

# Notes for Paper 1

## Paper 1

**Title: Data-driven decision making in power systems with probabilistic guarantees: Theory and applications of chance-constrained optimization**

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**Tags: chance-constraint; optimization; power system; data-driven**

### Notes:

#### 1. Problem description:

The problem of **Chance-constrained optimization(CCO)** can usually be expressed like this:

$$\begin{aligned} \min_x \quad & c^T x \\ \text{s.t.} \quad & P_\varepsilon(f(x, \varepsilon) \leq 0) \geq 1 - \epsilon \\ & x \in \chi \end{aligned}$$

This kind of problem can be converted into another type:

Violation Probability:  $V(x^o) := P_\varepsilon(f(x, \varepsilon) \geq 0)$

where  $x^o$  is a feasible solution to CCO.

Define:

$$F_\epsilon := \{x \in R^n : V(x) \leq \epsilon\} = \{x \in R^n : P_\varepsilon(f(x, \varepsilon) \geq 0) \leq 1 - \epsilon\}$$

Then,  $x^o$  is feasible to CCO if  $x \in \chi \cap F_\epsilon$ .

So CCO can be equivalently written as:

$$\begin{aligned} \min_x \quad & c^T x \\ \text{s.t.} \quad & V(x) \leq \epsilon \end{aligned}$$

$$x \in \chi$$

So, the first step is checking its **feasibility** ( $x^o$  satisfies the constraint), then is checking its **optimality** (the distance between the corresponding  $o^o$  and the best optimal value  $o^*$ ).

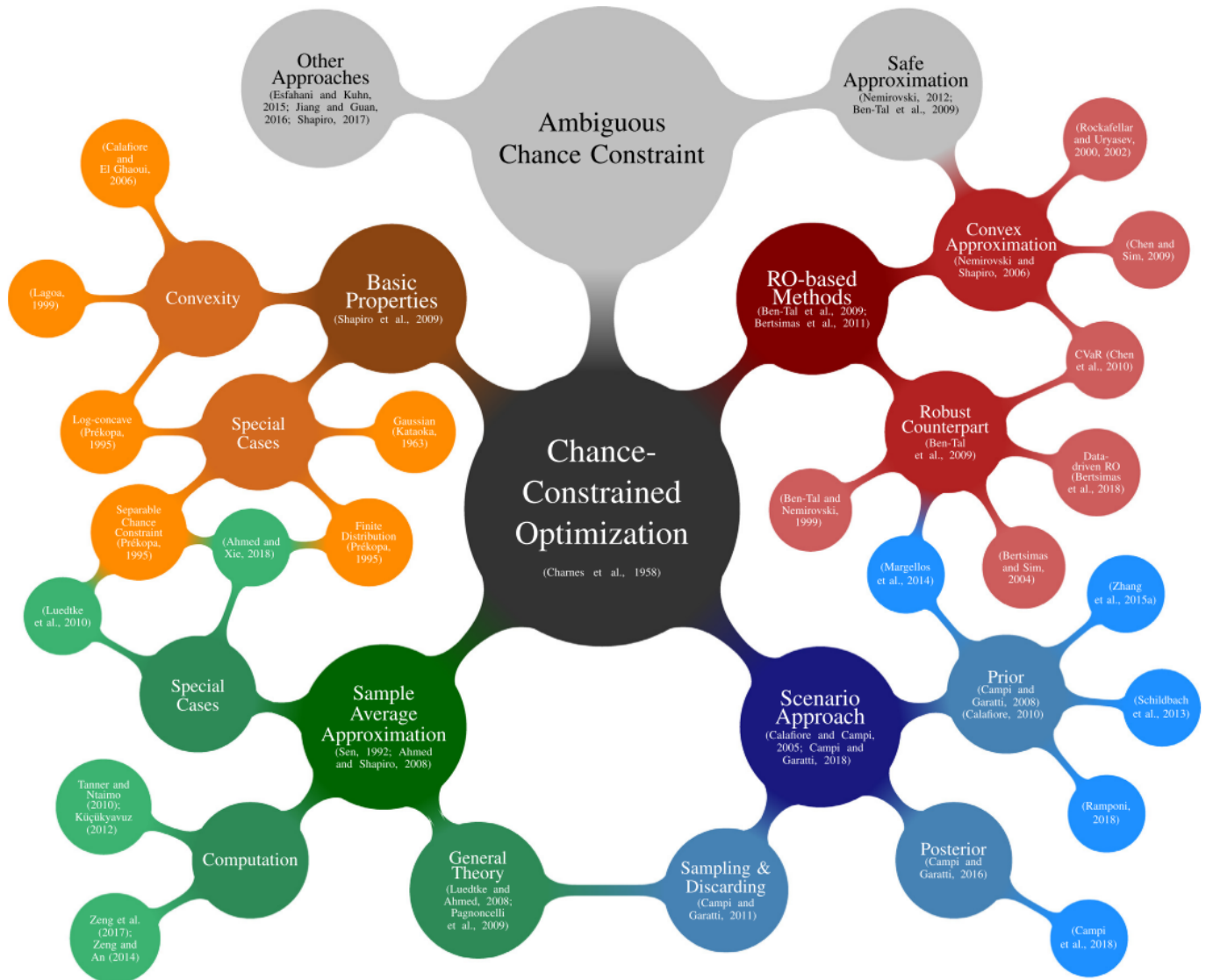
## 2. difficulties or issues:

