

CSCI 4061

Lecture 1

Course Mechanics & Introduction

Instructor: Jack Kolb

January 18, 2022

Instructor Information

- Name: Jack Kolb
- Email: jhkolb@umn.edu
- Office Hours
 - 10am – 12pm on Monday
 - 1pm – 2pm on Thursday
 - Can always email to schedule an appointment (including on-line)
- Office Location: 333 Shepherd Labs



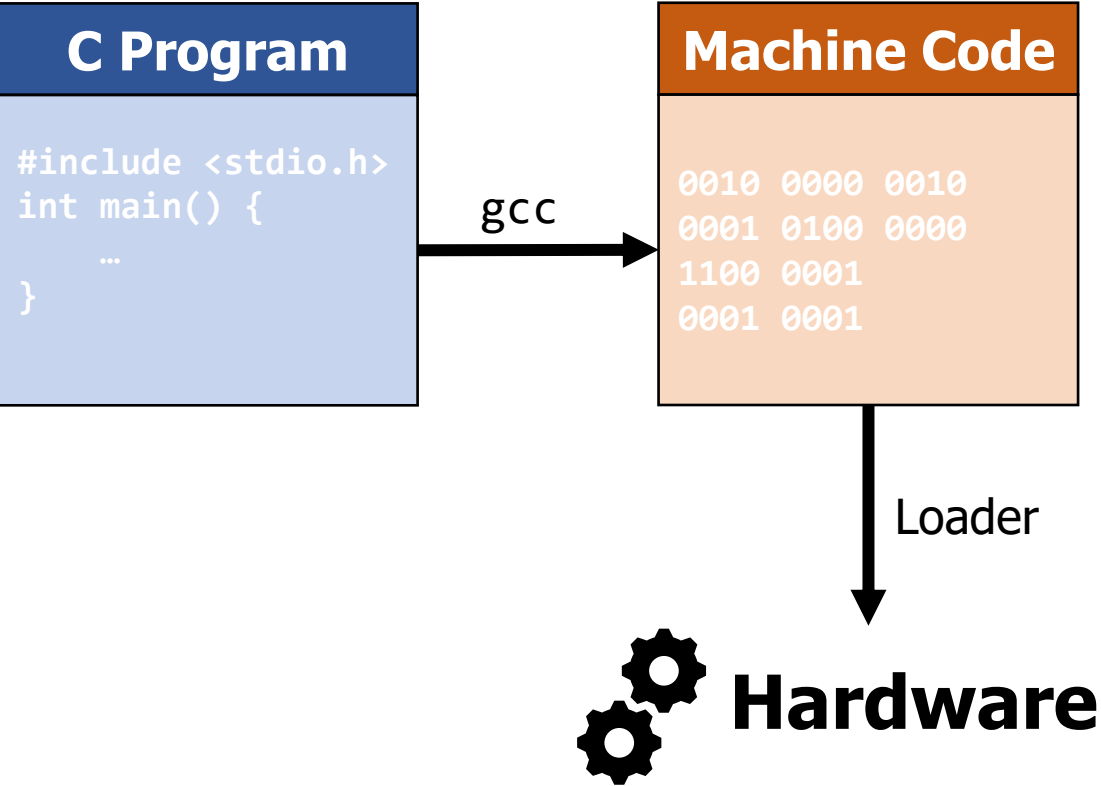
What's the point of this class?

- Consider: What is computer science?
- The study of **solving problems** with **computers**
 - One aspect: designing **algorithms** and **data structures** to form your solution
 - But a solution only comes to life when it executes on **real hardware**
- 2021 is a good start to learning how execution of code really works
 - C -> assembly -> machine code -> electrical signals in hardware
 - Hardware ingredients: CPU, memory, buses, I/O, disk storage
- But it gives you a very incomplete picture!
 - One Exception: Virtual Memory

