 0 | 3 | 0 | 4 | 2 | 3 | 0 | 3 | 2 | 1 | 2 | 0 | 1 | 0 | 7 |
|----|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| F1 | | | | | | | | | | | | | | | | | | | |
| F2 | | | | | | | | | | | | | | | | | | | |
| F3 | | | | | | | | | | | | | | | | | | | |
| PF | | | | | | | | | | | | | | | | | | | |

9. A computer provides the user with virtual address space of 2^32 (2 to the power 32) bytes. Pages of size 4096 (4K or 2^12) bytes are used for implementing virtual memory where the total physical memory is equal to 2^18 bytes. If the hexadecimal virtual

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10. A computer has four page frames. The time of loading, time of last access, and the R and M bits for each page are as shown below (the times are in clock ticks):

| Page | Loaded | Last Ref. | R M |
|------|--------|-----------|-----|
| 0 | 230 | 285 | 1 0 |
| 1 | 120 | 265 | 0 0 |
| 2 | 140 | 270 | 0 1 |
| 3 | 110 | 280 | 1 1 |

- (a) Which page will NRU replace? (1)
- (b) Which page will FIFO replace? (3)
- (c) Which page will LRU replace? (1)
- (d) Which page will second chance replace? (1)
- 11. Consider a logical address space of 64 pages of 2048 words each, mapped unto a physical memory of 32 frames.
 - a) How many bits are there in logical address?

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b) How many bits are there in physical address?

16

12. A system with 32 bit virtual address. If the page size is 4 KB and each table entry occupies 4 bytes, what is the size of the page table?

Consider a swapping system in which the memory consists of the following hole sizes: 10K, 4K, 20K, 15K, 9K. Assume worst fit algorithm is used. Which holes are taken for successive segment requests of 8K,

- 13. If there are 64 pages and the page size is 2048 words, what is the length of logical address?
- 14. A system with 32 bit virtual address. If the page size is 4 KB and each table entry occupies 4 bytes, what is the size of the page table?

```
total pages = 2^2
```

size of page table =
$$4B * 2 ^ 20 = 4MB$$

size of virtual memory = 2[^] 20 * 4KB = 4GB

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