There are **two main difficulties** to solve the problems:

(D1) calculating the probability involves multivariate integration, which is NP-Hard.

(D2) feasible region  $F_\epsilon$  is often non-convex

## 3. Methods



### Scenario approach

Scenario approach utilizes a dataset with  $N$  scenarios  $\{\epsilon_i\}_{i=1}^N$  to approximate CCO:

$$\begin{aligned} (SP_N) \min_{x \in \chi} \quad & c^T x \\ \text{s.t.} \quad & f(x, \epsilon^1) \leq 0, \dots, f(x, \epsilon^N) \leq 0 \end{aligned}$$

Because massive scenario number  $N$  will cause huge computation problem, so it involves another question: how to find the suitable scenario number (sample complexity  $N$ )?

### a. A-priori feasibility guarantees

steps:

1. exploring the problem structure and obtain the upper bound  $\bar{h}$  on the number of support scenarios;
2. choosing a good sample complexity  $N$
3. Solving  $SP_N$  and obtain  $x_N^*$  and  $o_N^*$

Many scholars make efforts on how to find the upper bound  $\bar{h}$  and the good sample complexity  $N$

### b. A-posteriori feasibility guarantees

steps:

1. given dataset  $\{\varepsilon_i\}_{i=1}^N$  solve  $SP_N$  and obtain  $x_N^*$ ;
2. find the support scenarios number (denoted as  $s_N^*$ )
3. calculate the posterior violation probability  $\epsilon(\beta, s_N^*, N)$
4. if  $\epsilon(\beta, s_N^*, N) \geq \epsilon$ , repeat step 1 to 3 with more scenarios until reaching  $\epsilon(\beta, s_N^*, N) \leq \epsilon$ . If all available scenarios are used but still fails to reach the condition, then it might be impossible to obtain a solution  $x_N^*$ .

$\epsilon(\beta, s_N^*, N)$  is the solution in the interval  $(0, 1)$  of the equality below:

$$\frac{\beta}{N+1} \sum_{i=k}^N \binom{i}{k} t^{i-k} - \binom{N}{k} t^{i-k} = 0$$

where  $\beta \in (0, 1)$

Many scholars make efforts on how to find the upper bound  $\bar{\epsilon}$ . If  $\bar{\epsilon} < \epsilon$ , then  $x^o$  is feasible.

### c. Optimality guarantees

The optimality checking is conducted by constructing lower bounds  $\underline{o}$  on  $o^*$

### Sample average approximation

Sample average approximation converts CCO into:

$$(SAA) : \min_x c^T x$$

$$s.t. \frac{1}{N} \sum_{i=1}^N 1_{\bar{f}(x, \varepsilon^i) > 0} \leq \varepsilon$$

where

$\bar{f}(x, \varepsilon^i) := \max \{f_1(x, \varepsilon), \dots, f_m(x, \varepsilon)\}$  and the violation probability  $\varepsilon$  differs from  $\epsilon$  in CCO.

**Data-driven:** SAA approximates the true distribution from  $N$  samples  $\{\varepsilon_i\}_{i=1}^N$ .

SAA further converts CCO into:

$$\begin{aligned} \min_{x, z} \quad & c^T x \\ \text{s.t.} \quad & f(x, \varepsilon^i) - M z_i \mathbf{1}_m \leq 0 \\ & \frac{1}{N} \sum_{i=1}^N z_i \leq \varepsilon \\ & x \in \chi, z_i \in \{0, 1\}, i = 1, 2, \dots, N \end{aligned}$$

Because  $M$  are big coefficients (weak formulations), it may cause numerical issues. Many scholars pay their attention on finding strong formulations without big coefficients  $M$ .