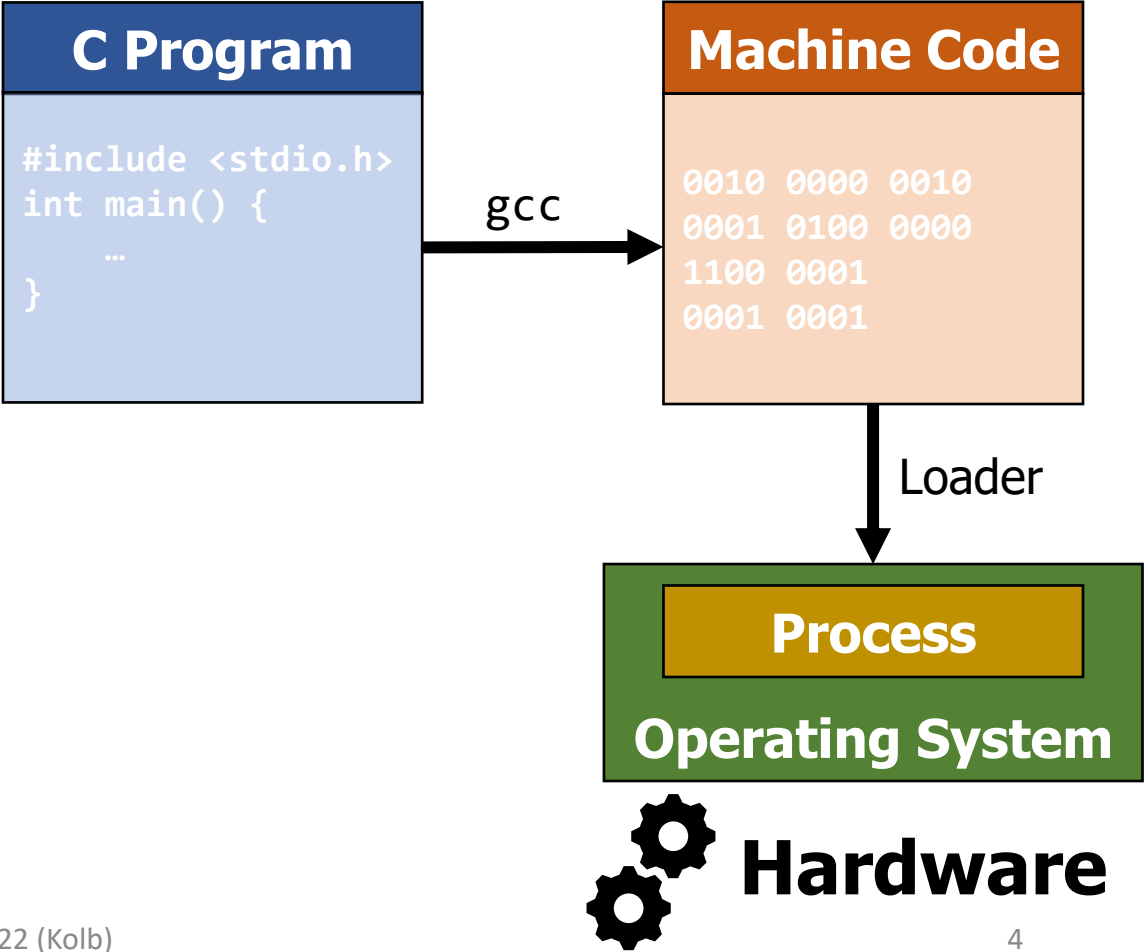


Enter Operating Systems

2021 View of the World



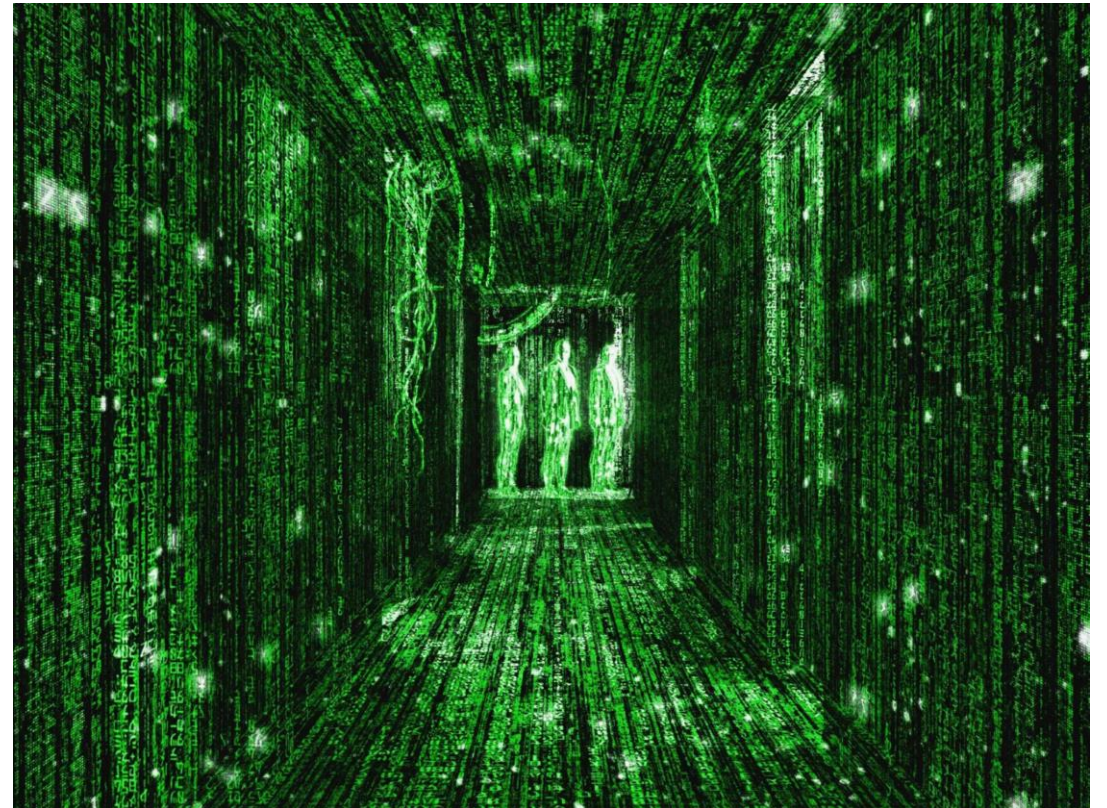
4061 (More Real) View of the World



Process Sees an OS-Constructed “Reality” -- A Virtual Machine

- When a process is running, it sees a consistent set of abstractions:
 - Virtual Memory Addresses
 - Threads
 - Files
 - Signals (like a seg fault)
 - Sockets (for networking)
- Regardless of what the underlying hardware looks like

All illusions created by the OS and hardware to isolate and control your program!



Why do this?

- Reason #1: **Hide diversity of hardware from programs**
- Example: Disk Storage
 - Spinning Metal or Solid-State
 - Made by Intel, Samsung, Toshiba, Micron, Western Digital, Seagate, ...
 - Different Capacities
- Programs just see files and same code “just works” for all disks
- Example: Networking
 - Wired: Ethernet and Cables
 - Wireless: Wi-Fi and Antennas
 - Made by Intel, Broadcom, Atheros, ...
- Programs just see sockets and same code “just works” for all devices



Why do this?

- Reason #2: **Protect different processes from each other**
- **Multiprogramming:** Several programs execute on a computer at same time
- Share CPU (*Time*)
- Share Memory (*Space*)
- The OS is responsible for enforcing safe sharing of resources among all processes
- Processes interact with each other only in restricted ways

Closer to Reality: Sharing Resources

- Thus far in the class, we've assumed a **uniprogramming** model
- One program runs on a computer at a time
- That program has exclusive access to the hardware

```
gcc -o fshow fshow.c
CFLAGS = -O -Wall -m32
LIBS = -lm
all: btest fshow
btest: btest.c bits.s decl.c tests.c btest.h bits.h
$(CC) $(CFLAGS) $(LIBS) -o btest btest.c decl.c tests.c
fshow: fshow.c
$(CC) $(CFLAGS) -o fshow fshow.c
install: fshow
$(CC) $(CFLAGS) -o fshow fshow.c
test: btest
./btest
clean:
rm -f *.o btest fshow *
```

Course Schedule

Week	Date	Lecture	Textbook	Homework	Lab	Project	Exam
1	9/6/2021	021 No Class - Labor Day Holiday		Entry Survey Released			
	9/8/2021	Intro, Course Mechanics 021	1.1-1.6		Lab 0		
	9/10/2021	More Intro, Start C Basics 021	C Reference				
2	9/13/2021	C Programming Basics, Pointers, Arrays 021	C Reference				
	9/15/2021	C Programming - Pointers, Memory Management 021	C Reference				

RIP apollo

Learning about Operating Systems

CSCI 4061: Introduction to OS

- Looking at operating systems through lens of *systems programming*
- Understand what operating systems offer us and how to intelligently use these features
- Discuss OS internals when motivated by practical applications

CSCI 5103: Operating Systems

- Looking at operating system internals for its own sake
- Understand theory and practice of designing/improving OS
- Actually implement the features that regular user programs depend on



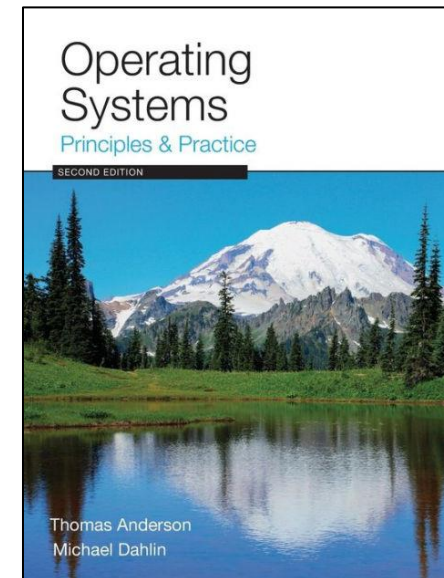
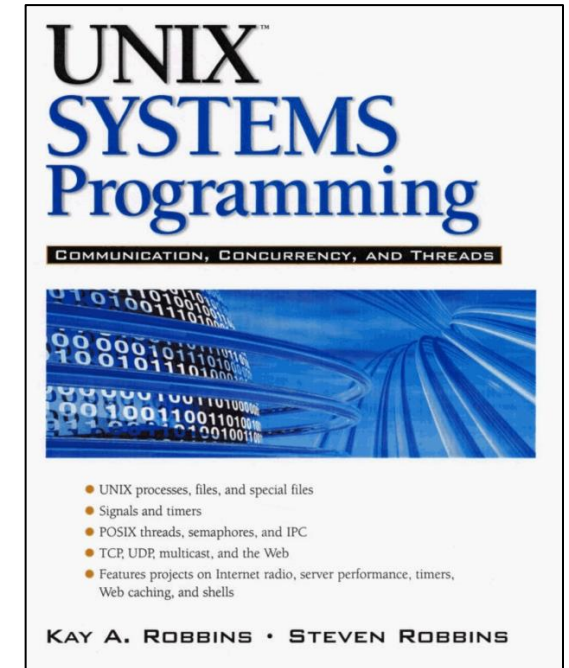
Course Mechanics