### a. feasibility guarantees

The feasible region of SAA is defined as:

$$F_{\varepsilon, \gamma}^N := \{x \in \chi : \frac{1}{N} \sum_{i=1}^N \mathbf{1}_{\bar{f}(x, \varepsilon^i) + \gamma \leq 0} \geq 1 - \varepsilon\}$$

Solutions of SAA is feasible to CCO with high probability  $1 - \beta$ :

$$P(F_{\varepsilon, \gamma}^N \subseteq F_{\epsilon}) \geq 1 - \beta$$

### b. optimality guarantees

use SAA to generate lower bounds of CCO  $o_L^*$  with probability at least  $1 - \delta$ .

### Robust optimization related methods

Robust optimization's typical form is:

$$\begin{aligned} (RC) : \min_{x \in \chi} \quad & c^T x \\ \text{s.t.} \quad & f(x, \varepsilon) \leq 0, \varepsilon \in U_{\epsilon} \end{aligned}$$

RC finds the optimal solution which is feasible all realizations of uncertainties. The key point is how to construct an good uncertainty set  $U_{\epsilon}$ .

Two points for an good uncertainty set:

1. RC is computationally tractable.
  2. The optimal solution of RC is not too conservative.
- Many scholars made efforts on the second point.

## safe approximation

$$(SA) : \min_{x \in \chi} c^T x \\ s.t. \quad x \in \underline{F}$$

where  $\underline{F} \subseteq F_\epsilon$ .  $F_\epsilon$  is the feasible region of CCO.

The authors used a large space to talk how to find the uncertainty sets by applying safe approximation in **individual chance constraints**. However, it is hard for a rookie like me to figure out the mathematical derivation in each step.

Authors introduced several approaches to apply safe approximation into joint chance constraints:

1. convert joint chance constraints to individual chance constraints. For example, use

$$P(f_i(x, \epsilon) \leq 1 - \epsilon_i, \quad i = 1, \dots, m)$$

or use the pointwise maximum

$$\bar{f}(x, \epsilon) := \max\{f_1(x, \epsilon), \dots, f_m(x, \epsilon)\}$$

2. directly deal with joint chance constraints. Authors introduced three typical approaches.

**(The mathematic part is over, next is application part.)**

## Applications in power systems

Pivotal task: maintain the real-time balance of supply and demand while ensuring the system is low-cost and reliable.

### Security-constrained economic dispatch(SCED)

**a. Deterministic SCED:** no uncertainties.

Example: direct current optimal power flow (DCOPF)

$$(det - DCOPF) : \min_g c(g) \\ s.t. \quad 1^T g = 1^T d - 1^T \hat{w} \\ f = H_g g + H_w \hat{w} - H_d d \\ \underline{f} \leq f \leq \bar{f} \\ \underline{g} \leq g \leq \bar{g}$$

where decision variables are generation output levels  $g \in R^{n_g}$ .  $c(g)$  is the total generation cost,  $1^T d - 1^T \hat{w}$  is net demands,  $H$  is power transfer distribution factor (PTDF) matrix,  $f$  is transmission line flows,  $\hat{w}$  is wind generation.

## b. Chance-constrained SCED

Treat wind generation  $w$  as a random vector.

$$\begin{aligned} (cc - DCOPF) : \min & c(g) \\ s.t. & 1^T g = 1^T d - 1^{g,\eta T} w \\ & f(\hat{w}, \tilde{w}) = H_g(\underline{g} - 1^T \tilde{w} \eta) - H_d d + H_w(\hat{w} + \tilde{w}) \\ & P_{\hat{w}}(\underline{f} \leq f(\hat{w}, \tilde{w}) \leq \bar{f}) \text{ and } \underline{g} \leq g - 1^T \tilde{w} \eta \leq \bar{g} \geq 1 - \epsilon \\ & 1^T \eta = 1 \\ & \underline{g} \leq g \leq \bar{g} \\ & -1 \leq \eta \leq 1 \end{aligned}$$

where  $\eta$  is affine control policy  $\eta \in [-1, 1]^{n_g}$  (proportionally allocate total wind fluctuations to each generate), it's also called participation factor or distribution vector.

## Security-constrained unit commitment(SCUC)

SCUC is used to minimize the generation cost.