- Lecture: Tuesday/Thursday, 75 minutes
 - Will feature in-class exercises to be done in small groups
 - Recordings will be posted on “Media Gallery” of our Canvas site
 - Usually at least one longer exercise/break
- Canvas Site: Main source of information for this class
 - Syllabus posted there has full details
 - Full Course Schedule, Lecture Slides, Assignment Documents
 - View grades
 - Office Hours Calendar: Bottom of Main Page
- Piazza: Q&A Forum
- Gradescope: Submit labs and projects, exam regrade requests



Course Textbooks

- *UNIX Systems Programming* (Second Edition), Robbins & Robbins.
 - A decent text, may want to consult Stephens and Rago (listed on Syllabus) after this class
- *Operating Systems: Principles and Practice* (Second Edition), Anderson & Dahlin.
 - Optional but useful
 - Focuses on OS internals and implementation
 - Particularly relevant for virtual memory



Labs

- Meets weekly on Monday, code due at end of next day (Tues)
- A mix of recitation and hands-on activities
- Lab assignments are 10% of course grade (drop two lowest)
- You will learn **a lot** by doing in this class!
 - Lecture will not be able to teach you all you need to know
- Encouraged to work in groups
- Submit to Gradescope



Weekly Quizzes

- 10% of Course Grade (drop two lowest scores)
- Weekly: Released on Monday, due following Monday
 - Posted on Canvas site
- Meant to give you a chance to review and apply lecture concepts
 - And give you practice for exams!
- Open Resource
 - Collaboration encouraged (as long as you understand your answers)
 - Unlimited submissions
- First quiz is out now, Due January 24



Projects

- 35% of Course Grade
- 4 Projects Planned: 8.75% of grade per project
- You will learn **a lot** by doing in this class
- To be completed individually or in pairs
- Submit to Gradescope
 - Grade determined by **both** test results and manual inspection
 - Score on Gradescope will not match final score on Canvas



Exams

- Midterm Exam 1: 12.5% of Course Grade
 - **Thursday, February 17**, During Normal Lecture Time
- Midterm Exam 2: 12.5% of Course Grade
 - **Thursday, March 31**, During Normal Lecture Time
- Final Exam: 15% of Course Grade
 - **Friday, May 6** (001) or **Thursday, May 5** (020)
 - Cumulative
 - Clear your schedule now



Participation

- 5% of Course Grade
- 3 Class Surveys
 - Entrance Survey (Out Now!)
 - Mid-Semester Survey
 - Exit Survey
- Engaging in Lecture Discussions
- Asking/Answering Questions on Piazza
- Don't stress this: expect to receive full credit if you submit all 3 surveys and participate reasonably in class/Piazza



Additional Policies

Our Goal: Create an environment where you can **learn** effectively and stay **safe**

1. Academic Integrity
2. Online Behavior
3. COVID-19



Academic Integrity

- **All work you submit (e.g., assignment answers, project code, exam answers) must be a result of your own effort and reflect your individual understanding of the material.**
- Collaboration is allowed in a few select cases (Quizzes, Labs)
- OK to talk about project ideas at a high level, but not to get into specifics about code
- Cheating cases will be referred to Office of Community Standards



Online Behavior

- Some elements of the class will take place online
 - Example: Piazza Discussions
- Be respectful to others and courteous at all times
 - Full University policy available [online](#)
- Bottom Line: Don't be a troll



COVID-19 Policies

- We will be following University policies
- **Wear a mask during all class meetings**
- Do not come to class if you are experiencing COVID symptoms
 - Watch lecture recordings
 - We can help you stay up to date with the course
- Things can always change if virus situation gets better or worse

We are in this together. Be ready to adapt if things change.



I know, we were basically in the same place at the start of last semester...



Advice

- This can be a challenging class (and these are challenging times!)
- Mastering the material will take **practice** and **persistence**
- If you are struggling, ask for help!
 - No one knows all this stuff, and you are not alone!
 - Office Hours
 - Piazza
- Respect and learn from each other



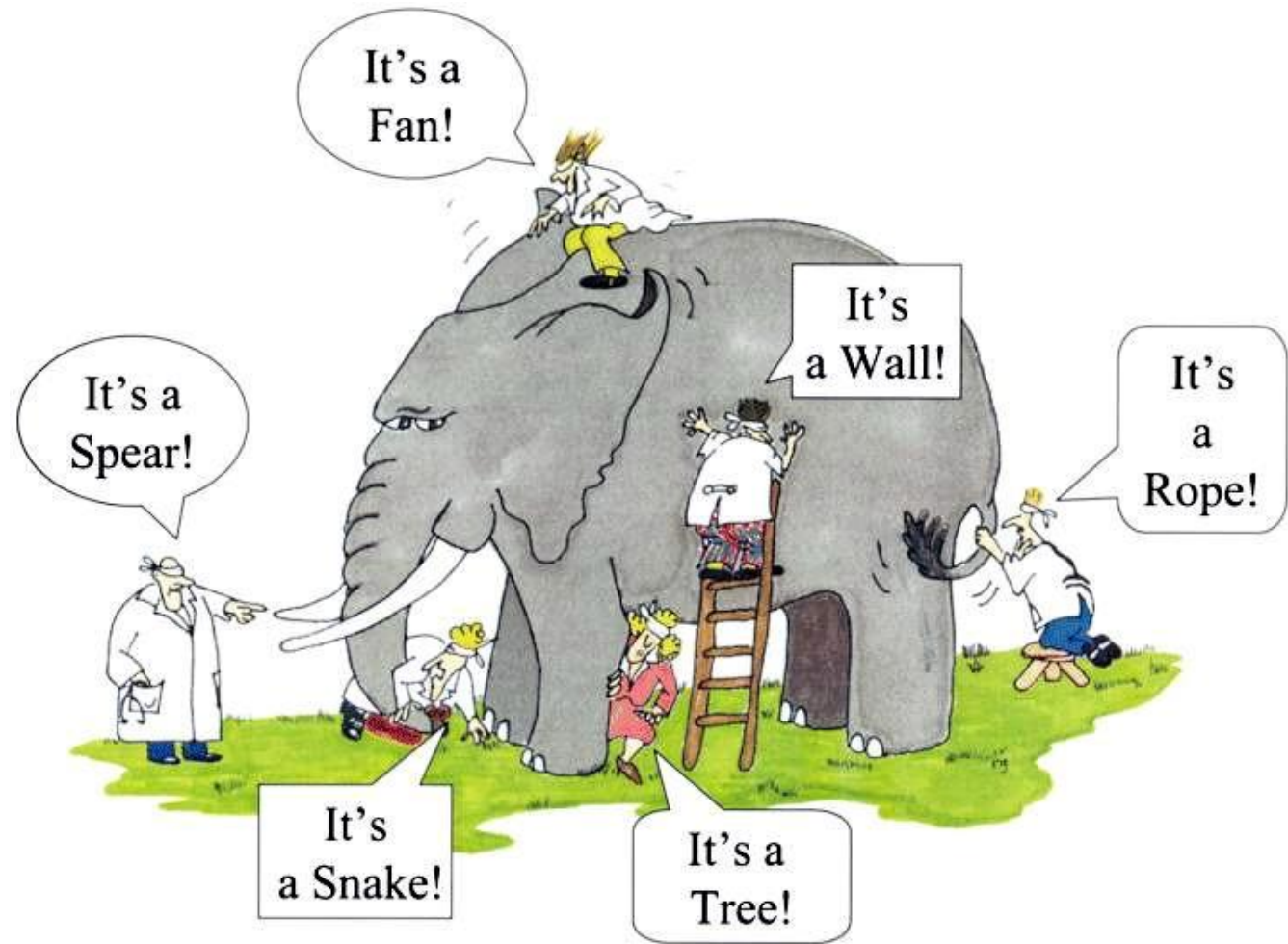
Advice

- It is our job as course staff to help you get the most out of this class
 - But we can't force you to put in the necessary effort and focus
 - Invest your time and attention, consistent and disciplined effort is the way to succeed
- Keep up with Canvas website, Piazza announcements
- Keep up with weekly lectures and readings
 - It becomes much harder to learn material when you are short on time
- Start on assignments early
 - Projects are going to take you significant time to complete
 - You will run into bugs, problems, and questions



What is an operating system?