### Deterministic SCUC

$$\begin{aligned} (det - SCUC) : \min_{z,u,v,g,s} & \sum_{t=1}^{n^t} c_n^T z^t + c_u^T u^t + c_v^T v^t + c_g^T g^{t,0} + c_s^T s^t \\ s.t. & 1^T g^{t,k} \geq 1^T \hat{d}^t - 1^T \hat{w}^t \end{aligned}$$

$$\underline{f} \leq H_g^{t,k} g^{t,k} - H_d^{t,k} \hat{d}^t + H_w^{t,k} \hat{w}^t \leq \bar{f} \quad (76c)$$

$$\underline{r} \leq g^{t,k} - g^{t-1,k} \leq \bar{r} \quad (76d)$$

$$a^k \circ (g^{t,0} - s^t) \leq g^{t,k} \leq a^k \circ (g^{t,0} + s^t) \quad (76e)$$

$$k \in [0, n_k], t \in [1, n_t]$$

$$\underline{g} \circ z^t \leq g^{t,0} \leq \bar{g} \circ z^t \quad (76f)$$

$$\underline{s} \circ z^t \leq s^t \leq \bar{s} \circ z^t \quad (76g)$$

$$\underline{g} \circ z^t \leq g^{t,0} - s^t \leq g^{t,0} + s^t \leq \bar{g} \circ z^t \quad (76h)$$

$$z^{t-1} - z^t + u^t \geq 0 \quad (76i)$$

$$z^t - z^{t-1} + v^t \geq 0 \quad (76j)$$

$$t \in [1, n_t]$$

$$z_i^t - z_i^{t-1} \leq z_i^t, \quad \iota \in [t+1, \min\{t + \underline{u}_i - 1, n_t\}] \quad (76k)$$

$$z_i^{t-1} - z_i^t \leq 1 - z_i^t, \quad \iota \in [t+1, \min\{t + \underline{v}_i - 1, n_t\}] \quad (76l)$$

$$i \in [1, n_g], \quad t \in [2, n_t]$$

Objective function is total operation costs, including no-load costs  $c_n^T z^t$ , startup costs  $c_u^T u^t$ , shutdown costs  $c_v^T v^t$ , generation costs  $c_g^T g^{t,0}$  and reserve costs  $c_s^T s^t$ .

The first constraint assures enough supply to meet net demand.

Constraints (76c), (76d) and (76g) are about transmission capacity, generation ramping capability and reserve limit in contingency scenario k at time t.

Constraints (76f) and (76g) are generation and reserve capacity.

(76i)-(76j) are logistic constraints about commitment status, startup and shutdown decisions.

Constraints (76k)-(76l) are minimum on/off time constraints for all generators.

Deterministic SCUC has no uncertain variables

## Chance-constrained SCUC

(cc-SCUC):

$$\min_{z,u,v,g,s} \sum_{t=1}^{n_t} c_n^T z^t + c_u^T u^t + c_v^T v^t + c_g^T g^{t,0} + c_s^T s^t \quad (77a)$$

$$\text{s.t. (76b), (76c), (76d), (76e), } k \in [0, n_k], t \in [1, n_t] \\ (76f), (76g), (76h)), (76i), (76j), t \in [1, n_t] \\ (76k), (76l), i \in [1, n_g], t \in [2, n_t]$$

$$\mathbb{P}\left(\mathbf{1}^T g^{t,k} \geq \mathbf{1}^T (\hat{d}^t + \tilde{d}^t) - \mathbf{1}^T (\hat{w}^t + \tilde{w}^t), \quad (77b)$$

$$\underline{f} \leq H_g^{t,k} g^{t,k} - H_d^{t,k} (\hat{d}^t + \tilde{d}^t) \\ + H_w^{t,k} (\hat{w}^t + \tilde{w}^t) \leq \bar{f}, \quad (77c)$$

$$k \in [0, n_k], t \in [1, n_t] \Big) \geq 1 - \epsilon \quad (77d)$$

Wind generation is modeled as a random vector consisting of a deterministic predicted component  $\hat{w} \in R^{n_w}$  and a stochastic error component  $\tilde{w} \in R^{n_w}$

The demand is also treated as a random vector with deterministic component  $\hat{d}$  and a stochastic error component  $\tilde{d}$ .

### solving chance-constrained SCUC

Sample average approximation is commonly used. There is no upper bound on the number of support scenarios for non-convex problems, so the scenario approach cannot be directly applied on cc-SCUC.

### Generation and transmission expansion Generation

#### a. Purposes:

1. when to invest on new elements such as transmission lines and generators;
2. what types of new elements are necessary;
3. how much capacity is needed and where the best location would be for those new elements.

#### b. Objective function:

1. total cost of investment in new generators and transmission line;
2. environmental impacts;
3. cost of generation.

#### c. Constrains:

1. total or individual costs within budget;
2. capacity constraint;
3. reliability requirement;



4. supply-demand balance;
5. power flow equations;
6. operation requirements such as generation or transmission limits.

**d. Uncertainties:**

1. demand;
2. generation;
3. transmission outages;
4. renewables.