- No universally accepted definition
- Depends on what you happen to care about at the moment
 - Computer User
 - Application Programmer
 - Systems Programmer
 - Hardware Engineer
- Example: Is the graphical user interface (window design and appearance, menus, etc.) part of the OS?



Possible OS Definition

- The operating system is the **lowest, foundational layer of software** in a computer system
- Easily the most complex software we've ever created
- The only software that is allowed to talk directly to hardware
- Manages all other software running on a machine
- Provides a consistent API that all other processes must go through to interact with the outside world: hardware, other processes
- Offers useful services to other processes (e.g., network management)



Roles of the Operating System

1. **Referee:** Supervise all running processes

- Resource sharing (Time on CPU, space in RAM)
- Maintain isolation and protection between competing processes

2. **Illusionist:** Clean, easy to use abstractions

- As coders (and as a process), we think we have a dedicated machine
- “Infinite” memory, no one else using CPU
- Build higher-level objects from hardware: Files, users, messaging
- Mask limitations, evolution of hardware

• **Glue:** Common Services

- Storage, graphical interface, networking, authorization



Short Exercise: Three Programs

#1

```
#include <stdio.h>

int main() {
    printf("Hello, World!\n");
    return 0;
}
```

#2

```
#include <string.h>
#include <unistd.h>

int main() {
    char *message = "Hello, World!\n";
    write(STDOUT_FILENO, message,
          strlen(message));
    return 0;
}
```

#3

```
.text
.global _start

_start:
    movq $1, %rax
    movq $1, %rdi
    movq $msg, %rsi
    movq $14, %rdx
    syscall

    movq $60, %rax
    movq $0, %rdi
    syscall

.data
msg:
    .ascii "Hello, world!\n"
```

1. What does each program do?
2. What are the similarities between these programs?
3. What are the differences between these programs?



System Calls

- User programs (e.g., our `hello_print` program) never interact directly with hardware like the screen
- Instead, a program makes a request to the OS to do something on its behalf by making a **system call**
- Like a function call - control transfers to some specified destination
- Not like a function call - OS takes over and runs code inside kernel
 - CPU operates in *supervisor mode* rather than *user mode*, kind of like `sudo`
- Return control back to where program left off when done
 - But with requested work all completed, or an error indicated



What do we mean by system call?

- Strictly speaking, system calls can *only* be done in assembly
 - That's the whole point of the `syscall` instruction in the x86_64 ISA
- But we won't be writing much assembly in this class
 - So you'll never write a system call for all of 4061?
 - In a way, yes!
- A function like `write` is a C library function that does the work of setting up and making a system call for us
- But, almost all programmers (and textbooks) refer to a call to `write` as a system call
 - And we'll follow this convention as well



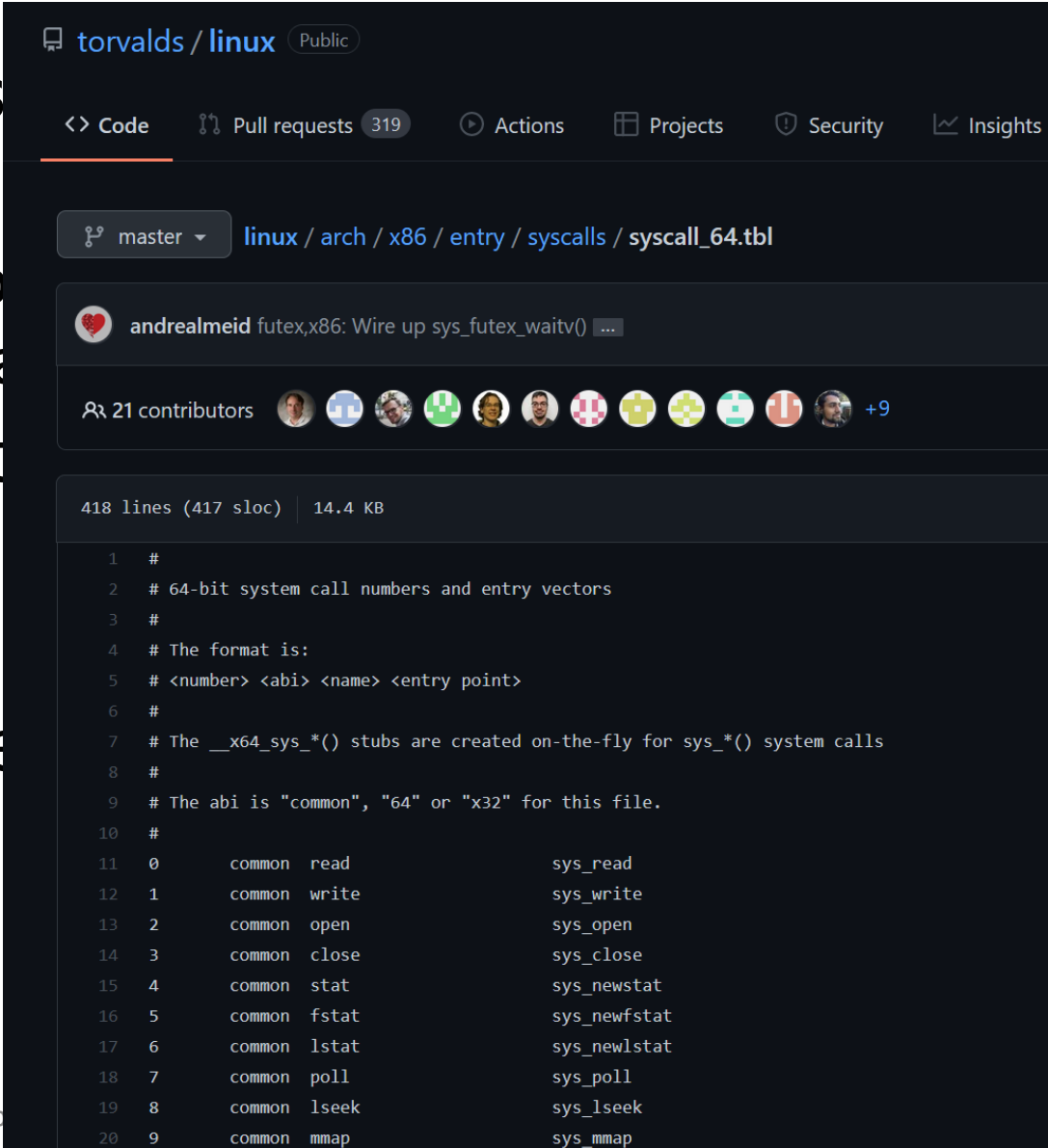
System Call Design Consideration

- Why is the series of steps to make a sys call so rigid?
- Put a specific number in a register
 - Mapping from number to operation to perform is part of OS standard
 - 64-bit Linux has ~300 different operations allowed, [each with a unique number](#)
- Transfer control to OS (syscall instruction in x86-64)
- Wouldn't it be easier to make this more like a regular function call?
 - Jump to address of first instruction in code want to execute
- Operating systems **do not trust** the user programs they supervise
 - That's part of their job
- Can only enter, start running OS code in very specific, controlled ways that are guaranteed to be safe



System Call Design Consideration

- Why is the series of steps to make a sys
- Put a specific number in a register
 - Mapping from number to operation to perform
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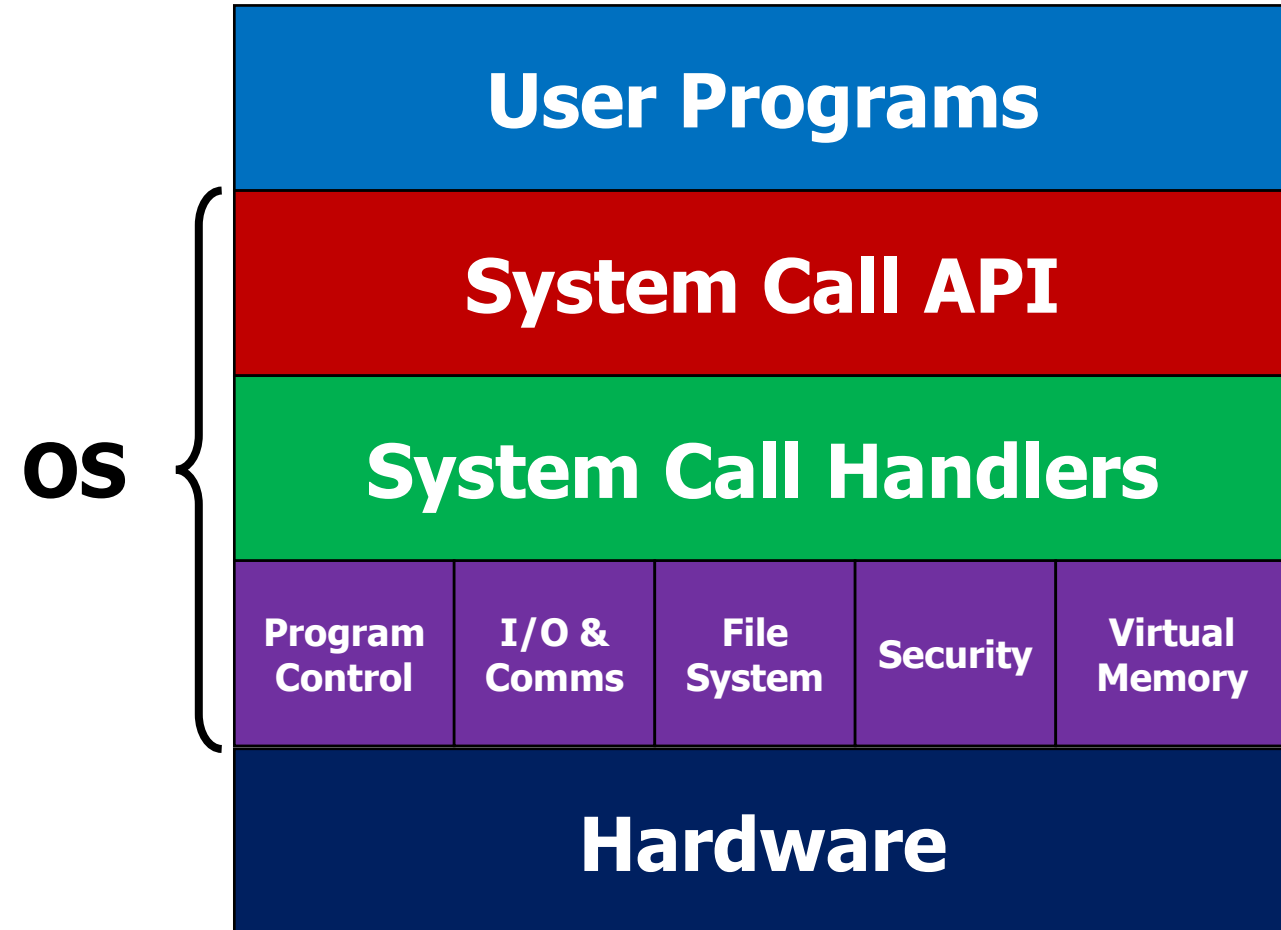
The screenshot shows a GitHub repository for the Linux kernel, specifically the file `syscall_64.tbl` in the `linux/arch/x86/entry/syscalls` directory. The repository is owned by `torvalds` and is public. It has 319 pull requests and 21 contributors. The file is 14.4 KB and contains 418 lines of code (417 sloc). The code defines the system call table for 64-bit Linux, listing system call numbers, their ABI, names, and entry points. The table is organized into columns for system call number, ABI, name, and entry point. The first 10 lines of the table are as follows:

Line	Number	ABI	Name	Entry Point
11	0	common	read	sys_read
12	1	common	write	sys_write
13	2	common	open	sys_open
14	3	common	close	sys_close
15	4	common	stat	sys_newstat
16	5	common	fstat	sys_newstat
17	6	common	lstat	sys_newstat
18	7	common	poll	sys_poll
19	8	common	lseek	sys_lseek
20	9	common	mmap	sys_mmap



OS Structure

- Operating systems are built out of layers, like many things in CS
- **4061**: Outer layers
- **5103**: Inner layers
- Hardware? EE and architecture



OS Structure

- Operating systems are built out of layers, like many things in CS
- **4061**: Outer layers
- **5103**: Inner layers
- Hardware? EE and architecture
- Reminder
 - OS unifies huge range of hardware under one API
 - And this API supports a huge range of software

OS



Systems Programming

- Like operating systems, has a loose definition
- **An area of study/expertise:** Writing and designing programs that make heavy use of the operating system's API
- **A point of contrast**
 - *Applications programming* is about software that provides services to the end user (e.g., word processor)
 - *Systems programming* is about software that acts a foundation for other software (e.g., Python interpreter, Web Server)



Systems Programming

- Also an **ethos/state of mind**
- Systems programmers have knowledge/awareness of hardware
 - It affects efficiency/correctness of your code!
 - Writing cache-aware code in 2021 is a good start
- Systems programmers are paranoid and meticulous
 - Almost all system calls can result in an error
 - What if you try to write a file to a full storage device?
 - True for library functions too: `malloc` can return `NULL` if error
 - Code must check for *any* possible error and address it
- Because they have to be: if your code fails, so does everything that depends on it



Systems Programming

A problem has been detected and Windows has been shut down to prevent damage to your computer.

DRIVER_IRQL_NOT_LESS_THAN_OR_EQUAL_TO

If this is the first time you've seen this Stop error screen, restart your computer. If this screen appears again, follow these steps:

Check to make sure any new hardware or software is properly installed. If this is a new installation, ask your hardware or software manufacturer for any Windows updates you might need.

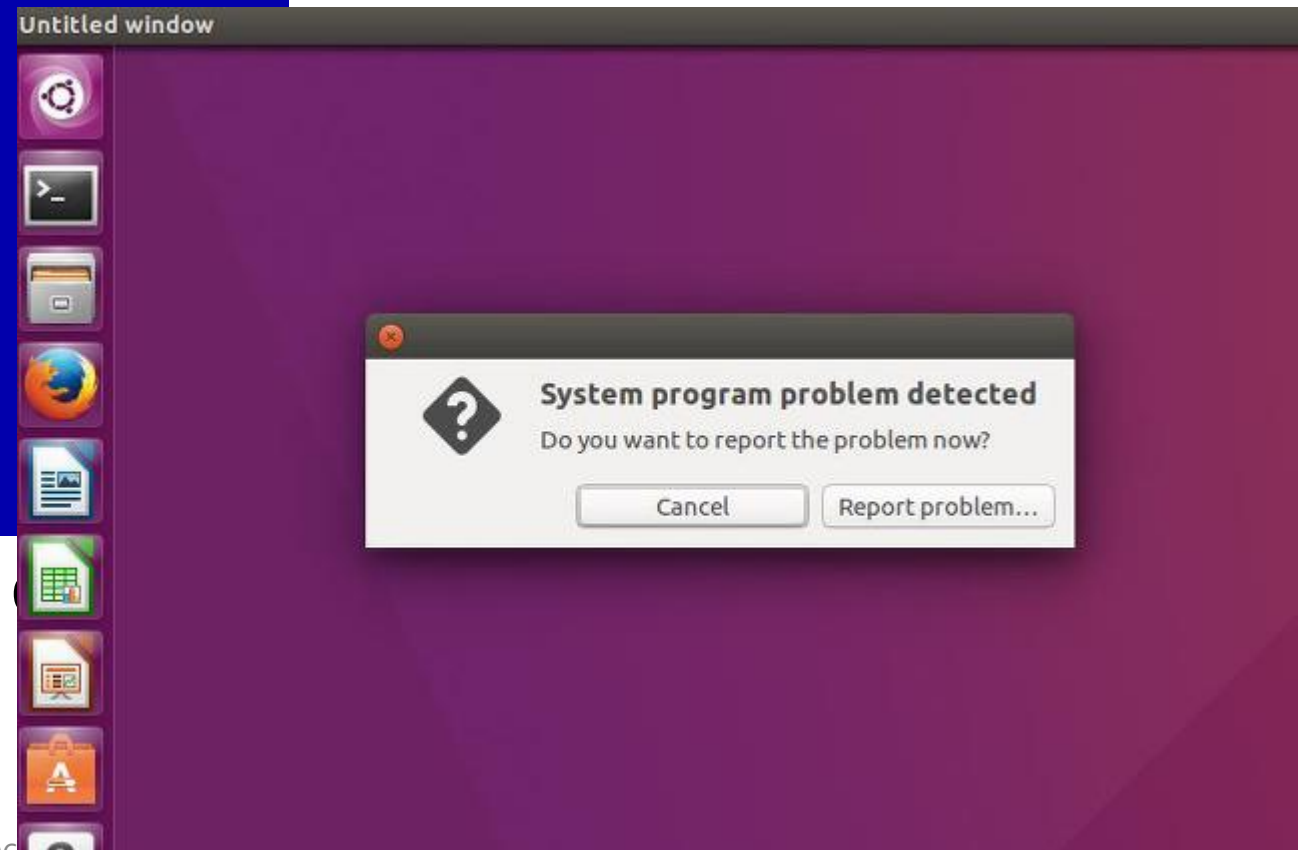
If problems continue, disable or remove any newly installed hardware or software. Disable BIOS memory options such as caching or shadowing. If you need to use Safe Mode to remove or disable components, restart your computer, press F8 to select Advanced Startup options, and then select Safe Mode.

Technical information:

*** STOP: 0x000000D1 (0x00000000, 0x00000000)

awareness of hardware

- Because they have to be: if your code depends on it



Core Concepts in Systems Programming

Concurrency	Multiple things are in progress at once, their order is unpredictable
Asynchrony	Events can happen at any point
Coordination	Competing processes/threads must avoid interfering with each other
Communication	Share info within a machine (processes/threads) or across the world (networking)
Security	Access to info, operations can be restricted
File Storage	Layout of data on disks, algorithms for efficient reading/writing
Memory	Illusion of independent address space for each process, safe and controlled sharing
Robustness	Handle unexpected events gracefully (unless you want to crash)
Efficiency	Use CPU, Memory, Disk to fullest potential, minimize overhead for other programs



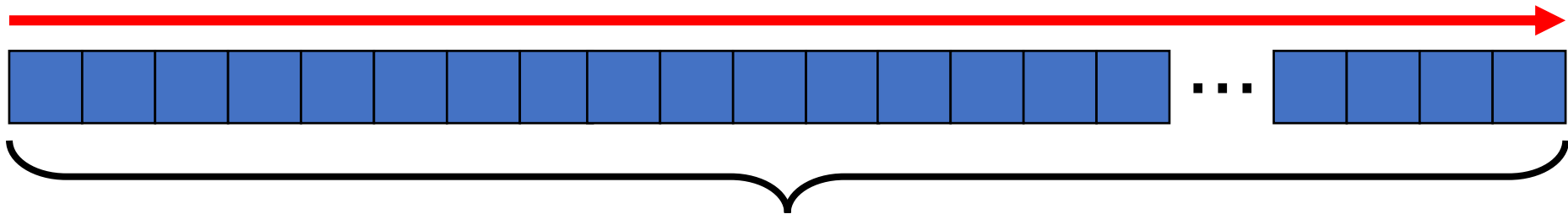
Major Course Topics

1. Quick C Review, C Program Structure (Some review from 2021)
2. Processes, Process Creation and Management
3. Low-Level I/O
4. File Systems
5. Virtual Memory
6. Signals and Event Handling
7. Interprocess Communication
8. Threads
9. Thread Synchronization
10. Network Communication and Programming
11. Networked Systems (Time Permitting)



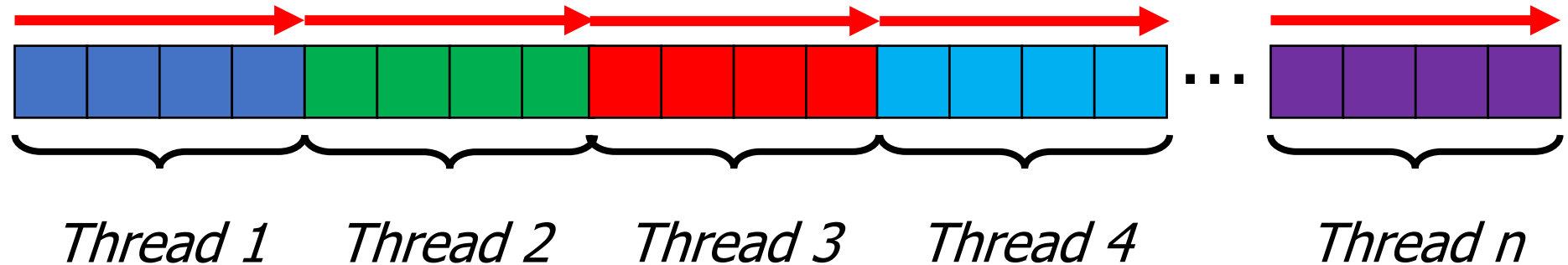
Preview: Multithreading

- Arguably the most important concept in this class
- One process can contain multiple **threads**, each executing its own sequence of instructions (code)
 - Lots of reasons we may want to do this
 - One obvious one: do work *in parallel* on a CPU with multiple cores
- Example: Summing all elements in an array
- Traditional Sequential Approach:



Preview: Multithreading

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- One process can contain multiple **threads**, each executing its own sequence of instructions (code)
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- Example: Summing all elements in an array
- Multithreaded Approach:



Exercise: Sketching Thread Code

- Assume we have some way to create threads, each executing a call to the `sum_elements` function in parallel with other threads
- We have two versions of `sum_elements`:

```
void sum_elements_v1(long *sum, int
*elements, int start, int stop) {
    for (int i = start; i < stop; i++) {
        *sum += elements[i];
    }
}
```

```
void sum_elements_v2(long *sum, int
*elements, int start, int stop) {
    long temp_sum = 0;
    for (int i = start; i < stop; i++) {
        temp_sum += elements[i];
    }
    *sum += temp_sum;
}
```

- If we have 8 threads and an array with 100 elements, what should the function arguments be for each thread?
- Generalize this to a case with n threads and an array with m elements
- Is one version of `sum_elements` better than the other? Why?

Exercise: Sketching Thread Code



Exercise: Sketching Thread Code

- Threads give you a lot of power as a programmer, but they can be very difficult to use
- You'll have to adjust your whole mental model of how code executes
- Understanding these concepts is a major advantage
 - Academic standpoint: Many advanced topics/courses rely on this knowledge
 - Career standpoint: Distinguishes you as a developer/engineer
- Point 1: Understanding hardware (CPU and memory) is even more crucial to writing efficient code when using multiple threads
- Point 2: Need to take extra care with multithreaded code to ensure it behaves correctly and consistently
 - We'll spend a lot of time on why this is and the tools we use to address it



Assumption: You Know Some C

- CSCI 2021 (or equivalent) is prereq, covers some hardware, basic C programming, and some interactions between C and hardware
- We'll assume you know C syntax and its major features, such as:
 - Loops
 - Conditionals
 - Functions
 - Data types, including arrays and structs
 - Pointers and memory addresses
 - Memory management with `malloc()` and `free()`
 - Stack vs. heap storage
 - How to edit, compile and execute code
 - How to debug code, particularly with `gdb`



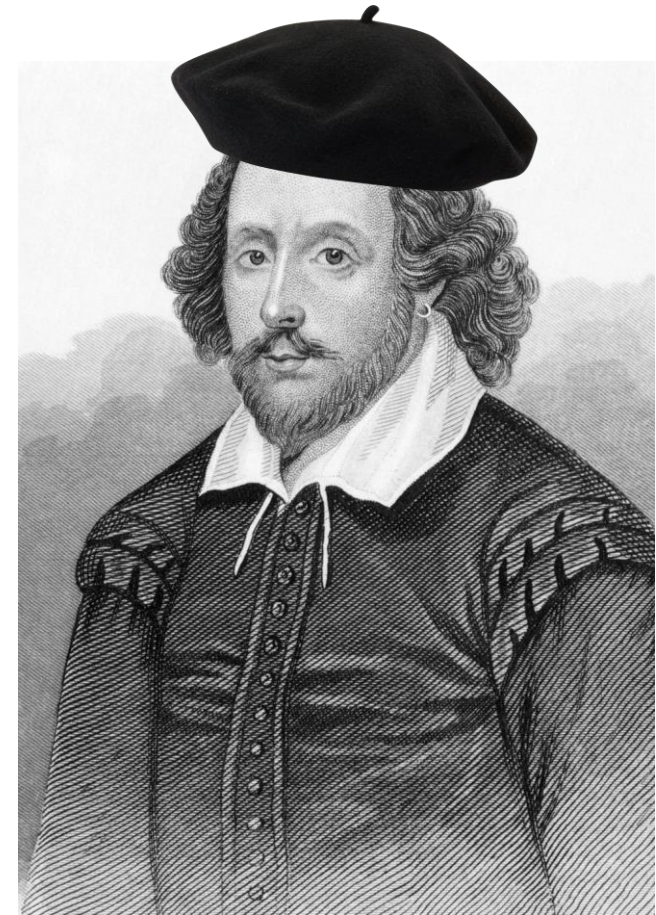
Getting up to Speed on C

- We'll do some review in our next lecture
 - Meant as a reminder, not a complete tutorial
- May want to locate a good C reference
 - *The C Programming Language* by Kernighan and Ritchie is a good one
- "C Programming Resources" [page](#) on our Canvas site
- For suggestions on useful tools, see "Programming Environment Guide" page on our Canvas site
- GDB [guide](#) linked to from our Canvas site
 - Originally written for CSCI 2021 but highly relevant
- **Make sure you invest the necessary effort here, otherwise this course will be very difficult for you**



Why use C in 4061?

- C can certainly be challenging
 - Memory management
 - Pointers
 - Strange bugs
- But it is **the** language for operating systems and systems programming
- Gives you most direct access to reality of real hardware
- Gives you direct and “unfiltered” access to the real OS API

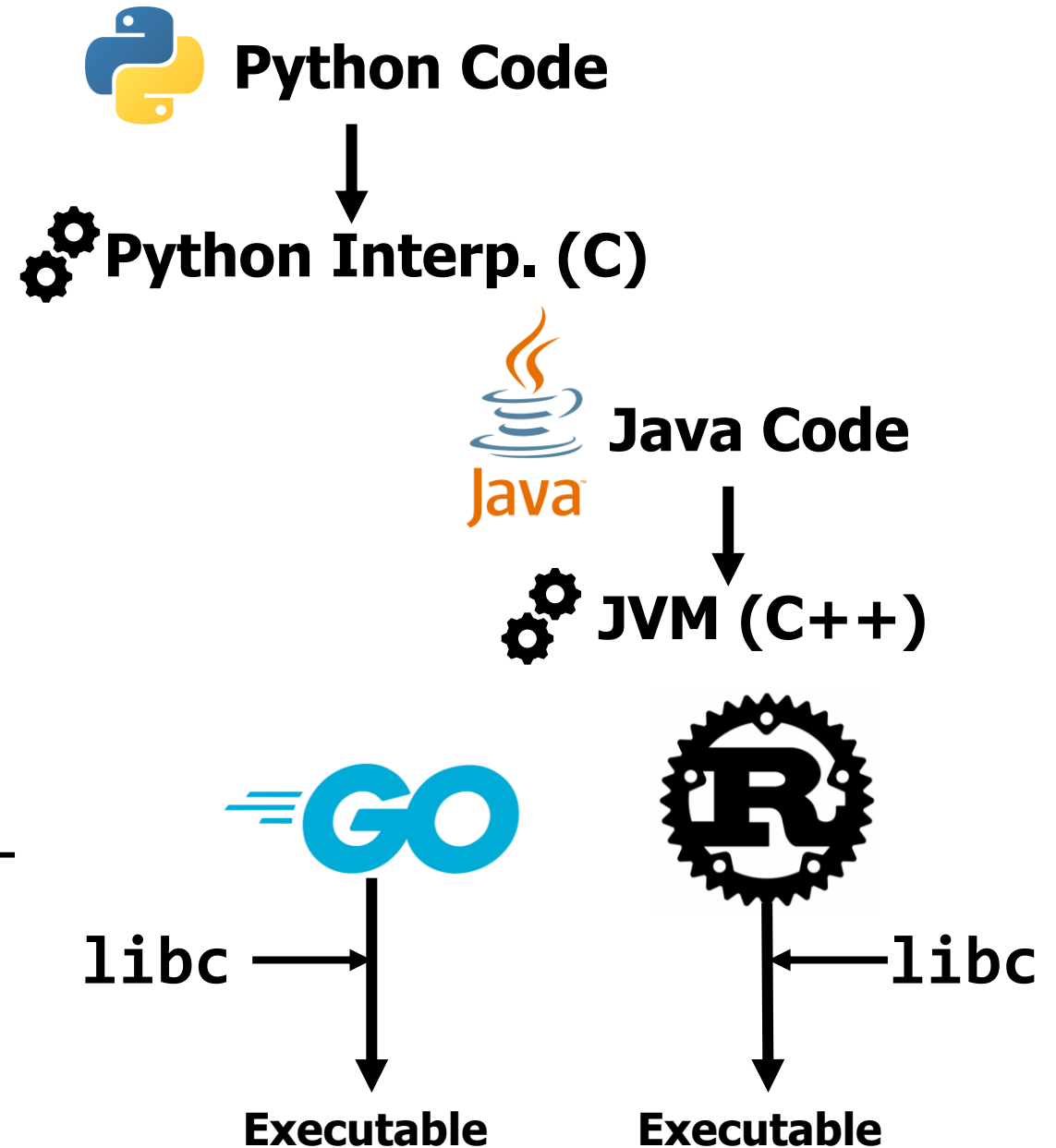


Same Reason We Don't
Study Shakespeare as
Translated to French



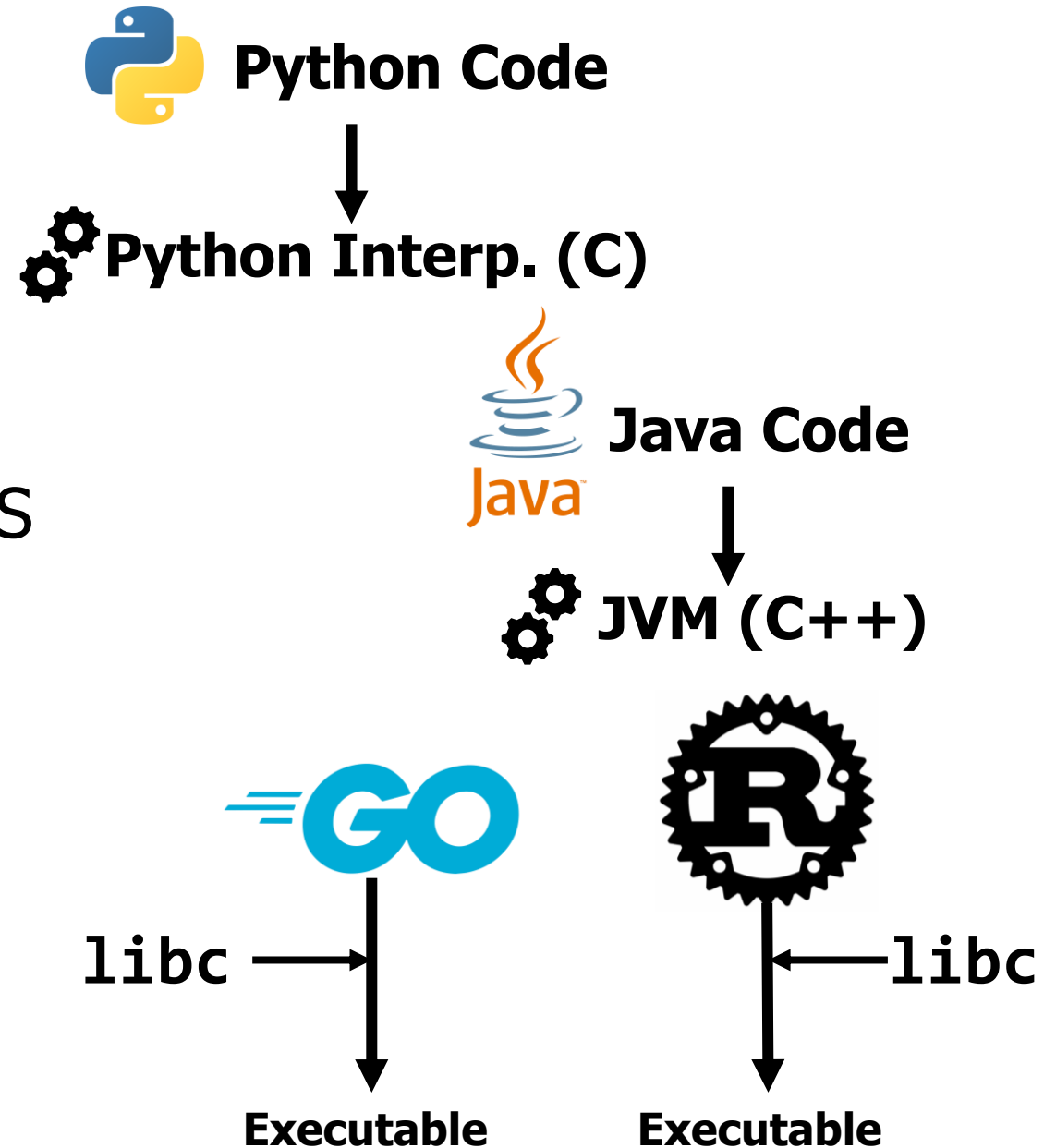
Why use C in 4061?

- C is **the** language for operating systems and systems programming
- Almost everything else is built on top of C and its low-level libraries (libc)
- Python, Java have interpreters written in C and C++
- JavaScript interpreted by your browser (probably) written in C++
- Even trendy languages like Go and Rust rely on libc



Why use C in 4061?

- C is **the** language for operating systems and systems programming
- These other languages have libraries that give you access to OS API
- They usually just wrap C code
- Sacrifice flexibility/expressiveness for convenience
 - Assumptions about common cases
 - Hard to work outside of these



Some (Hopefully) Clarifying History

- In the “beginning” there was **Multics**
 - Multiplexed Information and Computing Service
- This is back in the days of mainframes shared by many users
 - Time sharing among their programs was a key concern
- Started in 1964, heavily developed into the 1970s
- Introduced a lot of ideas we now take for granted
 - Hierarchical File System (Directories and Subdirectories)
 - Dynamic Linking (map code into memory at run time)
- Very large and complex, leading to problems



Huge Milestone in OS History: Unix

- Development began in 1969 by researchers who split from Multics
- Key Idea: Keep OS internals as simple as possible
 - Build small, sharply focused tools that combine flexibly
 - Keep Kernel implementation as small as possible while still useful
- Defined an API still in wide use today and has been hugely influential
 - The “language” in which we think about organizing computation on real systems
 - Processes, communication, files, etc.
 - Old is a **good** thing: well-tested and has remained useful and effective over decades of technological advances
- Also often **open** to inspection and community development
- Led to many “Unix-style” OS implementations that retain this API



Unix-Derived Operating Systems



Notable Exception:



Different API and history,
same concepts



QNX SOFTWARE SYSTEMS



POSIX: Modern, Unifying OS Standard

- By the 1990s and 2000s, many Unix-derived operating systems were being maintained and widely used
- All incorporated design and API ideas from original Unix
- Remember, “API” here means set of available system calls, their names, and their precise behavior (what are they guaranteed to do?)
- Different Unix descendants were largely the same but still had API differences, especially over time as new features/ideas added in
 - Makes it impossible to write code that would work on multiple systems
- POSIX defines a standard API that everyone adheres to
 - So, if you know POSIX, you can write code for many OS variants
 - Still not perfect, individual operating systems at times add own functionality



Summarizing All of This

- Our textbook would more accurately be titled *Posix Systems Programming*
- And in this class, we'll be focusing on **Linux** systems programming
- *Confused yet?*
- We aren't writing code for Unix itself, but we are using a lot of the API features it originally defined
- This is because Linux conforms to the Posix standard, and Posix is based on the ideas/conventions established in Unix
- Unfortunately, you'll see people say or write "Unix" when they really mean "Posix"
 - Not precise, but common enough that you'll have to get used to it

